

Greedy Algorithms

JOB SEQUENCING WITH DEADLINES

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Problem Statement

The problem is stated as below.

- There are n jobs to be processed on a machine.
- Each job i has a deadline $d_i \geq 0$ and profit $p_i \geq 0$.
- P_i is earned iff the job is completed by its deadline.
- The job is completed if it is processed on a machine for unit time.
- Only one machine is available for processing jobs.
- Only one job is processed at a time on the machine.

Problem Statement (Contd..)

- A feasible solution is a subset of jobs J such that each job is completed by its deadline.
- An optimal solution is a feasible solution with maximum profit value.

Example : Let $n = 4$, $(p_1, p_2, p_3, p_4) = (100, 10, 15, 27)$,
 $(d_1, d_2, d_3, d_4) = (2, 1, 2, 1)$

Example : Let $n = 4$, $(p_1, p_2, p_3, p_4) = (100, 10, 15, 27)$,
 $(d_1, d_2, d_3, d_4) = (2, 1, 2, 1)$

Solution

Sr.No.	Feasible Solution	Processing Sequence	Profit value
(i)	(1,2)	(2,1)	110
(ii)	(1,3)	(1,3) or (3,1)	115
(iii)	(1,4)	(4,1)	127 is the optimal one
(iv)	(2,3)	(2,3)	25
(v)	(3,4)	(4,3)	42 ↑
(vi)	(1)	(1)	100
(vii)	(2)	(2)	10
(viii)	(3)	(3)	15
(ix)	(4)	(4)	27

GREEDY ALGORITHM TO OBTAIN AN OPTIMAL SOLUTION

- Consider the jobs in the non increasing order of profits subject to the constraint that the resulting job sequence J is a feasible solution.

ALGORITHM

Algorithm 4.2: JobSequence($job[1..n]$, $p[1..n]$, $d[1..n]$, $list[1..n]$, n)

1. Initialize $list$ with zeros. // make all n time slots empty
 2. $profit = 0$; // $profit$ = total profit
 3. For $i = 1$ to n do
 4. $k = d[i]$; // set i th job's deadline to k
 5. While($k > 0$)
 6. If($list[k] = 0$) // if k th slot is vacant,
 7. $list[k] = job[i]$; // assign i th job to k th slot
 8. add $p[i]$ to $profit$; // accumulate total profit
 9. go to step 3; // exit from while-loop (break)
 10. Endif
 11. Decrement k by 1; // otherwise get preceding slot
 12. Endwhile
 13. Endfor
 14. Print $list$
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COMPLEXITY ANALYSIS OF JS ALGORITHM

- Sorting (sort p_i into decreasing order) and selection of jobs take $O(n \log n)$ time in total
- Consider n jobs in turn.
- For each job, inserting the job into partial solution using its deadline take $O(n)$ time.
- Checking whether the new solution is still feasible takes $O(n)$ time.
- Hence total running time, in worst case is $O(n^2)$

EXAMPLE:

let $n = 5$, $(p_1, \dots, p_5) = (20, 15, 10, 5, 1)$ and $(d_1, \dots, d_5) = (2, 2, 1, 3, 3)$.

Let $n=7$, $(p_1, \dots, p_7) = (35, 30, 25, 20, 15, 12, 5)$ and $(d_1 \dots d_7) = (3, 4, 4, 2, 3, 1, 2)$



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4/29/2025

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