LL(1) Parser

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Current lab: Repository Link

Implementation details

Class **Parser** is responsible for implementing the LL(1) parser. It contains the first and follow methods that are used to obtain the *first* and *follow* tables. These two methods are implemented in the same way they were specified in the lectures.

The class Parser keeps only the final version of the *first* and *follow* tables. That means, the steps that lead to the final version of *first* and *follow* are not stored because they are not relevant for the problem.

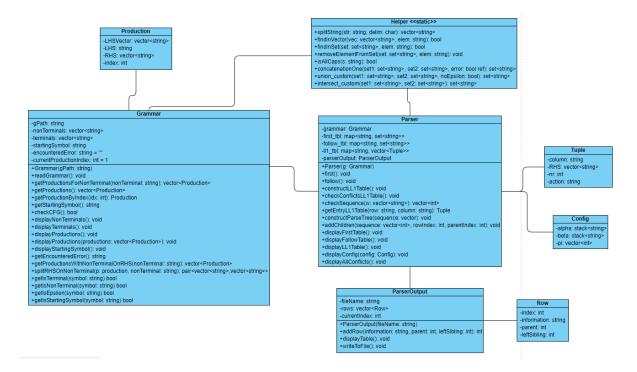
In order to write a cleaner implementation for the *first* and *follow* methods, some methods were added to the Grammar class. These methods are:

- *getIsTerminal*, *getIsNonTerminal*, *getIsEpsilon*, *getIsStartingSymbol*: these are used for checking whether a given symbol belongs to one of the above-mentioned classes
- *getProductionsWithNonTerminalOnRHS*: this method retrieves a vector of productions containing a given symbol in the right-hand side. This method is used in the *follow* algorithm.
- *splitRHSOnNonTerminal*: this method takes a production and a non-terminal (which must belong to the right-handside of the production). It splits the production into two parts, denoted by alpha and gamma. Alpha is the sequence of symbols before the non-terminal and gamma is the sequence after.

The set operations (union, concatenation of length one) were implented in the **Helper** class. The method for constructing the LL(1) parse table is *constructLl1Table*. Method *checkSequence* receives a sequence of terminal symbols and checks whether the sequence belongs to the language generated by grammar G or not. It returns a vector of integers representing the indexes of the productions that must be performed to solve the sequence.

Class **ParserOutput** is responsible for storing the parse tree (or syntax tree) generated by method *constructParseTree*. This method requires the sequence of indexes from *checkSequence*.

The class diagram becomes:



Experiments:

Example 1 (g1.in)

S A B C D a b c 5 S -> a A b | B A A -> a A | c A | c B -> D C D -> Epsilon | b C -> c S

[First table ...]
A: a c
B: b c
C: c
D: Epsilon b
S: a b c
a: a
b: b

[Follow table ...] A: b Epsilon B: a c C: a c D: c S: Epsilon

First table:

	F_0	F_1	F ₂	$F_3=F_2$
S	a	a	a, b, c	
A	a, c	a, c	a, c	
В	Ø	b, c	b, c	
С	С	С	С	
D	ε, b	ε, b	ε, b	

Follow table:

	L_0	L_1	L_2	$L_3=L_2$
S	3	3	ε	
A	Ø	b, ε	b, ε	
В	Ø	a, c	a, c	
С	Ø	Ø	a, c	
D	Ø	С	С	

Example 2 (g2.in)

S A
a b c
2
S -> A a | a
A -> b A | c
S

[First table ...] A: b c S: a b c a: a b: b c: c

[Follow table ...] A: a S: Epsilon

First table:

	F_0	F_1	$F_2=F_1$
S	a	a, b, c	
A	b, c	b, c	

Follow table:

	L_0	L_1	$L_2=L_1$
S	3	3	
A	Ø	a	

```
Example 3 (g3.in)
                                          [First table ...]
                                          A: b c
SABC
                                          B: d
a b c d
                                          C: d
4
                                          S: b c d
S \rightarrow A a \mid B B
                                          a: a
A \rightarrow b A \mid c
                                          b: b
B \rightarrow C a
                                          c: c
                                          d: d
C \rightarrow d
S
```

```
[Follow table ...]
A: a
B: d
C: a
S: Epsilon
```

First table:

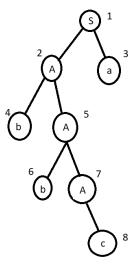
	F_0	F_1	F_2	$F_3 = F_2$
S	Ø	b, c	b, c, d	
A	b, c	b, c	b, c	
В	Ø	d	d	
С	d	d	d	

Follow table:

	L_0	L_1	$L_2=L_1$
S	3	ε	
A	Ø	a	
В	Ø	d	
С	Ø	a	

Sequence: bbca

```
[Checking sequence ...]
(bbca$ , S$ , )
[PUSH]
(bbca$ , Aa$ , 1)
PUSH]
(bbca$ , bAa$ , 13)
[POP]
(bca$ , Aa$ , 13)
[PUSH]
(bca$ , bAa$ , 133)
[POP]
(ca$´, Aa$ , 133)
[PUSH]
(ca$ , ca$ , 1334)
[POP]
(a$ , a$ , 1334)
[POP]
($ , $ , 1334)
[ACCEPT]
Sequence: 1 3 3 4
```



```
[Table ...]
          Parent Left-Sibling
   Info
    S
       0
          0
   Α
       1
          0
    а
       1
   b
          0
       2
          1
   b
       5
          0
       5
   Α
          1
          0
    done]
```

Example 4 (g4.in)

```
SBC
abcd
3
```

```
[First table ...]
B: a Epsilon
C: c Epsilon
S: a b c d
a: a
b: b
c: c
d: d
```

```
[Follow table ...]
B: b
C: d
S: Epsilon
```

```
\begin{array}{l} S \rightarrow B \ b \mid C \ d \\ B \rightarrow a \ B \mid Epsilon \\ C \rightarrow c \ C \mid Epsilon \\ S \end{array}
```

First table:

	F_0	F_1	$F_2=F_1$
S	Ø	a, b, c, d	
В	a, ε	a, ε	
С	ς, ε	ς, ε	

Follow table:

	L_0	L_1	$L_2=L_1$
S	3	3	
В	Ø	b	
С	Ø	d	

Sequence: aab

```
[Checking sequence ...]
(aab$ , S$ , )
[PUSH]
(aab$ , Bb$ , 1)
[PUSH]
aab$ , aBb$ , 13)
POP]
(ab$ , Bb$ , 13)
[PUSH]
ab$ , aBb$ , 133)
POP]
(b$ , Bb$ , 133)
[PUSH]
(b$ , b$ , 1334)
[POP]
($ , $ , 1334)
[ACCEPT]
Sequence: 1 3 3 4
```

```
Table ...]
Ιd
    Info Parent Left-Sibling
    S
        0
1
2
3
4
5
            0
    В
        1
            0
    b
        1
            1
        2
            0
    В
        2
            1
        5
    а
           0
                                  В
    В
           1
     Epsilon
                   0
                                        5
     done]
```

Example 8 (g8.in)

```
S \\ a b c \\ 1 \\ S \rightarrow a S b S \mid c \\ S
```

```
[LL1 table ...]
$: $: acc
S: a: (aSbS,1) c: (c,2)
a: a: pop
b: b: pop
c: c: pop
[... done]
```

```
[First table ...]
S: a c
a: a
b: b
c: c
[... done]
```

```
[Follow table ...]
S: Epsilon b
[... done]
```

Sequence: aacbcbc

```
[Checking sequence ...]
(aacbcbc$ , S$ , )
[PUSH]
(aacbcbc$ , aSbS$ , 1)
[POP]
(acbcbc$ , SbS$ , 1)
[PUSH]
(acbcbc$ , aSbSbS$ , 11)
[POP]
(cbcbc$ , SbSbS$ , 11)
[PUSH]
(cbcbc$ , cbSbS$ , 112)
[POP]
(bcbc$ , bSbS$ , 112)
[POP]
(cbc$ , SbS$ , 112)
[PUSH]
(cbc$ , cbS$ , 1122)
[POP]
(bc$ , bS$ , 1122)
[POP]
(c$ , S$ , 1122)
[PUSH]
(c$ , c$ , 11222)
[POP]
($ , $ , 11222)
[ACCEPT]
Sequence: 1 1 2 2 2
```

