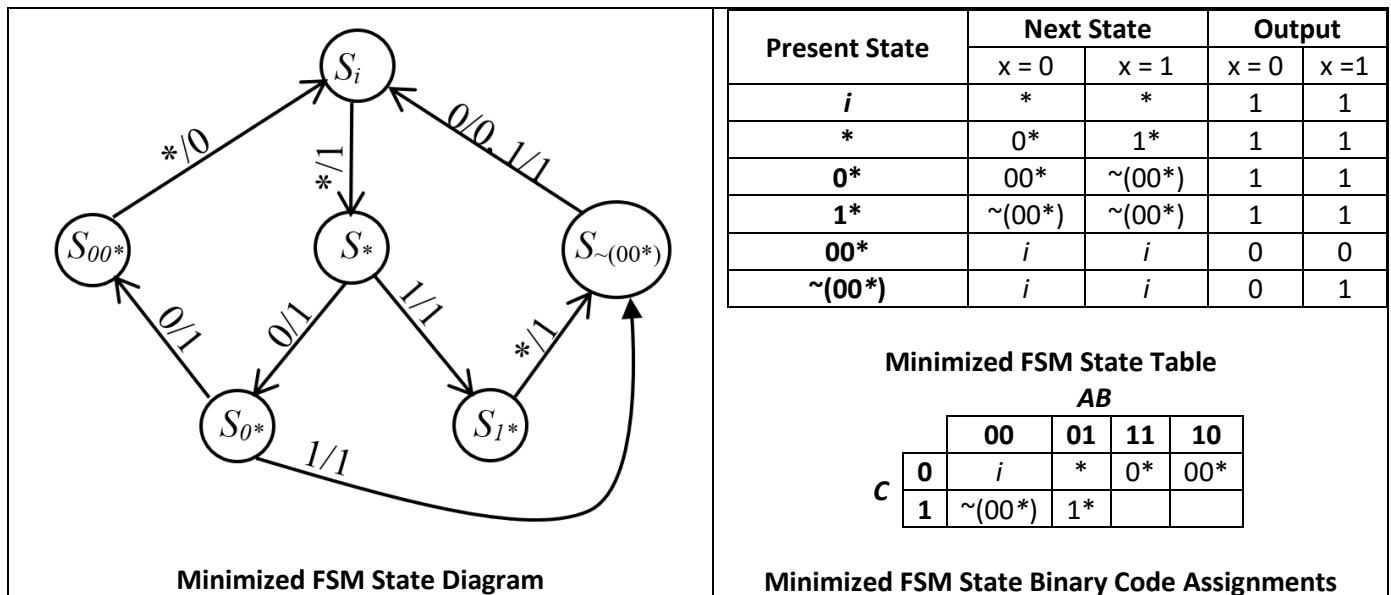


1. From last time
 - a. Created minimized FSM for a BCD code checker
 - i. BCD code checker outputs a 0 if the previous 4-bit sequence was between 0 – 9, and 1 otherwise
 - b. Filled out state table
 - i. Used minimized FSM to fill out new table
 - c. Assigned binary codes to minimize amount of logic
 - i. Try to assign states and their successors adjacent to each other (1 Hamming distance away)
 - ii. Create K-map to help us assign codes



- d. Make binary code table from the above
 - i. Be careful when assigning values!
 - ii. Table states don't line up exactly with table above

Present State	Binary Code	Present State			Input <i>x</i>	Next State			Output <i>z</i>
		<i>A</i>	<i>B</i>	<i>C</i>		<i>A'</i>	<i>B'</i>	<i>C'</i>	
<i>i</i>	000	0	0	0	0	0	1	0	1
<i>i</i>	000	0	0	0	1	0	1	0	1
~(00*)	001	0	0	1	0	0	0	0	0
~(00*)	001	0	0	1	1	0	0	0	1
*	010	0	1	0	0	1	1	0	1
*	010	0	1	0	1	0	1	1	1
1*	011	0	1	1	0	0	0	1	1
1*	011	0	1	1	1	0	0	1	1
00*	100	1	0	0	0	0	0	0	0
00*	100	1	0	0	1	0	0	0	0
	101	1	0	1	0	d	d	d	d
	101	1	0	1	1	d	d	d	d
0*	110	1	1	0	0	1	0	0	1
0*	110	1	1	0	1	0	0	1	1
	111	1	1	1	0	d	d	d	d
	111	1	1	1	1	d	d	d	d

- e. Create K-maps for each flip flop based on input and present state in table above
 - i. Be careful when entering values from the binary code table!
 - ii. Some states are missing and are don't cares, like 101
 - iii. Can insert missing inputs into table and write don't cares there to make filling K-maps easier

A'		AB			
		00	01	11	10
Cx	00	0	1	1	0
	01	0	0	0	0
	11	0	0	d	d
	10	0	0	d	d
		$A' = BCx$			

B'		AB			
		00	01	11	10
Cx	00	1	1	0	0
	01	1	1	0	0
	11	0	0	d	d
	10	0	0	d	d
		$B' = \overline{AC}$			

C'		AB			
		00	01	11	10
Cx	00	0	0	0	0
	01	0	1	1	0
	11	0	1	d	d
	10	0	1	d	d
		$C' = Bx + BC$			

