

# Selberg's Sieve - Bounding Twin Primes

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# Recall Setup

Let us recall that

$$S(A, \mathcal{P}, z) = \#\{a \in A : a \text{ is not divisible by any } p \leq z \text{ with } p \in \mathcal{P}\}$$

and Selberg's Sieve gives us

$$S(A, \mathcal{P}, z) \leq \frac{X}{V(z)} + \sum_{\substack{d_1, d_2 \leq z \\ d_1, d_2 | P_z}} |R_{[d_1, d_2]}|$$

where

$$V(z) = \sum_{\substack{d \leq z \\ d | P_z}} \frac{\mu^2(d)}{f_1(d)} \quad f(n) = \sum_{d|n} f_1(d) \quad \text{and} \quad |A_d| = \frac{X}{f(d)} + R_d$$

We want to find an upper bound for  $\frac{1}{V(z)}$ , thus a lower bound for  $V(z)$ , which motivates the following lemma:

# Lemma

## Lemma

Let  $\tilde{f}$  be a completely multiplicative function with  $\tilde{f}(p) := f(p)$  for all primes  $p$ . Then we have

$$V(z) \geq \sum_{\substack{e \leq z \\ p|e \Rightarrow p|P_z}} \frac{1}{\tilde{f}(e)} \quad \text{where} \quad P_z = \prod_{\substack{p \in \mathcal{P} \\ p \leq z}} p$$

**Note:** If  $\mathcal{P}$  is the set of all primes, then the second condition  $p \mid e \Rightarrow p \mid P_z$  is trivial.

# Twin Primes

## Definition

*A prime  $p$  is called a twin prime if  $p + 2$  is also a prime.*

Let

$$\pi_2(x) := \# \text{ of twin primes } \leq x$$

We would like to use Selberg's Sieve to obtain an upper bound for  $\pi_2(x)$  as  $x \rightarrow \infty$ .

# Outline of Steps

Once again, Selberg's Sieve gives us

$$S(A, \mathcal{P}, z) \leq \frac{X}{V(z)} + \sum_{\substack{d_1, d_2 \leq z \\ d_1, d_2 | P_z}} |R_{[d_1, d_2]}|$$

To use Selberg's Sieve, we will need to

- Find  $X$ , estimation of the size of  $A$
- Estimate  $|A_d|$  for  $d \mid P_z$  to find our multiplicative function,  $f$
- Find lower bound for  $V(z)$
- Estimate error term

# Understanding $S(A, \mathcal{P}, z)$

In the setting of this problem, we define

$$A = \{n(n+2) : n \leq x\} \quad \text{and} \quad \mathcal{P} = \text{set of all primes}$$

and for  $0 < z < x$ , we have

$$P_z = \prod_{\substack{p \in \mathcal{P} \\ p \leq z}} p = \prod_{p \leq z} p$$

# Understanding $S(A, \mathcal{P}, z)$

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$$P_z = \prod_{\substack{p \in \mathcal{P} \\ p \leq z}} p = \prod_{p \leq z} p$$

and so

$$\begin{aligned} S(A, \mathcal{P}, z) &:= \#\{n(n+2) : n \leq x, \ p \nmid n(n+2) \text{ for all } p \leq z\} \\ &= \#\{n(n+2) : n \leq x, \ p \nmid n \text{ and } p \nmid (n+2) \text{ for all } p \leq z\} \end{aligned}$$

# Understanding $S(A, \mathcal{P}, z)$ Cont'd

- If  $n \leq z$ ,  $n(n+2)$  is not counted in  $S(A, \mathcal{P}, z)$ .
- For all twin primes  $z < p \leq x$ ,  $p(p+2)$  is counted.

And so

$$\begin{aligned}\pi_2(x) &= \sum_{\substack{p \leq x \\ p+2 \in \mathcal{P}}} 1 = \pi_2(z) + \sum_{\substack{z < p \leq x \\ p+2 \in \mathcal{P}}} 1 \\ &\leq \pi_2(z) + S(A, \mathcal{P}, z) \\ &\leq z + S(A, \mathcal{P}, z)\end{aligned}$$



# Next Step

- Find  $X$ , estimation of the size of  $A$  ✓
- Estimate  $|A_d|$  for  $d \mid P_z$  to find our multiplicative function,  $f$
- Find lower bound for  $V(z)$
- Estimate error term

