

# Example Title

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## Section

### Some math

$$T(n) = \begin{cases} 0 & \text{if } n = 0 \\ 1 & \text{if } n = 1 \\ T(n-1) + T(n-2) + 1 & \text{if } n \geq 2 \end{cases} \quad (1)$$

$$T(n) = T(n-1) + T(n-2) + 1$$

$$T(n) = T(n-1) + T(n-2)$$

$$T(n) = 2 \cdot T(n-1)$$

$\Downarrow$

$$T(n) = 2 \cdot 2 \cdot T(n-2)$$

$$T(n) = 2^{k-1} \cdot T(n-k) \quad \{k = n-1\}$$

$$T(n) = 2^{k-1} \cdot T(1)$$

$$T(n) = O(2^n)$$

(2)

## A List

The two properties a problem must have to be solved with dynamic programming are:

- Optimal substructure
- Overlapping sub-problems

## A Table

Done after two iterations as there are no changes between iteration 1.5 and 2, see table 1. The shortest path from  $s$  to  $t$  is 6 and goes:  $s$ - $A$ - $B$ - $D$ - $t$ .

Iteration	s	A	B	C	D	t
0	0	$\infty$	$\infty$	$\infty$	$\infty$	$\infty$
1	0	1	3	3	4	6
2	0	1	3	3	4	6

Table 1: Iteratioins for finding the shortest path from  $s$  to  $t$

## Task 3: Maximum sum sub-array

```
1 def max_subarray_sum(A):
2     max_ending_here = max_so_far = A[0]
3
4     for num in A[1:]:
5         max_ending_here = max(num, max_ending_here + num)
6         max_so_far = max(max_so_far, max_ending_here)
7
8     return max_so_far
9
10 # Example usage
11 example_array = [-2, 1, -3, 4, -1, 2, 1, -5, 4]
12 print(max_subarray_sum(example_array))
```

The time complexity of this is  $O(n)$  as it only iterates through the array once.

## Task 4: Grid traveling

Traverse the grid from S to F by only moving right or down.

a)

S	
	F

When the grid is 2x2 there are 2 ways to get from S to F.

b)

S		
		F

When the grid is 3x3 there are 6 ways to get from S to F.

c)

S			
			F

When the grid is 4x4 there are 20 ways to get from S to F.

## Observation

The pattern can be recognized from table 2 where the number in the cells represents how many different ways one can traverse the grid from S to F with the specified rules.

1	1	1	1
1	2	3	4
1	3	6	10
1	4	10	20

Table 2: Pattern