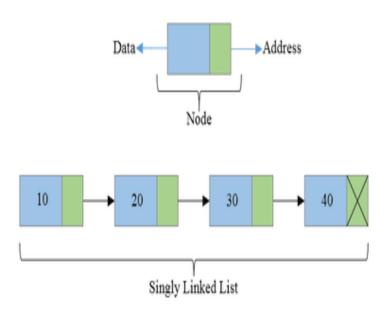
Module 2.1

Linear Data Structure: Linked List

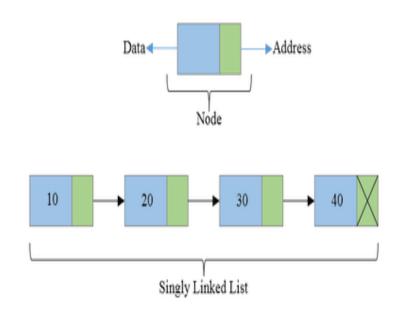
Outline

- Introduction and Representation of Linked List
- Linked List v/s Array
- Implementation of Linked List
- Circular Linked List
- Doubly Linked List,
- Application Polynomial Representation and Addition
- Other additional applications/Case study.

- Linear Collection of data elements called Nodes
- Linear order is given by means of pointers.

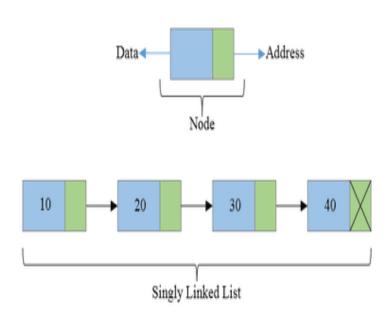


- Each node may be divided into atleast two fields for:
 - Storing Data
 - Storing Address of next element.



31-07-2024 5

- The Last node's Address field contains Null rather than a valid address.
- It's a NULL Pointer and indicates the end of the list.



Comparison between Array and Linked List

Advantages of Linked List

- Linked are Dynamic Data Structures
 - Grow and shrink during execution of the program
- Efficient Memory Utilization
 - As memory is not preallocated.
 - Memory can be allocated whenever required and deallocated when not needed.
- Insertion and deletions are easier and efficient
 - Provide flexibility in inserting a data item at a specified position and deletion of a data item from the given position
- Many complex applications can be easily carried out with linked lists

Disadvantages of Linked List

 Access to an arbitary data item is little bit cumbersome and also time consuming

More memory

• If the number of fields are more, then more memory space is needed.

Advantages of Arrays

- Simple to use and define
- Supported by almost all programming languages
- Constant access time
 - Array element can be accessed a[i]
- Mapping by compiler
 - Compiler maps a[i] to its physical location in memory.
 - This mapping is carried out in constant time, irrespective of which element is accessed

Disadvantages of Arrays

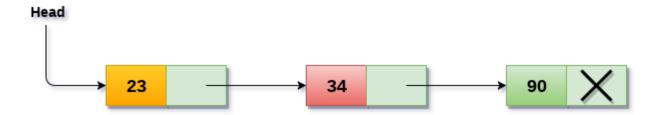
- Static Data Structure-
 - Size of an array is defined at the time of programming
- Insertion and Deletion is time consuming
- Requires Contiguous memory

Types of Linked List

- Singly Linked List
- Doubly Linked List
- Circular Linked List

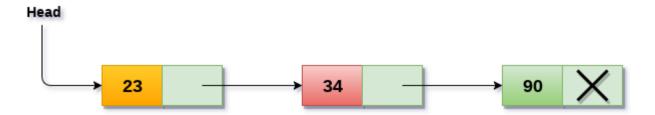
Singly Linked List

- All nodes are linked in sequential manner
- Linear Linked List
- One way chain
- It has beginning and end



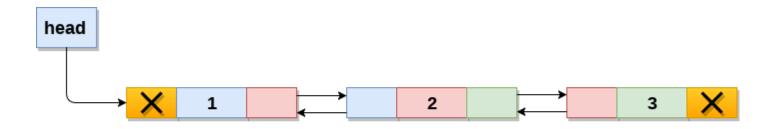
Singly Linked List

- Problem-
 - The predecessor of a node cannot be accessed from the current node.
 - This can be overcome in doubly linked list.