# Estimating $|A_d|$

Let  $d \mid P_z$ , say  $d = p_1 \cdots p_k$ . Then we have

$$\begin{aligned}|A_d| &= \#\{n(n+2) : n \leq x \text{ and } d \mid n(n+2)\} \\ &= \#\{n(n+2) : n \leq x \text{ and } n(n+2) \equiv 0 \pmod{d}\}\end{aligned}$$

## Notation

Let  $N(q) := \#$  of solutions to  $n(n+2) \equiv 0 \pmod{q}$ ,  $q \in \mathbb{N}$

## Notation

Let  $\omega(q) := \#$  the number of prime factors of  $q$ ,  $q \in \mathbb{N}$

# Estimating $|A_d|$ Cont'n

By the Chinese Remainder Theorem,  $n(n+2) \equiv 0 \pmod{d}$  has the same number of solutions as

$$n(n+2) \equiv 0 \pmod{p_1}$$

$$\vdots$$

$$n(n+2) \equiv 0 \pmod{p_k}$$

Since for each  $1 \leq i \leq k$ ,  $N(p_i) \leq 2$ , we have that

$$N(d) = N(p_1) \cdots N(p_k) \leq 2^k = 2^{\omega(d)}$$

# Estimating $|A_d|$ Cont'n

Further, since  $N(d)$  is only the number of solutions modulo  $d$  and we want all solutions  $\leq x$ , we can estimate the total number of solutions, ie. the size of  $A_d$  by

$$|A_d| = \frac{x}{d} \cdot N(d) + R_d, \quad \text{where } R_d \leq N(d) \leq 2^{\omega(d)}$$

Thus, we have our multiplicative function

$$f(d) = \frac{d}{N(d)}$$

which satisfies the conditions of Selberg's Sieve. And a simple fact for later:

$$f(p) = \frac{p}{N(p)} = \begin{cases} p, & \text{if } p = 2 \\ p/2, & \text{if } p > 2 \end{cases}$$

# Next Step

- Find  $X$ , estimation of the size of  $A$  ✓
- Estimate  $|A_d|$  for  $d \mid P_z$  to find our multiplicative function,  $f$  ✓
- Find lower bound for  $V(z)$
- Estimate error term

# Bounding $V(z)$ - Notations

First, let us define some notations

## Definition

For  $n \in \mathbb{N}$ , define

$$\tau_1(n) := \# \text{ odd divisors of } n$$

And so for  $n = 2^s p_1^{e_1} \cdots p_m^{e_m}$ , we have  $\tau_1(n) = (e_1 + 1) \cdots (e_m + 1)$

## Definition

For  $n \in \mathbb{N}$ , define

$$\tau(n) := \# \text{ divisors of } n$$

Note that if  $d$  is square free, then  $\tau(d) = 2^{\omega(d)}$ .

# Bounding $V(z)$

Let  $\tilde{f}$  be a completely multiplicative function with  $\tilde{f}(p) = f(p)$  for all primes  $p$ , as defined in our lemma. Then the lemma tells us that

$$\begin{aligned} V(z) &\geq \sum_{\substack{n \leq z \\ p|n \Rightarrow p|P_z}} \frac{1}{\tilde{f}(n)} = \sum_{n \leq z} \frac{1}{\tilde{f}(2)^s \tilde{f}(p_1)^{e_1} \cdots \tilde{f}(p_m)^{e_m}} \\ &= \sum_{n \leq z} \frac{1}{2^s (p_1/2)^{e_1} \cdots (p_m/2)^{e_m}} \\ &= \sum_{n \leq z} \frac{2^{e_1} \cdots 2^{e_m}}{n} \\ &\geq \sum_{n \leq z} \frac{(e_1 + 1) \cdots (e_m + 1)}{n} \\ &= \sum_{n \leq z} \frac{\tau_1(n)}{n} \end{aligned}$$

# Bounding $V(z)$ Cont'd

Next, we have

$$\sum_{n \leq z} \tau_1(n) = \sum_{n \leq z} \sum_{\substack{d|n \\ (d,2)=1}} 1 = \sum_{\substack{d \leq z \\ (d,2)=1}} \sum_{\substack{n \leq z \\ d|n}} 1 = \sum_{\substack{d \leq z \\ (d,2)=1}} \left[ \frac{z}{d} \right]$$



