## Knowledge Representation

Session in the course "Programmation Logique et Connaissances" at the École nationale supérieure des Télécommunications in Paris/France in spring 2011

by Fabian M. Suchanek

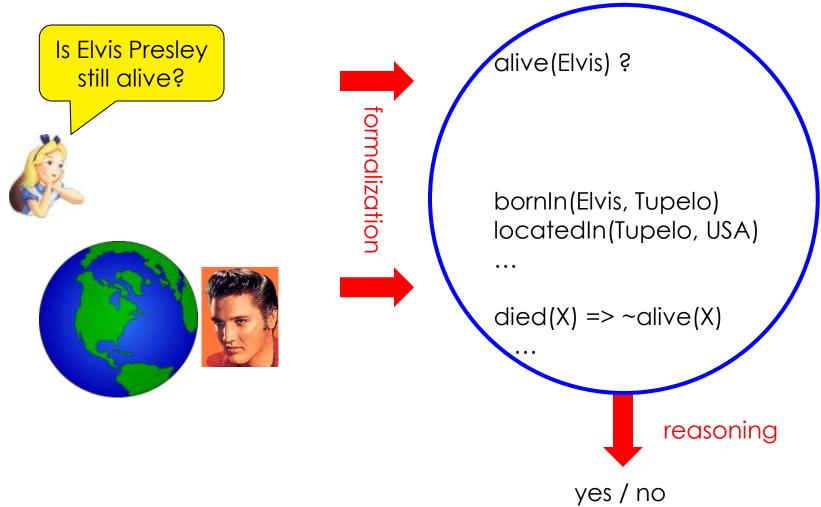


#### Organisation

- 3h class on Knowledge Representation
   15min break
- Slides are available at the Web-sites
   <u>http://suchanek.name/</u> → Teaching
   <u>http://icc.enst.fr/PLC</u> → Session of 2011-01-10
- Language of the Slides is English,
   lecturer will do his best to talk in French,
   questions are welcome throughout the session in French or English

#### Motivation

**Knowledge representation** (KR) is a research field of artificial intelligence that aims to formalize information so that it can be used for automated reasoning.

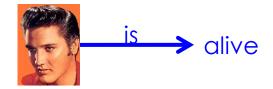


#### Motivation

**Knowledge representation** (KR) is a research field of artificial intelligence that aims to formalize information so that it can be used for automated reasoning.

There are different ways to represent knowledge:

- PROLOG-like: alive(Elvis), is(Elvis, alive)
- graphically



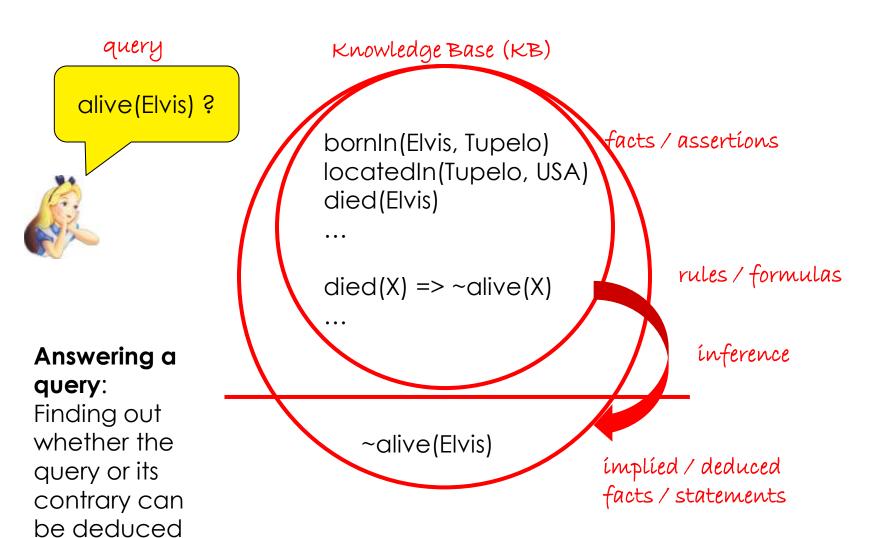
- in natural language
   "Flvis is alive"
- in propositional logic elvis\_alive .

- in first order logic
   \( \neg \) x: rocksinger(x) => alive(x)
- in a mathematical notation elvis ∈ Alive
- in a programming language elvis.alive=true
- in completely different formalism
   ⇔ elvis → ⊕ ⊕ ⊕

#### Overview

- Motivation
- Knowledge Representation Design
  - Knowledge Representation in the Semantic Web (RDF, RDFS, OWL, DL)

# KR Design: Big picture



from the KB

KR formalism: The language of the KB

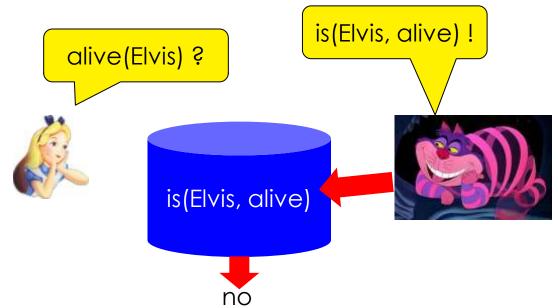
# KR Design: Canonicity

A KR formalism is **canonic** if one piece of knowledge can only be represented in one way

```
alive(Elvis)
is(Elvis, alive)
alive(Elvis, true)
vivant(Elvis)

not very canonic
```

Canonicity is in general desirable because it facilitates co-operation by different people



# KR Design: Canonicity

A KR formalism is **canonic** if one piece of knowledge can only be represented in one way.

A formalism can be made more canonic by

restricting it

```
e.g., by allowing only unary predicates: alive(Elvis)
is(Elvis, alive)
```

providing best practice guidelines

```
e.g., by prescribing certain conventions: alive(Elvis)

alice(elvis)
```

providing standard vocabularies

```
e.g., by listing predicates that should be used in a certain domain: {alive, dead, young, old, happy, sad}
```

## KR Design: Expressiveness

A KR formalism is **more expressive** than another one if we can say things in the first formalism that we cannot say in the second.

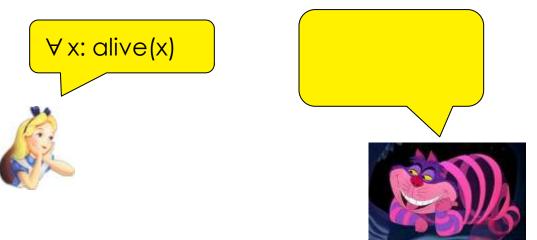
First order logic

Yx: rocksinger(x) => alive(x)

Propositional logic

?

In general, a higher expressiveness is desirable from a modeling point of view



... but it comes at a cost...

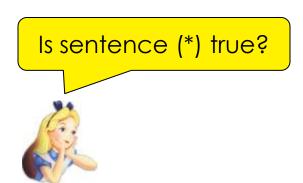
# KR Design: Decidability

A KR formalism is **decidable**, if there is an algorithm (computer program) that can answer any query on a knowledge base in that formalism.

(Technically: A logical system is decidable iff there is an effective method for determining whether arbitrary formulas are theorems of the logical system.)

Some formalisms are so expressive that they are undecidable:

Natural language



(\*) This sentence is false.

First order logic

First order logic is so powerful that it can express sentences of which it is impossible to determine whether they are true or not.

# KR Design: Decidability

A KR formalism is **decidable**, if there is an algorithm (computer program) that can answer any query on a knowledge base in that formalism.

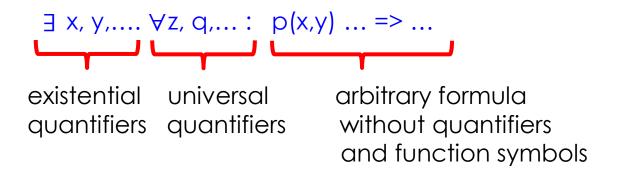
(Technically: A logical system is decidable iff there is an effective method for determining whether arbitrary formulas are theorems of the logical system.)

In general, decidability is desirable.

The more expressive a formalism is, the more likely it is to be undecidable.

Often, a formalism can be made decidable by restricting it

- Propositional logic is decidable
- First order logic is decidable if all formulae are of the following form:



#### KR Design: Closed world

A KR formalism follows the **closed world assumption (CWA)**, if any statement that cannot be proven is assumed to be false

```
PROLOG, e.g., follows the CWA:

?- assert(bornIn(Elvis, Tupelo)).

yes

?- alive(Elvis).

no
```

In many contexts, the open world assumption (OWA) is more appropriate. Under the OWA, a statement can be

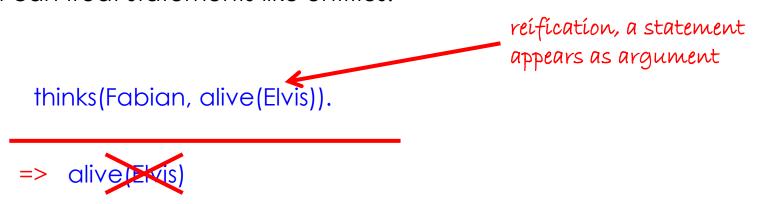
- provably false
- provable true
- unknown

```
?- alive(Elvis).

I have no clue
```

## KR Design: Reification

A KR formalism allows **reification**, if it can treat statements like entities.



