# Cooperative PSO with Spatially Meaningful Neighbours

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# Abstract

# Contents

1	Intr	roducti	ion	5
2	Bac	kgroui	$\mathbf{nd}$	5
	2.1	Partic	le Swarm Optimization	5
		2.1.1	Overview	5
		2.1.2	Particle	5
		2.1.3	Position Update Formula	6
		2.1.4	Velocity Update Formula	6
		2.1.5	Swarm	7
		2.1.6	Fitness Function	8
		2.1.7	Parameters	8
	2.2	Co-op	erative Particle Swarm Optimization	10
		2.2.1	CPSO-S	11
		2.2.2	Variant: $CPSO-S_k$	12
		2.2.3	Variant: $CPSO-R_k$	13
	2.3	Neighl	borhood Topologies	13
		2.3.1	Common Spatially Random Topologies	14
		2.3.2	Spatially Meaningful Neighbors	15
	2.4	Dynar	mic Connections	16
	2.5	Delau	nay Triangulation	19
		2.5.1	Triangulation	19
		2.5.2	Delaunay Triangulation	19
3	$\operatorname{Lit}\epsilon$	erature	e Review/Previous work	20
	3.1	CPSO		20
	3.2	Delau	nay Triangulation	22

4	$\mathbf{CP}$	SO using Delaunay Triangulation	23
	4.1	Experimental Setup	23
	4.2	CPSO configuration	24
	4.3	Benchmark functions	25
	4.4	Experiment 1: Robustness Test on Lower Dimensions	26
		4.4.1 Results	26
		4.4.2 Discussion	26
	4.5	Experiment 2: Robustness Test on Higher Dimensions	29
		4.5.1 Discussion	29
5	Cor	nclusion	31
J	Coi	iciusion	91
Bi	ibliog	graphy	32
$\mathbf{L}$	$\mathbf{ist}$	of Tables	
	1	PSO Pseudo Code	7
	2	CPSO-S Pseudo Code	11
	3	$\mathrm{CPSO}\text{-}\mathrm{S}_k$ Pseudo Code	12
	4	Choosing The Best Neighbour	18
	5	Functions, Stop Criteria and Domains	25
	6	Rastrigin Robustness Analysis (n=6)	27
	7	Rosenbrock Robustness Analysis (n=6)	27
	8	Griewanck Robustness Analysis (n=6)	28
	9	Ackley Robustness Analysis (n=6)	28
	10	Rastrigin Robustness Analysis (n=20)	29
	11	Rosenbrock Robustness Analysis (n=20)	30
	12	Griewanck Robustness Analysis (n=20)	30

CPSO	P. Wil	Ison
13	Ackley Robustness Analysis (n=6)	31
${f List}$	of Figures	
1	Visual Representation of the Neighborhoods[1]	14
2	Von Neumann[2]	16

How to know if the particles are working together [1]  $\ \ldots \ \ldots \ \ldots \ \ldots$ 

Circumcircles drawn amongst the Triangles [3]  $\ \ldots \ \ldots \ \ldots \ \ldots \ \ldots$  17

20

3

4

# 1 Introduction

# 2 Background

## 2.1 Particle Swarm Optimization

#### 2.1.1 Overview

Particle Swarm Optimization (PSO) is an iterative population-based computational intelligence algorithm which finds an optimal solution for non-linear continuous problems. Computational Intelligence is a domain of study in which the patterns and behaviours of intelligent agents (such as humans or animals) are used to help solve other types of problems [4]. Particle Swarm Optimization specifically draws from the flight patterns of birds in their search of food as well as the schooling pattern of fish. In both cases, the animals tend to follow each other from a distance which allows each to get a different vantage point. When one of the fish/birds find the best source of food (or the best place to land) they all tend to converge and head toward that point.

#### 2.1.2 Particle

PSO builds off this metaphor by representing the birds or fish as Particles. A Particle has a position, velocity and personal best value. It keeps track of all these values and updates them with each iteration. the position of the particle maps logically to the location of where the bird is on a map. The x,y,z,... values of the position also correlate to the input values in the problem which we are trying to optimize. The more input values (variables) the problem contains, the higher dimensions are needed to represent the problem. The initial position of the particle should be randomly generated to allow the swarm to effectively test the entire search space. If they all begin with similar solutions, it is unlikely that all potential solutions

are tested which leads to a higher possibility of converging on a local minimum as opposed to the best overall value. Velocity represents the direction and speed of the Particle. It determines how you update the position after each iteration. The velocity, too is updated at the end of each iteration and is influenced by the current velocity, the personal best value (or the best combination of values found previously by the particle which is maintained by the particle) and the global/network best (the best value found by all the particles in the swarm). The velocity and position are updated each iteration with the following functions.

#### 2.1.3 Position Update Formula

$$x_i = x_i + v_i$$

Updating the position of the particle is as simple as taking the current position and adding the current velocity to it. For multi-dimensional problems where the particles have multiple dimensions, this addition will need to be done for each dimension. In this case, the i represents the index of the position.

#### 2.1.4 Velocity Update Formula

$$v_i = \alpha v_i + \omega C_1 (y_i - x_i) + \lambda C_2 (\hat{y}_i - x_i)$$

The velocity update function has a little more to it. The new velocity is determined by a combination of it's current velocity, the distance it is from the particles personal best and the distance it is from the global best. In this case  $y_i$  represents the personal best for that index and  $\hat{y}_i$  represents the global best. The influence that each has is determined by the inertial weight  $(\alpha)$ , the cognitive weight  $(\omega)$  and social weight  $(\lambda)$  respectively. These

values typically add up to 1 and determine the percentage of influence they have on the new velocity. Additionally, the formula also makes use of a random component. The values  $C_1$  and  $C_2$  are randomly generated values in the range [0,1] that adds more randomness and aids in keeping the swarm from converging too quickly.

#### 2.1.5 Swarm

The Swarm is the heart that makes PSO work. It is the combination of a number of particles working in unison to solve the same problem. The Swarm needs to maintain the individual particles but it also needs to store a Global Best value. The global best is the position of the best solution found by any of the particles within the swarm. Each particle will in turn use that value to help influence the change in velocity for each iteration as seen in the velocity update formula above. This transfer of knowledge amongst the particles helps the swarm converge on the optimal solution. With each iteration, the swarm tries a number of different potential solutions and gains information to help it decide the overall best solution. The code for the swarm is relatively straight forward.

Table 1: PSO Pseudo Code

```
Initialize particles randomly
while current iteration < max iterations do
   for each particle j \in swarm do
       //Update the personal and global bests
       if f(j.\vec{x}) > f(j.\vec{y}) then
           j.\vec{(y)} \leftarrow j.\vec{(x)}
       end if
       if f(j.\vec{y}) > f(\hat{y}) then
           (y) \leftarrow j.(y)
       end if
   end for
   for all particles \in swarm do
       Update the Velocity
       Update the Position
   end for
end while
```

For every iteration, all the swarm needs to do it update the values of the personal best for each particle (represented by  $\vec{y}$ ) and update the global best if a new best value is found. This is accomplished by passing the solution set into the fitness function (see below) and getting a value. If you are working on a minimization problem, the smaller of the two will become the personal best, otherwise you will take the larger value. It then needs to iterate through each particle updating it's velocity and position. It is important to complete the global best update before updating the velocity since the global best has an influence on the new velocity. If the global best changes during an iteration, that new global best needs to be reflected in the velocity update in ALL particles, not just the particles that appear afterwards.

#### 2.1.6 Fitness Function

In order to find the best (or optimal) solution to a problem, we first need a way to evaluate those solutions. The Fitness Function is responsible for taking the list of inputs that make up the solution and determining how well it fits the problem. It will then return a single value that helps the swarm determine if it is a good solution or a bad solution, ultimately helping it converge on the best value. The Fitness Function itself is the representation of the problem that we are looking to optimize and is how we tailor the Particle Swarm Optimization algorithm to solve different problems.

#### 2.1.7 Parameters

There are a number of different parameters that can be altered to affect the success of the algorithm. Increasing or decreasing these values can lead to different results depending on the problem at hand. Experimentation is often required in order to get the best performance out of the algorithm.

Max Iterations: This helps determine how many iterations the swarm will go

through in to search for the optimal solution. Increasing this value can potentially lead to better results however is meaningless if the swarm were to have already converged to a local maxima.

Swarm Size: the size of the swarm determines the number of particles it employs to search for the optimal solution. Increasing this gives the swarm more information for each iteration and allows the swarm to have a higher likely hood of finding the overall best solution

Inertial Weight ( $\alpha$ ): the Inertial weight determines how much of an influence the current velocity has on future velocities. A large Inertial weight will lead to the particle carry on in it's current direction allowing it to better search the search space (though potentially heading off towards worse values). A low Inertial weight will allow it to converge towards the best values more quickly but may have a hard time searching the rest of the potential values

Cognitive Weight ( $\omega$ ): the Cognitive weight notes the effect the particles personal best solution has on the future velocity of the particle. Increasing this will keep the particle from venturing too far from this solution and only explore the neighboring space.

**Social Weight** ( $\lambda$ ): this is called the social weight because it determines the effect other particles have on the current particle. This is the influence the global best (or the best solution currently found by the entire swarm) has on the future velocity of each particle. A low value will have the particles mostly work on their own, whereas a large social weight will have all particles converge toward the best value.

## 2.2 Co-operative Particle Swarm Optimization

Although Particle Swarm Optimization is a very successful algorithm for solving continuous problems with reasonable dimensions, it often struggles as those dimensions increase. If the problem has a large enough number of inputs, it can be difficult for the swarm to hone in on an optimal solution given the breadth of the search space size, often leading it to get stuck in local optima. This issue is often referred to as the "curse of dimensionality" and affects most other stochastic optimization algorithms [5].

Much of the impetus of this problem spawns from having all the dimensions update at the same time and then calculating the fitness of the function. This can result in a situation where a a new solution is objectively worse for most dimensions but leads to a better fitness value due to the a better value for a single dimension.[6] For instance, if the intention of the algorithm is to minimize the sum of the values (represented as  $\sum_{n=0}^{d} x_n$ ) the position [5,5,5] will give a fitness value of 15, while the position [6,1,7] returns a better value of 14 even though the first and last dimension values are moving away from the optimal solution.[6] This can lead to the particles heading in the wrong direction, away from the optimal solution.