# Bounding $V(z)$ Cont'd

Next, we have

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# Bounding $V(z)$ Cont'd

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# Bounding $V(z)$ Cont'd

For the summation

$$\sum_{\substack{d \leq z \\ (d,2)=1}} \frac{1}{d}, \quad \text{we choose } c_d = \begin{cases} 1, & \text{if } (d,2)=1 \\ 0, & \text{otherwise} \end{cases} \quad \text{and } f(d) = \frac{1}{d}$$

Then by the partial summation technique, we have

$$\sum_{\substack{d \leq z \\ (d,2)=1}} \frac{1}{d} = \frac{1}{z} \sum_{\substack{d \leq z \\ (d,2)=1}} 1 + \int_1^z \left( \frac{1}{t^2} \sum_{\substack{d \leq t \\ (d,2)=1}} 1 \right) dt$$

# Bounding $V(z)$ Cont'd

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For the summation

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# Bounding $V(z)$ Cont'd

Hence we have that

$$\sum_{n \leq z} \tau_1(n) \geq z \sum_{\substack{d \leq z \\ (d,2)=1}} \frac{1}{d} - z \geq \frac{1}{2} z \log z - \underbrace{(c+1)}_D z$$

Now, for

$$\sum_{n \leq z} \frac{\tau_1(n)}{n}, \quad \text{we choose } c_n = \tau_1(n) \quad \text{and } f(n) = \frac{1}{n}$$

Apply partial summation again, and we get

$$V(z) \geq \sum_{n \leq z} \frac{\tau_1(n)}{n} \geq \frac{1}{4} \log^2(z) + \left( \frac{1}{2} - D \right) \log z - D \gg \log^2(z)$$

# Next Step

- Find  $X$ , estimation of the size of  $A$  ✓
- Estimate  $|A_d|$  for  $d \mid P_z$  to find our multiplicative function,  $f$  ✓
- Find lower bound for  $V(z)$  ✓
- Estimate error term

# Estimate Error Term

First, let us note that

$$\sum_{n \leq x} \tau(n) = \sum_{n \leq x} \sum_{\substack{d \leq x \\ d|n}} 1 = \sum_{d \leq x} \sum_{\substack{n \leq x \\ d|n}} 1 = \sum_{d \leq x} \left\lfloor \frac{x}{d} \right\rfloor \leq x \sum_{d \leq x} \frac{1}{d}$$

Taking  $c_n = 1$  and  $f(t) = \frac{1}{t}$ , we can use partial summation to get that

$$x \sum_{d \leq x} \frac{1}{d} = x \left( \frac{1}{x} \cdot [x] + \int_1^x \frac{[t]}{t^2} dt \right) \leq x(1 + \log x) \ll x \log x$$

Hence,

$$\sum_{n \leq x} \tau(n) \ll x \log x$$



# Estimate Error Term Cont'd

Note that our error term when estimating  $|A_d|$  satisfies

$$R(d) \leq N(d) \leq 2^{\omega(d)}$$

Thus, we have for the error term from Selberg's Sieve,

$$\sum_{\substack{d_1, d_2 \leq z \\ d_1, d_2 | P_z}} R([d_1, d_2]) \leq \sum_{\substack{d_1, d_2 \leq z \\ d_1, d_2 | P_z}} 2^{\omega([d_1, d_2])} \leq \sum_{\substack{d_1, d_2 \leq z \\ d_1, d_2 | P_z}} 2^{\omega(d_1)} 2^{\omega(d_2)}$$

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- Find  $X$ , estimation of the size of  $A$  ✓
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- Find lower bound for  $V(z)$  ✓
- Estimate error term ✓

# Finalé

We shall recall our bound on  $\pi_2(x)$  from before:

$$\pi_2(x) \leq x + S(A, \mathcal{P}, z)$$

As well, from Selberg and all the work we've done, we have

$$S(A, \mathcal{P}, z) \leq \frac{X}{V(z)} + \sum_{\substack{d_1, d_2 \leq z \\ d_1, d_2 | P_z}} |R_{[d_1, d_2]}| \ll \frac{x}{\log^2(z)} + (z \log z)^2$$

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We shall recall our bound on  $\pi_2(x)$  from before:

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And so

$$\pi_2(x) \ll z + \frac{x}{\log^2(z)} + (z \log z)^2$$

Now, if we pick  $z = x^{1/4}$ , we have

$$\pi_2(x) \ll x^{1/4} + 16 \cdot \frac{x}{\log^2(x)} + \frac{1}{16} \sqrt{x} \log^2(x) \ll \frac{x}{\log^2(x)}$$

# The End

*Thank You!*