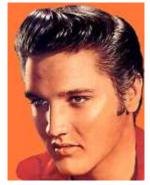
Modal logic allows talking about necessity and possibility

"It is possible that Elvis is alive"

♦ alive(Elvis)

#### KR Design: Unique Name

A KR formalism follows the **unique name assumption (UNA)**, if different names always refer to different objects.



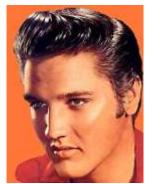




The King

The UNA is not useful if different people want to use different identifiers:







- ?- assert(alive(TheKing)).
   yes
- ?- assert (The King=Elvis). hmm, OK
- ?- alive(Elvis). yes

A KR formalism is **schema-bound**, if one has to decide upfront which entities can have which properties.

PROLOG is schema-free, any entity can have any property:

- ?- assert(alive(Elvis)).
  yes
- ?- assert(hasResolution(Elvis,5 megapixel))
   ves



In schema-bound formalisms, one has to decide a priori for classes of things and their properties:

#### person:

- alive/dead
- birthday
- profession

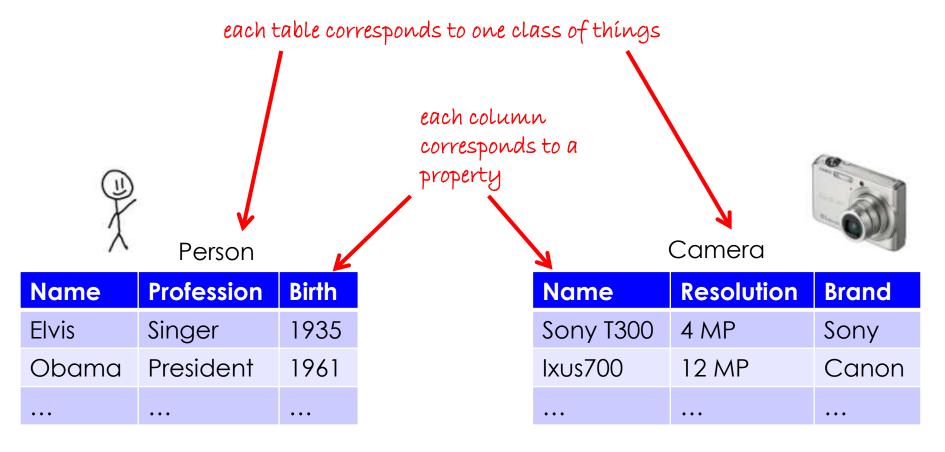
#### camera:

- resolution
- shutter speed
- weight

A schema-bound formalism puts more modeling constraints, but can exclude non-sensible statements.

**Databases** are a particular schema-bound KR formalism.

A database can be seen as a set of tables.



each row corresponds to a thing

**Databases** are a particular schema-bound KR formalism.

A database can be queried in the Structured Query Language (SQL).

SELECT name, birth FROM person WHERE profession='singer' AND birth>1930

		Person	
Name		fession	Birth
Elvis	h	ger	1935
Obama	Ξ	esident	1961
		_	•••

#### Camera

Name	Resolution	Brand
Sony T300	4 MP	Sony
lxus700	12 MP	Canon
•••	•••	•••

Elvis, 1935 JohnLennon, 1940

17

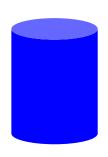
**Databases** are a particular schema-bound KR formalism.

A database can be queried in the Structured Query Language (SQL).

Databases are used in practically every major enterprise, with the main database systems being

- Oracle
- Microsoft SQL Server
- IMB's DB2
- Postgres
- MySQL

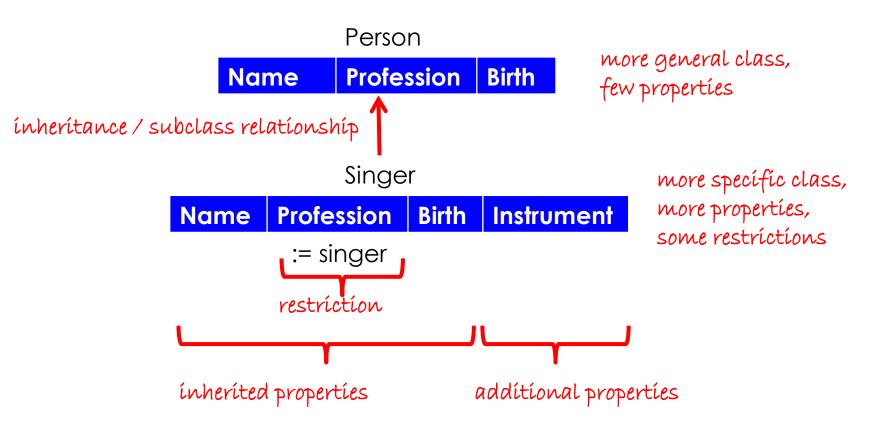
Name	Profession	Birth
Elvis	Singer	1935
Obama	President	1961
•••		•••



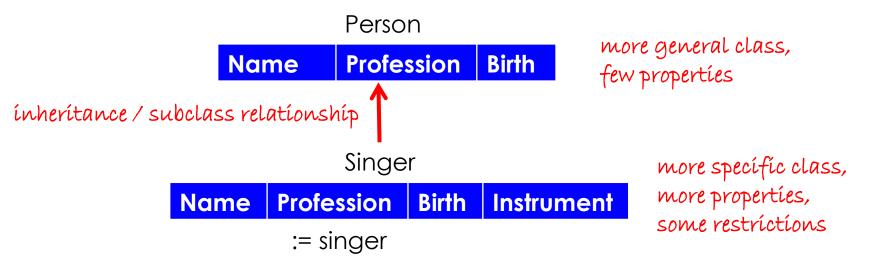


Headquarters of Oracle in Redwood Shores, CA, USA

A KR formalism supports **inheritance**, if properties specified for one class of things can be automatically transferred to a more specific class. A **class** is a set of entities with the same properties.



A KR formalism supports **inheritance**, if properties specified for one class of things can be automatically transferred to a more specific class. A **class** is a set of entities with the same properties.



Inheritance is useful, because

- it avoids duplication of work (no need to reinvent the wheel)
- it makes the KR formalism more canonic
   (the subclass automatically has the same properties as the superclass)

**Object-oriented programming languages** (such as e.g., Java) support inheritance.

```
super-class
public class Person {
                                       * declares two properties
     String name;
     String profession;
                                               sub-class
public class Singer extends Person {
                                               * overwrites a property
     String profession="Singer";
                                               * adds a property
     String instrument;
}
                                         Elvis is a singer
Singer elvis=new Singer();
                                         * but he inherited properties
elvis.name="Elvis";
elvis.instrument="guitar";
System.out.println(elvis.profession) * and has predefined values
     → Singer
```

**Object-oriented programming languages** (such as e.g., Java) support inheritance.

Most programming languages support object-orientation today, with the most important ones being

- Java
- C++
- C#
- Visual Basic
- Python
- Ruby



# KR Design: Monotonicity

A KR formalism is **monotonous**, if adding new knowledge does not undo deduced facts.

First order logic and propositional logic are monotonous:

```
elvis_is_person
elvis_is_alive
elvis_is_dead => ~ elvis_is_alive
elvis_is_dead => ~ elvis_is_dead

=> elvis_is_alive
elvis_is_dead => ~ elvis_is_alive
+ elvis_is_dead

=> elvis_is_dead
-> elvis_is_alive
elvis_is_dead
-> elvis_is_dead
-> elvis_is_alive
-> elvis_is_alive
-> elvis_is_dead
-> michael_jackson_alive
```

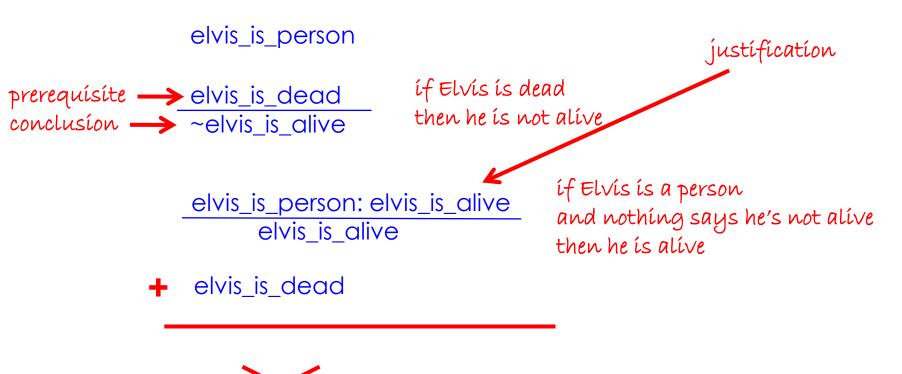
Monotonicity can be very counter-intuitive.