Doubly Linked List

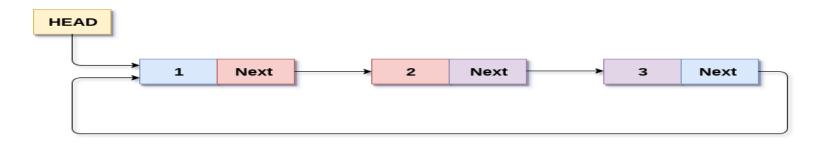
- Linked List holds two pointer fields
- Addresses of next as well as preceding elements are linked with current node.
- This helps to traverse in both Forward or Backward direction



Doubly Linked List

Circular Linked List

- The first and last elements are adjacent.
- A linked list can be made circular by
 - Storing the address of the first node in the link field of the last node.



Circular Singly Linked List

Linked List Operations

- Creation
- Insertion
- Deletion
- Traversal
- Searching

Implementation of Linked Lists

• Structures in C are used to define a node

 Address of a successor node can be stored in a pointer type variable

Struct Basics

Struct Syntax

Struct examples

Pointer to structure

Sample Programs for Pointer to structures

Initialization of Pointers to Structures

```
struct node
{
    int info;
    struct node *link;
}

Pointer that points to the structure itself,
Thus Linked List.
```

Singly Linked List

Creation of a new node

```
struct node{
               type1 member1;
               type2 member2;
                                    Info part
               ......
              struct node *link;
           };
struct node
                                    info
   int info;
   struct node *link;
};
struct node *start = NULL;
```

Creation of a new node

New node=temp

```
struct node *tmp;
tmp= (struct node *) malloc(sizeof(struct node));
tmp->info=data;
tmp->link=NULL;
```

Creating a Linked List

```
create_list(int data)
        struct node *q,*tmp;
        tmp= (struct node *) malloc(sizeof(struct node));
        tmp->info=data;
        tmp->link=NULL;
        if(start==NULL) /*If list is empty */
                 start=tmp;
        else
            /*Element inserted at the end */
                 q=start;
                 while(q->link!=NULL)
                          q=q->link;
                 q->link=tmp;
}/*End of create_list()*/
```

Explanation-

```
if(start==NULL) /*If list is empty */
   {
     start=tmp;
}
```

Explanation-

```
else
{    /*Element inserted at the end */
        q=start;
    while(q->link!=NULL)
        q=q->link;
        q->link=tmp;
}
```

31-07-2024 26

Traversing a Linked List

• Visit the node and print the data value

Traversing a Linked List

Assign the Value of start to another pointer say q

```
struct node *q=start;
```

- Now q also points to the first element of linked list.
- For processing the next element, we assign the address of the next element to the pointer q as-

```
q=q->link;
```

Traverse each element of the Linked list through this assignment until
pointer q has NULL address, which is link part of last element.

```
while(q!=NULL)
{
    q=q->link;
}
```

31-07-2024 28

Explanation-

```
while(q!=NULL)
     {
          q=q->link;
     }
```

Explanation-

Algorithm for Traversing a Linked List

Step 1:[INITIALIZE] SET PTR=START

Step 2: Repeat Steps 3 and 4 while PTR!=NULL

Step 3: Print PTR->INFO

Step 4: Set PTR=PTR->LINK

[End of Loop]

Step 5:EXIT

Algorithm for Counting the number of elements in a Linked List

Step 1:[INITIALIZE] SET COUNT=0

Step 2:[INITIALIZE] SET PTR=START

Step 3: Repeat Steps 4 and 5 while PTR!=NULL

Step 4: Set COUNT=COUNT +1

Step 5: Set PTR=PTR->LINK

[End of Loop]