CPSO attempts to solve the "curse of dimensionality" by instead dividing the problem into multiple sub problems. This allows you to solve a subset of the dimensions separately by converting a single high dimensional problem into a number of low dimensional problems. With each iteration, the particles are then combined together to get a total fitness value. It is important to not update all these particles before gathering the fitness value as you will lead to the same issue as previously. There are a number of ways of doing this split which have been implemented into different variant of the algorithm. A problem which arises from this separation is that it removes the dependencies of variables. For instance, if two dimensions are required to work together to improve the fitness, they will be no longer able to do so directly if they are being updated with different swarms.

#### 2.2.1 CPSO-S

The original version of CPSO is called CPSO-S. Developed by Van den bergh and Engelbrecht [5] it tackles the problem by dividing an n-dimensional problem in to n 1-dimensional swarms. It solves each swarm separately and then merges them to get the fitness value. It does this by creating a randomly generated solution and only updating the part that the current swarm is optimizing. It will continue using this randomly generated solution throughout the algorithm to ensure the other swarms do not affect the convergence of the current swarm. Apart from that, the rest is identical to the standard Particle Swarm Optimization algorithm, however on top of iterating through each particle per iteration, you also need to go through each swarm.

Table 2: CPSO-S Pseudo Code

```
create and initialize n one-dimensional PSO Swarms
Randomly initialize solution vector
while current iteration < max iterations do
   for each swarm i \in swarms do
       for each particle i \in i do
           //Update the personal and global bests
           if f(j.\vec{x}) > f(j.\vec{y}) then
              j.(y) \leftarrow j.(x)
           end if
           if f(j.\vec{y}) > f(\hat{y}) then
               (y) \leftarrow j.(y)
           end if
           Update the Velocity
           Update the Position
       end for
   end for
end while
```

Since 1 dimensional problems are trivial to solve through Particle Swarm Optimization, separating the problem into 1-dimensional problems, it significantly simplifies the end goal. However, since each swarm operates almost in a vacuum, it has no idea of the values obtained by the other problems. This means if one value is dependent or affected by another, there

is no way of that being taken into account. This leads to the swarms getting stuck in a position that is locally optimal for each dimension but is not the global best solution to the problem.

#### 2.2.2 Variant: CPSO- $S_k$

The first variant proposed by Van den bergh and Engelbrecht is  $CPSO-S_k$ . Instead of dividing the problem dimensions into n 1-dimensional swarms, it now evenly divides the problem into k swarms (or as close to even as possible). The basis of this variant is to attempt to combine dimensions that are dependent on each other in hopes to overcome getting stuck in local optima. The rest of the algorithm is identical to CPSO-S. It optimizes those multi-dimensional swarms separately and have them converge on a local optima, ultimately combining the individual solutions as it's final optimization value.

Table 3:  $CPSO-S_k$  Pseudo Code

```
K_1 \leftarrow nmodK
K_2 \leftarrow K - (nmodK)
Initialize K_1[n/K]-dimensional PSOs
Initialize K_2 | n/K |-dimensional PSOs
Randomly initialize solution vector
while current iteration < max iterations do
   for each swarm i \in swarms do
       for each particle j \in i do
           //Update the personal and global bests
           if f(j.\vec{x}) > f(j.\vec{y}) then
               j.(y) \leftarrow j.(x)
           end if
           if f(j.\vec{y}) > f(\hat{y}) then
               (y) \leftarrow j.(y)
           end if
           Update the Velocity
           Update the Position
       end for
   end for
end while
```

However, though this leads to a better chance of matching dependent dimensions, it still tends to get stuck in local optima. This is because the dimension split that exists in the CPSO solutions don't exist in the actual problem. By dividing the problem into sub problems, we are fundamentally altering the initial problem. It is also unlikely that dividing the algorithm into even sections will capture all the connected values and since there is no communication between swarms, it is impossible for the algorithm to break out of those local optima.

#### 2.2.3 Variant: CPSO- $R_k$

One potential solution to the local optima problem would be to randomly divide the problem into differently sized sub problems. This leads to an even larger likelihood of combining dependant variables though that is entirely left up to chance, however, if it does successfully make those connections it can very quickly converge on the optimal solution. CPSO- $R_k$  (also known as Random CPSO) tackles this methodology by dividing up the problem into k randomly sized PSOs. Once the swarms are created, it functions very similar to CPSO- $S_k$ . It solves all the swarms separately, placing them into a temporary solution vector to get the fitness. Once the max number of iterations is reached, it will take the global best of each swarm and combine them into a single solution.

# 2.3 Neighborhood Topologies

The Neighborhood Topology of a PSO dictates how the Swarms' individual particles are allowed to communicate with each other. This communication is the real strength of the PSO as it allows it numbers to work together in solving the problem. This particle communication is achieved through the use of 'Global Best' values. For each particle, when it is updating it's velocity for the next iteration, it will consult the global best in its network to aid in determining it's next move. In the standard PSO (Section ??), all particles are able to communicate with each other. This means that all particles in the swarm are 'led' by the

overall best particle in the entire swarm. This behaviour can sometimes lead to a naive swarm that quickly converges on an 'optimal' solution without full exploring the entire search space of the problem. To combat this, a local PSO [2] was proposed, which utilized different, more constrained, neighborhood topologies in hopes of creating a better PSO .

