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Stallions

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\mathcal{O}

cin.tie(0);

1 Contest	1	<pre>cout << fixed << setprecision(15);</pre>
2 Data Structures	2	return 0; }
3 Graphs	4	compilation.txt
_		2 lines
4 Strings	10	g++ -02 -std=c++17 -Wno-unused-result -Wshadow - Wall -0 %e %e.cpp
5 Geometry	13	g++ -std=c++17 -Wshadow -Wall -o %e %e.cpp - fsanitize=address -fsanitize=undefined - D_GLIBCXX_DEBUG -g
6 Math	16	
7 Convolutions	18	s.sh 6 lines
8 Various	22	for((i = 0; ; i++)) do echo \$i
9 Formulas	22	./gen \$i > in diff -w <(./a < in) <(./brute < in) break [\$? == 0] break
Contest (1) template.hpp		hash.sh 1 lines
$\frac{1}{//hash} = 7d0184$	27 lines	cpp -dD -P -fpreprocessed \$1 tr -d '[:space:]' md5sum cut -c-6
<pre>#include <bits stdc++.h=""> using namespace std;</bits></pre>		troubleshoot.txt 59 lines
#define FOR(i, a, b) for(int i = (a); i < (b) ++)	; i	Pre-submit: F9. Write a few manual test cases.
#define RFOR(i, a, b) for(int i = (a) - 1; i : b); i)	>= (Calculate time and memory complexity. Check limits.
<pre>#define SZ(a) int(a.size()) #define ALL(a) a.begin(), a.end() #define PB push_back</pre>		Check overflows, size of arrays, clearing mutitestcases, uninitialized variables.
#define MP make_pair		Wrong answer:
<pre>#define F first #define S second</pre>		F9. Print your solution!
<pre>typedef long long LL; typedef vector<int> VI; typedef pair<int, int=""> PII;</int,></int></pre>		Read your code. Check Pre-submit. Are you sure your algorithm works? Think about precision errors and hash collitions.
typedef double db;		Have you understood the problem correctly? Write brute and generator.
<pre>int main()</pre>		TC
<pre>{ ios::sync with stdio(0);</pre>		If you don't know what to do: Are you clearing all data structures between test

```
cout << fixed << setprecision(15);</pre>
return 0;
ompilation.txt
++ -O2 -std=c++17 -Wno-unused-result -Wshadow -
  Wall -o %e %e.cpp
++ -std=c++17 -Wshadow -Wall -o %e %e.cpp -
  fsanitize=address -fsanitize=undefined -
  D GLIBCXX DEBUG -q
.sh
or((i = 0; ; i++)) do
echo $i
./qen $i > in
                                                   F9.
diff -w < (./a < in) < (./brute < in) || break
[ $? == 0 ] || break
one
ash.sh
pp -dD -P -fpreprocessed $1 | tr -d '[:space:]'|
   md5sum |cut -c-6
oubleshoot.txt
re-submit:
9.
rite a few manual test cases.
alculate time and memory complexity. Check
  limits.
heck overflows, size of arrays,
learing mutitestcases, uninitialized variables.
rong answer:
9.
rint your solution!
ead your code.
heck Pre-submit.
re you sure your algorithm works?
hink about precision errors and hash collitions.
ave you understood the problem correctly?
rite brute and generator.
f you don't know what to do:
```

cases?

```
Can your algorithm handle the whole range of
   input?
Read the full problem statement again.
Do you handle all corner cases correctly?
Any uninitialized variables?
Any overflows?
Confusing N and M, i and j, etc.?
What special cases have you not thought of?
Are you sure the STL functions you use work as
   you think?
Add some assertions, maybe resubmit.
Go for a small walk, e.g. to the toilet.
Is your output format correct? (including
   whitespace)
Rewrite your solution from the start or let a
   teammate do it.
Runtime error:
Print your solution!
Read your code.
F9 with generator.
If you don't know what to do:
Invalidated pointers or iterators?
Are you reading or writing outside the range of
   any vector?
Any assertions that might fail?
Any possible division by 0? (mod 0 for example)
Have you tested all corner cases locally?
Any uninitialized variables?
Any possible infinite recursion?
Are you using too much memory?
Debug with resubmits (e.g. remapped signals, see
   Various).
Time limit exceeded:
What is the complexity of your algorithm?
Are you copying a lot of unnecessary data? (
   References)
Do you have any possible infinite loops?
How big is the input and output? (consider scanf)
Avoid vector, map. (use arrays/unordered_map)
What do your teammates think about your algorithm
Memory limit exceeded:
Calculate memory usage with stack in recurtion.
```

Data Structures (2)

```
dsu.hpp
                                           2de4ff, 34 lines
struct DSU
  int n;
 VI p;
 VI sz;
  void init(int _n)
   n = _n;
    sz.assign(n, 1);
    p.resize(n);
    iota(ALL(p), 0);
  int find(int v)
    if (v == p[v])
      return v;
    return p[v] = find(p[v]);
 bool unite(int u, int v)
    u = find(u);
    v = find(v);
    if (u == v)
      return false;
    if (sz[u] > sz[v])
      swap(u, v);
    p[u] = v;
    sz[v] += sz[u];
    return true;
};
Fenwick.hpp
                                          d4ebda, 43 lines
struct Fenwick
  int n;
 vector<LL> v;
```

void init(int _n)

v.assign(n, 0);

