# Manda o Double de Campeão CEFET-MG

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## 1 DP

## 1.1 Exemplo Sapo

```
// There are N stones, numbered 1, 2, \ldots, N.
// For each i (1<=i<=N), the height of Stone i is hi.
// There is a frog who is initially on Stone 1.
// He will repeat the following action some number of times to reach
   Stone N:
// If the frog is currently on Stone i, jump to one of the following:
   Stone i+1,i+2,...,i+K. Here, a cost of | hi - hj | is incurred,
   where j is the stone to land on.
// Find the minimum possible total cost incurred before the frog
   reaches Stone N.
dca int n, k;
// Top Down
4d3 int dp(int i) {
563
       if(i == 0) return 0;
7f9
       auto& ans = memo[i];
d64
       if (\simans) return ans;
5f9
       int ret = INF;
a7f
        f(j, max(0ll,i-k), i)
f97
            ret = min(ret, dp(j) + abs(h[j] - h[i]));
655
        return ans = ret;
641 }
// Bootom Up
e63 int dp_2(int x) {
        memo[0] = 0;
ecd
b85
        f(i,1,x) {
90b
            int best = INF:
203
            f(j, max(011, i-k), i) {
428
                best = min(best, memo[j] + abs(h[i] - h[j]));
818
            }
bc2
            memo[i] = best;
832
d56
        return memo[x-1];
```

```
63d void solve() {
0a1     cin >> n >> k;
3f0     f(i,0,n) cin >> h[i];
8e4     cout << dp(n-1) << endl;
1d6 }
```

6f3 }

## 1.2 Knapsack tradicional

```
// O(n * cap)
b94 const int MAXN = 110;
689 \text{ const int MAXW} = 1e5+10;
ba9 int n, memo[MAXN][MAXW];
310 int v[MAXN], w[MAXN];
74a int pego[MAXN] = {0};
// Retorna o lucro maximo
12c int dp(int id, int cap) {
        if(cap < 0) return -LLINF;</pre>
1bb
        if(id == n or cap == 0) return 0;
ecb
c1a
        int &ans = memo[id][cap];
d64
        if (\simans) return ans;
86f
        return ans = max(dp(id+1, cap), dp(id+1, cap-w[id]) + v[id]);
d95 }
// Armazena em pego os itens pegos
7d0 void recuperar(int id, int cap) {
efa
        if(id >= n) return;
fca
        if(dp(id+1, cap-w[id]) + v[id] > dp(id+1, cap)) { // se pegar}
   eh otimo
            pego[id] = true;
44 c
            recuperar(id+1, cap-w[id]);
3fd
        } else { // nao pegar eh otimo
4ee
884
            pego[id] = false;
45d
            recuperar(id+1, cap);
549
        }
845 }
63d void solve() {
311
        int cap; cin >> n >> cap;
457
        memset(memo, -1, sizeof memo);
```

```
03ъ
        f(i,0,n) \{ cin >> w[i] >> v[i]; \}
304
        int lucro_max = dp(0, cap);
ae7
        recuperar(0, cap);
4 c c
        int lucro = 0, peso = 0;
418
        f(i,0,n) {
ecd
            if(pego[i]) {
73f
                lucro += v[i];
20e
                peso += w[i];
c3f
            }
b7f
        }
d13
        assert(lucro_max == lucro and peso <= cap);</pre>
f6f }
```

## 2 Estruturas

#### 2.1 DSU

```
// Une dois conjuntos e acha a qual conjunto um elemento pertence por
   seu id
// find e unite: O(a(n)) \sim = O(1) amortizado
8d3 struct dsu {
825
        vector < int > id, sz;
b33
        dsu(int n) : id(n), sz(n, 1) { iota(id.begin(), id.end(), 0); }
0cf
        int find(int a) { return a == id[a] ? a : id[a] = find(id[a]);
  }
        void unite(int a, int b) {
440
605
            a = find(a), b = find(b);
            if (a == b) return;
d54
956
            if (sz[a] < sz[b]) swap(a, b);
            sz[a] += sz[b], id[b] = a;
640
ea7
8e1 };
// DSU de bipartido
```

```
//
// Une dois vertices e acha a qual componente um vertice pertence
// Informa se a componente de um vertice e bipartida
// find e unite: O(log(n))
8d3 struct dsu {
        vector<int> id, sz, bip, c;
5b4
        dsu(int n) : id(n), sz(n, 1), bip(n, 1), c(n) {
db8
            iota(id.begin(), id.end(), 0);
f25
        }
ef0
        int find(int a) { return a == id[a] ? a : find(id[a]); }
        int color(int a) { return a == id[a] ? c[a] : c[a] ^
   color(id[a]); }
440
        void unite(int a, int b) {
263
            bool change = color(a) == color(b);
            a = find(a), b = find(b);
605
            if (a == b) {
a89
4ed
                if (change) bip[a] = 0;
505
                return:
32d
            }
956
            if (sz[a] < sz[b]) swap(a, b);
efe
            if (change) c[b] = 1;
2cd
            sz[a] += sz[b], id[b] = a, bip[a] &= bip[b];
22b
118 };
// DSU Persistente
// Persistencia parcial, ou seja, tem que ir
// incrementando o 't' no une
//
// find e unite: O(log(n))
8d3 struct dsu {
33c
        vector<int> id, sz, ti;
733
        dsu(int n) : id(n), sz(n, 1), ti(n, -INF) {
db8
            iota(id.begin(), id.end(), 0);
        }
aad
5e6
        int find(int a, int t) {
```

```
6ba
            if (id[a] == a or ti[a] > t) return a;
ea5
            return find(id[a], t);
6cb
       }
        void unite(int a, int b, int t) {
fa0
            a = find(a, t), b = find(b, t);
84f
d54
            if (a == b) return;
956
            if (sz[a] < sz[b]) swap(a, b);</pre>
35d
            sz[a] += sz[b], id[b] = a, ti[b] = t;
513
       }
6c6 };
// DSU com rollback
//
// checkpoint(): salva o estado atual de todas as variaveis
// rollback(): retorna para o valor das variaveis para
// o ultimo checkpoint
// Sempre que uma variavel muda de valor, adiciona na stack
// find e unite: O(log(n))
// checkpoint: 0(1)
// rollback: O(m) em que m e o numero de vezes que alguma
// variavel mudou de valor desde o ultimo checkpoint
8d3 struct dsu {
825
        vector < int > id, sz;
27 c
        stack<stack<pair<int&, int>>> st;
98d
        dsu(int n) : id(n), sz(n, 1) {
            iota(id.begin(), id.end(), 0), st.emplace();
1 c c
8cd
        }
bdf
        void save(int &x) { st.top().emplace(x, x); }
30d
        void checkpoint() { st.emplace(); }
5cf
        void rollback() {
ba9
            while(st.top().size()) {
6bf
                auto [end, val] = st.top().top(); st.top().pop();
149
                end = val:
            }
f9a
            st.pop();
25a
       }
3c6
        int find(int a) { return a == id[a] ? a : find(id[a]); }
ef0
```