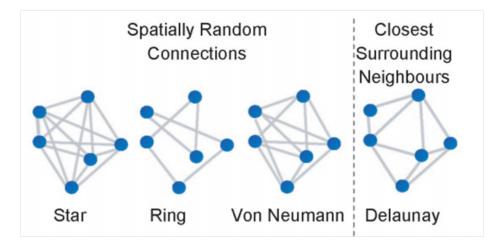


Figure 1: Visual Representation of the Neighborhoods[1]

#### 2.3.1 Common Spatially Random Topologies

There are a number of Neighborhood Topologies that are currently being used and tested in the PSO community. Many of these topologies have their own pros and cons and can lead to different or better results depending on the problem in question. These topologies are known as 'Spatially Random' since the location of the particles at any point in the calculation has no bearing on what particles it is connected to.

#### Star

The Star Topology is the most common since it is used in the standard PSO. As mentioned in the intro, Star Topology connects every particle with each other. That means the 'neighborhood best' for every particle is the same as the 'global best' position in the swarm. The Star Topology uses this information to quickly converge onto the optimal best without a thorough exploration of the search space.

#### Ring

The Ring Topology changes the network topology by reducing the connections per particle to just two. All the particles are ultimately still connected to each other, but through multiple intermediary particles. This topology slows the transfer of information among the particles, decreasing the impact the global best has on the swarm itself[2]. This helps to slow down the convergence of the swarm, allowing a more thorough investigation of the search space. However, as a result, it requires a larger number of iterations to find an optimal solution. The connections are often determined randomly during the initialization of the particles.

#### Von Neumann

The Von Neumann Topology expands on the Ring Topology by increasing the number of connections each particles has while still maintaining the slow information transfer rate within the network. This helps to speed up convergence while allowing for exploration. The particles are connected in a cubic-lattuce arrangement where particles are connected to particles on the left, right, above and below. This isn't, however, determined by actual location of the particle, but often based on the index value that particle is in the Swarm array, where above and below are determined by some sort of offset based on the swarm size (See Figure ??). This means it does not take advantage of positioning information and the connections are generated only in the initialization phase.

#### 2.3.2 Spatially Meaningful Neighbors

Where the updates to the Spatially Random Neighborhood Topologies allowed for more time in searching the search space, the actual searching ultimately ends up being mostly random

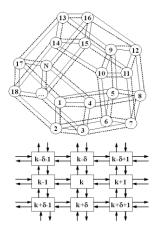


Figure 2: Von Neumann[2]

movements. This is mostly due to the fact that these connections are both randomly selected and only determined once during the life cycle of the swarm. If the connections were dynamically chosen based on information within the swarm, that information could be used to better direct the collaboration between particles. In the paper "Particle Swarm Using Spatially Meaningful Neighbors" [1] Lane, Engelbrecht and Gain proposed using the distance between particles to determine which particle is connected to which. It was theorized that particles that are close together have a high probability that they are searching the same subspace (a pocket of space that has it's own local max/min). The proposed a new Neighborhood Topology called 'Spatially Meaningful Neighbors' that uses the Delaunay Triangulation algorithm to connect all particles with the neighbors that are closest to them. Those connections then represent which particles they can talk to and draw information from. Decisions are then made to avoid or proceed in searching a space based on how the particles are moving (See Section ??). This led to a more guided search and a higher success rate in their tests. [1]

# 2.4 Dynamic Connections

Just having the list of the two closest neighbours for each particle isn't in itself inherently helpful for Particle Swarm Optimization. Of this list, we now want to determine which of

these particles are are working together so we can have the particles cooperate to find a local optima. Setting up connections between these particles will allow them to communicate and share information, allowing them to more quickly converge on that optima. A particle is considered to be working with another if:

- Particle  $P_1$  is following behind Particle  $P_2$ . i.e. both are heading in the same direction. This this case, we'd want to establish a directed connection so the particle in front can send information to the trailing particle
- Both Particle  $P_1$  and Particle  $P_2$  are heading towards each other but not passing each other. Here we would want to set up an undirected connection so that both particles can share information.

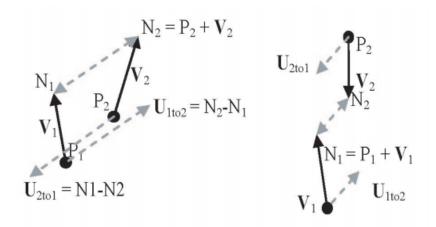


Figure 3: How to know if the particles are working together[1]

You can test both cases with a simple equation:

$$V_1 \cdot U_{1to2} > 0$$

In the function above, the  $V_1$  represents the current velocity of  $P_1$  and  $U_{1to2}$  is the vector representing the distance from  $P_1$  to  $P_2$  after the two Particles next position update. Using the dot product on both these vectors allows us to determine the angle between them. If

the dot product is greater than 0, we know the angle between the two vectors is less than 90 degrees. This outcome can only occur if it satisfies one of the two.

A caveat to this would be that we still wish to encourage local exploration. Delaunay Triangulation does not guarantee that the particles are close together. If two neighbouring Particles are working together, but are still reasonably far away, having them communicate will inhibit local exploration. Therefore we need to set a threshold in which we want the particles to search within before they start communicating. 4 details the process of finding the best neighbour to communicate with.