Step 6:Print COUNT

Step 7:EXIT

Searching a Linked List

- First traverse the linked list
- While traversing compare the info part of each element with the given element

31-07-2024 33

Searching a Linked List

```
search(int data)
{
       struct node *ptr = start;
       int pos = 1;
       while(ptr!=NULL)
       {
               if(ptr->info==data)
               {
                       printf("Item %d found at position %d\n",data,pos);
                       return;
               ptr = ptr->link;
               pos++;
       if(ptr == NULL)
               printf("Item %d not found in list\n",data);
}/*End of search()*/
```

Algorithm for Searching a Linked List

```
Step 1:[INITIALIZE] SET POSITION=1
```

Step 2:[INITIALIZE] SET PTR=START

Step 3: Repeat Steps 4 while PTR!=NULL

Step 4: If DATA=PTR->INFO

Print POSITION

Exit

[End of If]

Set PTR=PTR->LINK

Set POSITION=POSITION +1

[End of Loop]
Step 5:If PTR=NULL

Print Search Unsuccessful

[End of If]

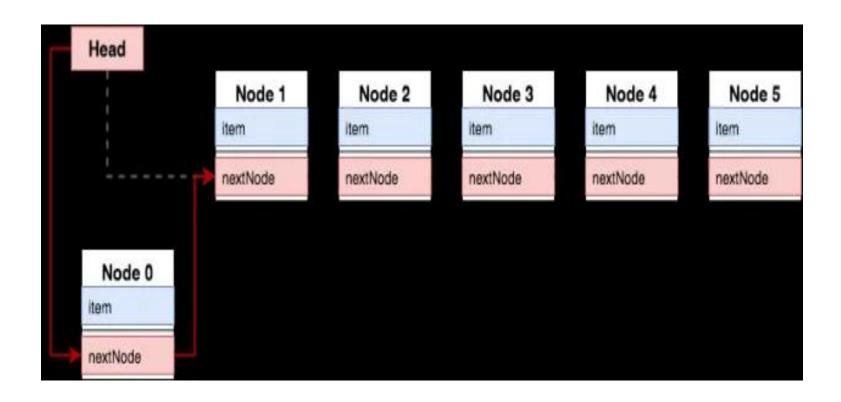
Step 6: Exit

31-07-2024 35

Insertion into a Linked List

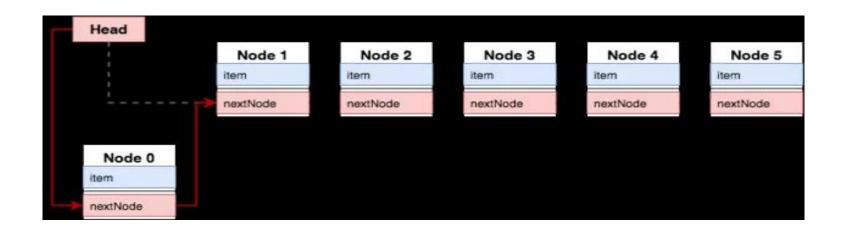
- Insertion is possible in two ways:
 - Insertion at Beginning
 - Insertion in Between

Case 1- Insertion at Beginning



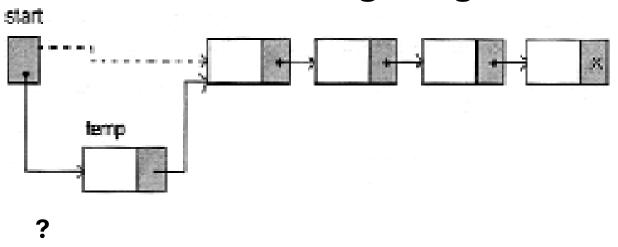
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Case 1-Insertion at Beginning



