OptCuts: Joint Optimization for Seam Placement and Parameterization of 3D Surfaces

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Parameterizing 3D surfaces to the 2D plane with low mapping distortion is a critical problem in computer graphics with many applications including texture mapping, remeshing, and detail transfer. In most practical cases, it is not possible to map 3D surfaces to the 2D plane without introducing discontinuities (seams) and distortion. Most existing techniques follow a two step process - first placing seams and then minimizing distortion - that usually leads to suboptimal results. In this paper, we attack seam placement and parameterization jointly in a single alternating optimization framework, where seams are optimally and progressively introduced or removed (in topology steps) in between distortion minimization (in descent steps). A linear combination of symmetric Dirichlet energy and seam length are taken as our objective, of which the stationary w.r.t. both UV topology and UV coordinates are guaranteed to be reached within a bounded number of alternating iterations per balancing factor, input model, and initial embedding.

Specifically, in descent steps, we minimize symmetric Dirichlet energy using projected Newton method given the current UV topology. In topology steps, we search for a nearby UV topology that locally decrease the objective the most by querying a filtered set of basic topological operations. To be appropriately aggressive on searching in the topological space, we develop an analogous line search method as in continuous settings and allow multiple fracture initiation. Since in application scenarios, an upper bound for distortion or seam length is more intuitive than picking a balancing factor, we also provide a constrained optimization view of this broader problem that seeks stationary w.r.t. both primal (distortion and seams) and dual (balancing factor) variables subject to user specified distortion or seam length upper

Our method automatically produces both visually pleasing and high quality UV maps with low mapping distortion and an optimally sparse set of seams without any user assistance. We also show that given a UV configuration by other methods, our method can improve the distortion and seam placement, and that our framework has the potential to also handle bijectivity, seamlessness, and user preferences jointly within it as well.

CCS Concepts: • Computing methodologies → Mesh geometry mod-

Additional Key Words and Phrases: geometry processing, mesh parameterization, seam placement, numerical optimization, ...

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INTRODUCTION

RELATED WORKS

related methods: AutoCuts [Poranne et al. 2017], Seamster [Sheffer and Hart 2002], geometry images [Gu et al. 2002], Multi-chart geometry images [Snyder et al. 2003], D-Chart [Julius et al. 2005]

components: Bijective parameterization with free boundaries [Smith and Schaefer 2015], projected Newton [Teran et al. 2005], MIPS [Hormann and Greiner 2000],

AN ALTERNATING FRAMEWORK OF CONTINUOUS AND DISCRETE OPTIMIZATION FOR MESH **PARAMETERIZATION**

The most basic and intuitive mesh parameterization objective regarding both seams and distortion is minimizing distortion with as-sparse-as-possible seams introduced. However, seam sparsity usually leads to discontinuous energies w.r.t. UV coordinates $U \in \mathcal{R}^{2n_v}$, (n_{τ}) is the number of vertices on the input mesh), which is non-trivial to be considered into existing distortion minimization routines. Instead of progressively approximating seam sparsity energy with a continuous counterpart applying homotopy optimization method as [Poranne et al. 2017], we handle this discrete energy in a combinatoric way - searching in the topological space.

3.1 Formulation

This topological space is a directed graph G_T with its vertices $v_T \in$ V_T being all possible UV topologies of a given 3D surface, and its edges $e_T \in E_T$ are the basic topological operations on a mesh such as vertex split, edge merge, etc, that can transform one UV topology to a nearby topology.

Now, if we consider both distortion and seam in one objective E_w , we can define the value f_v of vertex $v_{T,i}$ as

$$f_v(v_{T,i}) = \min_{U_i} E_w$$

and the weights f_w of edge $e_{T,m}$ from $v_{T,i}$ to $v_{T,j}$ could just be

$$f_{\mathcal{W}}(e_{T,m}) = f_{\mathcal{V}}(v_{T,i}) - f_{\mathcal{V}}(v_{T,i})$$

Thus our problem could be written as

$$\min_{U,v_T} E_w$$

which could be stated as to search for a $v_{T,i}$ on G_T where all edges connected to it satisfies $f_w \ge 0$.

However, computing f_v for one UV topology requires a whole continuous optimization process, and even the number of neighbors of one UV topology is in the scale of n_v^2 . Consequently, we construct a search path on G_T by progressively introducing or removing seams, and we only estimate f_w on a local stencil of U for a filtered set of neighbors so that the whole process of continuous optimization is only conducted while necessary.