## 2.2 Fenwick Tree (BIT)

```
// Operacoes O-based
// query(1, r) retorna a soma de v[1..r]
// update(1, r, x) soma x em v[1..r]
//
// Complexidades:
// build - O(n)
// query - O(log(n))
// update - O(log(n))
e04 namespace bit {
06d
        int bit[2][MAX+2];
1a8
        int n;
727
        void build(int n2, vector<int>& v) {
1e3
            n = n2:
            for (int i = 1; i <= n; i++)</pre>
535
                 bit [1] [min(n+1, i+(i\&-i))] += bit[1][i] += v[i];
a6e
d31
1a7
        int get(int x, int i) {
7c9
            int ret = 0:
360
            for (; i; i -= i&-i) ret += bit[x][i];
edf
            return ret;
        }
a4e
920
        void add(int x, int i, int val) {
503
            for (; i <= n; i += i&-i) bit[x][i] += val;</pre>
fae
        }
        int get2(int p) {
3d9
            return get(0, p) * p + get(1, p);
c7c
33c
9e3
        int query(int 1, int r) { // zero-based
ff5
            return get2(r+1) - get2(1);
25e
7ff
        void update(int 1, int r, int x) {
e5f
            add(0, 1+1, x), add(0, r+2, -x);
f58
            add(1, 1+1, -x*1), add(1, r+2, x*(r+1));
```

```
5ce
       }
17a }:
63d void solve() {
        vector <int > v {0,1,2,3,4,5}; // v[0] eh inutilizada
97a
c7b
        bit::build(v.size(), v);
67f
        int a = 0, b = 3;
        bit::query(a, b); // v[a] + v[a+1] + ... + v[b] = 6 | 1+2+3 =
9b0
   6 | zero-based
        bit::update(a, b, 2); // v[a...b] += 2 | zero-based
b3d
7b4 }
2.3 SegTree
// Recursiva com Lazy Propagation
// Query: soma do range [a, b]
// Update: soma x em cada elemento do range [a, b]
// Pode usar a seguinte funcao para indexar os nohs:
// f(1, r) = (1+r) | (1!=r), usando 2N de memoria
//
// Complexidades:
// build - O(n)
// query - 0(log(n))
// update - O(log(n))
0d2 const int MAX = 1e5+10;
fb1 namespace SegTree {
098
        int seg[4*MAX], lazy[4*MAX];
052
        int n, *v;
b90
        int op(int a, int b) { return a + b; }
2c4
        int build(int p=1, int l=0, int r=n-1) {
3c7
            lazv[p] = 0;
6cd
            if (1 == r) return seg[p] = v[1];
            int m = (1+r)/2;
ee4
```

return seg[p] = op(build(2\*p, 1, m), build(2\*p+1, m+1, r));

317

985

0d8

680

6f2

acb

}

}

void build(int n2, int\* v2) {

n = n2, v = v2:

build();

```
void prop(int p, int 1, int r) {
ceb
            seg[p] += lazv[p]*(r-l+1);
cdf
2c9
            if (1 != r) lazy[2*p] += lazy[p], lazy[2*p+1] += lazy[p];
3c7
            lazy[p] = 0;
        }
c10
        int query(int a, int b, int p=1, int l=0, int r=n-1) {
04a
6b9
            prop(p, 1, r);
527
            if (a <= l and r <= b) return seg[p];</pre>
786
            if (b < 1 \text{ or } r < a) \text{ return } 0;
ee4
            int m = (1+r)/2:
19e
            return op(query(a, b, 2*p, 1, m), query(a, b, 2*p+1, m+1,
   r));
1c9
       }
f33
        int update(int a, int b, int x, int p=1, int l=0, int r=n-1) {
6b9
            prop(p, 1, r);
9a3
            if (a \le 1 \text{ and } r \le b) {
b94
                 lazy[p] += x;
6b9
                prop(p, 1, r);
534
                return seg[p];
821
            }
e9f
            if (b < l or r < a) return seg[p];</pre>
ee4
            int m = (1+r)/2:
a8f
            return seg[p] = op(update(a, b, x, 2*p, 1, m), update(a,
   b, x, 2*p+1, m+1, r));
08f
       }
        // Se tiver uma seg de max, da pra descobrir em O(log(n))
        // o primeiro e ultimo elemento >= val numa range:
        // primeira posicao >= val em [a, b] (ou -1 se nao tem)
        int get left(int a, int b, int val, int p=1, int l=0, int
119
   r=n-1) {
6b9
            prop(p, 1, r);
            if (b < l or r < a or seg[p] < val) return -1;</pre>
f38
205
            if (r == 1) return 1;
            int m = (1+r)/2:
ee4
753
            int x = get_left(a, b, val, 2*p, l, m);
            if (x != -1) return x:
50e
            return get_left(a, b, val, 2*p+1, m+1, r);
сЗс
68 c
        }
        // ultima posicao >= val em [a, b] (ou -1 se nao tem)
992
        int get_right(int a, int b, int val, int p=1, int l=0, int
   r=n-1)
```

```
6b9
            prop(p, 1, r);
            if (b < l or r < a or seg[p] < val) return -1;</pre>
f38
205
            if (r == 1) return 1;
            int m = (1+r)/2;
ee4
            int x = get_right(a, b, val, 2*p+1, m+1, r);
1 b 1
50e
            if (x != -1) return x;
6a7
            return get_right(a, b, val, 2*p, 1, m);
1b7
        }
        // Se tiver uma seg de soma sobre um array nao negativo v, da
            pra
        // descobrir em O(log(n)) o maior j tal que
           v[i]+v[i+1]+...+v[i-1] < val
89b
        int lower_bound(int i, int& val, int p, int l, int r) {
6b9
            prop(p, 1, r);
6e8
            if (r < i) return n;</pre>
b5d
            if (i <= l and seg[p] < val) {</pre>
bff
                val -= seg[p];
041
                return n;
            }
634
Зсе
            if (1 == r) return 1;
            int m = (1+r)/2;
ee4
514
            int x = lower_bound(i, val, 2*p, 1, m);
            if (x != n) return x;
ee0
8b9
            return lower_bound(i, val, 2*p+1, m+1, r);
0.1d
a15 };
63d void solve() {
213
        int n = 10;
        int v[] = {1, 2, 3, 4, 5, 6, 7, 8, 9, 10};
89e
2d5
        SegTree::build(n, v);
        cout << SegTree::query(0, 9) << endl; // seg[0] + seg[1] + ...</pre>
   + seg[9] = 55
310
        SegTree::update(0, 9, 1); // seg[0, ..., 9] += 1
6d9 }
```

## 2.4 Sparse Table Disjunta

```
// Description: Sparse Table Disjunta para soma de intervalos
// Complexity Temporal: O(n log n) para construir e O(1) para consultar // Permite atribuir valores a linhas e colunas, e consultar a posicao
// Complexidade Espacial: O(n log n)
2b7 #include <bits/stdc++.h>
ca4 using namespace std;
```

```
005 #define MAX 100010
352 #define MAX2 20 // log(MAX)
82d namespace SparseTable {
        int m[MAX2][2*MAX], n, v[2*MAX];
9bf
b90
        int op(int a, int b) { return a + b; }
        void build(int n2, int* v2) {
8b0
1e3
            n = n2;
            for (int i = 0; i < n; i++) v[i] = v2[i];</pre>
df4
a84
            while (n&(n-1)) n++;
3d2
            for (int j = 0; (1<<j) < n; j++) {
1c0
                int len = 1<<i:</pre>
d9b
                for (int c = len; c < n; c += 2*len) {
332
                     m[j][c] = v[c], m[j][c-1] = v[c-1];
                     for (int i = c+1; i < c+len; i++) m[j][i] =</pre>
668
   op(m[i][i-1], v[i]);
                    for (int i = c-2; i >= c-len; i--) m[j][i] =
432
   op(v[i], m[j][i+1]);
eda
f4d
            }
ce3
9e3
        int query(int 1, int r) {
f13
            if (1 == r) return v[1];
e6d
            int j = __builtin_clz(1) - __builtin_clz(1^r);
d67
            return op(m[j][1], m[j][r]);
a7b
        }
258 }
63d void solve() {
        int n = 9:
1a3
        int v[] = \{1, 2, 3, 4, 5, 6, 7, 8, 9\};
3f7
        SparseTable::build(n, v);
        cout << SparseTable::query(0, n-1) << endl; // sparse[0] +</pre>
   sparse[1] + ... + sparse[n-1] = 45
241 }
```