It requires everything to be known upfront.

# KR Design: Monotonicity

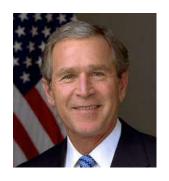
A KR formalism is **monotonous**, if adding new knowledge does not undo deduced facts.

**Default logic** is not monotonous:



# KR Design: Fuzziness

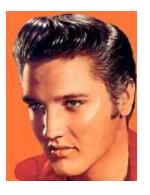
A KR formalism is **fuzzy**, if certain statements can hold to a certain degree. The opposite of fuzzy is **crisp**.



fantastic (Bush) 0.1



fantastic (Madonna) 0.8



fantastic (Elvis)
1.0

First order logic, PROLOG and propositional logic are all crisp.

#### KR Design: Fuzziness

A KR formalism is **fuzzy**, if certain statements can hold to a certain degree. **Fuzzy logic** is a fuzzy KR formalism.

```
rainy (0.8)

rainy => bad_weather

bad_weather (0.8)

rainy (0.8)

windy(1.0)

rainy \/ windy => bad_weather

bad_weather (??)

1.0
```

Fuzzy logic defines how to compute fuzzy values for complex combinations of fuzzy predicates, e.g.

- OR is computed as maximum
- AND is computed as minimum
- NOT is computed as 1-x

## KR Design: Contradictions

A KR formalism is **tolerant to contradictions** if it allows contradictions in the knowledge base.

First order logic and propositional logic are not tolerant to contradictions:

```
alive(Elvis). ~alive(Elvis).
```

=> life\_is\_beautiful

ex falso quod libet...

Some domains require handling of contradictions (e.g. information extraction from the Web)

Elvisfans.com: Elvis is alive Wikipedia: Elvis is dead.

## KR Design: Contradictions

A Markov Logic Network is a set of propositional logic formulas with weights:

```
f1: elvis_alive [0.8]
f2: elvis_dead => ~elvis_alive [0.5]
f3: elvis_dead [0.6]
```

A **possible world** is an assignment of truth values to the predicates:

```
w1: w2: w3: w4:
elvis_alive ✓ elvis_alive × elvis_alive × elvis_dead ✓ elvis_dead ✓

0.3

0.1

0.5

0.9
```

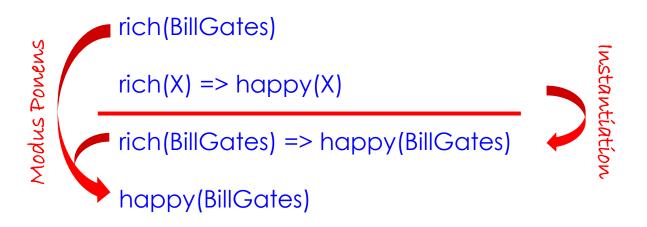
A Markov Logic Network assigns a probability to each possible world:

```
P(w) \sim P(w4) \sim e^{0.8} \cdot e^{0.6} = 4.0
satisfied formula f f1 f2 f3
(normalized by the sum of the values for all worlds)
```

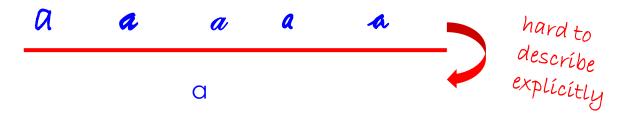
Finding the most probable world boils down to the **Weighted Maximum Satisfiability Problem**.

A KR formalism is **explicit**, if it can explain why a certain statement was deduced.

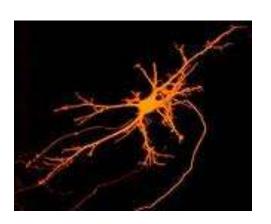
Logic-based formalisms are usually explicit:



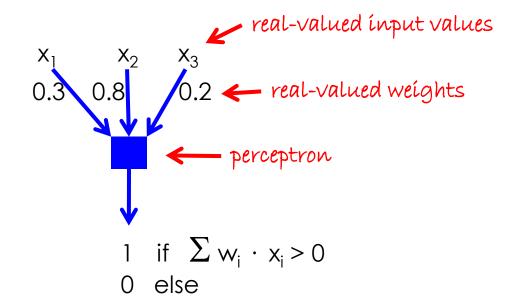
Explicitness may not always be useful, consider pattern recognition:



A perceptron is a simplified model of a human nerve cell (neuron)







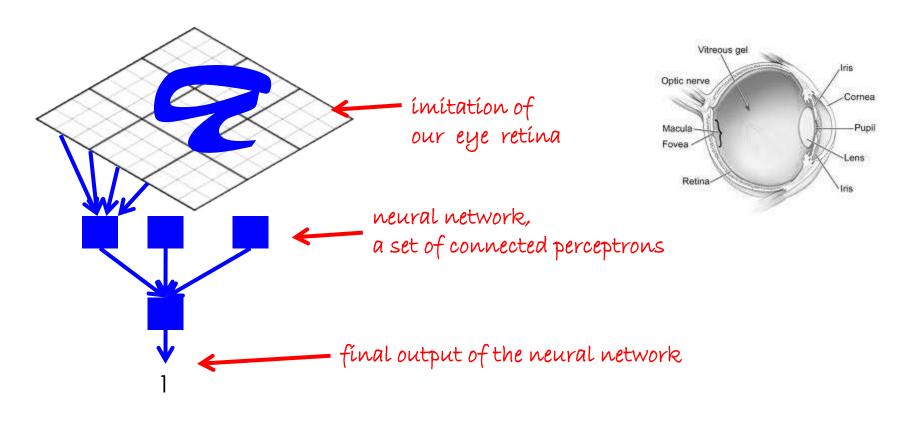
#### **Examples:**

input: 
$$x_1 = 1$$
,  $x_2 = 2$ ,  $x_3 = 1$   

$$\sum w_i \cdot x_i = 0.3 \cdot 1 + 0.8 \cdot 2 + 0.2 \cdot 1 = 2.1 > 0$$

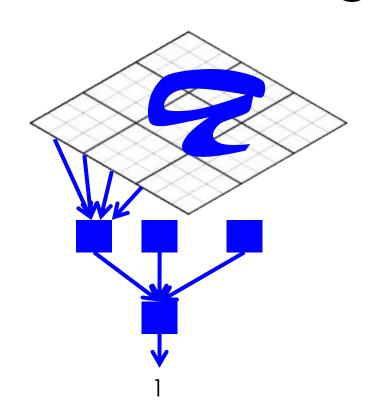
=> output: 1

input: 
$$x_1 = 0$$
,  $x_2 = -2$ ,  $x_3 = 0$ 



**Training a neural network**: Adjusting the weights so that the network says 1 for all positive examples.





A Neural Network is a non-explicit knowledge representation.

Neural Networks are used, e.g,

- to recognize postal codes on letters
- to optimize search engine results

This network "knows" what the letter "a" is.

But this knowledge is not explicit.

# KR Design: Distributedness

A KR formalism is **distributed**, if it encourages use and co-operation by different people / systems across different places / organizations.

Every rock singer is alive!

The King is a rock singer.



I am Elvis, the King.



