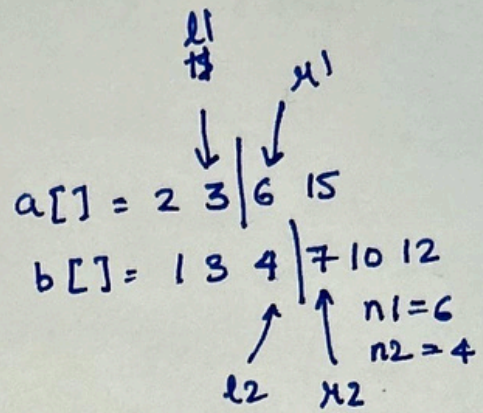


Q1. $a[]$, $b[]$
 if ($a.length > b.length$) {
 int $n1 = a.length$;
 $n2 = b.length$;
 }
 else {
 ~~int~~ $n1 = b.length$;
 $n2 = a.length$;
 }



$$4 + 6 = 10 / 2 = 5$$

$n = n1 + n2$;

$left = (n1 + n2 + 1) / 2$; \rightarrow left half

$low = 0$, $high = n1$;

while ($low \leq high$) {

$mid1 = (low + high) / 2$;

$mid2 = left - mid1$;

$l1 = (mid1 > 0) ? a[mid1 - 1] : Integer.MIN_VALUE$;

$l2 = (mid2 > 0) ? b[mid2 - 1] : Integer.MIN_VALUE$;

$h1 = (mid1 < n1) ? a[mid1] : Integer.MAX_VALUE$;

$h2 = (mid2 < n2) ? b[mid2] : Integer.MAX_VALUE$;

 if ($l1 \leq h2$ & & $l2 \leq h1$) {

 if ($n \% 2 == 1$) return $Math.max(l1, l2)$;

 else

 return $((double) (Math.max(l1, l2) + Math.min(h1, h2))) / 2.0$;

 }

 else if ($l1 > h2$)

$high = mid1 - 1$;

 else

$low = mid1 + 1$;

}

Q2. $\text{int cnt} = 0;$
 $\text{el} = 0;$

$n = \text{arr.length};$

for ($i = 0 \rightarrow n$) {

if ($\text{cnt} == 0$) {

$\text{cnt} = 1;$

$\text{el} = \text{arr}[i];$

}

else if ($\text{el} == \text{arr}[i]$) {

$\text{cnt}++;$

}

else {

$\text{cnt}--;$

}

}

Complexity = $O(1)$

2 2 1 1 1 2 2

$\text{el} = 2$

$\rightarrow \text{cnt} = \cancel{0} \times \cancel{2} \times \cancel{0} \times \cancel{0} \times \cancel{0} \times \cancel{0} \times 1$

Q3. $\text{cnt1} = 0; \text{cnt2} = 0;$

$\text{el1} = \text{Integer.MIN_VALUE}; \text{el2} = \text{Integer.MIN_VALUE}; n = \text{arr.length};$

for ($i = 0 \rightarrow n$) {

if ($\text{cnt1} == 0 \ \&\& \ \text{el2} \neq \text{arr}[i]$) {

$\text{cnt1} = 1;$

$\text{el1} = \text{arr}[i];$

}

else if ($\text{cnt2} == 0 \ \&\& \ \text{el1} \neq \text{arr}[i]$) {

$\text{cnt2} = 1;$

$\text{el2} = \text{arr}[i];$

}

else if ($\text{arr}[i] == \text{el1}$) $\text{cnt1}++;$

else if ($\text{arr}[i] == \text{el2}$) $\text{cnt2}++;$

else {

$\text{cnt1}--;$

$\text{cnt2}--;$

}

}

2 1 1 3 1 4 5 6

$\text{cnt1} = 1$

$\text{el1} = 2$

$\text{cnt2} = \cancel{0} \times 2$

$\text{el2} = \text{min}$

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Complexity = $O(1)$