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Method 3.2

Let's consider a simple situation, minimizing normalized symmetric Dirichlet energy [Smith and Schaefer 2015]

$$E_{SD} = \frac{1}{n_t |A|} \sum_t |A_t| (\sigma_{t,1}^2 + \sigma_{t,2}^2 + \sigma_{t,1}^{-2} + \sigma_{t,2}^{-2})$$

and normalized total seam length

$$E_{se} = \frac{1}{\sqrt{n_t |e|}} \sum_{i \in \mathcal{S}} 2|e_i|$$

where a balancing factor $\lambda \in [0, 1]$ is controlling the ratio between

$$E_{w} = \lambda E_{SP} + (1 - \lambda)E_{SD}$$

We minimize E_w by iteratively alternate between continuous optimization (in descent steps) and discrete optimization (in topology

• In descent steps, we compute $f_v(v_{T,i})$ via projected Newton method [Teran et al. 2005]:

$$f_{v}(v_{T,i}) = E_{se,i} + \min_{U_{i}} E_{SD}$$

• In topology steps, we estimate $f_v(v_{T,j})$ for a filtered set of neighbors on a local stencil of U as \hat{f}_v and move onto the neighbor $v_{T,i+1}$ with smallest \hat{f}_v .

If in a descent step, $f_v(v_{T,i}) \ge f_v(v_{T,i-1})$ is detected, we stop the process by rolling back to $v_{T,i-1}$, which is the stationary of E_w w.r.t. both UV topology and UV coordinates that we are searching for.

3.3 Convergence

As our method is defined to guarantee convergence, we now analyze the convergence rate. First, it's easy to see that E_w is monotonically decreasing looking at each end of descent steps. Now we look at descent step *i* and i + 1, from $E_w^i \ge E_w^{i+1}$ we have

$$E_{SD}^i - E_{SD}^{i+1} \geq \frac{\lambda}{1-\lambda}(E_{se}^{i+1} - E_{se}^i) \geq \frac{\lambda}{1-\lambda}\frac{1}{\sqrt{n_t|e|}}2|e|_{min}$$

if we now only consider splitting operations that keep increasing E_{se} . It's obvious that E_{SD} 's theoretical lower bound is defined to be 4, so we have

$$n_{alter} \leq \frac{(1-\lambda)\sqrt{n_t}\overline{|e|}}{2\lambda|e|_{min}}(E_{SD}^0-4)$$

The most important hint we can read from this is, to accelerate convergence, we can move through multiple vertices on G_T in each topology step to increase $E_{se}^{i+1} - E_{se}^{i}$. Consequently, we build an anologous line search method and allow multiple fracture initiation to be appropriately agressive when searching in the topological space and ensure that we won't fall into bad locally optimal UV topologies.

Merge operations should be defined carefully to ensure convergence, and the proof will need updates. For example, it needs to decrease E_w in order to be considered. What will be the filtering measurement? How do we prevent from element inversion potentially caused by merge? How will we choose among merge and split?

Is their energy decrease comparable? Will we need to also propagate merge like "zippering"? Do we need merge to be performed on non-corner edge pairs like the sense of interior splits?

3.4 Potential Extensions

It will be interesting to replace E_{SD} with other types of distortion energies, especially conformal energies like MIPS [Hormann and Greiner 2000] to see how seams that benefits conformality will be different from seams that benefits isometry. Besides, bijectivity could potentially be achieved by augmenting distortion energy with a penalty-based collision handling energy, possibly also assisted with air mesh method [?]. Similarly, seamless properties could also potentially be achieved by augmenting distortion energy with the correspondingly developed new differentiable objectives, and our alternating framework stays the same.

If an objective derived from an application is discontinuous and it could be expressed using mesh topology, then we can simply augment it into E_{se} and tackle it in the topology steps. For example, the smoothness of seams, user preferences on regional seam placement, and properties related to charts should all be able to be considered in this way.

Besides, SIMD type of parallelism could easily accelerate the \hat{f}_v evaluations in the topology steps, and it also has the potential to improve the quality of the topology search by directly evaluating f_v for neighbors and track multiple branches.