Elvis is alive

~~~~> We will see a very popular distributed formalism shortly

# KR Design: Summary

In general, a KR formalism serves to

- represent knowledge
- infer facts from that knowledge
- answer queries on the knowledge

There are many KR formalisms with different properties:

- Canonicity (does the formalism allow only one form to represent a statement?)
- Expressiveness (how powerful is the formalism)
- Decidability (can every query be answered?)
- Closed World (does the formalism assume that everything unknown is false?)
- Unique Name (do two names always mean two different things?)
- Schema-bound (do we have to decide upfront on properties of things?)
   Databases/SQL is a schema-bound KR formalism
- Inheritance (can properties be transferred from one class to another?)

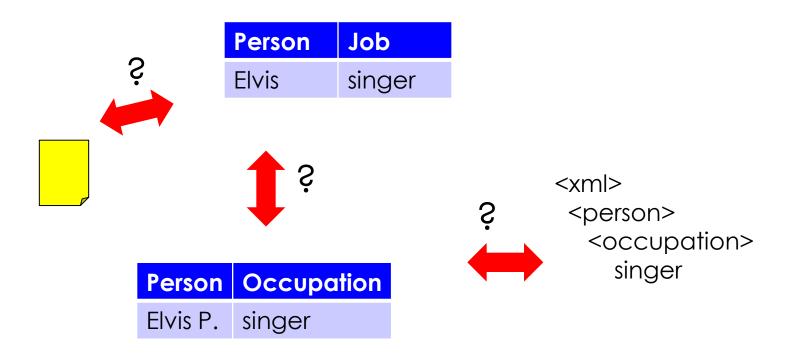
  Object-oriented programming languages support inheritance
- Monotonicity (will new knowledge never undo existing knowledge?)
   Default logic allows non-monotonicity
- Fuzziness (can properties be fufilled to a certain degree?)
   Fuzzy logic is a fuzzy KR formalism
- Tolerance to contradictions (is it OK to have a contradiction in the KB?)
   Markov Logic Networks can deal with contradictions
- Explicitness (is knowledge always explainable?)
   Neural Networks store knowledge only implicitly

#### Overview

- Motivation
- Knowledge Representation Design
- Knowledge Representation in the Semantic Web

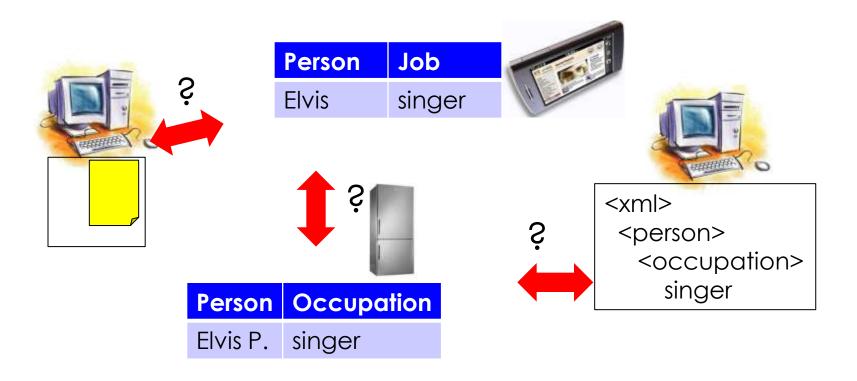
#### Motivation

Interaction between data on the Web is difficult, in particular if the data is in different formats



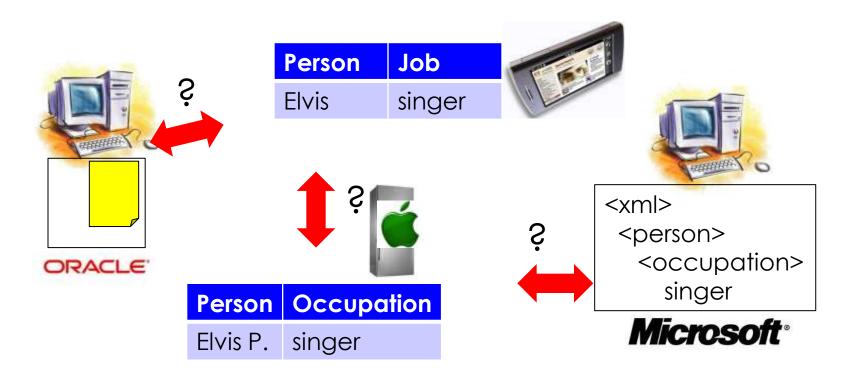
#### Motivation

Interaction between data on the Web is difficult, in particular if the data is in different formats, on different machines or devices



#### Motivation

Interaction between data on the Web is difficult, in particular if the data is in different formats, on different machines or devices or in different companies



PYCOPHILIBE

### Motivation: Use cases

#### Examples:

- Booking a flight
   Interaction between office computer, flight company, travel agency, shuttle services, hotel, my calendar
- Finding a restaurant Interaction between mobile device, map service, recommendation service, restaurant reservation service
- Web search Interaction between client, search service, Web page content provider
- Web service composition
   Interaction between client and Web services and Web services themselves
- Intelligent home
   Fridge knows my calendar, orders food if I am planning a dinner
- Intelligent cars
   Car knows my schedule, where and when to get gas, how not to hit other cars, what are the legal regulations

# Motivation: Merging

#### Examples:

- Adding data to a database
   From XML files, from other databases
- Merging data after company mergers (e.g. Apple buys Microsoft)
   Different terminology has to be bridged, accounts to be merged
- Merging data in research
   e.g. biochemical, genetic, pharmaceutical research data

(Less exciting, but probably more frequent)

#### Motivation: Semantic Web

We need a Knowledge Representation that

- allows machines to process data from other machines
- ensures interoperability between different schemas, devices and organizations
- allows data to describe data
- allows machines to reason on the data
- allows machines to answer semantic queries

This is what the Semantic Web aims at

The **Semantic Web** is an evolving extension of the World Wide Web, which promotes a new distributed knowledge representation.

#### Semantic Web

The Semantic Web provides KR standards to

- Identify entities (URIs)
- Express facts (RDF)
- Express concepts (RDFS)
- Share vocabularies
- Reason on facts (OWL)

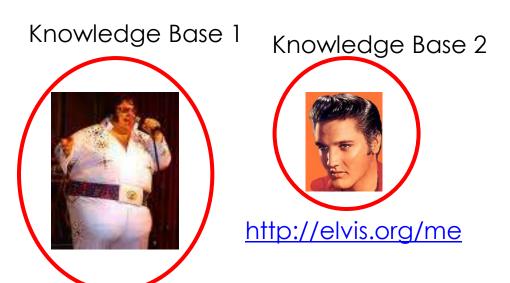
... and it even works.

These standards are produced and endorsed by the Word Wide Web Consortium (W3C)

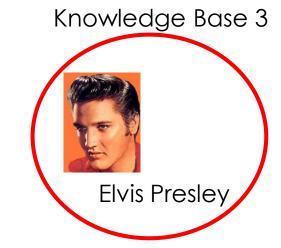
The **Semantic Web** is an evolving extension of the World Wide Web, which promotes a new distributed knowledge representation.

#### **URIs**

A **Uniform Resource Identifier** (URI) is a string of characters used to identify an entity on the Internet







http://onto.com/people/singers/EP

Entity: anything A person, a city, a relationship ("loves"), a class of things ("singer")...

### **URIs**

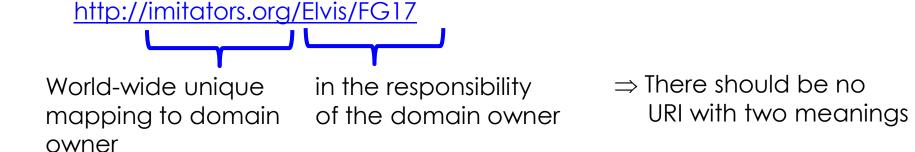
A **Uniform Resource Identifier** (URI) is a string of characters used to identify an entity on the Internet

A URI is a generalization of a URL (the stuff you type in the address bar in a browser)

|                                                   | Age                              |
|---------------------------------------------------|----------------------------------|
| http://elvis.org/me Identifies the not Internet-o | ·                                |
|                                                   | tifies a file,<br>net-accessible |

#### **URIs**

A **Uniform Resource Identifier** (URI) is a string of characters used to identify an entity on the Internet



- ⇒ People can invent all kinds of URIs
- a company can create URIs to identify its products
- an organization can assign sub-domains and each sub-domain can define URIs
- individual people can create URIs from their homepage
- people can create URIs from any URL for which they have exclusive rights to create URIs

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... and it even works.

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The **Semantic Web** is an evolving extension of the World Wide Web, which promotes a new distributed knowledge representation.

### RDF: Statements

The **Resource Description Framework** (RDF) is a distributed KR formalism.

Assume we have the following URIs:

A URI for Elvis: <a href="http://elvis.org/himself">http://elvis.org/himself</a>

A URI for "winning a prize": <a href="http://inria.fr/rdf/dta#won">http://inria.fr/rdf/dta#won</a>

A URI for the Grammy award: <a href="http://g-a.com/prize">http://g-a.com/prize</a>

An **RDF statement** is a triple of 3 URIs: The subject, the predicate and the object.