#### 2.5 Tabuleiro

```
// Description: Estrutura que simula um tabuleiro M x N, sem realmente
    criar uma matriz
    mais frequente
// Complexidade Atribuir: O(log(N))
// Complexidade Consulta: O(log(N))
// Complexidade verificar frequencia geral: O(N * log(N))
```

```
9a0 #define MAX_VAL 5 // major valor que pode ser adicionado na matriz
   + 1
8ee class BinTree {
d9d
        protected:
ef9
            vector < int > mBin;
673
        public:
            explicit BinTree(int n) { mBin = vector(n + 1, 0); }
            void add(int p, const int val) {
e44
dd1
                for (auto size = mBin.size(); p < size; p += p & -p)</pre>
174
                     mBin[p] += val:
b68
            }
e6b
            int query(int p) {
                int sumToP {0};
e1c
b62
                for (; p > 0; p -= p \& -p)
                     sumToP += mBin[p];
ec1
838
                return sumToP;
793
            }
a5f };
b6a class ReverseBinTree : public BinTree {
673
        public:
83e
            explicit ReverseBinTree(int n) : BinTree(n) {};
e44
            void add(int p, const int val) {
850
                BinTree::add(static_cast < int > (mBin.size()) - p, val);
705
            }
            int query(int p) {
e6b
                return BinTree::query(static_cast < int > (mBin.size()) -
164
   p);
a21
            }
6cf };
952 class Tabuleiro {
673
        public:
177
            explicit Tabuleiro (const int m, const int n, const int q)
   : mM(m), mN(n), mQ(q) {
                mLinhas = vector < pair < int , int8_t >> (m, {0, 0});
958
                mColunas = vector<pair<int, int8_t>>(n, {0, 0});
d68
                mAtribuicoesLinhas = vector(MAX_VAL,
   ReverseBinTree(mQ)); // aARvore[51]
                mAtribuicoesColunas = vector(MAX VAL.
9e5
   ReverseBinTree(mQ)):
```

```
13b
            }
            void atribuirLinha(const int x, const int8_t r) {
bc2
                 mAtribuirFileira(x, r, mLinhas, mAtribuicoesLinhas);
e88
            }
062
ca2
            void atribuirColuna(const int x, const int8 t r) {
689
                 mAtribuirFileira(x, r, mColunas, mAtribuicoesColunas);
a40
            }
d10
            int maxPosLinha(const int x) {
f95
                 return mMaxPosFileira(x. mLinhas. mAtribuicoesColunas.
   mM):
8ba
            }
            int maxPosColuna(const int x) {
ff7
                 return mMaxPosFileira(x, mColunas, mAtribuicoesLinhas,
b95
   mN):
            }
252
80e
            vector < int > frequenciaElementos() {
a35
                vector < int > frequenciaGlobal(MAX_VAL, 0);
45a
                for(int i=0; i<mM; i++) {</pre>
                     vector<int> curr = frequenciaElementos(i,
ebd
   mAtribuicoesColunas):
97f
                     for(int j=0; j<MAX_VAL; j++)</pre>
ef3
                         frequenciaGlobal[j] += curr[j];
094
01e
                return frequenciaGlobal;
            }
b7a
bf2
        private:
69d
            int mM. mN. mQ. mMoment {0}:
0a6
            vector < ReverseBinTree > mAtribuicoesLinhas,
   mAtribuicoesColunas:
f2d
            vector < pair < int , int8_t >> mLinhas , mColunas ;
            void mAtribuirFileira(const int x, const int8_t r,
e7a
   vector<pair<int, int8_t>>& fileiras,
                                 vector < ReverseBinTree > & atribuicoes) {
1d7
                if (auto& [oldQ, oldR] = fileiras[x]; oldQ)
224
bda
                     atribuicoes[oldR].add(oldQ, -1);
914
                 const int currentMoment = ++mMoment:
b2c
                fileiras[x].first = currentMoment:
```

```
80b
                fileiras[x].second = r;
                atribuicoes[r].add(currentMoment, 1);
f65
            }
5de
            int mMaxPosFileira(const int x, const vector<pair<int,</pre>
2b8
   int8_t>>& fileiras, vector < ReverseBinTree >&
   atribuicoesPerpendiculares, const int& currM) const {
                auto [momentoAtribuicaoFileira, rFileira] =
   fileiras[x];
8d0
                vector < int > fileiraFrequencia(MAX_VAL, 0);
729
                fileiraFrequencia[rFileira] = currM;
85a
                for (int8_t r {0}; r < MAX_VAL; ++r) {</pre>
8ca
                     const int frequenciaR =
   atribuicoesPerpendiculares[r].query(momentoAtribuicaoFileira + 1);
                     fileiraFrequencia[rFileira] -= frequenciaR;
04a
                     fileiraFrequencia[r] += frequenciaR;
72e
                }
6b0
b59
                return MAX_VAL - 1 -
   (max_element(fileiraFrequencia.crbegin(),
   fileiraFrequencia.crend()) - fileiraFrequencia.crbegin());
372
7c4
            vector < int > frequenciaElementos(int x,
   vector < ReverseBinTree > & atribuicoesPerpendiculares) const {
8d0
                vector < int > fileiraFrequencia(MAX_VAL, 0);
                auto [momentoAtribuicaoFileira, rFileira] = mLinhas[x];
583
                fileiraFrequencia[rFileira] = mN;
                for (int8_t r {0}; r < MAX_VAL; ++r) {</pre>
85a
8ca
                     const int frequenciaR =
   atribuicoesPerpendiculares[r].query(momentoAtribuicaoFileira + 1);
                     fileiraFrequencia[rFileira] -= frequenciaR;
04a
72e
                     fileiraFrequencia[r] += frequenciaR;
6b0
                }
2e6
                return fileiraFrequencia;
            }
15d
20c };
63d void solve() {
```

```
e29
        int L, C, q; cin >> L >> C >> q;
56c
        Tabuleiro tabuleiro (L, C, q);
a09
        int linha = 0, coluna = 0, valor = 10; // linha e coluna sao 0
   based
        tabuleiro.atribuirLinha(linha, static_cast<int8_t>(valor)); //
b68
   f(i,0,C) matriz[linha][i] = valor
        tabuleiro.atribuirColuna(coluna, static_cast <int8_t>(valor));
   // f(i,0,L) matriz[i][coluna] = valor
        // Freuencia de todos os elementos, de O a MAX VAL-1
155
        vector<int> frequenciaGeral = tabuleiro.frequenciaElementos();
        int a = tabuleiro.maxPosLinha(linha); // retorna a posicao do
176
   elemento mais frequente na linha
        int b = tabuleiro.maxPosColuna(coluna); // retorna a posicao
981
   do elemento mais frequente na coluna
9b5 }
```