CREATE THE NEW NODE, CONNECT NEW NODE TO THE OLD FIRST NODE CONNECT THE START POINTER TO THE NEW NODE,

Case 1- Insertion at Beginning



31-07-2024

Case 1- Insertion at Beginning

- Lets say tmp is the pointer which points to the node that has to be inserted
- Assign data to the new node

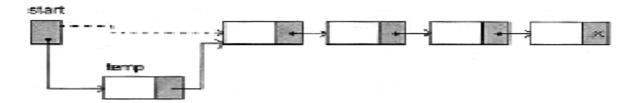
tmp->info=data;

- Start points to the first element of linked list
- Assign the value of start to the link part of the inserted node as

tmp->link=start;

- Now inserted node points beginning of the linked list.
- To make the newly inserted node the first node of the linked list:

start=temp



Algorithm for Insertion at Beginning in a Linked List

- First check whether Memory is available for the new node.
- If the memory has exhausted then an Overflow message is printed
- Else We allocate memory for the new node

Algorithm for Insertion at Beginning in a Linked List

Step 1:If AVAIL=NULL Then

WRITE OVERFLOW

Go to Step 7

[End of If]

Step 2: Set TEMP=AVAIL

Step 3: Set AVAIL=AVAIL->LINK

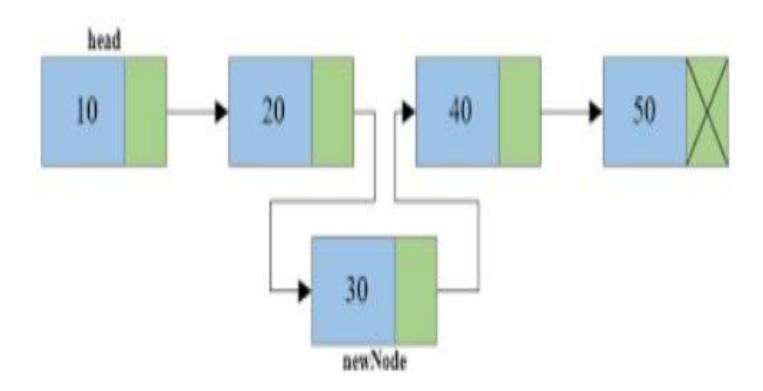
Step 4: Set TEMP->INFO=DATA

Step 5: Set TEMP->LINK=START

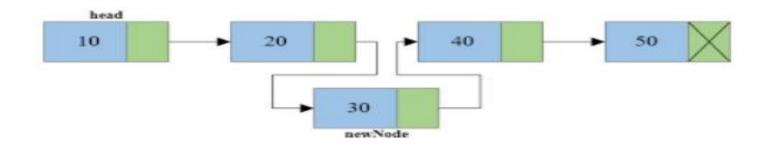
Step 6: Set START=TEMP

Step 7:EXIT

Case 2- Insertion in Between



Case 2- Insertion in Between



CREATE THE NEW NODE, CONNECT THE NEW NODE TO THE NEXT NODE CONNECT THE PREVIOUS TO THE NEW NODE,

.....

Case 2-Insertion in Between

q=start;

• First we traverse the linked list for obtaining the node after which we want to insert the element

```
for(i=0;i<pos-1;i++)
            q=q->link;
            if(q==NULL)
                   printf("There are less than %d elements",pos);
                   return;
                                      Case 2-
   }/*End of for*/
                                          start
```

Explanation-

Explanation-

31-07-2024

Algorithm for Insertion in Between

```
Step 1:If AVAIL=NULL Then
      WRITE OVERFLOW
      Go to Step 13
    [End of If]
Step 2: Set TEMP=AVAIL
Step 3: Set AVAIL=AVAIL->LINK
Step 4: Set TEMP->INFO=DATA
Step 5: Set TEMP->LINK=NULL
Step 6: Read POSITION from User
Step 7: Set Q=START
Step 8 : Set I=0
Step 9: Repeat step 10 till I<POS-1
Step 10:
                     Set Q=Q->LINK
             If Q=NULL
                     Print Less Number of Elements Case 2-
```

