

Contiki-NG Protothreads

Luca Mottola
`luca.mottola@polimi.it`

Contiki-NG

- Why?
 - A **unique design point** when it comes to concurrency and networking
- Fork of original Contiki
- Focus on
 - RFC-compliant **IPv6 networking**
 - Modern **32-bit platforms**
- Mature: almost 15 years of development up to now (including original Contiki)
- Key features:
 - Configurability, ease of porting to new platforms
 - Efficient, configurable IPv6-based networking
 - Event-driven kernel, **protothread** concurrency model
 - Small memory footprint
- Used both in academia and industry



Contiki-NG Programming

- Use a dialect of the C language
 - Development environments and toolchains are essentially those of C
 - Debugging is a different ballgame though..
- Programs run on real hardware but also
 - On the **native** platform, that is, your host
 - In the COOJA simulator and MSPSim emulator

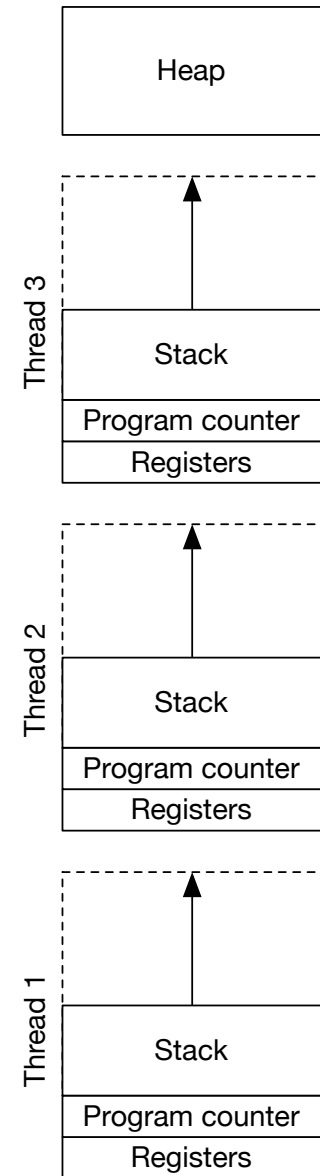


Protothreads



Existing Models: Multi-thread

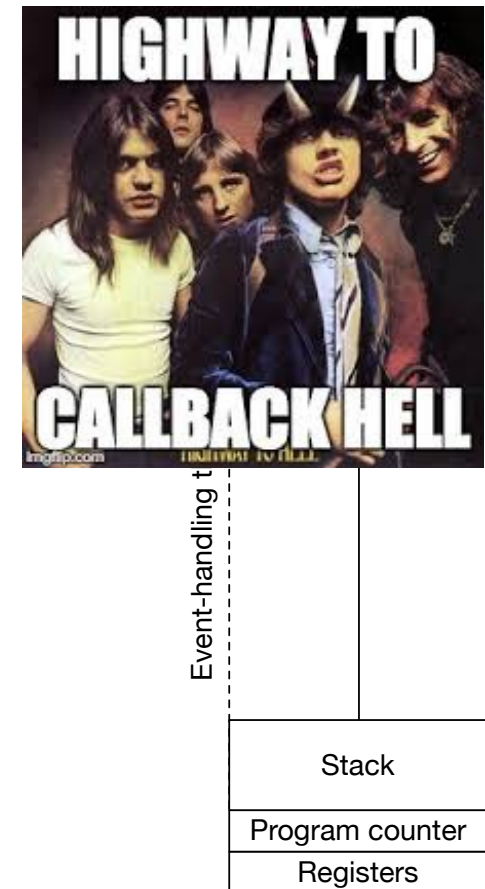
- Similar to **multi-threading** in C/C++
 - A subset of traditional synchronization primitives available
- Not many embedded OSes provide it
 - RIoT, ...
- Pros:
 - **Familiar to programmers**
 - Facilitates porting non-embedded code
- Cons:
 - Heavy on memory: each thread requires its own stack
 - Difficult to tune: hard to anticipate every thread's max stack size



Existing Models: Event-driven

Heap

- **Non-blocking calls** and (possibly asynchronous) **callbacks**
- Provided by many
 - ARM mBed, Contiki, Zephyr, Apache MyNewt...
- Pros:
 - Facilitates implementing reactive code
 - **Memory-efficient**: only one stack
- Cons:
 - Difficult to program: code as state machines, no sequential semantics