#### Union-Find (Disjoint Set Union)

```
f3b const int MAX = 5e4+10;
074 int p[MAX], ranking[MAX], setSize[MAX];
Ocd struct UnionFind {
        int numSets;
c55
02d
        UnionFind(int N) {
680
            iota(p,p+N+1,0);
340
            memset(ranking, 0, sizeof ranking);
            memset(setSize, 1, sizeof setSize);
f0a
            numSets = N:
0bd
        }
142
        int numDisjointSets() { return numSets; }
c59
        int sizeOfSet(int i) { return setSize[find(i)]; }
a5b
8ee
        int find(int i) { return (p[i] == i) ? i : (p[i] =
   find(p[i])): }
        bool same(int i, int j) { return find(i) == find(j); }
da3
92e
        void uni(int i, int j) {
ea5
            if (same(i, j))
505
                return;
```

```
c56
            int x = find(i), y = find(j);
e4f
            if (ranking[x] > ranking[y])
9dd
                swap(x, y);
ae9
            y = [x]q
            if (ranking[x] == ranking[y])
6e9
                ++ranking[y];
3cf
223
            setSize[y] += setSize[x];
92a
            --numSets;
e3f
       }
b6b };
63d void solve() {
f98
        int n, ed; cin >> n >> ed;
f4e
        UnionFind uni(n);
        f(i,0,ed) {
31c
602
            int a, b; cin >> a >> b; a--, b--;
45e
            uni.uni(a,b);
c0f
        }
350
        cout << uni.numDisjointSets() << endl;</pre>
01b }
```

## 3 Grafos

## 3.1 Emparelhamento Max Grafo Bipartido (Kuhn)

```
6c6 struct kuhn {
14e
        int n, m;
789
        vector < vector < int >> g;
d3f
        vector<int> vis, ma, mb;
        kuhn(int n_, int m_) : n(n_), m(m_), g(n),
40e
8af
            vis(n+m), ma(n, -1), mb(m, -1) {}
        void add(int a, int b) { g[a].push_back(b); }
ba6
        bool dfs(int i) {
caf
29a
            vis[i] = 1;
29b
            for (int j : g[i]) if (!vis[n+j]) {
8c9
                vis[n+i] = 1;
2cf
                if (mb[j] == -1 or dfs(mb[j])) {
bfe
                    ma[i] = j, mb[j] = i;
8a6
                    return true;
                }
b17
82a
d1f
            return false;
        }
4ef
bf7
        int matching() {
1ae
            int ret = 0, aum = 1;
5a8
            for (auto& i : g) shuffle(i.begin(), i.end(), rng);
392
            while (aum) {
618
                for (int j = 0; j < m; j++) vis[n+j] = 0;
c5d
                aum = 0:
830
                for (int i = 0: i < n: i++)
01f
                    if (ma[i] == -1 and dfs(i)) ret++, aum = 1;
085
edf
            return ret;
2ee
        }
b0d }:
63d void solve() {
be0
        int n1; // Num vertices lado esquerdo grafo bipartido
        int n2; // Num vertices lado direito grafo bipartido
e4c
761
        kuhn K(n1, n2);
732
        int edges;
6e0
        while(edges--) {
b1f
            int a, b; cin >> a >> b;
3dc
            K.add(a.b): // a -> b
5b7
        }
```

```
69b
        int emparelhamentoMaximo = K.matching();
76b }
3.2 Fluxo - Dinitz (Max Flow)
// Encontra fluxo maximo de um grafo
// O(min(m * max flow, n^2 m))
// Grafo com capacidades 1: O(\min(m \text{ sqrt}(m), m * n^2(2/3)))
// Todo vertice tem grau de entrada ou saida 1: O(m sqrt(n))
472 struct dinitz {
        const bool scaling = false; // com scaling -> 0(nm
   log(MAXCAP)),
                                     // com constante alta
206
        int lim;
670
        struct edge {
358
            int to, cap, rev, flow;
7f9
            bool res:
d36
            edge(int to_, int cap_, int rev_, bool res_)
a94
                : to(to_), cap(cap_), rev(rev_), flow(0), res(res_) {}
f70
        };
002
        vector < vector < edge >> g;
216
        vector < int > lev, beg;
a71
        11 F;
190
        dinitz(int n) : g(n), F(0) {}
087
        void add(int a, int b, int c) {
            g[a].emplace_back(b, c, g[b].size(), false);
bae
4c6
            g[b].emplace_back(a, 0, g[a].size()-1, true);
5c2
        }
123
        bool bfs(int s, int t) {
90f
            lev = vector<int>(g.size(), -1); lev[s] = 0;
64c
            beg = vector<int>(g.size(), 0);
8b2
            queue < int > q; q.push(s);
402
            while (q.size()) {
                int u = q.front(); q.pop();
be1
bd9
                for (auto& i : g[u]) {
                    if (lev[i.to] != -1 or (i.flow == i.cap)) continue;
dbc
                    if (scaling and i.cap - i.flow < lim) continue;</pre>
b4f
185
                    lev[i.to] = lev[u] + 1;
8ca
                    q.push(i.to);
f97
                }
            }
e87
0de
            return lev[t] != -1;
742
        }
```

```
dfb
        int dfs(int v, int s, int f = INF) {
50b
            if (!f or v == s) return f;
88f
            for (int& i = beg[v]; i < g[v].size(); i++) {</pre>
                auto& e = g[v][i];
027
                if (lev[e.to] != lev[v] + 1) continue;
206
                 int foi = dfs(e.to, s, min(f, e.cap - e.flow));
ee0
749
                if (!foi) continue:
                e.flow += foi, g[e.to][e.rev].flow -= foi;
3c5
45c
                return foi;
            }
618
bb3
            return 0;
4b1
        }
ff6
        11 max flow(int s. int t) {
a86
            for (lim = scaling ? (1 << 30) : 1; lim; lim /= 2)
9d1
                 while (bfs(s, t)) while (int ff = dfs(s, t)) F += ff;
4ff
            return F;
8b9
        }
86f };
// Recupera as arestas do corte s-t
dbd vector<pair<int, int>> get_cut(dinitz& g, int s, int t) {
        g.max_flow(s, t);
f07
68c
        vector < pair < int , int >> cut;
        vector < int > vis(g.g.size(), 0), st = {s};
1b0
321
        vis[s] = 1:
3c6
        while (st.size()) {
            int u = st.back(); st.pop_back();
b17
322
            for (auto e : g.g[u]) if (!vis[e.to] and e.flow < e.cap)</pre>
                vis[e.to] = 1, st.push_back(e.to);
c17
d14
        }
481
        for (int i = 0; i < g.g.size(); i++) for (auto e : g.g[i])
9d2
            if (vis[i] and !vis[e.to] and !e.res) cut.emplace_back(i,
   e.to);
d1b
        return cut:
1e8 }
63d void solve() {
1a8
        int n; // numero de arestas
b06
        dinitz g(n);
732
        int edges;
        while(edges--) {
6e0
1e1
            int a, b, w; cin >> a >> b >> c;
f93
            g.add(a,b,c); // a -> b com capacidade c
        }
fa1
```

```
07a    int maxFlow = g.max_flow(SRC, SNK); // max flow de SRC -> SNK
a7b }
```