http://elvis.org/himself http://inria.fr/rdf/dta#won http://g-a.com/prize

We can understand an RDF statement as a First Order Logic statement with a binary predicate

won(Elvis, Grammy award)

## RDF: Namespaces

A **namespace** is an abbreviation for the prefix of a URI.

@prefix elvis: http://elvis.org/

@prefix inria: http://inria.fr/rdf/dta# @prefix grammy: <a href="http://g-a.com/">http://g-a.com/</a>

An **RDF statement** is a triple of 3 URIs: The subject, the predicate and the object.

http://elvis.org/himself http://inria.fr/rdf/dta#won

http://g-a.com/prize

... with the above namespaces, this becomes...

elvis:himself

inria:won

grammy:prize

The **default name space** is indicated by ":"

:himself

inria:won

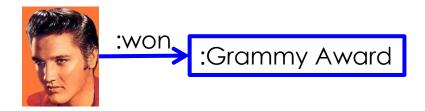
grammy:prize

## RDF: Graphs

A set of triples is **isomorphic to a labeled directed multi-graph**:

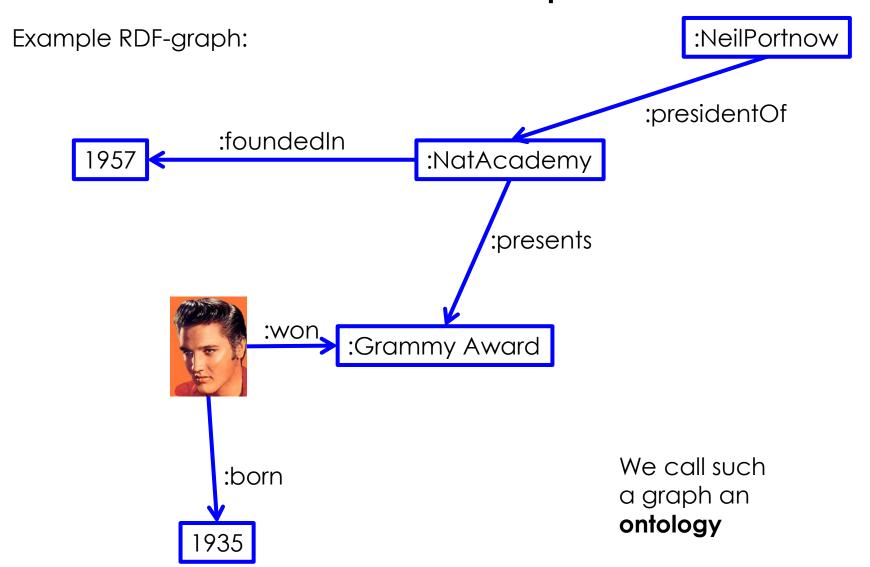
The subject and object of a triple correspond to nodes, the predicate corresponds to directed edge from subject to object with a label given by the predicate.

:elvis :won :GrammyAward



In the following, we will use this notation, assuming some default name space.

## RDF: Graphs



#### RDF: Event entities

The **Resource Description Framework** (RDF) is a KR formalism that allows only binary predicates.

All tabular data and n-ary relations can be expressed in RDF by event entities.

| Person  | Prize            | Year     |
|---------|------------------|----------|
| Elvis   | Grammy<br>Award  | 1967     |
| won(Elv | vis, GrammyAward | d, 1967) |
|         |                  |          |

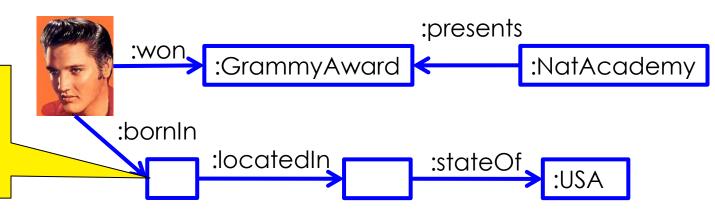
**Event entities** are artificial entities (nodes) that represent a complex constellation

#### RDF: Blank nodes

The **Resource Description Framework** (RDF) is a KR formalism that allows only binary predicates.

A **blank node** is an RDF node that has no name.

A blank node can be unified with another node, so that a blank node does not violate uniqueness constraints.



#### **Semantics:**

A triple <s,p,o> is interpreted as a First Order Logic fact p(s,o).

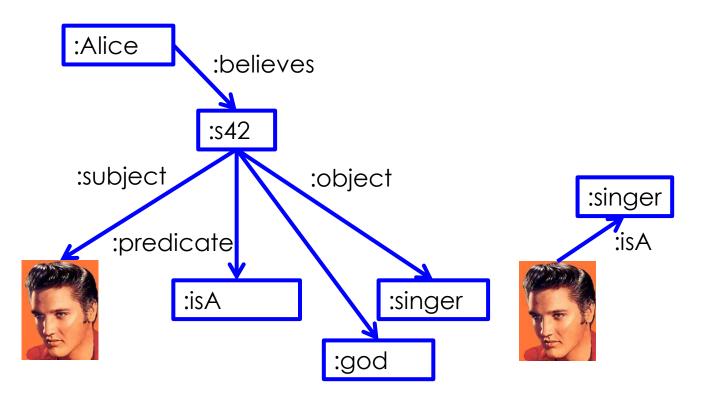
A blank node is interpreted as an existential variable.

won(Elvis, GrammyAward)
presents(NatAcademy, GrammyAward)

 $\exists x,y: bornIn(Elvis, x) / locatedIn(x, y) / stateOf(y,USA)$ 

#### RDF: Reification

A **reified statement** is an entity that represents a statement.



The reified statement can be part of the ontology or not,

i.e., it does not have to be true.

### Semantic Web

The Semantic Web provides KR standards to

- Identify entities (URIs)
- Express facts (RDF)
- Express concepts (RDFS)
- Share vocabularies
- Reason on facts (OWL)

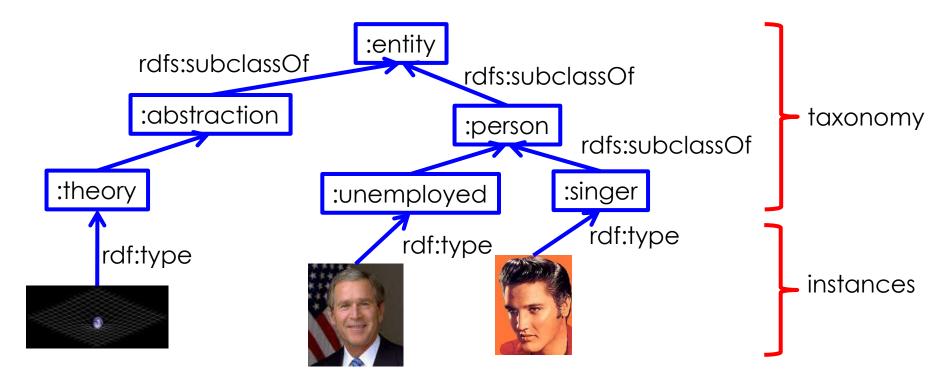
... and it even works.

These standards are produced and endorsed by the **Word Wide Web Consortium** (W3C)

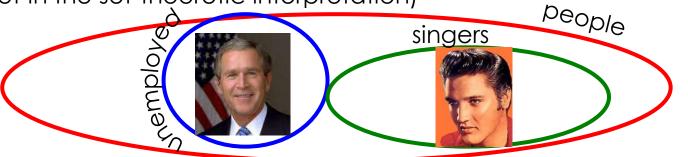
The **Semantic Web** is an evolving extension of the World Wide Web, which promotes a new distributed knowledge representation.

## RDFS: Classes

A class (also called concept) can be understood as a set of similar entities.



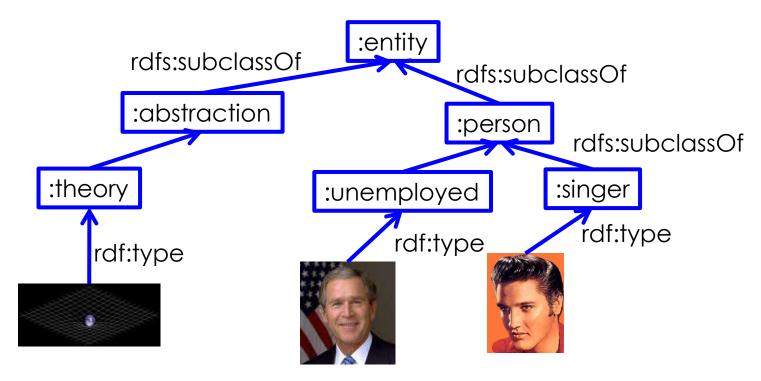
A **super-class** of a class is a class that is more general than the first class (a super-set in the set-theoretic interpretation)



55

#### RDFS: Classes

A **class** (also called concept) can be understood as a set of similar entities.



The fact that an entity belongs to a class is expressed by the **type** predicate from the standard namespace rdf (<a href="http://w3c.org/">http://w3c.org/</a>...).

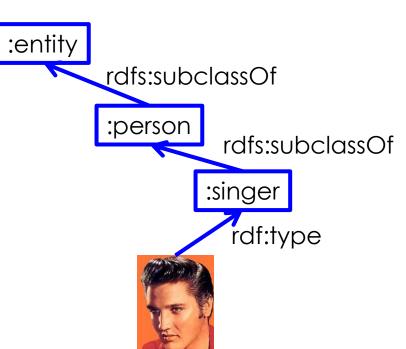
The fact that a class is a sub-class of another class is expressed by the **subclassOf** predicate from the standard namespace rdfs (<a href="http://w3c.org/">http://w3c.org/</a>...).

RDFS knows inferences (called **entailment** in RDFS).

Entailment rule:

If the graph contains such and such triples

then add this triple



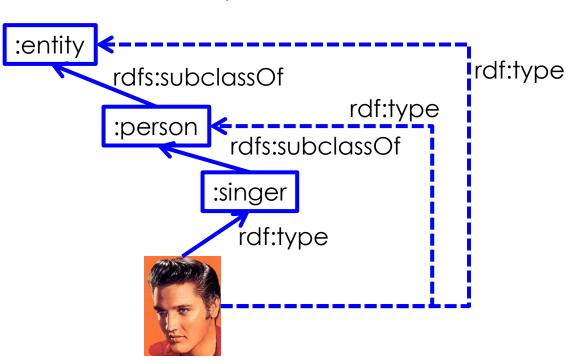
RDFS knows inferences (called **entailment** in RDFS).

Type entailment rule:

X rdf:type Y

Y rdfs:subclassOf Z

X rdf:type Z



In first order logic:

 $\forall$  x, y, z: type(x,y) /\ subclassOf(y,z) => type(x,z)

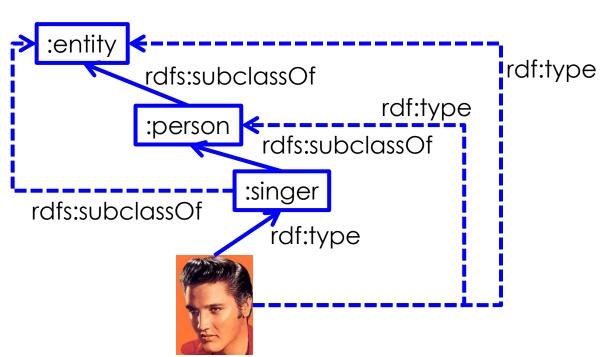
RDFS knows inferences (called **entailment** in RDFS).

Subclass entailment rule:

X rdfs:subclassOf Y

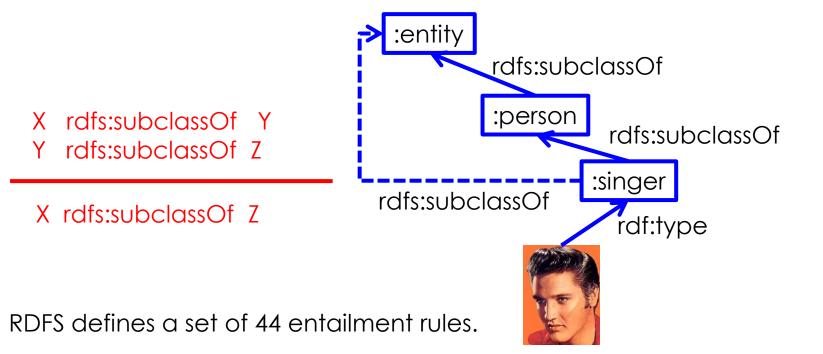
Y rdfs:subclassOf Z

X rdfs:subclassOf Z



In first order logic:

 $\forall$  x, y, z: subclassOf(x,y) /\ subclassOf(y,z) => subclassOf(x,z)



The entailment rules are applied recursively until the graph does not change any more.

This can be done in polynomial time. Whether this is done physically or deduced at query time is an implementation issue.

### Semantic Web

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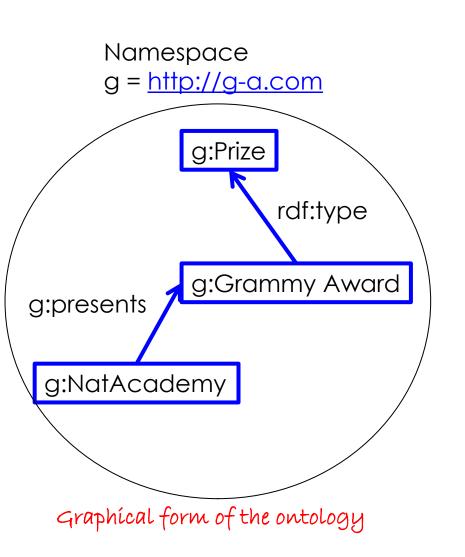
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# Storing data

RDF data is usually stored on a server (=internet accessible computer)





The server at <a href="http://g-a.com">http://g-a.com</a> stores:

@prefix g: http://g-a.com

@prefix rdf: http://www.w3.org/...

g:GrammyAward rdf:type g:Award

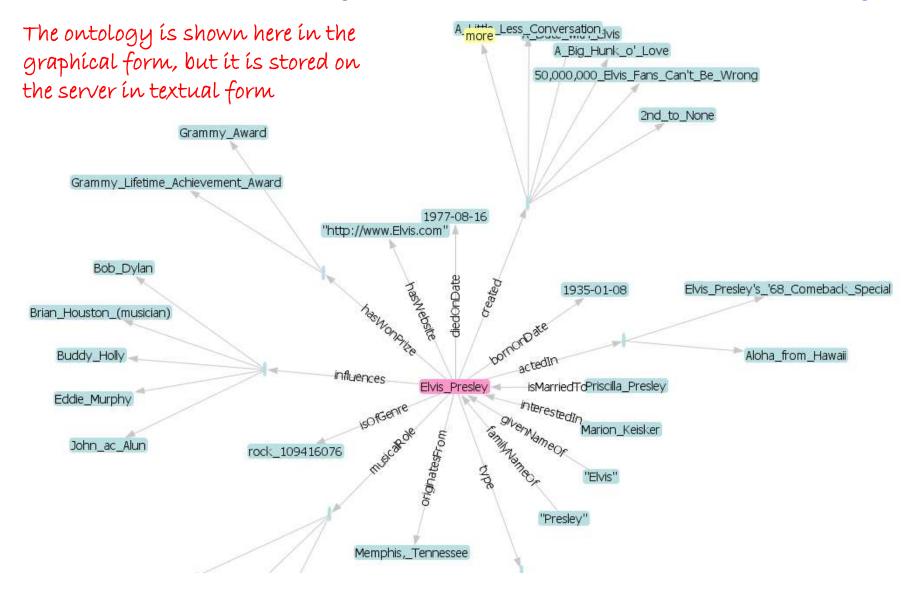
g:NatAcademy g:presents g:GrammyAward

Textual form of the ontology

Try this

# Storing data

Example: The YAGO ontology is stored at the server at <a href="http://mpii.de/yago">http://mpii.de/yago</a>



# Sharing: Cool URIs

A URI is not necessarily **dereferenceable** (i.e., it cannot be accessed online)

http://g-a.com/GrammyAward => HTTP not found error

... but it can be dereferenceable. This means that if I access the URL, the server responds with an RDF snippet:

@prefix g: http://g-a.com

@prefix rdf: http://www.w3.org/1999/02/22-rdf-syntax-ns#

g:GrammyAward rdf:type g:Award

http://elvis.com/elvis g:won g:GrammyAward

Try this out:

rdf:type = <a href="http://www.w3.org/1999/02/22-rdf-syntax-ns#type">http://www.w3.org/1999/02/22-rdf-syntax-ns#type</a>

# Sharing: Cool URIs

#### Server at <a href="http://elvis.com">http://elvis.com</a>