Exit

[End of If]

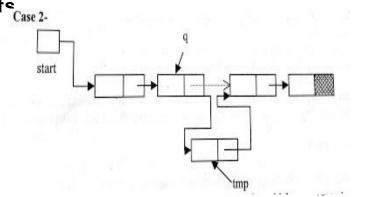
SET I=I+1

Step 11:Set TEMP->LINK=Q->LINK

[End of Loop]

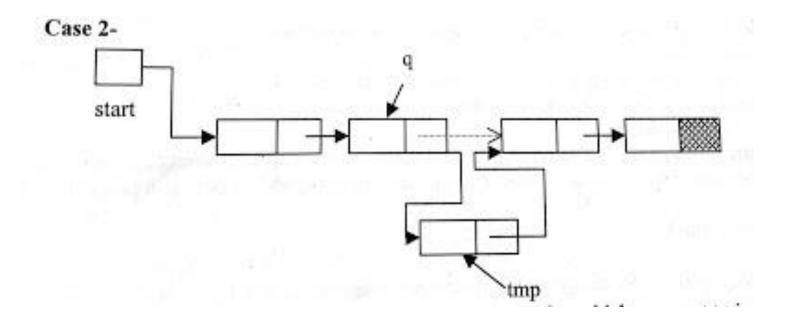
Step 12:Set Q->LINK=TEMP





Case 2-Insertion in Between

• Then we add the new node by adjusting address fields



Case 2- Insertion at the end

- Without Using Position?
- At the end?

Algorithm for Insertion at the end of a Linked List

```
Step 1: If AVAIL=NULL Then
WRITE OVERFLOW
Go to Step 10
[End of If]
```

Step 2: Set TEMP=AVAIL

Step 3: Set AVAIL=AVAIL->LINK

Step 4: Set TEMP->INFO=DATA

Step 5: Set TEMP->LINK=NULL

Step 6: Set Q=START

Step 7: Repeat step 8 while PTR->LINK!=NULL

Step 8: SET Q=Q->LINK

[End of Loop]

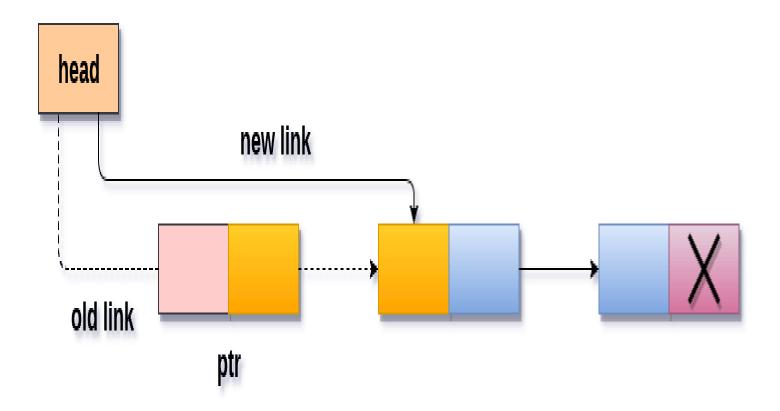
Step 9: SET Q->LINK=TEMP

Step 10 : EXIT

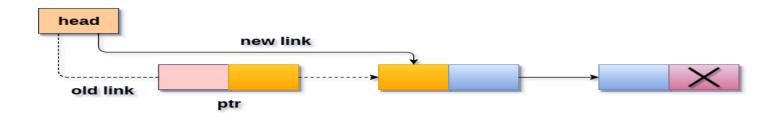
Deletion from a Linked List

- For deleting the node from a linked list, first we traverse the linked list and compare with each element.
- After finding the element there may be two cases for deletion-
 - Deletion in beginning
 - Deletion in between