#### 3.3 Fluxo - MinCostMaxFlow

```
// min_cost_flow(s, t, f) computa o par (fluxo, custo)
// com max(fluxo) <= f que tenha min(custo)</pre>
// min_cost_flow(s, t) -> Fluxo maximo de custo minimo de s pra t
// Se for um dag, da pra substituir o SPFA por uma DP pra nao
// pagar O(nm) no comeco
// Se nao tiver aresta com custo negativo, nao precisa do SPFA
// O(nm + f * m log n)
123 template < typename T > struct mcmf {
670
        struct edge {
b75
            int to, rev, flow, cap; // para, id da reversa, fluxo,
   capacidade
7f9
            bool res: // se eh reversa
            T cost: // custo da unidade de fluxo
635
            edge(): to(0), rev(0), flow(0), cap(0), cost(0),
892
   res(false) {}
1d7
            edge(int to_, int rev_, int flow_, int cap_, T cost_, bool
   res_)
f8d
                : to(to_), rev(rev_), flow(flow_), cap(cap_),
   res(res_), cost(cost_) {}
723
        };
002
        vector < vector < edge >> g;
168
        vector < int > par_idx, par;
f1e
        T inf;
a03
        vector <T> dist;
        mcmf(int n) : g(n), par_idx(n), par(n),
b22
   inf(numeric_limits <T>::max()/3) {}
91c
        void add(int u, int v, int w, T cost) { // de u pra v com cap
   w e custo cost
            edge a = edge(v, g[v].size(), 0, w, cost, false);
2fc
234
            edge b = edge(u, g[u].size(), 0, 0, -cost, true);
b24
            g[u].push_back(a);
c12
            g[v].push_back(b);
0ed
        }
```

```
8bc
        vector<T> spfa(int s) { // nao precisa se nao tiver custo
   negativo
871
             deque < int > q;
3d1
             vector < bool > is_inside(g.size(), 0);
             dist = vector<T>(g.size(), inf);
577
a93
             dist[s] = 0;
a30
            q.push_back(s);
             is_inside[s] = true;
ecb
14d
             while (!q.empty()) {
b1e
                 int v = q.front();
ced
                 q.pop_front();
48d
                 is_inside[v] = false;
76e
                 for (int i = 0; i < g[v].size(); i++) {</pre>
9d4
                     auto [to, rev, flow, cap, res, cost] = g[v][i];
                     if (flow < cap and dist[v] + cost < dist[to]) {</pre>
e61
943
                          dist[to] = dist[v] + cost;
                         if (is_inside[to]) continue;
ed6
020
                          if (!q.empty() and dist[to] > dist[q.front()])
   q.push_back(to);
b33
                          else q.push_front(to);
b52
                          is_inside[to] = true;
2d1
                     }
8cd
                 }
f2c
            }
8d7
            return dist;
96c
        }
2a2
        bool dijkstra(int s, int t, vector<T>& pot) {
489
             priority_queue < pair < T, int > , vector < pair < T, int > > ,
   greater<>> q;
577
             dist = vector<T>(g.size(), inf);
             dist[s] = 0;
a93
115
            q.emplace(0, s);
             while (q.size()) {
402
91b
                 auto [d, v] = q.top();
833
                 q.pop();
68b
                 if (dist[v] < d) continue;</pre>
                 for (int i = 0; i < g[v].size(); i++) {</pre>
76e
9d4
                     auto [to, rev, flow, cap, res, cost] = g[v][i];
e8c
                     cost += pot[v] - pot[to];
                     if (flow < cap and dist[v] + cost < dist[to]) {</pre>
e61
943
                         dist[to] = dist[v] + cost;
441
                         q.emplace(dist[to], to);
88b
                         par_idx[to] = i, par[to] = v;
```

```
873
de3
9d4
1d4
             return dist[t] < inf;</pre>
c68
3d2
        pair < int, T > min_cost_flow(int s, int t, int flow = INF) {
             vector <T> pot(g.size(), 0);
3dd
9e4
             pot = spfa(s); // mudar algoritmo de caminho minimo aqui
d22
            int f = 0;
            T ret = 0:
4a0
             while (f < flow and dijkstra(s, t, pot)) {</pre>
bda
                 for (int i = 0; i < g.size(); i++)</pre>
d2a
                     if (dist[i] < inf) pot[i] += dist[i];</pre>
                 int mn_flow = flow - f, u = t;
71b
                 while (u != s){
045
                     mn_flow = min(mn_flow,
90f
07d
                         g[par[u]][par_idx[u]].cap -
   g[par[u]][par_idx[u]].flow);
                     u = par[u];
3d1
935
                 ret += pot[t] * mn_flow;
1f2
476
                 u = t:
045
                 while (u != s) {
                     g[par[u]][par_idx[u]].flow += mn_flow;
e09
                     g[u][g[par[u]][par_idx[u]].rev].flow -= mn_flow;
d98
                     u = par[u];
                 }
bcc
04d
                 f += mn flow:
36d
15b
             return make_pair(f, ret);
        }
cc3
        // Opcional: retorna as arestas originais por onde passa flow
            = cap
182
        vector < pair < int , int >> recover() {
             vector < pair < int , int >> used;
24a
2a4
             for (int i = 0; i < g.size(); i++) for (edge e : g[i])</pre>
                 if(e.flow == e.cap && !e.res) used.push_back({i,
587
   e.to}):
f6b
             return used:
```

```
390  }
697 };
63d void solve(){

1a8    int n; // numero de vertices
4c5    mcmf < int > mincost(n);

ab4    mincost.add(u, v, cap, cost); // unidirectional
983    mincost.add(v, u, cap, cost); // bidirectional

073    auto [flow, cost] = mincost.min_cost_flow(src, end/*, initialFlow*/);
```