Table 4: Choosing The Best Neighbour

```
input: N_i Particle i's closest neighbours
procedure CHOOSEBESTNEIGHBOUR(N_i)
   hasConnectedNeighbours = false
   P_f \leftarrow min(N_k)
   for k \leftarrow 1 toneighbourset_size do
       if working_t ogether(P_i, P_k) and dist(X_i - P_c) < dist(X_i - P_k) and f(P_k) < f(X_i)
then
           P_c \leftarrow P_k
           hasConnectedNeighbours \leftarrow true
       end if
   end for
   localExploitationRation \leftarrow swarm.diameter/200
   if hasConnectedNeighbours and
distance(X_i - P - c)/distance(X_i - P - f) > localExploitationRatio and
V_i < 2 * distance(X_i - P_x) then
       P_n \leftarrow P_c
   else
       P_n \leftarrow P_f
   end if
return P_n
end procedure
```

## 2.5 Delaunay Triangulation

#### 2.5.1 Triangulation

Triangulation in Computational Geometry is the decomposition of a polygon or set of points into a set of pairwise non-intersecting triangles in which all points in the set or edge in the polygon are an edge in at least one of newly the formed triangles. In n-dimensions, the triangles are instead represented as simplices (the n-dimensional equivalent). Triangulation is mainly used in computer graphics to simplify complex shapes into easily drawn triangles, though the information gathered from this process can be utilized in a number of different problems. An optimal triangulation is one that minimizes the total length of all the edges, otherwise known has having the minimum weight. What this means is, an optimal triangulation will, on average, connect all the points to their nearest neighbours. Definitively finding the optimal triangulation of a set of points is considered to be an NP-Hard problem, however, there are a number of triangulation algorithms that do a very good job of coming close to optimal.

#### 2.5.2 Delaunay Triangulation

Delaunay Triangulation (DT) is considered to be one of the most efficient triangulation methods in computational geometry and often leads to optimal or close to optimal triangulations. This triangulation methodology creates a set of triangles such that the circumcircle of each triangle contains no other points. What this means is that if one were to draw a circle around every triangle where each point of the triangle lies on the circumference of that circle, no points will fall inside another triangles circle. This guarantees the triangles are evenly spaces out and often results in a lower number of 'thin' triangles.

The issue with Delaunay Trinagulation is one of dimensionality. In 2-dimensions, Delaunay Triangulation can be done with a worst-case time complexity of just O(nlogn) which is one of the quickest methods for pseudo-optimal triangulation. However, it becomes in-

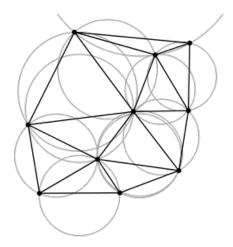


Figure 4: Circumcircles drawn amongst the Triangles[3]

creasingly difficult to compute as dimensions increase making it a less viable solution to all versions of problem. The majority of algorithms that compute Delaunay Triangulation are typically limited to just 2 and 3 dimensions. One way of computing DT in higher dimensions would be to increase the dimension by 1 (giving the value of the new dimension as  $|p|^2$  and then computing the convex hull on those sets of points. However, as the dimensions increase even that becomes increasingly complex to compute. In 4-dimensions the best case algorithm will have a worst-case  $O(n^3)$  run time and in d-dimensions, it takes  $O(n^{\lceil d/2 \rceil + 1})$  time.

# 3 Literature Review/Previous work

#### 3.1 CPSO

In van den Bergh and Engelbrecht's 2004 paper "A Cooperative Approach to Particle Swarm Optimization", they introduced a new concept called Cooperative Particle Swarm Optimization (CPSO). They noted that, like other stochastic optimization algorithms, PSO suffered from the "curse of dimensionality". This references the increasing difficulty of finding an optimal solution as dimension increases as the potential search space increases exponentially. Taking note of Potter's dimensional decomposition to help solve this problem in Genetic Algorithms (another stochastic optimization algorithm), they attempted to do the same for

Particle Swarm Optimization and in the process, solving the issue of pseudominima, in which each swarm gets stuck in a local optima which isn't globally optimal.

Their approach was to create two new PSO algorithm variants. The first was CPSO- $S_k$ , which took the idea of dividing the problem into smaller dimensions, but instead of n 1-dimensional sub problems, it is divided into k problems of equal size. This is in direct response to the issue of pseudominima, which spawns from the issue of each swarm optimizing completely separately. The division is now done in hopes that the dependent variables are grouped together in the same smaller sub swarms. The second proposed variant was CPSO- $H_k$  (Hybrid CPSO). This took advantage of PSOs ability to overcome local optima and CPSO's ability for quick convergence by having both CPSO- $S_k$  and a standard PSO run simultaneously. The two algorithms would share information and help each other find a global optima

Their tests revealed that both variants of CPSO vastly outperformed the standard Particle Swarm Optimization algorithm in higher dimension problems. Each algorithm was tested against a number of standard optimization formulas with 30 dimensions each and not only did they see an increase in terms of the quality of solutions returned, but also in the robustness of the solutions (meaning they were able to reach an optimum solution prior to hitting the max loops). CPSO algorithms were able to converge onto solutions significantly faster than the standard PSO while still leading to global optimums. The problem of pseudominima still persisted in the CPSO- $S_k$ , though at a lesser degree, however the other proposed variant CPSO- $H_k$  performed much better in this case. While the Hybrid variant required more memory space and processing power to compute, it overall was able to overcome local minima better than its counter part and overall led to the best solutions.