DESCENT STEPS FOR CONTINUOUS OPTIMIZATION

Newton-type Iterations

for each descent step inner iteration *j*:

compute E_{SD} Hessian proxy P^j using projected Newton;

compute E_{SD} gradient g^j ;

solve for search direction p^j ($P^jp^j=-g^j$) using PARDISO symmetric indefinite solver;

compute initial step size α_0^J by avoiding element inversion;

backtracking line search with Armijo rule; update $U^{j+1} = U^j + \alpha^j p^j$;

record energy decrease $(1 - \lambda_t)\Delta E_{SD}^j$;

4.2 Potential Accelerations for Practical Use

Since our topological operations only change the mesh locally both on connectivity and coordinates, we could also update the Hessian or the decomposition locally after topology changes to save time. Besides, it's also interesting to try other Hessian approximation methods like L-BFGS or Majorization to explore further acceleration by finding a balance between computational cost and convergence

For convergence tolerance of descent steps, $||\nabla E_{SD}||^2 \le 10^{-6}$ (note that our energy is normalized) works generally well for all input models judging from the initiated fracture in the following topology step. In fact more inexact solve performs well on most of the models with even $||\nabla E_{SD}||^2 \le 10^{-4}$, but some may result even better with $||\nabla E_{SD}||^2 \le 10^{-8}$. Since we are conducting non-convex optimization, $||\nabla E_{SD}||^2$ is not always decreasing, which is also why we don't use Wolfe conditions for line search. The argument here for tolerance issue is that, it depends on whether we are truly in

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the infinitesimal region of a stationary. Some configuration with $||\nabla E_{SD}||^2 \le 10^{-6}$ may still not inside the infinitesimal region of a stationary, where if optimization goes on, the $||\nabla E_{SD}||^2$ will go up and then fall down again to a real stationary, which is understandable in non-convex optimization.

TOPOLOGY STEPS FOR DISCRETE OPTIMIZATION

Evaluating Topological Operations via Optimization on **Local Stencils**

Candidate Filtering: for each vertex compute divergence of local gradients independently picking $\sqrt{n_{v,b}^i}$ boundary vertices and $\sqrt{n_{v,i}^i}$ interior vertices with largest divergence as candidates

Local Evaluation: for each candidate vertex if on boundary for each interior incident edge split and compute $\Delta E_{SD,I}$ locally compute $\Delta E_{w,l} = (1 - \lambda_t) \Delta E_{SD,l} + \lambda_t \Delta E_{se}$ else for each pair of incident edges forming a smooth path split and compute $\Delta E_{SD,l}$ locally compute $\Delta E_{w,l} = 0.5((1 - \lambda_t)\Delta E_{SD,l} + \lambda_t \Delta E_{se})$

split the vertex with largest $|\Delta E_{w,l}|$ turn on fracture propagation Since we are using the local estimation to approximate true global energy decrease, it would be necessary to find a balance between accuracy and efficiency by weighing between size of local stencils and number of optimization iterations to run.

Enable merge operation.

Line Search in Topological Space

Once a fracture is initiated, it almost always be extended further in the later topology steps, which justifies the robustness of our method. To speed up the process, instead of waiting for another descent step and query all the filtered candidates again, we propagate this newly initiated fracture further in between the first several inner iterations of the following descent step.

After the fracture has been initiated, we first go to descent step to run an inner iteration and record the energy decrease ΔE_{w}^{J} and energy $E_w^{j,0}$. Then we evaluate \hat{f}_v 's for splitting the tail vertex of the newly initiated fracture along its incident edges. If the one with largest \hat{f}_v satisfies $\hat{f}_v - E_w^{j,0} \le \Delta E_w^j$, we propagate the fracture along this edge, and run another inner iteration to do another propagation query. If there's no propagation that could benefit more than running the inner iteration, we stop query propagation in the current descent step and run inner iterations till convergence:

for each fracture tail vertex k for each interior incident edge of ksplit and compute $\Delta E_{SD,l}$ locally compute $\Delta E_{w,l} = (1-\lambda_t)\Delta E_{SD,l} +$ $\lambda_t \Delta E_{se}$ if the largest $|\Delta E_{w,l}|$ is larger than $|(1-\lambda_t)\Delta E_{SD}^j|$ propagate fracture by splitting the vertex else turn off fracture propagation for the rest of the current descent step

Improve current propagation scheme by developing analogous line search method

Multiple Fracture Initiation

Multiple fracture initiation would be redundant on a single conelike region since only one fracture would be enough to release the distortion. However, initiating one fracture for each cone-like region

within a single topology step would certainly accelerate the whole

This could be done by partitioning the UV space according to distortion or filtering measurement and initiate a valid fracture (if any) in each region in every topology step. The partition criteria could be separating the domain so that on each subdomain, the function is convex (has exactly one stationary).