#### 3.4 Fluxo - Problemas

```
// 1: Problema do Corte
7a9 - Entrada:
bc1
       - N itens
388
       - Curso Wi Inteiro
       - M restricoes: se eu pegar Ai, eu preciso pegar Bi...
387 - Saida: valor maximo pegavel
ac2 - Solucao: corte maximo com Dinitz
       - dinitz(n+m+1)
019
593
       - f(i,0,n): i \rightarrow SNK com valor Ai
       - f(i,0,m):
9eb
9e2
           * SRC -> n+i com valor Wi
           * ParaTodo dependente Bj: n+i -> Bj com peso INF
a9e
8a0
       - ans = somatorio(Wi) - maxFlow(SRC, SNK);
/* ============ */
```

## 4 Matematica

## 4.1 Numero de Digitos

```
// Calcula o numero de digitos de n // 1234 = 4; 0 = 1
```

```
09c int numDigits(int n) {
       if (n == 0) return 1;
662
        n = std::abs(n);
146
       return static_cast < int > (std::floor(std::log10(n))) + 1;
d2b }
    Primos - Lowest Prime Factor
// Menor fator primo de n
// O(sqrt(n))
074 int lowestPrimeFactor(int n, int startPrime = 2) {
9d5
        if (startPrime <= 3) {</pre>
            if (not (n & 1)) return 2:
fb4
5a0
            if (not (n % 3)) return 3;
            startPrime = 5;
72a
43a
       }
        for (int i = startPrime; i * i <= n; i += (i + 1) % 6 ? 4 : 2)</pre>
c94
dcb
            if (not (n % i))
d9a
                return i;
041
        return n;
6c5 }
4.3 Primos - Primo
// Verifica se um numero eh primo
// O(sqrt(n))
5f7 bool isPrime(int n) {
        return n > 1 and lowestPrimeFactor(n) == n;
822 }
    String
5.1 Longest Common Subsequence 1 (LCS)
// Retrorna a LCS entre as string S e T.
// Armazena em memo[i][j] o LCS_SZ de s[i...n] e t[j...m].
// Implementacao recursiva
```

//

```
// Temporal: O(n*m)
// Espacial: O(n*m)
da5 const int MAXN = 1e3+10;
dd0 int memo[MAXN][MAXN];
// Calcula tamanho do LCS recursivamente
28f int lcs_sz(string& s, string& t, int i, int j) {
        if(i == s.size() or j == t.size()) return 0;
e80
        auto& ans = memo[i][j];
d64
        if (\simans) return ans;
1a9
        if(s[i] == t[i])
            ans = 1 + lcs_sz(s,t,i+1, j+1);
176
295
        else
3af
            ans = max(
c19
                    lcs_sz(s,t,i+1,j),
364
                    lcs_sz(s,t,i,j+1)
616
                );
ba7
        return ans:
// Armazena em ans a LCS entre S e T
10e void lcs(string& ans, string& s, string& t, int i, int j) {
        if(i >= s.size() or j >= t.size()) return;
524
        if(s[i] == t[j]) {
b80
            ans.push_back(s[i]);
081
            return lcs(ans, s, t, i+1, j+1);
a00
4cb
        if(lcs_sz(s,t,i+1,j) > lcs_sz(s,t,i,j+1)) return lcs(ans, s,
   t, i+1, j);
4f2
        return lcs(ans, s, t, i, j+1);
bc5 }
63d void solve() {
bfb
        string s, t; cin >> s >> t;
457
        memset(memo, -1, sizeof memo);
a4d
        string ans; lcs(ans, s, t, 0,0);
        cout << ans << endl;</pre>
886
030 }
```

