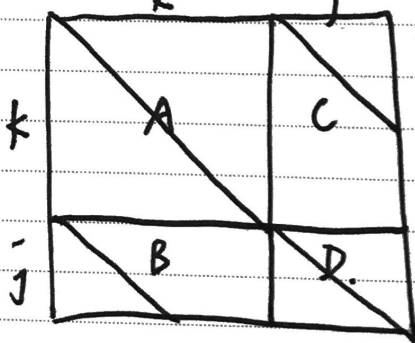


P7.1.8.

Compute the eigenvalues of  $M$ , where  $A, B, C, D$  are diagonal.

$$M = \begin{bmatrix} A & C \\ B & D \end{bmatrix} \begin{matrix} k \\ j \end{matrix}$$

If  $k > j$ , we have



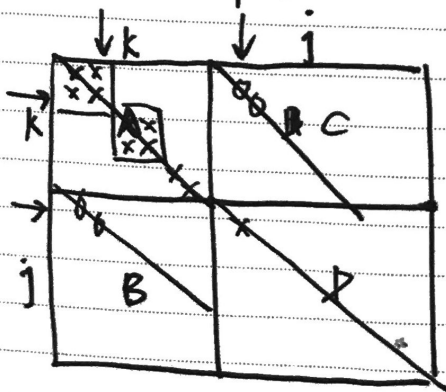
We need to move the diagonal entries of  $B$  and  $C$  to the off-diagonal entries in  $D$  using permutation matrices on both sides of  $M$ , changing the rows and cols at the same time.

$$N = P^{-1}MP \text{ and } M \text{ are similar, } \lambda(M) = \lambda(N).$$

In this way, we can zero out  $B, C$  and fill those entries in  $A, D$ .

And in  $A, D$ , the diagonal will be either  $1 \times 1$  or  $2 \times 2$ , we can compute them easily.

If  $k < j$ , we have swap them



The same method works for  $k < j$ .

After we do the swaps, we can get a nice almost diagonal matrix and compute the  $2 \times 2$  block ones and get the  $1 \times 1$  block ones as eigenvalues.

P3.

$H$  is the sparse connectivity matrix of websites. ( $n \times n$ )

$A$  is the transition possibility matrix. ( $n \times n$ )

$\alpha = 0.85$ .

Let  $C_j$  be the column sums of  $H$ .

$$C_j = \sum_i H_{ij}$$

So if  $C_j \neq 0$ , the probability of row  $i$  page is chosen is

$$\alpha \cdot \frac{H(i,j)}{C_j} + (1-\alpha) \cdot \frac{1}{n}$$

And if  $C_j = 0$ , the probability becomes  $\frac{1}{n}$ .

So  $A$  matrix can be formed as

$$A(i,j) = \begin{cases} \frac{\alpha H(i,j)}{C_j} + (1-\alpha) \frac{1}{n} & , C_j \neq 0 \\ \frac{1}{n} & , C_j = 0 \end{cases}$$

which can be also written as

$$A = \alpha \cdot H \cdot C + \text{ones}(n,1) \cdot u^T$$

where  $C$  is a sparse matrix with the reciprocals of column sums in its diagonal.  $u$  is  $n \times 1$  vector with  $(1-\alpha) \frac{1}{n}$  and  $\frac{1}{n}$  in its entries.

$$u_i = \begin{cases} (1-\alpha) \frac{1}{n} & , C_j \neq 0 \\ \frac{1}{n} & , C_j = 0 \end{cases}$$

And because  $x = Ax = \alpha \cdot \overset{\text{sparse}}{H \cdot C} \cdot x + \overset{n \times 1}{\text{ones}(n,1)} \cdot \left( \overset{n \times n}{H^T} \cdot \overset{n \times 1}{x} \right)$

We can use power method and avoid using huge dense matrix. ↓ scalar