



## Problem 2 Check if Sorted Without Using Sort (20 points)

Dr X claims to have an algorithm that takes an input sequence  $S$  and produces an output sequence  $T$  that is a sorting of the  $n$  elements in  $S$ . However, you are in doubt about its correctness. Write a python function `is_sorted(T, S)`, that tests in  **$O(n \log n)$**  time in the worst case if  $T$  is a sorting of  $S$  (8pts for the runtime).

- Assume both  $S$  and  $T$  are lists of integers, and they have the same length  $n$
- Assume that there are no duplicated elements in  $S$  or  $T$ .
- No sorting functions are allowed in your solution.
- Hint: write your own binary search function.

Example 1:

Input: `is_sorted([1, 2, 3, 4, 5], [2, 4, 3, 1, 5])`  
Return: True

Example 2:

Input: `is_sorted([1, 2, 3, 5, 9], [2, 4, 3, 1, 5])`  
Return: False

Example 3:

Input: `is_sorted([1, 2, 3, 5, 4], [2, 4, 3, 1, 5])`  
Return: False

Example 4:

Input: `is_sorted([5, 4, 3, 2, 1], [2, 4, 3, 1, 5])`  
Return: True

### Problem 3 All Possible Sum Finder (28 points)

You are given two Python lists  $A$  and  $B$  both of length  $n$ , each storing integers **in increasing order**. Given an integer  $m$ , implement a Python function `find_sum_m(A, B, m)` such that it returns **all** pairs of elements  $(a, b)$ :  $a \in A$  and  $b \in B$ , such that  $a + b = m$ , as a python list of tuples.

- If there are duplicates of  $a \in A$  that all match to one (or duplicates of)  $b \in B$ , then return one pair  $(a, b)$  (see example below):
- If there is no such pair, then return an empty list.
- Your program must have **runtime complexity:  $O(n)$**  in the worst case (8pts for the runtime),
- and **memory complexity:  $O(1)$**  (8pts for the memory complexity).

Example 1:

```
A = [-1, 4, 5, 6, 8, 10, 12], B = [0, 1, 2, 4, 9, 10, 20]
Input: find_sum_m(A, B, 14)
Return: [(4, 10), (5, 9), (10, 4), (12, 2)]
```

Example 2:

```
A = [-1, 4, 5, 6, 8], B = [0, 1, 2, 4, 10]
Input: find_sum_m(A, B, 100)
Return: []
```

Example 3:

```
A = [1, 2, 2, 3, 5], B = [1, 98, 98, 99, 99]
Input: find_sum_m(A, B, 100)
Return: [(1, 99), (2, 98)]
```

## Problem 4: Logarithm Calculator (28 points)

If  $m^k = n$  for integers  $m, n > 1$ , and  $k \geq 0$ , we say that  $k$  is the logarithm base  $m$  of  $n$ . In general, this logarithm need not necessarily be an integer. Let us define the **integer logarithm base  $m$  of  $n$**  to be the greatest integer  $k \geq 0$ , such that  $m^k \leq n$ , and denote it by  $k = \lg_m n$ . In the following exercise, we will calculate the integer logarithm function.

The only arithmetical operations you may use are  $+$  and  $*$ . When analyzing the time complexity of these algorithms, you are to count arithmetical operations only, assuming that every operation takes a single time unit.

a) Write a python function **LG1(m,n)** with input integers  $m, n > 1$ , that calculates  $\lg_m n$  by repeatedly calculating the powers  $m^0, m^1, \dots, m^k$ , until a number  $k$  is found satisfying  $m^k \leq n < m^{k+1}$ .

There is no runtime complexity limit for this question. However, you should analyze your implementation's runtime complexity by yourself. Write your program's big **O** complexity **within your .py file in term of  $k$** . (As comment, tightest big O in worst case) (8pts for the runtime),

b) It is well known that each positive integer  $k$  can be written uniquely as a sum of integer powers of 2, i.e., in the form  $k = 2^{l_1} + 2^{l_2} + \dots + 2^{l_j}$ , where  $l_1 > l_2 > \dots > l_j \geq 0$ . For example,

$$12 = 2^3 + 2^2$$

$$31 = 2^4 + 2^3 + 2^2 + 2^1 + 2^0$$

Hence,  $m^k = m^{2^{l_1}} \cdot m^{2^{l_2}} \dots m^{2^{l_j}}$ , and if we need to calculate  $k = \lg_m n$ , it is enough to find the appropriate exponents  $l_1, l_2, \dots, l_j$ .

Design an iterative (i.e., non-recursive) algorithm **LG2(m,n)** to calculate  $\lg_m n$  by first finding an integer  $l_1$  satisfying  $m^{2^{l_1}} \leq n < m^{2^{(l_1+1)}}$ , then finding an integer  $l_2 < l_1$ , satisfying  $m^{2^{l_1}} \cdot m^{2^{l_2}} \leq n < m^{2^{l_1}} \cdot m^{2^{(l_2+1)}}$ , and so on.

- Your program must have **runtime complexity:  $O(\log k)$**  in the worst case (8pts for the runtime), in which  $k = \log n$
- and **memory complexity:  $O(\log k)$**  (8pts for the memory complexity).

### **Problem 5 Why is $O(n^2)$ faster than $O(n \log n)$ sometimes? (8 points)**

Al and Bob are arguing about their algorithms. Al claims his  $O(n \log n)$ -time method is always faster than Bob's  $O(n^2)$ -time method. To settle the issue, they perform a set of experiments. To Al's dismay, they find that if  $n < 100$ , the  $O(n^2)$ -time algorithm runs faster, and only when  $n \geq 100$ ,  $O(n \log n)$ -time one runs faster. Explain how this is possible.

Important:

- You should submit a .txt file for this question.
- If you are submitting this question on gradescope, simply upload the text file, the text file will not get auto graded, TA will manually grade your text file after the deadline.