```
@prefix e: <a href="http://elvis.com">http://elvis.com</a>
```

@prefix rdf: http://www.w3.org/1999/02/22-rdf-syntax-ns#

```
e:elvis rdf:type e:singer
```

e:elvis e:born 1935

#### Server at <a href="http://g-a.com">http://g-a.com</a>

```
@prefix g: http://g-a.com
```

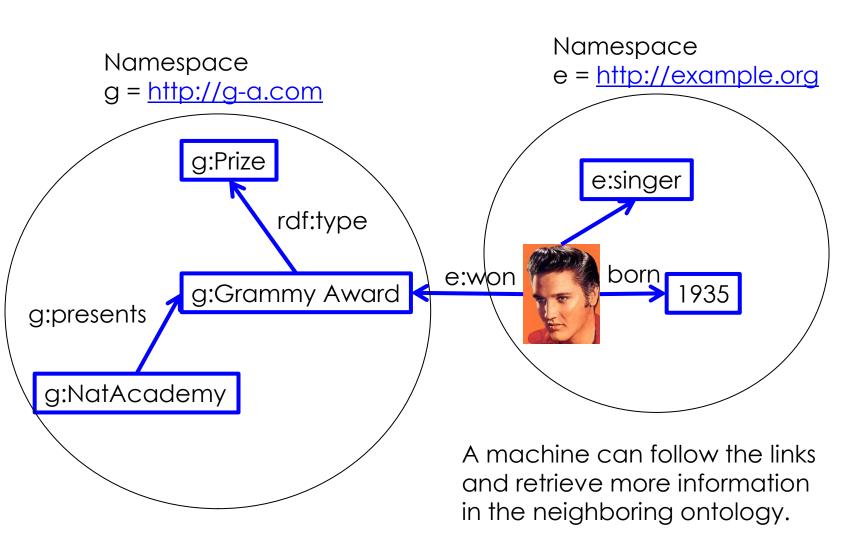
@prefix rdf: http://www.w3.org/1999/02/22-rdf-syntax-ns#