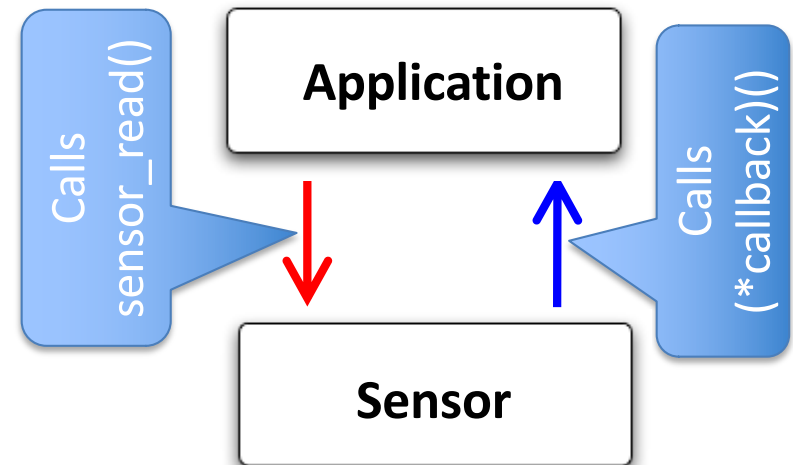


Event-driven: Example

```
// Sensor interface: sensor.h  
void sensor_read(void (*callback)(error_t, val_t));
```

```
#include "sensor.h"  
  
void my_sensor_read_done(error_t err,  
                        val_t value) {  
    // ...  
    if (err == OK) {  
        do_something(value);  
    }  
}  
  
void foo() {  
    // ... do something  
    sensor_read(&my_sensor_read_done);  
    // ... do something more  
}
```

The corresponding `sensor.c` is provided by sensor-specific implementations



If callbacks are asynchronous, the execution is preempted unless in an **atomic** block



Protothreads

- "Best of both worlds"
- Maintain sequential semantics...
- ...with a single stack
- Provided in Contiki and Contiki-NG

```
#include "contiki.h"
#include <stdio.h>

PROCESS(test_proc, "Test process");
AUTOSTART_PROCESSES(&test_proc);

PROCESS_THREAD(test_proc, ev, data)
{
    PROCESS_BEGIN();

    printf("Hello, world!\n");

    PROCESS_END();
}
```

We will see later
where displayed...



Protothreads: Handling Events

- Protothreads process incoming events
 - A timer expiring, a packet received...
- They are scheduled cooperatively
 - Individual protothreads decide when to release the MCU

```
PROCESS_THREAD(test, ev, data)
{
    static struct etimer et;
    PROCESS_BEGIN();

    etimer_set(&et, CLOCK_SECOND); /* Trigger a 1s timer */
    for(;;) {
        PROCESS_WAIT_EVENT_UNTIL(etimer_expired(&et));
        printf("Hello, world!\n");
        etimer_set(&et, CLOCK_SECOND);
    }

    PROCESS_END();
}
```

Set a new timer,
triggering the next event
for this protothread

Every time the
protothread resumes,
this condition is checked



Protothreads: Implementation

- Uses **local continuations** as result of macro expansion
 - When **set**, allow one to tag specific places in the code
 - When **resumed**, allow the execution to restart from within the code
- Use of macros spares the need for a dedicated preprocessor



Protothreads: Implementation

```
PROCESS_THREAD(test, ev, data)
{
    static struct etimer et;
    PROCESS_BEGIN();

    etimer_set(&et, CLOCK_SECOND);
    for(;;) {
        PROCESS_WAIT_EVENT_UNTIL(etimer_expired(&et));
        printf("Hello, world!\n");
        etimer_set(&et, CLOCK_SECOND);
    }

    PROCESS_END();
}
```

```
static
char test (struct pt* pt,
           event_t ev, void* data)
{
    static struct etimer et;
    switch(pt->lc) { case 0:

        etimer_set(&et, CLOCK_SECOND);
        for(;;) {
            pt->lc = 15; case 15:
                if (!(etimer_expired(&et))) {
                    return 0;
                }
                printf("Hello, world!\n");
                etimer_set(&et, CLOCK_SECOND);
            }

        pt->lc = 0; return 2;
    }
```

- What really happens when resuming:
 - The protothread function is called again, and starts over
 - A **switch** statement reads the value of **pt->lc**
 - The execution jumps to the **case**: corresponding to where the execution yielded
- Therefore, local variables need to be **static**, or they get re-initialized every time the protothread resumes
 - **No free lunches**: they consume memory also when the protothread is not executing!
- Do not use **switch** inside protothreads, or you will mess things up



Protothreads: Events

- Events used to
 - Exchange data among protothreads
 - Achieve cooperative scheduling

Find code under
examples/events!


```
PROCESS_THREAD(ping_process, ev, data)
{
    static struct etimer timer;

    PROCESS_BEGIN();

    /* Start the ping pong... */
    ping_event=process_alloc_event();
    printf("Sent the first ping!\n");
    process_post(&pong_process, ping_event, &count);

    // ...
}
```

```
PROCESS_THREAD(pong_process, ev, data)
{
    // ...
    PROCESS_WAIT_EVENT();
    if (ev == ping_event) {
        printf("Got a ping: %d!\n", *(int*)data);
        process_post(&ping_process, pong_event, &count);
        printf("Sent a pong event to ping!\n");
    } else {
        printf("Got an unknoww event!\n");
    }
    // ...
}
```



Protothreads: APIs

- **PROCESS_PAUSE()**: releases control, rescheduled immediately
- **PROCESS_YIELD()**: releases control, wait for next event
- **process_poll(&process_name)**: polls another protothread
- **PROCESS_EXIT()**: stops a protothread
- **process_exit(&process_name)**: kills another protothread