Overall, they found that CPSO is an essential tool for solving optimization problems in higher dimensions and leads to objectively better results than it's non-cooperative counterparts

## 3.2 Delaunay Triangulation

Particle Swarm Optimization is a powerful algorithm for optimizing continuous non-linear function by having Particles swarm around the search space to help converge on the global optimum value. One way PSO can be improved upon is by introducing communication techniques to allow the swarms to more directly work together and 'team up' to better evaluate a local search space. A number of Particle knowledge-transfer topologies have been introduced in the past, however, in Lane, Engelbrecht and Gain's paper titled "Particle Swarm Optimization with Spatially Meaningful Neighbours" They propose a different topology utilizing Delaunay Triangulation's ability to easily locate each particles closest neighbour. In this paper, they propose the idea that trading information amongst close particles allow for better use of the local search space.

Delaunay Triangulation is a quick and effective way of finding neighbours in low dimensions that typically represent the closest neighbours to each particle. The algorithm returns triangles which make use of all the particles and also has approximately the minimum-weight (or shortest overall edge lengths). Lane, Engelbrecht and Gain use this information along with a number of other heuristics to determine which particle should communicate with which for best result. They noted that having particles that are 'working together' share information allowed them to collaborate to better explore the smaller local area. Particles were defined as working together if the were headed either in the same direction or towards each other. They also noted that communication with the optimal neighbor was also effective as long as the neighbor wasn't too far way.

After testing their theory, they were able to conclude that Delaunay Triangulation for information sharing between particles led to much better results in low dimensions. Its success rates were almost uni-formally higher than the other knowledge transfer approaches, meaning it was more likely to find the global optimum solution. The problem arose with larger dimension problems. As the dimension rose, the Delaunay Triangulation increasingly became the algorithms bottleneck. In 2 or 3 dimensions, Delaunay Triangulation is a quick

way of determining the particle's nearest neighbours. In larger dimensions, however, it becomes significantly harder to compute and is therefore less helpful for use in Particle Swarm Optimization

Lane, Engelbrecht and Gain go on to conclude that the use of spatially significant neighbours to knowledge transfer is comparatively successful to the other common knowledge transfer topologies. The use of Delaunay Triangulation, however limits the approach to just 2 and 3 dimensions.

# 4 CPSO using Delaunay Triangulation

Building off of the previous work, it is clear the major downside to the use of Delaunay Triangulation spawns from its issues in higher dimensions. In the past, such issues were solved by utilizing CPSO's ability to break the problem into a number of smaller sub problems, as seen with PSO's "curse of dimensionality". Knowing that it has been known to lead to better results, it would be worth trying Delaunay Triangulation in conjunction with CPSO algorithms to see if they can both benefit from each others skills to better solve problems in n-dimensions. It was stated that the tipping point for this benefit is for problems greater than 3 dimensions, as such, all CPSO iterations have been initialized such that no one swarm is handling more than 3 dimensions.

# 4.1 Experimental Setup

To test this theory, a number of known CPSO algorithms were implemented with the added functionality of the communication via spatially significant neighbours using the Delaunay Triangulation approach. The following algorithms were tested for this experiment:

• **PSO**: For comparison's sake, all CPSO variants will be compared to the standard PSO algorithm to ensure either of the changes lead to a meaningful benefit. All the values will be identical to their CPSO counterparts.

• **CPSO-S**: This variant restricts each swarm to just one dimension each particle. Since Delaunay Triangulation does not exist in 1-dimension, it was not used (nor required) for finding the closest neighbours. The rest of the algorithm, however, was still applicable.

- CPSO-S<sub>k</sub>: Due to Delaunay Triangulation being limited to just 2 and 3 dimensions, a
  k needed to be selected in which the swarm is divided into 1, 2 or 3 dimension chunks.
  For the purposes of this test, swarms were divided to maximize the dimensions for each swarm. (k = dimensions/3)
- CPSO- $\mathbf{R}_k$ : The Algorithm was altered to ensure all the random swarms are selected to only be between 1 and 3 dimensions to ensure the spatially significant neighbours algorithm works. The number of swarms is set to half that of the overall dimensions to allow for an even distribution of 1 dimension and 3 dimensional swarms.

# 4.2 CPSO configuration

The CPSO and PSO parameters were selected to match those tested in both the Delaunay Triangulation and the CPSO paper in attempt to replicate the results achieved in those experiments. As per the other experiments, the number of iterations were capped at 10,000. To keep the playing field even, the particles allotted to each algorithm will be distributed evenly amongst each swarm. This means that each iteration for the swarm will contain the exact same number of fitness calls, allowing us to directly compare their results. Each experiment was performed 50 times; the reported results are the average of the global/network bests after all of the 50 runs. The experiments were also repeated for each type of swarm using 15, 20, and 25 particles for 6 dimension problems and 75,100,125 for 30 dimension problem also with and without the use of spatially significant neighbours. All variants of CPSO used the following input values:

- Social Weight ( $\lambda$ ): 1.49
- Cognitive Weight ( $\omega$ ): 1.49

• Inertial Weight ( $\alpha$ ): starts at 1 and linearly decreases with each iteration, ultimately hitting 0 on the last iteration

- Max Velocity  $(V_{max})$ : clamped to the domain of the function
- Max Position ( $P_{max}$ ): also clamped to the domain of the function
- Particle Count: particles are evenly distributed among the swarms in the algorithm. for example, if the particle count is set to 100 and the CPSO has 10 different swarms, each swarm will only have 10 particles to solve its particular subproblem.

#### 4.3 Benchmark functions

Again, in an attempt to replicate the results of the previous experiments, for comparisons sake, the same benchmark functions were selected here as were in the two previous papers.