```
g:GrammyAward rdf:type g:Award
```

http://elvis.com/elvis g:won g:GrammyAward

# Sharing

If two RDF graphs share one node, they are actually one RDF graph.



## Sharing: Standard Vocabulary

A number of standard vocabularies have evolved

rdf: The basic RDF vocabulary

http://www.w3.org/1999/02/22-rdf-syntax-ns#

rdfs: RDF Schema vocabulary

http://www.w3.org/1999/02/22-rdf-syntax-ns#

dc: Dublin Core (predicates for describing documents)

http://purl.org/dc/elements/1.1/

foaf: Friend Of A Friend (predicates for relationships between people)

http://xmlns.com/foaf/0.1/

cc: Creative Commons (types of licences)

http://creativecommons.org/ns#

## Sharing: Standard Vocabulary

A number of standard vocabularies have evolved

rdf: The basic RDF vocabulary

http://www.w3.org/1999/02/22-rdf-syntax-ns#

rdfs: RDF Schema vocabulary

http://www.w3.org/1999/02/22-rdf-syntax-ns#

Standard vocabulary provided by the W3C:

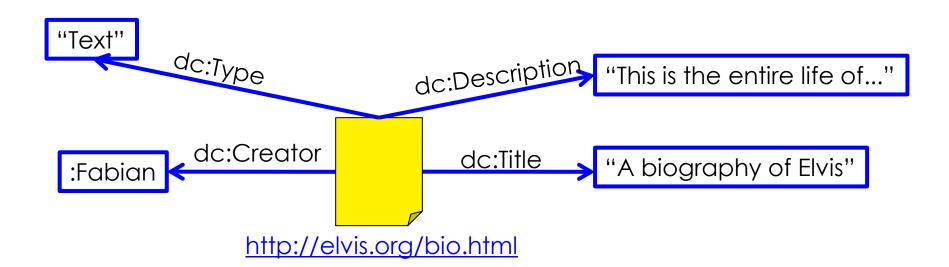
- rdf:type
- rdfs:subclassOf
- rdf:List
- around a dozen others

## Sharing: Dublin Core

A number of standard vocabularies have evolved

dc: Dublin Core (predicates for describing documents)

http://purl.org/dc/elements/1.1/

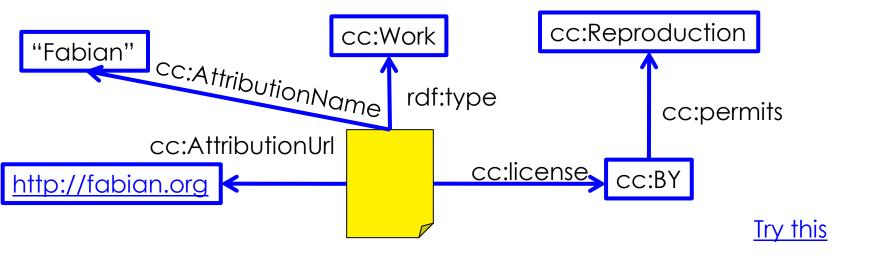


# Sharing: Creative Commons

A number of standard vocabularies have evolved

cc: Creative Commons (types of licences)

http://creativecommons.org/ns#



**Creative Commons** is a non-profit organization, which defines very popular licenses, notably

- CC-BY: Free for reuse, just give credit to the author
- CC-BY-NC: Free for reuse, give credit, non-commercial use only
- CC-BY-ND: Free for reuse, give credit, do not create derivative works

#### What we have seen so far

RDF is a knowledge representation formalism

that has a graphical form and a textual form



that uses URIs to identify entities (and abbreviates them with name spaces)
 @prefix yago: <a href="http://mpii.de/yago#">http://mpii.de/yago#</a>

yago:elvis = <a href="http://mpii.de/yago#elvis">http://mpii.de/yago#elvis</a>

i.e., two ontologies can talk about exactly the same entity or relation

that is stored on servers on the Internet and can be accessed by a machine



### Semantic Web

The Semantic Web provides KR standards to

- Identify entities (URIs)
- Express facts (RDF)
- Express concepts (RDFS)
- Share vocabularies

  ✓
- Reason on facts (OWL)

... and it even works.

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### **OWL:** Goal

RDFS just allows us to define classes and subclasses with very limited inference.

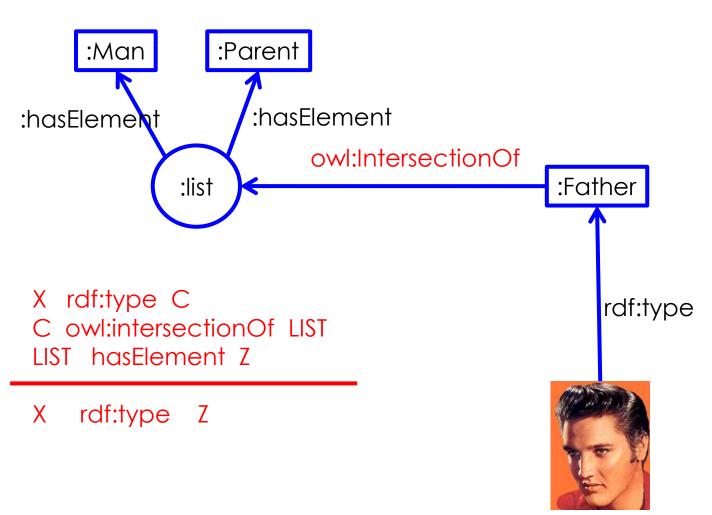
Can we go further?

- Reasoning
   If X is left of Y and Y is left of Z
   then X is left of Z
- Class definitions
   The class of husbands is
   the class of married men
- Class properties
   People and tables are two disjoint classes

Goal of the Web Ontology Language (OWL)

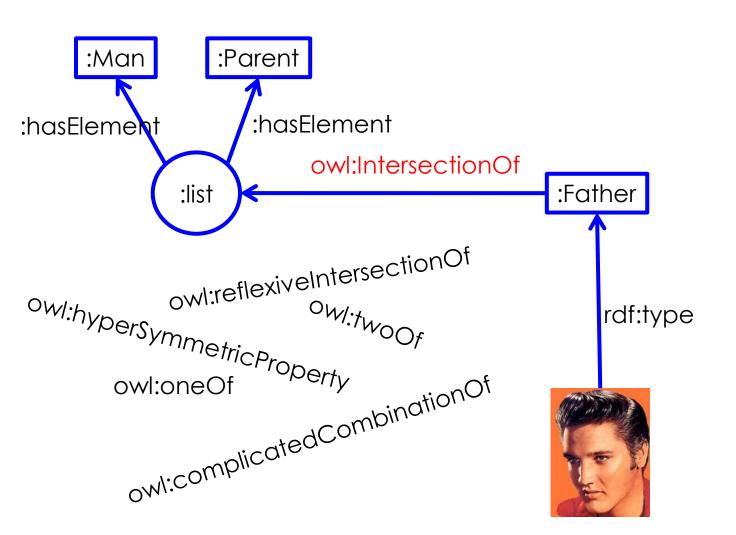
## OWL: Vocabulary

OWL is a namespace that defines predicates with certain semantic rules.



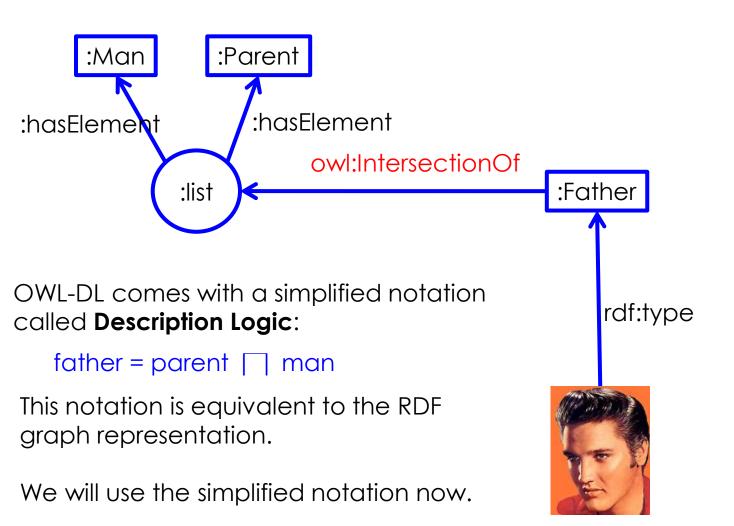
# **OWL: Undecidability**

OWL is a namespace that defines predicates with certain semantic rules. OWL defines so powerful predicates that it is **undecidable**.



## **OWL-DL:** Goal

OWL-DL is a subset of OWL that is decidable.



### OWL-DL: Class constructors

#### Class constructors:

```
The class of things that are in X or in Y
X \mid Y \mid Y
                The class of things that are in both X and Y
                The class of things that are not in X
```

#### Example:

Assume that we have the following classes:

person, parent, hardRockSinger, softRockSinger, happyPerson, marriedPerson, malePerson

```
father = parent | malePerson
rockSinger = hardRockSinger | softRockSinger
unmarried-rock-singing-father
= (parent | man) | (hardRockSinger | | softRockSinger) | ~married
non-rock-singing-person
= person □ ~(hardRockSinger □ softRockSinger)
```

#### Class constructors:

```
    X ☐ Y
    X ☐ Y
    The class of things that are in X or in Y
    ~X
    The class of things that are not in X
    V R.C
    The class of things where all R-links lead to a C
    I R.C
    I R.C
    I R-link to a C
    I R-link to a C
```

R: A predicate/role

C: a class

```
has-happy-child = 3 hasChild.happyPerson
```

This corresponds to the First Order Logic formula:

```
\forall x: has-happy-child(x) <=> \exists y: hasChild(x,y) => happyPerson(y)
```

#### Class constructors:

