

type expr = | Integer of int | Pair of expr * expr | Apply of string * expr

type value = | INT of int | PAIR of value * value

let rec eval = function

| Integer n \rightarrow INT n

| PAIR (e1, e2) \rightarrow PAIR (eval e1, eval e2)

| Apply (f, e) \rightarrow eval-function(f, eval e)

let v = eval e in

C(eval-function(f, v))

1. Add a continuation parameter c to each function, return value.

(a)

let rec eval-cps c = function

| Integer n \rightarrow C (INT n)

| Pair (e1, e2) \rightarrow eval-cps (fun v1 \rightarrow (*PAIR 1*) v1, e2, c)
eval-cps (fun v2 \rightarrow (*PAIR 2*) v2, v1, c) (PAIR (v1, v2)) e2
e1) \rightarrow C (PAIR (v1, v2)) e2

| Apply (f, e) \rightarrow eval-cps (fun v \rightarrow C (eval-function(f, v)) e apply-cnt
(*FUNC*) - 2 f, c

Thus

let eval-2 e = eval-cps (fun x \rightarrow x) e. (* ID *) ... part b.

eval: expr \rightarrow value.

eval-cps: expr \rightarrow value

(IH) C (eval e) = eval-cps c e

(b) Eliminate high-order continuations.

1. Add a constructor to cnt for each fun (* CNT *) $\lambda \dots$ (free variables)

type cnt =

| ID

| PAIR1 of expr * cnt

| PAIR2 of value * cnt

| FUNC of string * cnt

Call apply-cnt at every application of continuation.

let rec eval-cps-dfn c = function

| Integer n \rightarrow apply-cnt C (INT n)

| Pair (e1, e2) \rightarrow eval-cps-dfn (PAIR1 (e2, c)) e1

| Apply (f, e) \rightarrow eval-cps-dfn (FUNC (f, c)) e

and apply-cnt = function

| (ID, v) \rightarrow v

| (PAIR1 (e2, c), v1) \rightarrow eval-cps-dfn (PAIR2 (v1, c)) e2

| (PAIR2 (v1, c), v2) \rightarrow apply-cnt (c, PAIR (v1, v2))

| (FUNC (f, c), v) \rightarrow apply-cnt (c, eval-function(f, v))

let eval-3 e = eval-cps-dfn ID e.

fun v1 \mapsto eval-cps ... e2
fun v2 \mapsto C (PAIR (v1, v2))
fun v \mapsto C (eval-function(f, v))

Mutually recursive.

eval-cps-dfn: cnt \rightarrow expr \rightarrow value

apply-cnt: cnt * value \rightarrow value

eval-3: expr \rightarrow value