The functions which were selected are as follows:

Function	Domain	Criterion
Rastrigin	[-5.12;5.12]	0.01
Rosenbrock	[-30;30]	100
Griewanck	[-600;600]	0.05
Ackley	[-32;32]	0.01

Table 5: Functions, Stop Criteria and Domains

The "Success Rate" evaluation indicates the percentage of runs that reached the end criterion and the "Iterations to Criterion" value returns the average number of iterations it took to reach that stopping point.

The functions themselves are as follows:

#### Rastrigin:

$$\sum_{i=1}^{n} x_i^2 + 10 - 10\cos(2\pi x_i)$$

Rosenbrock:

$$\sum_{i=1}^{n-1} 100(x_{i+1} - x_i^2)^2 + (x_i - 1)^2$$

Griewanck:

$$\frac{1}{4000} \sum_{i=1}^{n} x_i^2 - \prod_{i=1}^{n} \cos(\frac{x_i}{\sqrt{i}}) x_i + 1$$

Ackley:

$$20 + e - 20e^{-0.2\sqrt{\frac{\sum_{i=1}^{n} x_i^2}{n}} - e^{\frac{\sqrt{\sum_{i=1}^{n} \cos(2\pi x_i)}}{n}}$$

## 4.4 Experiment 1: Robustness Test on Lower Dimensions

The mark of a good PSO algorithm is one that can hit or get very close to the optimal result on a consistent basis. As such, the robustness of the algorithm needed to be tested as well. In this case, "Robustness" refers to the number of times and the speed at which the algorithm hit the exit threshold over the course of 50 runs. The tables presented list the algorithm being tested first, followed by the number of particles used in that calculation. Following that is the percentage of successful runs with the average iterations without Spatially Meaningful Neighbors followed by the algorithm with. The best runs are those that we successful the most times while taking the least number of iterations to achieve it. Note: if the run was not successful, it's iteration count does not factor into the average iterations. The tests were run against problems of 6-dimensions to see if Delaunay Triangulation led to noticeable improvements with those results.

#### **4.4.1** Results

#### 4.4.2 Discussion

The Robustness metric shows how successful the swarms are a finding optimal solutions and in comparison, the effect the Spatially Meaningful Neighborhood Topology has on this success value. In most cases, we can see that the number of iterations to reach the threshold (note: unsuccessful searches are not counted toward this) is larger in the new topology. This has to do with the idea that local search spaces are investigated more thoroughly before advancing. This difference can be a single digit to multiple hundreds of extra iterations as

Algorithm	s	Standard		Spatially Meaningful Neighbors	
	Б	Succeeded	Iterations	Succeeded	Iterations
	15	54	1291		
PSO	20	63	1264		
	25	60	1185		
	15	100	837	100	772
CPSO-S	20	100	597	100	624
	25	100	462	100	465
	15	90	1091	92	1191
$\mathbb{CPSO-S}_2$	20	98	1033	100	1165
	25	98	969	100	1145
	15	58	1173	74	1295
$\mathbb{CPSO-R}_3$	20	80	1110	80	1312
	25	88	1032	84	1314

Table 6: Rastrigin Robustness Analysis (n=6)

Algorithm	S	Stand	dard	Spatially M	eaningful Neighbors
		Succeeded	Iterations	Succeeded	Iterations
	15	98	346		
PSO	20	100	308		
	25	100	350		
	15	100	287	100	356
CPSO-S	20	100	128	100	116
	25	100	106	100	96
	15	100	284	100	316
$\mathbb{CPSO}$ -S <sub>3</sub>	20	100	153	100	156
	25	100	90	100	85
	15	94	301	98	386
${f CPSO-R_3}$	20	98	392	100	415
	25	100	260	100	275

Table 7: Rosenbrock Robustness Analysis (n=6)

it very infrequently less than it's standard PSO counterpart. As a result, however, we can see that in general, the success ratio benefits from this longer search. In a number of cases, especially for the Griewanck function ??, there is a large boost of 10%-20% from the change in the neighborhood topologies. This, however is not always the case. There are a number of instances where the new topology leads to around 2% fewer successes (or one extra failure). These discrepancies can often potentially be explained by the random nature of the initial

Algorithm	G	Standard		rd Spatially Meaningful Neighbor	
rigoriumi		Succeeded	Iterations	Succeeded	Iterations
	15	40	1323		
PSO	20	28	1269		
	25	22	1222		
	15	98	980	100	976
CPSO-S	20	100	724	96	759
	25	98	647	100	733
	15	82	1220	98	1465
$\mathbb{CPSO}$ -S <sub>3</sub>	20	84	1208	100	1440
	25	94	1104	100	1396
	15	62	1301	82	1531
$\mathbb{CPSO-R}_3$	20	66	1247	86	1581
	25	80	1230	98	1700

Table 8: Griewanck Robustness Analysis (n=6)

Algorithm	s	Standard		Spatially Meaningful Neighbors	
	Б	Succeeded	Iterations	Succeeded	Iterations
	50	100	1375		
PSO	75	100	1287		
	100	100	1262		
	50	94	1223	94	1272
CPSO-S	75	100	1115	100	1137
	100	100	1014	100	1013
	50	100	1261	100	1263
$\mathbb{CPSO-S}_3$	75	100	1202	100	1237
	100	100	1148	100	1189
	50	96	1351	94	1365
$\mathbb{CPSO-R}_3$	75	100	1309	100	1299
	100	100	1211	100	1267

Table 9: Ackley Robustness Analysis (n=6)

particle generation. What is more concerning is it seems to have the exact opposite effect for the Rastrigin function. Runs using the new Topology performed constantly much worse. A Hypothesis on why this might be could be because the search space domain is so small as is that the proximity benefit becomes mostly moot. The algorithm tends to struggle with a large number of floating point calculations and with particles as close as those, it requires a lot of floating point calculations to form the Delaunay Triangulation.

