Sistemas Distribuídos

José Orlando Pereira

Departamento de Informática
Universidade do Minho

Mutex with Peterson's/...

- Does it really work?

```
try {
    l.lock();
    c=c+1;
} finally {
    l.unlock();
}
```

while(flag[...]) { }

flag[i] = false;

Consequence

```
try {
    l.lock();
    c=c+1;
} finally {
    l.unlock();
}
```

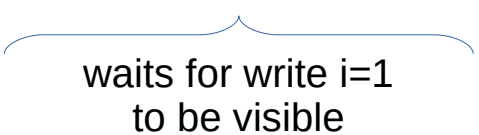
while(flag[...]) { }

flag[i] = false;

- Initially c=10
- One thread:
  - read c
  - write c = 11
  - write flag[i] = false
- Other thread:
  - read flag[...] = false
  - read c = 10!!!!

Solution: Memory barriers

- Declare: volatile int j; or AtomicInteger j;
- Reading from a volatile j waits for all writes preceding the observed value on j to be also visible
  - Writer code:
    - i=1; j=1;
  - Reader code:
    - rj=j; ri=i; System.out.println(rj+", "+ri);



Should I use it?

- Synchronization operations also act as memory barriers:
  - java.util.concurrent.* primitives
  - synchronized keyword
- Volatile variables also impact performance:
  - A volatile access takes as much time as using a lock
  - Therefore… **use the lock**!

The real world

- Is j.u.c.ReentrantLock implemented with Bakery or something similar?
  - **No!**
  - Needs *O(n)* space for each lock (with *n* threads)
  - Consumes CPU time / power when busy
- These algorithms are a good for:
  - Understanding concurrent programming and races
  - Preparing for distributed algorithms

The real world

- j.u.c.ReentrantLock is implemented with atomic operations and scheduling:
  - Atomic operations (testAndSet, compareAndSwap, …) are used to quickly check if mutex is available
    - For more on atomic operations, see Chapter 7 of TAOMP (not part of the program!)
  - The thread is suspended by the OS when waiting for longer periods
    - Remember scheduling from Operating Systems
    - man futex on Linux

j.u.c Locks vs Monitors

```
class C {
    private int i;

    synchronized public void m() {

        i++;

    }
}
```

There is a hidden “lock” in each object used by “synchronized”

```
class C {
    private int i;
    private Lock l =
        new ReentrantLock();

    public void m() {
        try { l.lock();
        i++;
        } finally { l.unlock(); }
    }
}
```

Equivalent code (aproximately...)

j.u.c. Locks vs Monitors

- Main differences, for now:
  - Synchronized blocks are nested in LIFO order

    vs.

    j.u.c. Locks can be unlocked in any order
    - To take advantage of two phase locking
  - Threads waiting for a synchronized block enter in any order

    vs.

    j.u.c. Lock can be configured for threads waiting to enter in a fair (not FIFO!) order

j.u.c Conditions vs Monitors

- class C {

There is a hidden “condition” in each object used by “wait()/notifyAll()”

```
synchronized public void m1() {
        while(…) wait();
}

synchronized public void m2() {
        notifyAll();
}
}
```

- class C {

```
private Lock l =
        new ReentrantLock();
private Condition c =
        l.newCondition();
public void m1() {
        try { l.lock();
            while(…) c.await()
        } finally { l.unlock(); }
}
public void m2() {
        try { l.lock();
            c.signalAll();
        } finally { l.unlock(); }
}
}
```

Equivalent code (aproximately...)

j.u.c. Conditions vs Monitors

- Main differences:
  - One implicit condition for each lock
    vs.
    Many j.u.c. conditions for the same lock
    - Avoids signalAll()
  - Threads waiting for a condition wakeup in any order
    vs.
    Threads waiting for a j.u.c. Condition obtained from a ReentrantLock wakeup in FIFO order (but may not aquire lock in FIFO order...)