

The Why and How of Nonnegative Matrix Factorization

Topic Presentation

Group 02

LINMA2380 — Matrix computations

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Summary

1 Introduction

2 Applications

3 Connections to other problems

Agenda

Use : Analysis of high-dimensional data by automatically extracts sparse and meaningful features from a set of nonnegative data vectors

- 1 What : Definitions and properties
- 2 Why : Applications
- 3 How : Algorithms
- 4 What next : Connections with Problems in Mathematics and Computer Science
- 5 Conclusion

What : Definitions and properties

Nonnegative matrix factorization (NMF) is a Linear dimensionality reduction (LDR)

LDR :

- From a set of data points $x_j \in R^p$ for $1 \leq j \leq n$
- To a set of dimension $r < \min(p, n)$
- Thanks to $w_k \in R^p$ for $1 \leq k \leq r$
- Such that : $\forall j, x_j \approx \sum_{k=1}^r w_k h_j(k)$, for some weights $h_j \in R^r$

Equivalent to **low-rank matrix approximation** : $X \approx WH$

- $X \in R^{p \times n}$: $X(:, j) = x_j$ for $1 \leq j \leq n$
- $W \in R^{p \times r}$: $W(:, k) = w_k$ for $1 \leq k \leq r$
- $H \in R^{r \times n}$: $H(:, j) = h_j$ for $1 \leq j \leq n$

NMF : decomposing a given nonnegative data matrix X as $X \approx WH$

Applications - Image processing

Goal : Facial Feature Extraction



Data matrix : $X \in \mathbb{R}_+^{p \times n}$

- p : total number of pixels
- n : number of faces
- $X(i, j)$: the gray-level of the i -th pixel in the j -th face

Applications - Image processing

$$\underbrace{X(:, j)}_{j\text{th facial image}} \approx \sum_{k=1}^r$$


$$\underbrace{W(:, k)}$$

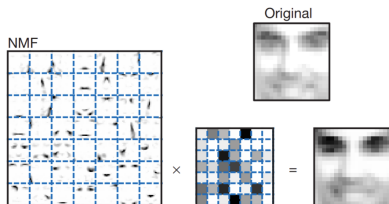
facial features



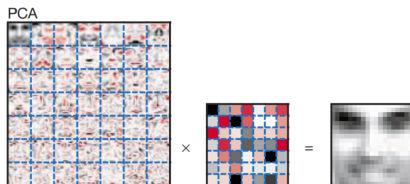
$\underbrace{H(k, j)}$	$=$	$\underbrace{WH(:, j)}$
importance of features in j th image		approximation of j th image



Applications - Image processing



NMF decomposition



PCA decomposition

Applications - Text Mining

Goal : Topic Recovery and Document Classification

Data matrix : $X \in \mathbb{R}_+^{n \times m}$

- each column : a document
- each line : a word
- $X(i, j)$: number of times the i -th word appears in the j -th document

$$\underbrace{X(:, j)}_{j\text{th document}} \approx \sum_{k=1}^r \underbrace{W(:, k)}_{k\text{th topic}} \underbrace{H(k, j)}_{\substack{\text{importance of } k\text{th topic} \\ \text{in } j\text{th document}}}, \quad \text{with } W \geq 0 \text{ and } H \geq 0.$$

Applications - Hyperspectral Unmixing

Goal :

- 1 Identify the constitutive materials present in an image
- 2 Classify the pixels according to their constitutive materials

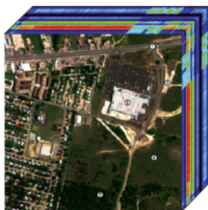
Spectral signature of a pixel: fraction of incident light being reflected by that pixel at different wavelengths

Applications - Hyperspectral Unmixing

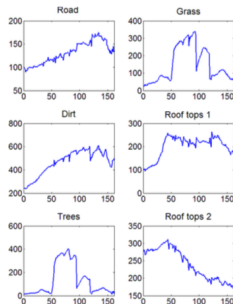
Data matrix : $X \in \mathbb{R}^{n \times m}$

- each column : spectral signature of a pixel

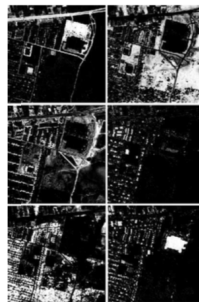
$\underbrace{X(:, j)}$
spectral signature
of j th pixel



$\approx \sum_{k=1}^r \underbrace{W(:, k)}$
spectral signature
of k th endmember



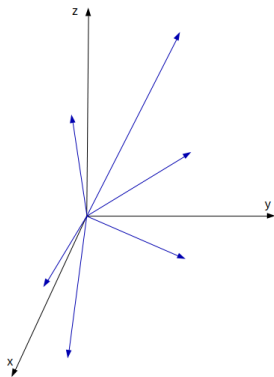
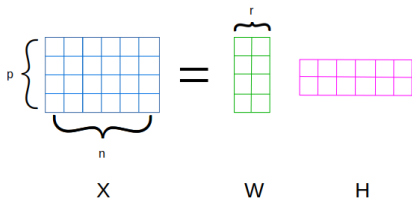
$\underbrace{H(k, j)}$
abundance of k th endmember
in j th pixel



Nonnegative rank

Definition (Nonnegative rank)

Given $X \in \mathbb{R}_+^{p \times n}$, the nonnegative rank of X , denoted $\text{rank}_+(X)$ is the minimum r s.t. $\exists W \in \mathbb{R}_+^{p \times r}, H \in \mathbb{R}_+^{r \times n}$ with $X = WH$.



