

# y2017p394. CPS & Defunctionalisation

type expr = | Integer of int

type value = | INT of int

| Pair of expr \* expr

| PAIR of value \* value

(\* apply a named function: APP \*)

Apply of string \* expr

let rec eval = function

| Integer n → INT n

| PAIR (e1, e2) → PAIR (eval e1, eval e2)

| Apply (f, e) → eval-function(f, eval e)

let v = eval e in

C(eval-function(f, v))

1. Add a continuation parameter c to each function, return value.

(a)

let rec eval-cps C = function

| Integer n → C (INT n)

| Pair (e1, e2) → eval-cps (fun v1 → (\* PAIR 1 \*) v1, e2, C) e1  
eval-cps (fun v2 → (\* PAIR 2 \*) v2, v1, C (PAIR (v1, v2))) e2

| Apply (f, e) → eval-cps (fun v → C(eval-function(f, v)) e, (\* FUNC \*) - λ f, c.

Thus

let eval-2 e = eval-cps (fun x → x) e. (\* ID \*)

(b) Eliminate high-order continuations.

1. Add a constructor to cnt for each fun (\* CNTN \*) λ... (free variables)

type cnt =

| ID

| PAIR1 of expr \* cnt

| PAIR2 of value \* cnt

| FUNC of string \* cnt

Call apply-cnt at every application of continuation.

let rec eval-cps-dfn c = function

| Integer n → apply-cnt C (INT n)

| Pair (e1, e2) → eval-cps-dfn (PAIR1 (e2, c)) e1

| Apply (f, e) → eval-cps-dfn (FUNC (f, c)) e

And apply-cnt = function

| (ID, v) → v

| (PAIR1 (e1, c), v1) → eval-cps-dfn (PAIR2 (v1, c)) e1

| (PAIR2 (v1, c), v2) → apply-cnt (c, PAIR (v1, v2))

| (FUNC (f, c), v) → apply-cnt (c, eval-function(f, v))

let eval-3 e = eval-cps-dfn ID e.

Mutually recursive.