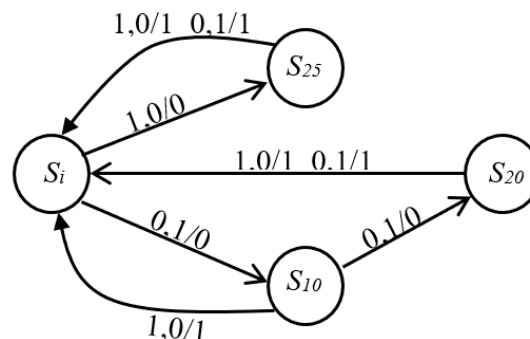
- f. Use derivations from the K-maps to design initial combinational circuit
- g. Create a K-Map based on flip-flops to determine the output combinational circuit
 - i. Since this is a Mealy model, we also use the input
 - ii. Don't cares are in same position as K-maps above

$$z = B + \overline{AC} + \overline{A}x$$

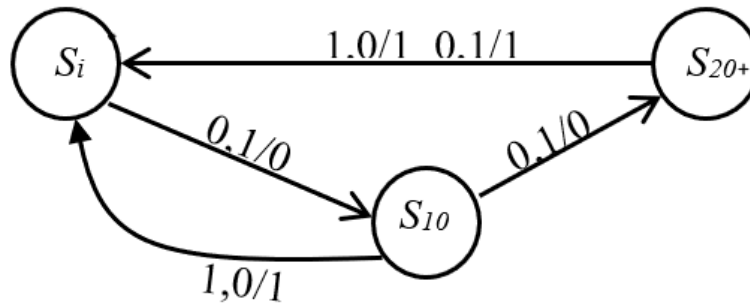
		AB			
		00	01	11	10
Cx	00	1	1	1	0
	01	1	1	1	0
	11	1	1	d	d
	10	0	1	d	d

2. Vending machine example

- a. Design a vending machine that only takes dimes and quarters
 - i. Merchandise is dispensed ($z = 1$) when the sum of the inputs ≥ 30 cents
 - ii. Machine does not give change
- b. $x_1 = \text{quarter}$, $x_2 = \text{dime}$
 - i. Assume that it is not possible to input both quarters and dimes simultaneously
 - ii. Two input, single output
- c. Will use a Mealy model
 - i. Provides for simpler logic in the end
- d. First, create state transition diagram
 - i. Inputs are (x_1, x_2) for (quarter, dime)
 - ii. Input (0, 0) is omitted, would just cause machine to stay in its current state



- e. Next, minimize the number of states using the Partition Minimization Procedure
 - i. $P_1 = (i, 10, 20, 25)$
 - ii. $P_2 = (i) (10) (20, 25)$
 1. 20 and 25 have same k-successors (i for both) so they stay together
- f. Draw new state transition diagram, with new state called 20+



- g. Assign code words next
 - i. $\text{ceil}(\log_2 3) = 2$ flip flops
 - ii. S_i starts in 00
 - iii. No way to place all adjacent states 1 Hamming distance away
 1. Do the best we can, though

		A	
		0	1
B	0	i	20+
	1	10	

- h. Next, create binary code table
 - i. Don't have to create state transition table since this is simple enough
 - ii. Will do the same as previous problem and add empty rows to the table for don't cares

Present State	Binary Code	Present State		Inputs		Next State		Output z
		A	B	x_1	x_2	A'	B'	
i	00	0	0	0	0	0	0	0
i	00	0	0	0	1	0	1	0
i	00	0	0	1	0	1	0	0
		0	0	1	1	d	d	d
10	01	0	1	0	0	0	1	0
10	01	0	1	0	1	1	0	0
10	01	0	1	1	0	0	0	1
		0	1	1	1	d	d	d
20+	10	1	0	0	0	1	0	0
20+	10	1	0	0	1	0	0	1
20+	10	1	0	1	0	0	0	1
		1	0	1	1	d	d	d
		1	1	0	0	d	d	d
		1	1	0	1	d	d	d
		1	1	1	0	d	d	d
		1	1	1	1	d	d	d

- i. Finally, create K-maps from table above
 - i. Be careful when entering values into K-map!
 1. For example, $(A, B, x_1, x_2) = 0011$ is missing since we can't have $(x_1, x_2) = (1, 1)$
 2. Make those inputs don't cares like normal
 3. We added extra rows to the binary code table
 - a. This way, we know exactly where the don't cares go

A'		AB			
		00	01	11	10
x₁x₂	00	0	0	d	1
	01	0	1	d	0
	11	d	d	d	d
	10	1	0	d	0
$A' = A\bar{x}_1\bar{x}_2 + Bx_2 + \bar{A}\bar{B}x_1$					

B'		AB			
		00	01	11	10
x₁x₂	00	0	1	d	0
	01	1	0	d	0
	11	d	d	d	d
	10	0	0	d	0
$B' = B\bar{x}_1\bar{x}_2 + \bar{A}\bar{B}x_2$					

z		AB			
		00	01	11	10
x₁x₂	00	0	0	d	0
	01	0	0	d	1
	11	d	d	d	d
	10	0	1	d	1
$z = Ax_2 + Bx_1 + Ax_1$					

3. Debugging an FSM
 - a. One good way of seeing if your FSM you drew was right is to give a stream of inputs into your FSM
 - i. For example, with the vending machine, give a dime, then a dime, then a quarter
 - ii. See what states you land at, and if those are what you expect
 - iii. Also see what outputs you get, and if those are what you expect as well