WEIGHTING THE OBJECTIVE AUTOMATICALLY BY INTRODUCING A DUAL PROBLEM

At each inner iterate k + 1, we fix some λ^{k+1} and minimize the bi-objective

$$\min_{T,V} E_{SE}(V,T) + \lambda^{k+1} E_{SD}(V,T)$$

How do we get λ^{k+1} ?

Our overall minimization is inequality constrained with a specified upper bound $b \in \mathbb{R}_+$ on distortion. (L2 norm on SD energy for now - pretty easy to modify to an extremal measure if we want later

Our model problem minimization is then

$$\min_{T,V} E_{SE}(V,T) : b - E_{SD}(V,T) \ge 0$$

Or, equivalently,

$$\min_{T, V} \max_{\lambda \ge 0} E_{SE}(V, T) + \lambda \left(E_{SD}(V, T) - b \right)$$

Of course this is nonsmooth in λ since it does not take into account very nicely the fact that per-iteration we will start away from feasibility and want to iteratively improve both our primal variables $\{V, T\}$ and our dual variable λ . So to smoothly update to a current λ^{k+1} from a previous estimate λ^k we will add a regularizer $R(\lambda, \lambda^k)$ to make sure λ iterates behave themselves reasonably. For now lets stick with something simple: a quadratic regularizer should do the trick $R = \frac{1}{2\kappa} (\lambda - \lambda^k)^2$. For iteration k + 1 this gives us

$$\min_{T,V} \max_{\lambda \geq 0} E_{SE}(V,T) + \lambda \left(E_{SD}(V,T) - b \right) - \frac{1}{2\kappa} (\lambda - \lambda^k)^2$$

And now we can first solve closed form for λ as

$$\lambda^{k+1} \leftarrow argmax_{\lambda \geq 0} E_{SE}(V,T) + \lambda \left(E_{SD}(V,T) - b \right) - \frac{1}{2\kappa} (\lambda - \lambda^k)^2$$

giving us

$$\lambda^{k+1} \leftarrow \max\left(0, \kappa(E_{SD}(V, T) - b) + \lambda^k\right)$$

We then can solve the inner iteration (with both discrete topology steps and smooth steps) with the energy

$$\min_{T,V} E_{SE}(V,T) + \lambda^{k+1} E_{SD}(V,T)$$

Followed by the next update of dual variable λ .

(Notice that throughout the above we can define a progressive λ without needing to employ subgradients to reason about nonsmoothness in our sparsity energy.)

R is to iteratively solving for λ so that it could have intermediate values between 0 and ∞ . Then starting from full mesh, λ will first increase as bound is not reached, and then it will decrease when bound is reached. But currently we don't have merge to increase E_{SE} and decrease E_{SD} , so our process will stop right after it reaches

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53 54

56 57 58 the bound. The bound is obvious to be reached, how do we know the path of λ is great? If we have merge, then will the optimization converge on all T, V, λ ?

7 RESULTS AND DISCUSSION

quality and timing comparison with previous methods, deal with closed surfaces by either starting from random rectangle embedding or basic heuristics like farthest points.

show improvements starting from results given by previous methods $% \left(1\right) =\left(1\right) \left(1\right) \left$

how does triangulation affect our result? try same shape with different triangulation.

given a symmetric shape, whether symmetrically triangulated or not, does our method preserve symmetry in UV space?

8 CONCLUSIONS AND FUTURE WORKS

take advantage of basic SIMD type of parallelism for accelerating query and improving results' quality by directly evaluating f_v for neighbors and track multiple branches, very useful for practical implementations

if the user won't mind getting a slightly different triangulation, we could also create fractures in the interior of an element and locally remesh the stencil

start and solve in 3D by reducing curvature so that the need for locally injective initial embedding in parameterization problems could be eliminated, and the result is only "biased" by it's 3D shape, which is the most reasonable bias

try conformal energy like MIPS

bijectivity, seamless, and other augmentation of continuous energy?

handle user preferences on seam placement seam smoothness, patch related discrete energy augmentation?

9 ACKNOWLEDGEMENTS

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