

Grammar

- Ambiguity
- Precedence and associativity

Equivalent

How to prove $L(G1) = L(G2)$?

- Simplify
- Chomsky normal form
- Push down automata

Parsing

(Terminal) Left to right

	Leftmost	Rightmost Derivation
Order	Top-down	Bottom-up
Example	Recursive Descent, LL(1)	LR(0), SLR(1), LR(1)
Strategy	Predict-match	Shift-reduce [deterministic]

LL(1)

- Why M?
 - First Set, Follow Set
- And how to write the above formally?
- Why it has non-determinism for some Grammar?
 - Left recursion

$$A \Rightarrow^+ A\alpha$$

- Add EOF to S
- Add extra $D = \text{then } C \mid \text{epsilon}$

LR(0)

- Why it has non-determinism for some Grammar?
 - shift-reduce conflict
 - both *shift* $A \rightarrow \alpha \bullet a \beta$ and *reduce* $B \rightarrow \beta \bullet$, $A, B \in N$
 - and no matter whether $a \in FOLLOW(B)$ in an NFA state set.

$$S \rightarrow L \bullet := R$$

$$R \rightarrow L \bullet$$

- reduce-reduce conflicts
 - see below
- NFA

SLR(1)

- LR(0) table structure
 - same parser operation (shift/ reduce)
 - **one token of lookahead**
 - to arbitrate among shift-reduce conflicts
 - **DFA (less non-determinism allowed)**
- How to construct a Full DFA?
 - Do it directly with epsilon closure. ✓
 - Power set / subset construction for converting NFAs to DFAs. (Time consuming)
- Why it has non-determinism for some Grammar?
 - shift-reduce
 - both *shift* $A \rightarrow \alpha \bullet a \beta$ and *reduce* $B \rightarrow \beta \bullet$, $A, B \in N$
 - and $a \in FOLLOW(B)$ in an NFA state set.

$$S \rightarrow L \bullet := R$$

$$R \rightarrow L \bullet, (:=) \in FOLLOW(R)$$

- reduce-reduce
 - both *reduce* $A \rightarrow \alpha \bullet$ and *reduce* $B \rightarrow \beta \bullet$, $A, B \in N$
 - and $\exists a. a \in FOLLOW(A)$ and $a \in FOLLOW(B)$ in an NFA state.
 - In particular, hold if $A = B$.

LR(1)

- *shift* $A \rightarrow \alpha \bullet a \beta, b$ and *reduce* $B \rightarrow \beta \bullet, b$
 - **reduce** only if the next token is exactly terminal b as $b \in first(B)$ rather than any of those in $Follow(B)$ in SLR(1) parsing.

$$R \rightarrow L \bullet, b$$

$$S \rightarrow L \bullet := R, b$$

- but causing more complex DFA

Formal Model of Language

CSG

- Earley Parser

- Chart Parsing
- Dependency Grammar/Parsing
- Categorical Grammar/Parsing