# 4.5 Experiment 2: Robustness Test on Higher Dimensions

The issue with the Delaunay Triangulation method as originally pitched was that it did not scale well to higher dimensions. At dimensions higher than 3 it became too slow for the benefit gained to be worth it. By in limiting the Delaunay Triangulation calculations to just 3-dimensions, we can keep the run time of said calculation to just  $O(n^2)$ . In the test, be use the Delaunay Triangulation method with CPSO to solve problems of dimension 20. This drastically reduces the run time of calculating 20 dimensions, even if the calculation needs to be performed 20 times per iteration. This test aims to see if it still has the same effect on the success rate as experienced in the previous section.

Algorithm	S	Standard		Spatially Meaningful Neighbors	
	5	Succeeded	Iterations	Succeeded	Iterations
	50	28	1434		
PSO	75	44	1351		
	75	62	1236		
	50	94	986	94	1012
CPSO-S	75	100	800	100	827
	100	100	639	100	624
	50	40	1377	36	1397
$\mathbb{CPSO}$ -S <sub>3</sub>	75	40	1373	32	1412
	100	62	1342	62	1350
	50	34	1324	26	1458
$\mathbf{CPSO} ext{-}\mathbf{R}_3$	75	56	1244	56	1548
	100	74	1152	48	1485

Table 10: Rastrigin Robustness Analysis (n=20)

#### 4.5.1 Discussion

After promising results in lower dimensions, tests were performed to see how well the algorithm scales to higher dimension problems. The problem set was increased from 6-dimensions to 20-dimensions to see the effect this has on the algorithms 'robustness'. The number of particles and swarms were uniformly increased with the dimension to ensure that all CPSO variants have enough particles to adequately search the search space. CPSO-S, for

Algorithm	S	Standard		Spatially Meaningful Neighbors	
	Б	Succeeded	Iterations	Succeeded	Iterations
	50	100	366		
PSO	75	100	300		
	100	100	250		
	50	100	653	94	605
CPSO-S	75	100	382	100	393
	100	100	259	100	187
	50	100	355	100	412
$\mathbb{CPSO-S}_2$	75	100	286	100	295
	100	100	254	100	272
	50	100	625	100	704
$\mathbb{CPSO-R}_3$	75	100	542	100	708
	100	100	419	100	611

Table 11: Rosenbrock Robustness Analysis (n=20)

Algorithm	s	Stand	dard	Spatially M	leaningful Neighbors
Aigorithin	٥	Succeeded	Iterations	Succeeded	Iterations
	50	92	1240		
PSO	75	100	1118		
	100	98	1064		
	50	88	1016	89	1052
CPSO-S	75	100	382	100	393
	100	74	705	80	714
	50	96	1210	100	1213
$\mathbb{CPSO}$ -S <sub>2</sub>	75	98	1069	98	714
	100	100	1247	100	1325
	50	42	1588	44	1588
$\mathbb{CPSO-R}_3$	75	50	1286	52	1570
	100	40	1240	64	1689

Table 12: Griewanck Robustness Analysis (n=20)

instance, would then need to be divided into 20 swarms so ensuring each swarm had at the very least, 2 particles was key. The results show something of a decrease in the benefit from the Topology. Part of the reason for this is the increase in particles led to all the variants have an adequate number to optimized the solution. The cases where it wasn't 100% saw a small increase in effectiveness but not a large enough boost the definitively conclude that it was solely due to the use of spatially meaningful neighbors. Additionally, the problem with

Algorithm	s	Standard		Spatially Meaningful Neighbors	
	Б	Succeeded	Iterations	Succeeded	Iterations
	50	96	1507		
PSO	75	100	1387		
	100	100	1295		
	50	100	1254	100	1262
CPSO-S	75	100	1168	100	1147
	100	100	997	100	992
	50	96	1524	96	1650
$\mathbb{CPSO-S}_2$	75	100	1376	100	1530
	100	100	1324	100	1422
	50	94	1426	100	1426
$\mathbb{CPSO-R}_3$	75	100	1290	100	1318
	100	100	1217	100	1246

Table 13: Ackley Robustness Analysis (n=6)

the Rastrigin algorithm is also present in higher dimensions and leads to the direct opposite effect of the benefit given to other functions. It must also be noted that the run-time of the Delaunay Triangulation version of the algorithms also ran exceedingly slow. This is mostly due to the number of times we are calling that Delaunay Triangulation (DT) algorithm. Since the swarm is divided in n/3 or n/3 sub swarm, each of which needs to calculate the DT with each iteration. This leads to an extra  $O(n*n^2)$  algorithm being run every iteration of the CPSO. Though it is not as time consuming as  $O(n^{\lceil d/2 \rceil + 1})$ , especially on significantly large dimensions, it is also not an insignificant use of resources.

# 5 Conclusion

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