Deletion in beginning



Deletion in beginning



CONNECT START POINTER TO THE SECOND NODE...... DELETE THE FIRST NODE

Deletion in beginning

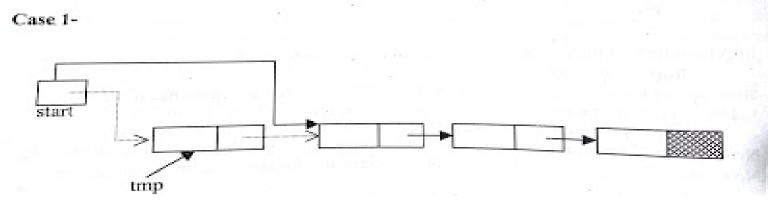
- Start points to the first element of linked list.
- If element to be deleted is the first element of linked list then we assign the value of start to tmp as-

tmp = start;

- So tmp points to the first node which has to be deleted.
- Now assign the link part of the deleted node to start as-

start=start->link;

- Since start points to the first element of linked list, so start->link will point to the second element of linked list.
- Now we should free the element to be deleted which is pointed by tmp.
 free(tmp);



31-07-2024

Algorithm for Deletion in the beginning of the Linked List

Step 1:If START=NULL

Step 2: Write UNDERFLOW

Step 3: Go to Step 7

[End of If]

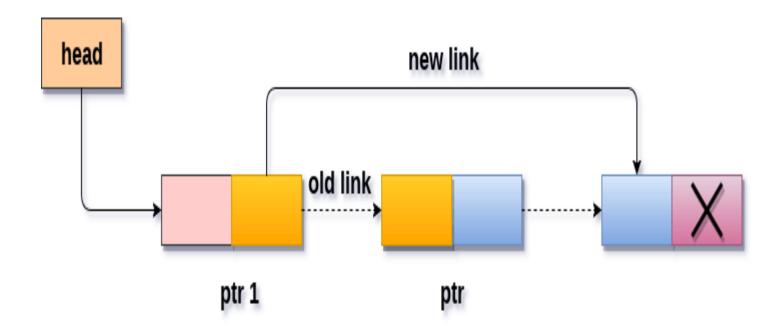
Step 4: Set TEMP=START

Step 5: Set START=START->LINK

Step 6:FREE TEMP

Step 7:EXIT

Deletion in between



Deletion a node from specified position

Deletion in between

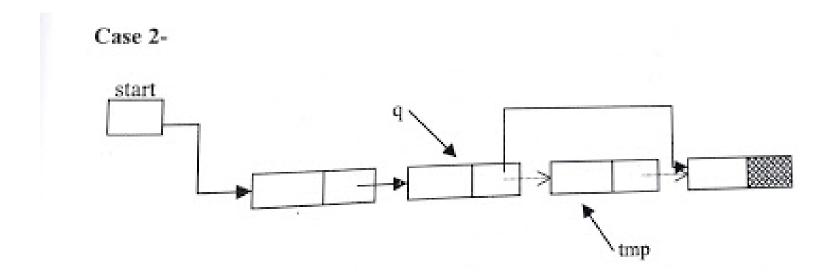


DELETE THE NODE AND CONNECT THE PREVIOUS AND THE NEXT NODE......

Deletion in between

- If the element is other than the first element of linked list then
 - we give the link part of the deleted node to the link part of the previous node.
 - This can be as-

```
tmp =q->link;
q->link = tmp->link;
free(tmp);
```



Deletion at the end

• If node to be deleted is last node of linked list then statement 2 will be as-

```
tmp =q->link;
q->link = NULL;
free(tmp);
```

Circular Linked List

Why Circular?

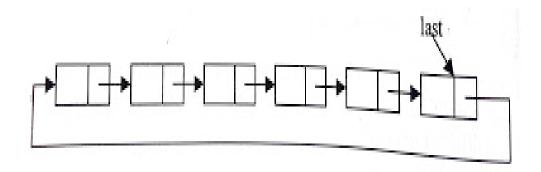
- In a singly linked list,
 - If we are at any node in the middle of the list, then it is not possible to access nodes that precede the given node.
 - This problem can be solved by slightly altering the structure of singly linked list.