## 5.2 Split de String

```
// O(|s| * |del|).
5a6 vector<string> split(string s, string del = " ") {
       vector < string > retorno;
cd5
0f4
      int start, end = -1*del.size();
016
      do {
a3b
           start = end + del.size();
257
        end = s.find(del, start);
36 c
         retorno.push_back(s.substr(start, end - start));
3a7
      } while (end != -1);
       return retorno;
5fa
f80 }
```

```
6 Extra
6.1 fastIO.cpp
int read_int() {
    bool minus = false;
    int result = 0;
    char ch;
    ch = getchar();
    while (1) {
       if (ch == '-') break;
       if (ch >= '0' && ch <= '9') break;
        ch = getchar();
    if (ch == '-') minus = true;
    else result = ch-'0';
    while (1) {
        ch = getchar();
       if (ch < '0' || ch > '9') break;
        result = result *10 + (ch - '0');
    if (minus) return -result;
    else return result;
6.2 hash.sh
# Para usar (hash das linhas [11, 12]):
# bash hash.sh arquivo.cpp 11 12
sed -n $2','$3' p' $1 | sed '/^#w/d' | cpp -dD -P -fpreprocessed | tr
   -d '[:space:]' | md5sum | cut -c-6
6.3 stress.sh
P = a
make ${P} ${P}2 gen || exit 1
for ((i = 1; ; i++)) do
    ./gen $i > in
    ./${P} < in > out
    ./${P}2 < in > out2
    if (! cmp -s out out2) then
        echo "--> entrada:"
        cat in
```

```
echo "--> saida1:"
        cat out
        echo "--> saida2:"
        cat out2
        break;
    echo $i
done
6.4 pragma.cpp
// Otimizacoes agressivas, pode deixar mais rapido ou mais devagar
#pragma GCC optimize("Ofast")
// Auto explicativo
#pragma GCC optimize("unroll-loops")
// Vetorizacao
#pragma GCC target("avx2")
// Para operacoes com bits
#pragma GCC target("bmi,bmi2,popcnt,lzcnt")
6.5 timer.cpp
// timer T; T() -> retorna o tempo em ms desde que declarou
using namespace chrono;
struct timer : high_resolution_clock {
    const time_point start;
    timer(): start(now()) {}
    int operator()() {
        return duration_cast < milliseconds > (now() - start).count();
    }
};
6.6 vimrc
d79 set ts=4 sw=4 mouse=a nu ai si undofile
7c9 function H(1)
        return system("sed '/^#/d' | cpp -dD -P -fpreprocessed | tr -d
   '[:space:]' | md5sum", a:1)
Obe endfunction
329 function P() range
dd8
       for i in range(a:firstline, a:lastline)
```

```
ссс
             let 1 = getline(i)
139
             call cursor(i, len(1))
             echo H(getline(search('{}'[1], 'bc', i) ? searchpair('{',
7c9
    '', '}', 'bn') : i, i))[0:2] 1
         endfor
bf9
Obe endfunction
90e vmap \langle C-H \rangle : call P()\langle CR \rangle
de2 "" }
6.7 debug.cpp
void debug_out(string s, int line) { cerr << endl; }</pre>
template < typename H, typename ... T>
void debug_out(string s, int line, H h, T... t) {
    if (s[0] != ',') cerr << "Line(" << line << ") ";</pre>
    do { cerr << s[0]; s = s.substr(1);</pre>
    } while (s.size() and s[0] != ',');
    cerr << " = " << h;
    debug_out(s, line, t...);
#ifdef DEBUG
#define debug(...) debug_out(#__VA_ARGS__, __LINE__, __VA_ARGS__)
#else
#define debug(...) 42
#endif
6.8 makefile
CXX = g++
CXXFLAGS = -fsanitize=address, undefined -fno-omit-frame-pointer -g
    -Wall -Wshadow -std=c++17 -Wno-unused-result -Wno-sign-compare
    -Wno-char-subscripts #-fuse-ld=gold
clearexe:
    find . -maxdepth 1 -type f -executable -exec rm {} +
6.9 temp.cpp
#include <bits/stdc++.h>
using namespace std;
#define _ ios_base::sync_with_stdio(0);cin.tie(0);
```

```
#define all(a)
                      a.begin(), a.end()
#define int
                      long lo
                                 ng int
#define double
                     long double
#define f(i,s,e)
                     for(int i=s;i<e;i++)</pre>
#define dbg(x) cout << #x << " = " << x << " ";
#define dbgl(x) cout << #x << " = " << x << endl;</pre>
#define vi
                      vector<int>
#define pii
                      pair < int , int >
                      "\n"
#define endl
#define print_v(a)
                      for(auto x : a)cout << x << " "; cout << endl</pre>
#define print_vp(a) for(auto x : a)cout << x.first << " "<< x.second << endl</pre>
#define rf(i,e,s)
                      for(int i=e-1:i>=s:i--)
                     ((a) + (b - 1))/b
#define CEIL(a, b)
#define TRUNC(x, n) floor(x * pow(10, n))/pow(10, n)
#define ROUND(x, n) round(x * pow(10, n))/pow(10, n)
const int INF = 1e9; // 2^31-1
const int LLINF = 4e18; // 2^63-1
const double EPS = 1e-9;
const int MAX = 1e6+10; // 10^6 + 10
void solve() {
}
int32_t main() { _
    int t = 1; // cin >> t;
    while (t--) {
        solve();
    }
    return 0;
}
6.10 rand.cpp
mt19937 rng((int)
    chrono::steady_clock::now().time_since_epoch().count());
int uniform(int 1, int r){
    uniform_int_distribution < int > uid(1, r);
    return uid(rng);
}
```