

ELEC5650 Homework2

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May 22, 2019

Problem 1

According to the conclusion we derived in the lecture, we have $R_1 \leq R_2 \Rightarrow g_1 \leq g_2$; and $\forall X, Y$, and $0 \leq X \leq Y$, then $g(X) \leq g(Y)$ for both g_1 and g_2 . We denote

$$\phi_k = E[P_k] = \alpha_k \phi_{k-1} \circ g_1 + (1 - \alpha_k) \phi_{k-1} \circ g_2$$

where $\phi_0 = I$ is the identity map. We define $\delta_k = \alpha_k - \alpha \geq 0$. We use mathematical induction to prove the problem, together with the statement $\phi_k(X) \leq \phi_k(Y)$, $\forall 0 \leq X \leq Y$.

When $k = 0$, $\phi_0 = f_0$, both statements holds.

When $k = 1$, $\phi_1 - f_1 = \delta_1(g_1 - g_2) \leq 0$;

$$\phi_1(X) = \alpha_1 g_1(X) + (1 - \alpha_1) g_2(X) \leq \alpha_1 g_1(Y) + (1 - \alpha_1) g_2(Y) = \phi_1(Y)$$

so both statements holds.

When $k \geq 2$, we assume $\phi_{k-1} \leq f_{k-1}$, and $\forall X \leq Y$, $\phi_{k-1}(X) \leq \phi_{k-1}(Y)$, obviously we have

$$\begin{aligned} \phi_k(X) &= \alpha_k \phi_{k-1} \circ g_1(X) + (1 - \alpha_k) \phi_{k-1} \circ g_2(X) \\ &\leq \alpha_k \phi_{k-1} \circ g_1(Y) + (1 - \alpha_k) \phi_{k-1} \circ g_2(Y) = \phi_k(Y) \end{aligned}$$

$$\begin{aligned} \phi_k - f_k &= \alpha_k \phi_{k-1} \circ g_1 + (1 - \alpha_k) \phi_{k-1} \circ g_2 - \alpha f_{k-1} \circ g_1 - (1 - \alpha) f_{k-1} \circ g_2 \\ &= \alpha(\phi_{k-1} \circ g_1 - f_{k-1} \circ g_1) + (1 - \alpha)(\phi_{k-1} \circ g_2 - f_{k-1} \circ g_2) + \delta_k(\phi_{k-1} \circ g_1 - \phi_{k-1} \circ g_2) \\ &\leq 0 \end{aligned}$$

First two term hold because of the first assumption, and the last term holds due to our second assumption.

Therefore, we have $\phi_k \leq f_k$, for all k .

Problem 2

1) In our lecture, we proved that $\forall X \leq Y, h(X) \leq h(Y), g(X) \leq g(Y)$. To prove $g^i h^i \leq h^i g^i$, we just need to prove $g^2 h \leq h g^2$ and $g h^2 \leq h^2 g$, omitting X .

$$gh \leq hg \implies gh \circ h \leq hg \circ h \implies gh \circ h \leq h \circ gh \leq h \circ hg \implies gh^2 \leq h^2 g$$

$$gh \leq hg \implies gh \circ g \leq hg \circ g \implies h \circ gg \geq g \circ hg \geq g \circ gh \implies g^2 h \leq h g^2$$

2) To find a sufficient condition of the previous statement, we should relax both size towards the statement and irrelevant to X .

$$gh(X) = \tilde{g} \circ hh(X) = [hh(X)^{-1} + BR^{-1}B']^{-1} \leq [BR^{-1}B']^{-1}$$

$$hg(X) = A'g(X)A + Q \geq Q$$

Noted that the matrix inversion lemma is used in \tilde{g} . After eliminating X related terms, we can easily find a sufficient condition for $gh \leq hg$, which is $[BR^{-1}B']^{-1} \leq Q$

Problem 3

Omitting θ , we first define

$$J_k(x) = x'Qx$$

$$f(x_k, u_k, \gamma_k) = x'_k Q x_k + \gamma_k u'_k R u_k$$

$$g(x_k, u_k, \gamma_k) = A x_k + \gamma_k B u_k$$

Then we can solve the schedule optimal by solving the equivalent dynamic programming problem

$$\begin{aligned} J_{k-1}^*(x_{k-1}) &= \min_{u_{k-1}} \{f(x_{k-1}, u_{k-1}, \gamma_{k-1}) + J_k^*(g(x_{k-1}, u_{k-1}, \gamma_{k-1}))\} \\ &= \min_{u_{k-1}} \{x'_{k-1} Q x_{k-1} + \gamma_{k-1} u'_{k-1} R u_{k-1} + (A x_{k-1} + \gamma_{k-1} B u_{k-1})' Q (A x_{k-1} + \gamma_{k-1} B u_{k-1})\} \\ &= \min_{u_{k-1}} \{x'_{k-1} Q x_{k-1} + \gamma_{k-1} u'_{k-1} R u_{k-1} + x'_{k-1} A' Q A x_{k-1} + \gamma_{k-1} x'_{k-1} A' Q A B u_{k-1} \\ &\quad + \gamma_{k-1} u'_{k-1} B' Q A x_{k-1} + \gamma_{k-1}^2 u'_{k-1} B' Q B u_{k-1}\} \end{aligned}$$

When $\gamma_{k-1} = 0$, we have

$$J_{k-1}^*(x_{k-1}) = \min_{u_{k-1}} \{x'_{k-1} Q x_{k-1} + x'_{k-1} A' Q A x_{k-1}\} = x'_{k-1} (A' Q A + Q) x_{k-1}$$

When $\gamma_{k-1} = 1$, according matrix calculus we have

$$\begin{aligned} u_{k-1}^* &= (R + B' Q B)^{-1} B' Q A x_{k-1} \\ J_{k-1}^*(x_{k-1}) &= x'_{k-1} Q x_{k-1} - x'_{k-1} A' Q B [B' Q B + R]^{-1} B Q A x_{k-1} \\ &= x'_{k-1} (Q - A' Q B [B' Q B + R]^{-1} B Q A) x_{k-1} \end{aligned}$$

If we define $P_k = Q$, then we can have

$$\begin{aligned}
J_{k-1}^*(x_{k-1}) &= x'_{k-1} P_{k-1} x_{k-1} \\
u_{k-1}^* &= -L_{k-1} x_{k-1} = (R + B'QB)^{-1} B'QA x_{k-1} \\
P_{k-1} &= \begin{cases} A'P_k A + Q = h(P_k) & , \text{if } \gamma_{k-1} = 0 \\ Q - A'P_k B [B'P_k B + R]^{-1} B P_k A = g(P_k) & , \text{if } \gamma_{k-1} = 1 \end{cases}
\end{aligned}$$

So the problem is equivalent to solve

$$\begin{aligned}
J_0^*(x_{k-1}) &= x'_0 P_0 x_0 \\
P_0 &= f_0 \circ f_1 \circ \dots \circ f_k(P_k) \\
f_i &= \begin{cases} h & , \text{if } \gamma_i = 0 \\ g & , \text{if } \gamma_i = 1 \end{cases}
\end{aligned}$$

From Problem 2, we proved that for d times g and $T - d$ times h

$$g^d h^{T-d} \leq f_0 \circ f_1 \circ \dots \circ f_k \leq h^{T-d} g^d$$

so it is trivial to find the optimal schedule is

$$\gamma_i(\theta) = \begin{cases} 0 & , \text{if } i = d, d+1, \dots, T-1 \\ 1 & , \text{if } i = 0, 1, \dots, d-1 \end{cases}$$