Computational Geometry : Nested polytopes problem

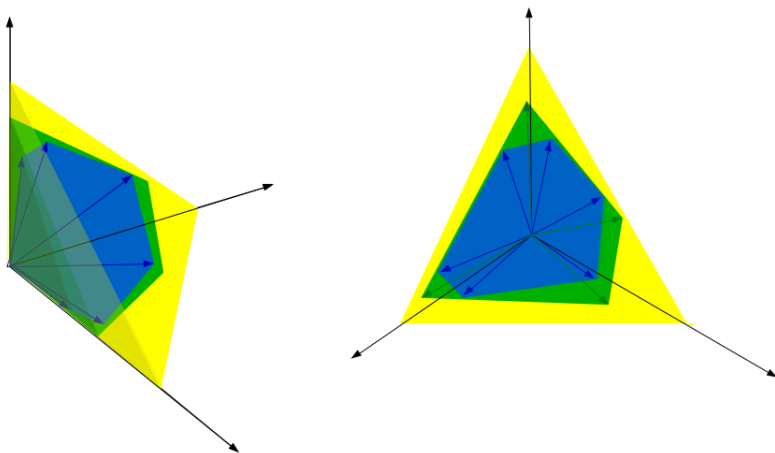


Figure: Finding a polytope with minimum nb of vertices nested between 2 polytopes

Graph Theory : Bipartite dimension

Let $G(X) = (V_1 \cup V_2, E)$ be a bipartite graph induced by X (i.e. $(i, j) \in E \Leftrightarrow X_{ij} \neq 0$).

Definition (Biclique and bipartite dimension)

- A *biclique* (or a *complete bipartite graph*) is a bipartite graph s.t. every vertex in V_1 is connected to every vertex in V_2 .
- The *bipartite dimension* (or the *minimum biclique cover*) $bc(G(X))$ is the minimum number of bicliques needed to cover all edges in E .

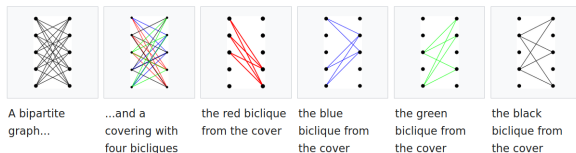


Figure: Example for biclique edge cover [biclique]

For any $(W, H) \geq 0$ s.t. $X = WH = \sum_{k=1}^r W_{:k}H_{k:}$, we have

$$G(X) = \cup_{k=1}^r G(W_{:k}H_{k:})$$

where $G(W_{:k}H_{k:})$ are complete bipartite subgraphs ($\text{bc}(G(W_{:k}H_{k:}) = 1 \forall k$).

Theorem (Rectangle covering bound)

$$\text{bc}(G(X)) \leq \text{rank}_+(X)$$

Communication complexity



Alice and Bob want to compute :

$$f : \{0, 1\}^m \times \{0, 1\}^n \rightarrow \{0, 1\} : (x, y) \rightarrow f(x, y)$$

While minimizing the nb of bits exchanged (i.e. communication complexity (CC)).

Nondeterministic comm. complex. of f (NCC): CC of f with oracle/message before starting the communication.

The communication matrix $X \in \{0, 1\}^{2^n \times 2^m}$ is equal to the function f for all possible combinations of inputs.

Theorem (Yannakis)

$$NCC \text{ of } f \leq \log_2(\text{rank}_+(X))$$

Linear Optimization : Extended formulation

$$\begin{aligned}
 (\text{LP}) \quad & \max && c^T x \\
 & \text{s.t.} && Ax \leq b \\
 & && x \in \mathbb{R}^n \geq 0
 \end{aligned}$$

Definition (Extended formulation)

The extended formulation of a polytope P is a higher dimensional polytope Q and a linear projection π s.t. $\pi(Q) = P$.

In our LP problem, an extended formulation of the polytope $P \subset \mathbb{R}^n$ defined by the constraints $Ax \leq b$, is a polytope $Q \subset \mathbb{R}^{n+r}$ defined by $Cx + Dy \leq d$ with $y \in \mathbb{R}^r$, s.t. $\pi(Q) = P$.

The slack matrix $X(i, j) = b_i - A_i v_j$.

With v_j , the j^{th} vertex of P and $\{x \in \mathbb{R}^n | b_i - A_i x \geq 0\}$ its i^{th} facet.

The (i, j) entry measures the slack of the i^{th} inequality for the vertex j .

Theorem (Yannakis)

The minimum size of an extended formulation Q of P is equal to $\text{rank}_+(X)$.



