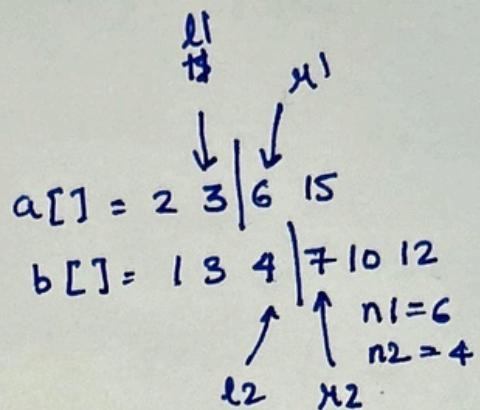


Q1. $a[], b[]$

```
if (a.length > b.length) {
    int n1 = a.length;
    n2 = b.length;
} else {
    n1 = b.length;
}
n2 = a.length;
```



$n = n1 + n2;$

$left = (n1 + n2 + 1) / 2;$ → left half

$low = 0, high = n1;$

while ($low \leq high$) {

$mid1 = (low + high) / 2;$

$mid2 = left - mid1;$

$l1 = (mid1 > 0) ? a[mid1 - 1] : Integer.MIN_VALUE;$

$l2 = (mid2 > 0) ? b[mid2 - 1] : Integer.MIN_VALUE;$

$x1 = (mid1 < n1) ? a[mid1] : Integer.MAX_VALUE;$

$x2 = (mid2 < n2) ? b[mid2] : Integer.MAX_VALUE;$

 if ($l1 \leq x2 \ \&\ l2 \leq x1$) {

 if ($n \% 2 == 1$) return Math.max(l1, l2);

 else .

 return ((double) (Math.max(l1, l2) + Math.min(x1, x2))) / 2.0;

}

 else if ($l1 > x2$)

 high = mid1 - 1;

 else

 low = mid1 + 1;

}

Q2. int cnt = 0;
 el = 0;
 $n = arr.length$;
 for ($i = 0$; $\rightarrow n$) {
 if ($cnt == 0$) {
 cnt += 1;
 el = arr[i];
 }
 else if ($el == arr[i]$) {
 cnt++;
 }
 else {
 cnt--;
 }
 }

$\begin{array}{cccccc} 2 & 2 & 1 & 1 & 1 & 2 & 2 \end{array}$
 $el = \cancel{X} 2$
 $\rightarrow cnt = \cancel{0} \cancel{X} \cancel{X} \cancel{0} \cancel{X} \cancel{0} 1$

$S \cdot \text{Complexity} = O(1)$

Q3. cnt1 = 0; cnt2 = 0;
 $el1 = \text{Integer.MIN_VALUE}$; $el2 = \text{Integer.MAX_VALUE}$; $n = arr.length$;

for ($i = 0 \rightarrow n$) {
 if ($cnt1 == 0 \& el1 != arr[i]$) {
 cnt1 = 1;
 el1 = arr[i];
 }
 else if ($cnt2 == 0 \& el2 != arr[i]$) {
 cnt2 = 1;
 el2 = arr[i];
 }
 else if ($arr[i] == el1$) cnt1++;
 else if ($arr[i] == el2$) cnt2++;
 else {
 cnt1--;
 cnt2--;
 }
 }

$\begin{array}{ccccccc} 2 & 1 & 1 & 3 & 1 & 4 & 5 & 6 \end{array}$
 $cnt1 = 1$
 $el1 = 2$
 $cnt2 = \cancel{0} \cancel{X} 2$
 $el2 = \min 1$

$S \cdot \text{Complexity} = O(1)$