```
    X ☐ Y
    X ☐ Y
    The class of things that are in X or in Y
    ~X
    The class of things that are not in X
    V R.C
    It also that are in X or in Y
    The class of things that are not in X
    V R.C
    It also that are in both X and Y
    X or in Y
    X he class of things that are not in X
    It also that are in both X and Y
    X or in Y
    X he class of things where all R-links lead to a C
    It also that are in both X and Y
    X or in Y
    X or in Y
    X or in Y
    X or in Y
    Y also that are in both X and Y
    X or in Y
    X or in Y
    X or in Y
    X or in Y
    Y or in Y
```

R: A predicate/role

C: a class

```
has-only-happy-children = ∀ hasChild.happyPerson
```

This corresponds to the First Order Logic formula:

```
\forall x: hohc(x) <=> \forall y: hasChild(x,y) => happyPerson(y)
```

#### Class constructors:

```
    X ☐ Y
    X ☐ Y
    The class of things that are in X or in Y
    ~X
    The class of things that are not in X
    V R.C
    It also that are in X or in Y
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    V R.C
    It also that are in both X and Y
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    X he class of things that are not in X
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    X or in Y
    X he class of things where all R-links lead to a C
    It also that are in both X and Y
    X or in Y
    X or in Y
    X or in Y
    X or in Y
    Y also that are in both X and Y
    X or in Y
    X or in Y
    X or in Y
    X or in Y
    Y or in Y
```

R: A predicate/role

C: a class

```
person-with-happy-child = person ☐ ∃ hasChild.happyPerson

person-with-only-happy-children = person ☐ ∀ hasChild.happyPerson
```

#### Class constructors:

∃ R.C The class of things where there is a R-link to a C

#### Assertions:

```
X Y X is a subclass of Y (everything in X is also in Y)
```

```
singer person
```

This corresponds to the First Order Logic formula:

```
\forall x: singer(x) => person(x)
```

#### Class constructors:

```
X | Y The class of things that are in both X and Y X | Y The class of things that are in X or in Y The class of things that are not in X
```

 $\forall$  R.C The class of things where all R-links lead to a C  $\exists$  R.C The class of things where there is a R-link to a C

#### Assertions:

elvis: singer

This corresponds to the First Order Logic formula: singer(elvis)

#### Class constructors:

```
X | Y The class of things that are in both X and Y X | Y The class of things that are in X or in Y \times X The class of things that are not in X
```

∀ R.C The class of things where all R-links lead to a C
 ∃ R.C The class of things where there is a R-link to a C

#### Assertions:

```
X  X is a subclass of Y (everything in X is also in Y)
a:C a is a thing in the class C
```

elvis: person ☐ ∃ hasChild.happyPerson

### This corresponds to:

elvis: specialClass

specialClass = person | ∃ hasChild.happyPerson

#### Class constructors:

X | Y The class of things that are in both X and Y

X | Y The class of things that are in X or in Y

~X The class of things that are not in X

∀ R.C The class of things where all R-links lead to a C

∃ R.C The class of things where there is a R-link to a C

#### Assertions:

 $X \sqsubseteq Y$  X is a subclass of Y (everything in X is also in Y)

a:C a is a thing in the class C

(a,b):R a and b stand in the relation R, i.e., R(a,b)

(elvis,lisa): hasChild This corresponds to: hasChild(elvis,lisa)

(elvis,priscilla): marriedTo This corresponds to: marriedTo(elvis,priscilla)

#### Class constructors:

The class of things that are in both X and Y
The class of things that are in X or in Y

The class of things that are not in X

∀ R.C ∃ R.C

The class of things where all R-links lead to a C The class of things where there is a R-link to a C

#### Assertions:

X [ Y

X is a subclass of Y (everything in X is also in Y)

a:C

a is a thing in the class C

(a,b):R

a and b stand in the relation R, i.e., R(a,b)

#### Examples:

assume the classes: male, person, happyPerson

and the predicates: marriedTo, hasChild

- build the class of married people
- build the class of people married to at least one happy person
- build the class of happy male married people
- say that Elvis is such a person

#### OWL-DL assertions **entail** other assertions:

elvis: singer singer person

(elvis,lisa): hasChild lisa: happyPerson

(elvis,priscilla): marriedTo

elvis: ∀ marriedTo.happyPerson

Elvis is a singer Every singer is a person

Elvis has child Lisa Lisa is a happy person Elvis is married to Priscilla

All people that Elvis is married to are happy

elvis: person

elvis: ∃ hasChild.happyPerson

priscilla: happyPerson

OWL-DL assertions entail other assertions.

Analogue in First Order Logic

elvis: sing<u>er</u>

singer \_ person

(elvis,lisa): hasChild

lisa: happyPerson

singer(elvis)

 $\forall$  x: singer(x) => person(x)

hasChild(elvis,lisa)

happyPerson(lisa)

elvis: person

elvis: ∃ hasChild.happyPerson

person(elvis)

∃ x: hasChild(elvis,x) /\ happyPerson(x)

OWL-DL is decidable

Thus, OWL-DL corresponds to a decidable subset of First Order Logic

**OWL-DL** is a decidable KR formalism.

Every OWL-DL statement has an analogue

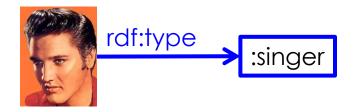
- in first order logic
- and in RDFS

singer(elvis) elvis:singer

notation

Description logic First Order Logic

notation



RDF graphical notation

Thus, OWL-DL corresponds to a decidable fragment of First Order Logic and to a decidable fragment of OWL.

There are a number of free OWLDL reasoners available online:

- Pellet
- FaCT++
- Prova

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  ✓
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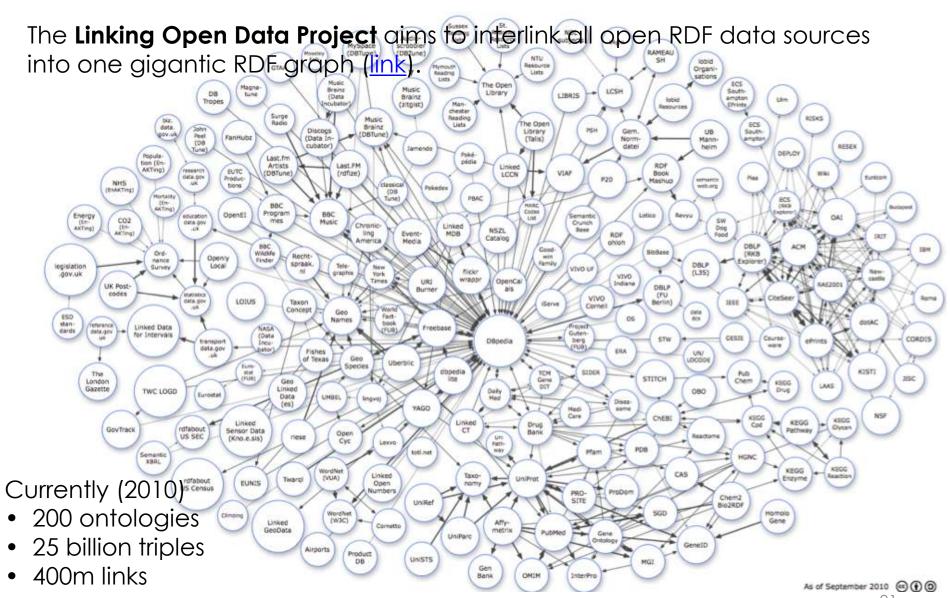
The **Semantic Web** is an evolving extension of the World Wide Web, which promotes a new distributed knowledge representation.

## Existing ontologies

There are already hundreds of RDF ontologies on the Web (<a href="http://www4.wiwiss.fu-berlin.de/lodcloud/">http://www4.wiwiss.fu-berlin.de/lodcloud/</a>)

- US census data
- BBC music database
- Gene ontologies
- DBpedia general knowledge (and hub vocabulary), + YAGO, + Cyc etc.
- UK government data
- geographical data in abundance
- national library catalogs (Hungary, USA, Germany etc.)
- publications (DBLP)
- commercial products
- all Pokemons
- ...and many more

## Linking Data: The Cloud

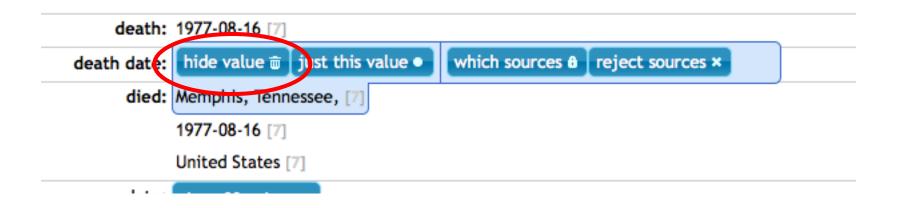


## Sigma

**Sigma** is a Semantic Web search engine developed at the DERI Ireland. (<a href="http://sig.ma">http://sig.ma</a>)



Sigma also allows the user to correct factually wrong information (such as the urban legend that Elvis would be dead).



### Semantic Web

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- ... and it even works. 🗸

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## Summary

We have seen RDF and OWL, the knowledge representation formalisms of the Semantic Web

- RDF is a distributed knowledge representation formalism
- OWL-DL is a decidable fragment of First Order Logic
- The Semantic Web already contains a sizeable number of ontologies

