Introduction to Artificial Intelligence and Problem Formulations

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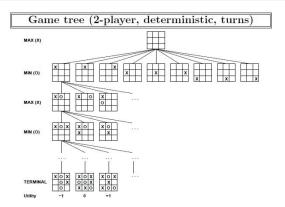
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Game Tree - Tic-Tac-Toe



- From the initial state, MAX has **nine** possible moves.
- 9! = 362880 terminal nodes
- Play alternates between MAX'S placing an X and MIN'S placing an 0
 until we reach leaf nodes corresponding to terminal states such that one
 player has three in a row or all the squares are filled.



Introduction to Artificial Intelligence and Problem Formulations

Game Tree - Tic-Tac-Toe...

Different functions

- S₀: The initial state, which specifies how the game is set up at the start.
- ullet PLAYER(s): Defines which player has the move in a state.
- ACTIONS(s): Returns the set of legal moves in a state.
- RESULT(s, a): The transition model, which defines the result of a move.
- TERMINAL-TEST(s): A terminal test, which is true when the game is over and false
 otherwise. States where the game has ended are called terminal states.
- UTILITY (s, p): A utility function (also called an objective function or payoff function), defines the final numeric value for a game that ends in terminal state s for a player p. In chess, the outcome is a win, loss, or draw, with values +1, 0, or \(\frac{1}{2}\). Some games have a wider variety of possible outcomes; the payoffs in backgammon range from 0 to +192. A zero-sum game is (confusingly) defined as one where the total payoff to all players is the same for every instance of the game. Chess is zero-sum because every game has payoff of either 0 + 1, 1 + 0 or \(\frac{1}{2}\) + \(\frac{1}{2}\). "Constant-sum" would have been a better term, but zero-sum is traditional and makes sense if you imagine each player is charged an entry fee of \(\frac{1}{2}\).





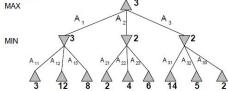
Optimal Decisions in Games - The minimax Algorithm

Minimax

Perfect play for deterministic, perfect-information games

Idea: choose move to position with highest minimax value = best achievable payoff against best play

E.g., 2-ply game:



- MAX: Upper triangle and MIN: Lower triangle
- MAX'S best move at the root is A_1 , it leads to the highest minimax value and MIN'S best reply is A_{11} , because it leads to the successor with the lowest minimax value.



Optimal Decisions in Games - The minimax Algorithm

• Given a game tree, the optimal strategy can be determined by examining the **minimax value** of each node, which we write as MINIMAX(n).

```
\begin{aligned} & \text{MINMAX-VALUE}(n) = \\ & & \text{UTLITY}(n) \\ & \text{max}_{n \in Successors(n)} & \text{MINMAX-VALUE}(s) & \text{if } n \text{ is a MIN node} \\ & \text{min}_{n \in Successors(n)} & \text{MINMAX-VALUE}(s) & \text{if } n \text{ is a MIN node}. \end{aligned}
```

Properties of minimax

Complete?? Yes, if tree is finite (chess has specific rules for this)

Optimal?? Yes, against an optimal opponent. Otherwise??

Time complexity?? $O(b^m)$

Space complexity?? O(bm) (depth-first exploration)

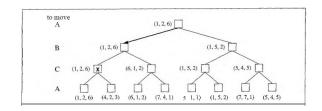
For chess, $b \approx 35$, $m \approx 100$ for "reasonable" games \Rightarrow exact solution completely infeasible

But do we need to explore every path?

- The first MIN node, has three successors with values 3, 12, and 8, so its minimax value is 3.
- The root node is a MAX node; its successors have minimax values 3, 2, and 2; so it has a minimax value of 3.

Optimal Decisions in Multiplayer Games

- Three Players: A, B, C, Vector: $\langle v_A, v_B, v_C \rangle$
- (1, 2, 6) and (4, 2, 3) = (1, 2, 6) (Player C)







Optimal Decisions in Multiplayer Games

- Suppose A and B are in weak positions and C is in a stronger position.
- Then it is often optimal for both A and B to attack C rather than each other.
- Let C destroy each of them individually.
- In this way, collaboration emerges from purely selfish behavior. Of course, as soon as C weakens under the joint onslaught, the alliance loses its value, and either A or B could violate the agreement.

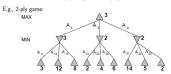


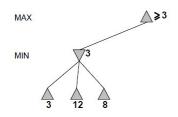


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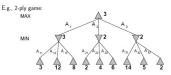


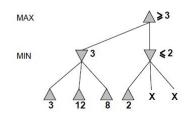


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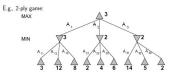


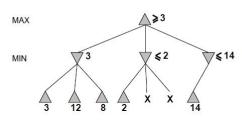


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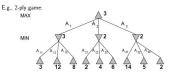


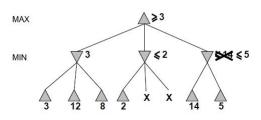
Minimax

Perfect play for deterministic, perfect-information games

Idea: choose move to position with highest minimax value

— best achievable payoff against best play





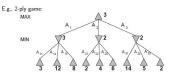


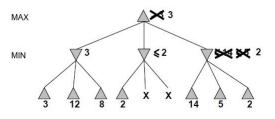


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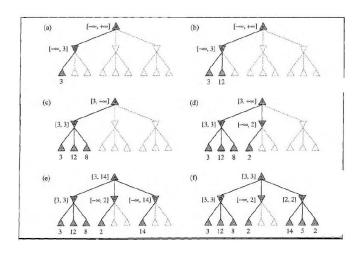
• When we apply alpha-beta pruning to a standard minimax tree, it returns the same move as minimax would, but prunes away branches that cannot possibly influence the final decision.

Example

```
MINIMAX(root) = max(min(3, 12, 8), min(2, x, y), min(14, 5, 2))
= \max(3, \min(2, x, y), 2)
= \max(3, z, 2) \text{ where } z = \min(2, x, y) \le 2
= 3
```





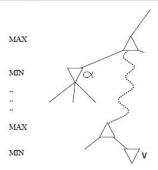


• [at least, at most]





Why is it called $\alpha - \beta$?



 α is the best value (to MAX) found so far off the current path If V is worse than α , MAX will avoid it \Rightarrow prune that branch Define β similarly for MIN



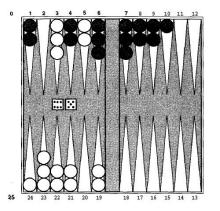


- Alpha-beta needs to examine only $O(b^{m/2})$ nodes to pick the best move, instead of $O(b^m)$ for minimax.
- This means that the effective branching factor becomes \sqrt{b} instead of b-for chess, about 6 instead of 35.





- Backgammon is a typical game that combines luck and skill.
- Two players, 24 narrow triangles, 15 checkers (white/black) per player
- ullet 0 \sim 6: Black's home board, 7 \sim 12: outer board
- ullet 19 \sim 25: White's home board, 13 \sim 18: outer board



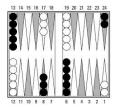




- The aim of the game is to move all one's pieces off the board.
- White moves clockwise toward 25, and black moves counterclockwise toward 0.
- A piece can move to any position except one where there are two or more of the opponent's pieces.
- If it moves to a position with one opponent piece, that piece is captured and has to start its journey again from the beginning.
- In this position, white has just rolled 6-5 and has four legal moves: (5-10,5-11), (5-11,19-24), (5-10,10-16), and (5-11,11-16).







Initial arrangement of checkers

- 1st player: 2 on first, 5 on twelve and ninteen, and 3 on seventeen point.
- 2nd player: 2 on twenty-fourth point, 5 on thirteen and sixth point, and 3 on eigth point.
- Both players have their own pair of dice and a dice cup used for shaking.
- To start the game, each player throws a single die. If equal numbers come up then both players roll again until they roll different numbers.
- The player throwing the higher number now moves his checkers according to the numbers showing on both dice.
- After the first roll, the players throw two dice and alternate turns.





Players begin with two checkers on their 24-point, three checkers on their 8-point, and five checkers each on their 13-point and their 6-point.

Rules

- A checker may be moved only to an open point one, i.e., not occupied by 2 or more opposing checkers.
- The numbers on the 2 dice constitute separate moves. For example, if a player rolls 5 and 3, he may move one checker 5 spaces to an open point and another checker 3 spaces to an open point or he may move the one checker a total of 8 spaces to an open point, but only if the intermediate point (either 3 or 5 spaces from the starting point) is also open.

Means-end Analysis

Means-end Analysis

- **①** Comparing the current state S_i to a goal state S_g and computing the difference D_{ig}
- ② An operator O_k is then selected to reduce the difference D_{ig} .
- The operator O_k is applied if possible. If not, the current state is saved, a subgoal is created and means-end analysis is applied recursively to reduce the subgoal.
- If the subgoal is solved, the saved state is restored and work is resumed on the original problem.

Example

Let the initial PL object $L_i=R$ & $(\sim\!P
ightarrow Q)$ and goal object $L_g=(Q\ V\ P)$

$$L_i' = (\sim P \rightarrow Q) \& R = (\sim \sim P \lor Q) \& R = (P \lor Q) \& R$$

$$L_g = (Q \lor P) \& R$$

NOTE: A comparison of the difference that R is on the left in L_i but on the right in L_g . This causes a subgoal to be set up to reduce this difference.