How?

- In a singly linked list, next part (pointer to next node) of the last node is NULL,
 - if we utilize this link to point to the first node then we can reach preceding nodes.

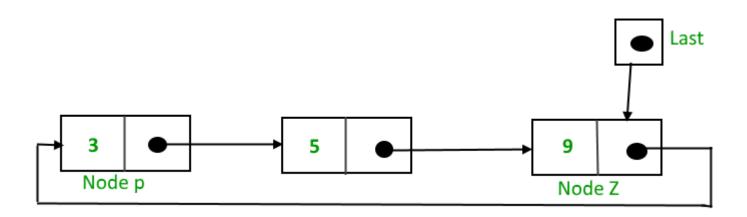
Implementation of circular linked list

- Creation of circular linked list is same as single linked list.
- Last node will always point to first node instead of NULL.



Implementation of circular linked list

- One pointer last,
 - which points to last node of list and link part of this node points to the first node of list.



Advantages of a Circular linked list

• In circular linked list, we can easily traverse to its previous node, which is not possible in singly linked list.

- Entire list can be traversed from any node.
 - If we are at a node, then we can go to any node. But in linear linked list it is not possible to go to previous node.

Advantages of a Circular linked list

- In Single Linked List, for insertion at the end, the whole list has to be traversed.
- In Circular Linked list,
 - with pointer to the last node there won't be any need to traverse the whole list.
 - So insertion in the begging or at the end takes constant time irrespective of the length of the list i.e O(1).
 - It saves time when we have to go to the first node from the last node.
 - It can be done in single step because there is no need to traverse the in between nodes

Disadvantages of Circular linked list

- Circular list are complex
 - as compared to singly linked lists.
- Reversing of circular list is a complex
 - as compared to singly or doubly lists.
- If not traversed carefully,
 - then we could end up in an infinite loop.

 Like singly and doubly lists circular linked lists also <u>doesn't</u> <u>supports direct accessing of elements.</u>

31-07-2024

Insertion into a circular linked list:-

Insertion in a circular linked list may be possible in two ways-

- Insertion in an empty list
- Insertion at the end of the list
- Insertion at beginning
- Insertion in between

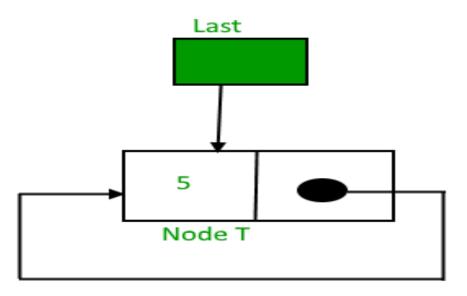
Insertion in an empty list

New element can be added as-

• If linked list is empty:

```
If (last==NULL)
{
last=tmp;
tmp->link=last
}
```





31-07-2024

Insertion at the end of circular linked list

• If linked list is not empty: Insertion at the end of the list tmp->link = last->link; /* added at the end of list*/ last->link = tmp; last = tmp; Last Node T 10 Last Node T 6

Insertion at the end of circular linked list

• If linked list is not empty: Insertion at the end of the list tmp->link = last->link; /* added at the end of list*/ last->link = tmp; last = tmp; Last Node T 10 Last Node T 6

31-07-2024

Insertion at the beginning of circular linked list

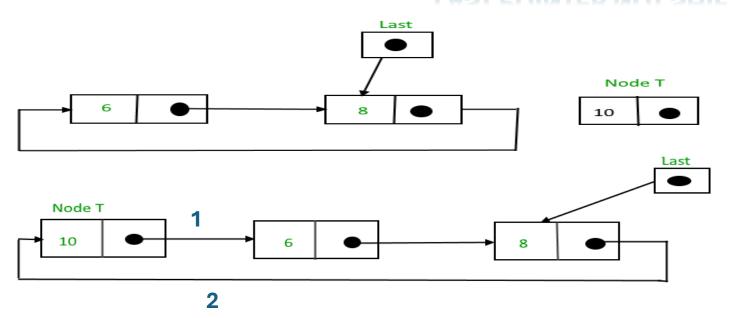
• If linked list is not empty:

Insertion at the beginning of the list

Follow these step:

- 1. Create a node, say tmp.
- 2. Make tmp-> next = last -> next.
- 3. $last \rightarrow next = tmp$.

LAST POINTER NOT SHIFTED!



Insertion at the beginning of circular linked list

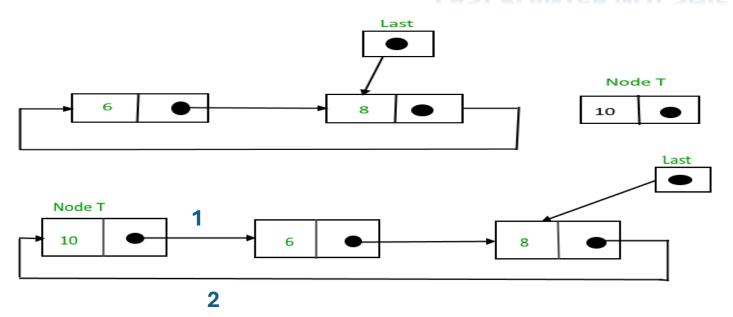
• If linked list is not empty:

Insertion at the beginning of the list

Follow these step:

- 1. Create a node, say tmp.
- 2. Make tmp-> next = last -> next.
- 3. $last \rightarrow next = tmp$.

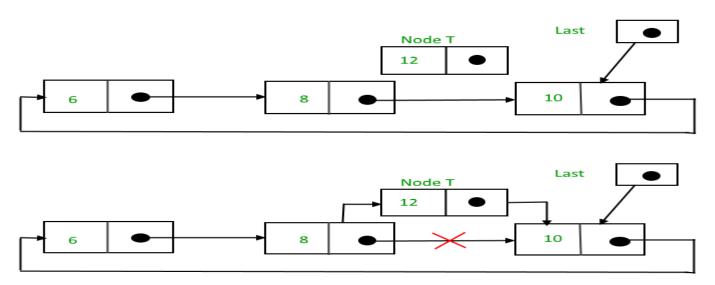
LAST POINTER NOT SHIFTED!



Insertion in between of circular linked list

• Insertion in between is same as in single linked list. This can be as-

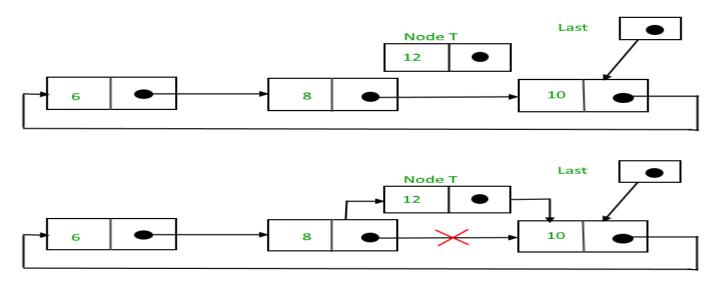
• Here q points to the node after which new node will be inserted.



Insertion in between of circular linked list

• Insertion in between is same as in single linked list. This can be as-

• Here q points to the node after which new node will be inserted.



Creation of CLL

```
create list(int num)
        struct node *q,*tmp;
        tmp= malloc(sizeof(struct node));
        tmp->info = num;
        if(last == NULL)
                last = tmp;
                tmp->link = last;
        else
                tmp->link = last->link; /*added at the end of list*/
                last->link = tmp;
                last = tmp;
}/*End of create_list()*/
```

31-07-2024