

LMAT2450 – First Homework

Question 1: Pseudorandom Generator.

Let G be a pseudorandom generator (PRG) such that, if $s \in \{0,1\}^n$, then $G(s) \in \{0,1\}^{l(n)}$ with $l(n) > n$.

We define, for $s_1, s_2 \in \{0,1\}^n$, $G'(s_1||s_2) := G(s_1)||s_2$ where $G'(s_1||s_2) \in \{0,1\}^{l(n)+n}$.

Show that either G' is a PRG by offering a reduction to the security of G , or that G' is not a PRG by exhibiting an attack (building a distinguisher with non negligible advantage).

Solution:

Question 2: Pseudorandom Function.

Consider F a pseudorandom function (PRF) such that $F : \{0, 1\}^n \times \{0, 1\}^n \longrightarrow \{0, 1\}^n$.

We define $F' : \{0, 1\}^n \times \{0, 1\}^{2n} \longrightarrow \{0, 1\}^n$ such that $F'(k, x||y) := F(k, x) \oplus F(k, y)$.

Show either that F' is a PRF by providing a reduction to the security of F , or that F' is not a PRF by exhibiting an attack (building a distinguisher with non negligible advantage).

Solution:

Question 3: How not to derive a PRG.

Let G be a PRG such that, if $s \in \{0, 1\}^n$, then $G(s) \in \{0, 1\}^{l(n)}$ with $l(n) > n$. Define also G' such that $G'(s) := G(s) \oplus (0^{l(n)-n} || s)$.

Show that G' may not be a PRG.

Solution: