Introduction to Artificial Intelligence: AO* and Hill Climbing algorithms

Dr. Tapas Kumar Mishra

Assistant Professor
Department of Computer Science and Engineering
SRM University
Amaravati-522502, Andhra Pradesh, India
Email: tapaskumar.m [at] srmap [dot] edu [dot] in





Heuristic Function

Admissible heuristics

E.g., for the 8-puzzle:

 $h_1(n) = \text{number of misplaced tiles}$

 $h_2(n) = \text{total Manhattan distance}$

(i.e., no. of squares from desired location of each tile)





Start State

Goal State

$$h_1(S) = ?? 6$$

 $h_2(S) = ?? 4+0+3+3+1+0+2+1 = 14$





Heuristic Function

- A typical solution is about 20 steps, although this of course varies depending on the initial state.
- The branching factor is about 3 (when the empty tile is in the middle, there are four possible moves; when it is in a corner there are two; and when it is along an edge there are three).
- \bullet An exhaustive search to depth $\bf 20$ would look at about $3^{20}=3.5\times 10^9$ states.
- There are only 9! = 362,880 different arrangements of 9 squares.
- Manhattan/City block distance: $|x_1 x_2| + |y_1 y_2|$





Heuristic Function

Dominance

If $h_2(n) \ge h_1(n)$ for all n (both admissible) then h_2 dominates h_1 and is better for search

Typical search costs:

$$\begin{array}{ll} d=14 & \mathsf{IDS}=3,\!473,\!941 \; \mathsf{nodes} \\ & \mathsf{A}^*(h_1)=539 \; \mathsf{nodes} \\ & \mathsf{A}^*(h_2)=113 \; \mathsf{nodes} \\ d=24 & \mathsf{IDS}\approx 54,\!000,\!000,\!000 \; \mathsf{nodes} \\ & \mathsf{A}^*(h_1)=39,\!135 \; \mathsf{nodes} \\ & \mathsf{A}^*(h_2)=1,\!641 \; \mathsf{nodes} \end{array}$$

Given any admissible heuristics h_a , h_b ,

$$h(n) = \max(h_a(n), h_b(n))$$

is also admissible and dominates $h_a,\ h_b$





Relaxed Problems

Relaxed problems

Admissible heuristics can be derived from the exact solution cost of a relaxed version of the problem

If the rules of the 8-puzzle are relaxed so that a tile can move anywhere, then $h_1(n)$ gives the shortest solution

If the rules are relaxed so that a tile can move to any adjacent square, then $h_2(n)$ gives the shortest solution

Key point: the optimal solution cost of a relaxed problem is no greater than the optimal solution cost of the real problem





Local Search Algorithms and Optimization Problems

- Local search algorithms operate using a single current state (rather than multiple paths) and generally move only to neighbors of that state.
- Typically, the paths followed by the search are not retained.
- Although local search algorithms are not systematic, they have two key advantages:
 - 1 They use very little memory-usually a constant amount
 - They can often find reasonable solutions in large or infinite (continuous) state spaces for which systematic algorithms are unsuitable.



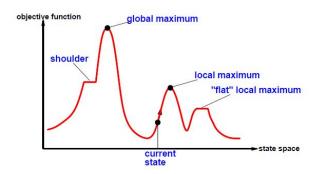


Local Search Algorithms and Optimization Problems

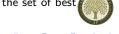
- Local search algorithms are useful for solving pure optimization problems, in which the aim is to find the best state according to an objective function.
- A landscape has both "location" (defined by the state) and "elevation" (defined by the value of the heuristic cost function or objective function).
 - If elevation corresponds to cost, then the aim is to find the lowest valley-a global minimum
 - If elevation corresponds to an **objective function**, then the aim is to find the **highest peak**-a **global maximum**.



Hill-Climbing Search



- It is simply a loop that continually moves in the direction of increasing value-that is, uphill.
- It terminates when it reaches a "peak" where no neighbor has a higher value.
- Hill-climbing algorithms typically choose randomly among the set of best successors, if there is more than one.



Hill-Climbing Search

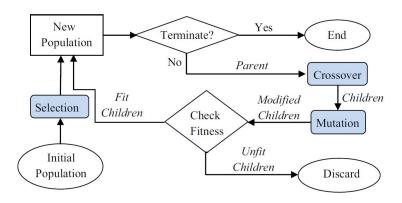
Hill climbing often gets stuck for the following reasons

- 1 Local maxima: a local maximum is a peak that is higher than each of its neighboring states, but lower than the global maximum.
- Q Ridges: Ridges result in a sequence of local maxima that is very difficult for greedy algorithms to navigate.
- Plateaux: A plateau is an area of the state space landscape where the evaluation function is flat. It can be a flat local maximum, from which no uphill exit exists, or a shoulder, from which it is possible to make progress.





Genetic Algorithm



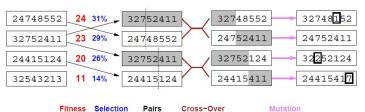




Genetic Algorithm

Genetic algorithms

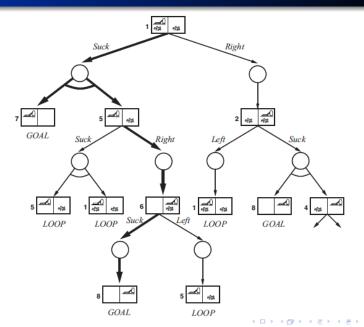
= stochastic local beam search + generate successors from **pairs** of states





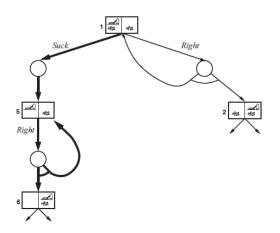


AND OR Search





AND OR Search







AND/OR graphs

- Some problems are best represented as achieving subgoals, some of which achieved simultaneously and independently (AND)
- Up to now, only dealt with OR options

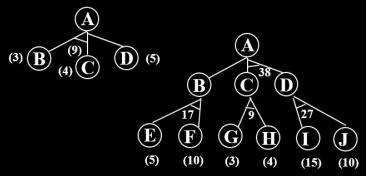






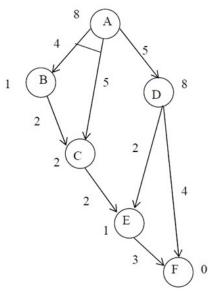
AND/OR search

 We must examine several nodes simultaneously when choosing the next move



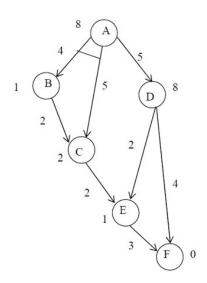


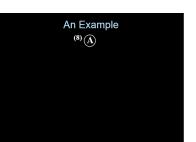




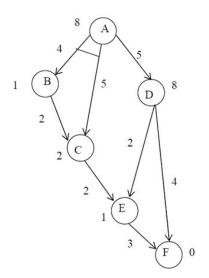


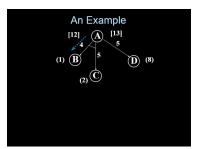






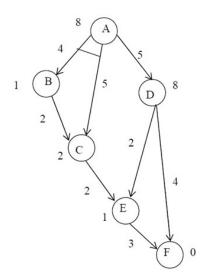


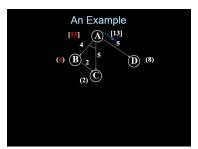






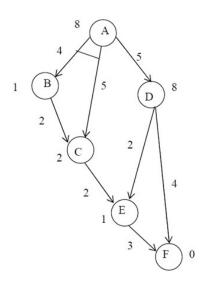


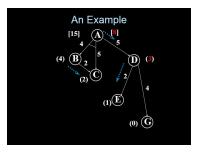






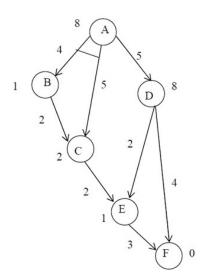


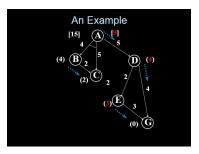






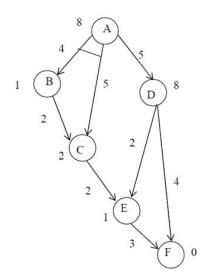


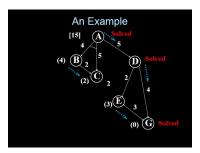
















- 1 Place the starting node s on open.
- ② Using the search tree constructed thus far, compute the most promising solution tree T_0 .
- **3** Select a node n that is both on open and a part of T_0 . Remove n from open and place it on closed.
- **4** If n is a terminal goal node, label n as solved. If the solution of n results in any of n's ancestors being solved, label all the ancestors as solved. If the start node s is solved, exit with success where T_0 is the solution tree. Remove from open all nodes with a solved ancestor.
- If n is not a solvable node (operators cannot be applied), label n as unsolvable. If the start node is labeled as unsolvable, exit with failure. If any of n's ancestors become unsolvable because n is, label them unsolvable as well. Remove from open all nodes with unsolvable ancestors.

- 6 Otherwise, expand node n generating all of its successors. For each such successor node that contains more than one subproblem, generate their successors to give individual subproblems. Attach to each newly generated node a back pointer to its predecessor. Compute the cost estimate h^* for each newly generated node and place all such nodes that do not yet have descendants on open. Next, recompute the values of h^* at n and each ancestor of n.
- Return to step 2.