 $n = _n;$

```
void add(int i, int x)
    for (; i < n; i = (i + 1) | i)
      v[i] += x;
  LL sum(int i)
   LL ans = 0;
    for (; i \ge 0; i = (i \& (i + 1)) - 1)
      ans += v[i];
    return ans;
  int lower bound(LL x)
    LL sum = 0;
    int i = -1;
    int lq = 31 - __builtin_clz(n);
    while (lq >= 0)
      int j = i + (1 << lq);
      if (j < n \&\& sum + v[j] < x)
        sum += v[j];
        i = j;
      lg--;
    return i + 1;
};
Fenwick.txt
Minimum on segment:
```

upd(r, -x, x * r);

```
1) Use two Fenwick trees with n = 2^k.
n = 1 \ll (32 - \underline{builtin_clz(n - 1))};
2) One tree for normal array and one for reversed
3) When querying for minimum on the segment
only consider segments [(i & (i + 1)), i]
from trees that are COMPLETELY inside [1, r]
Fenwick tree for adding on segment (prefixes):
1) Use 2 arrays: mult and add
2) upd(int i, int updMult, int updAdd)
default Fenwick update.
3) add x on segment [1, r]:
 upd(1, x, -x * (1 - 1));
```

```
4) to calculate sum on prefix r:
  sumAdd and sumMult - default Fenwick sum
  st - initial value of r
  ans = st * sumMult + sumAdd
treap.hpp
Description: uncomment in split for explicit key or in merge for
implicit priority.
                                          925cdb, 145 lines
mt19937 rng;
struct Node
  int 1, r;
  int x;
  int y;
  int cnt;
  int par;
  int rev;
  int mn;
  void init(int value)
    1 = r = -1;
    x = value;
    y = rnq();
    cnt = 1;
    par = -1;
    rev = 0;
    mn = value;
};
struct Treap
  Node A[N];
  int sz = 0;
  int getCnt(int v)
```

if (v == -1)

int getMn(int v)

if (v == -1)

return INF;

return A[v].mn;

int newNode(int val)

return 0;

return A[v].cnt;

```
A[sz].init(val);
  return sz++;
void upd(int v)
 if (v == -1)
    return;
 A[v].cnt = getCnt(A[v].1) +
 getCnt(A[v].r) + 1;
 A[v].mn = min(A[v].x,
 min(getMn(A[v].l), getMn(A[v].r)));
void reverse(int v)
 if ( \lor == -1 )
    return;
 A[v].rev ^= 1;
void push(int v)
 if (v == -1 | | A[v].rev == 0)
    return;
 reverse(A[v].1);
  reverse(A[v].r);
 swap(A[v].l, A[v].r);
 A[v].rev = 0;
PII split(int v, int cnt)
 return {-1, -1};
 push(v);
  int left = getCnt(A[v].1);
  PII res;
  // if (val \ll A/v | x)
  if (cnt <= left)</pre>
    if (A[v].1 != -1)
     A[A[v].1].par = -1;
    res = split(A[v].l, cnt);
    A[v].l = res.second;
    if (res.second !=-1)
      A[res.second].par = v;
    res.second = v;
  }
  else
    if (A[v].r != -1)
```

```
A[A[v].r].par = -1;
   // split(v, val)
    res = split(A[v].r, cnt - left - 1);
    A[v].r = res.first;
    if (res.first !=-1)
     A[res.first].par = v;
    res.first = v;
 upd(v);
 return res;
int merge(int v, int u)
 if (v == -1) return u;
 if (u == -1) return v;
 int res;
 // if (rng()\% (getCnt(v) + getCnt(u)) <
     getCnt(v)
 if (A[v].y > A[u].y)
   push(v);
   if (A[v].r != -1)
     A[A[v].r].par = -1;
    res = merge(A[v].r, u);
   A[v].r = res;
    if (res !=-1)
     A[res].par = v;
   res = v;
 }
 else
   push(u);
   if (A[u].l != -1)
    A[A[u].1].par = -1;
    res = merge(v, A[u].1);
   A[u].l = res;
    if (res !=-1)
     A[res].par = u;
    res = u;
 upd(res);
 return res;
int getIdx(int v, int from = -1)
 if (v == -1)
   return 0;
 int x = getIdx(A[v].par, v);
 if (from == -1 || A[v].r == from)
    x += getCnt(A[v].1) + 1;
```

```
push(v);
    return x;
};
ordered-set.hpp
#include <ext/pb_ds/assoc_container.hpp>
using namespace __gnu_pbds;
using namespace std;
typedef tree<int, null_type, less<int>,
    rb_tree_tag,
   tree_order_statistics_node_update>
   ordered_set;
// example: ordered_set s; s.insert(47);
// s.order\_of\_key(k); -- returns number of
    elements less then k
// s.find_by_order(k); - returns iterator to k-th
     element or s.end()
// s.count() does not exist.
sparse-table.hpp
                                         7fdd30, 30 lines
int lq[N + 1];
struct SparseTable
  int t[N][LOG];
  void init(const VI& v)
   lq[1] = 0;
    FOR (i, 2, N + 1) lg[i] = lg[i / 2] + 1;
    FOR (i, 0, N) FOR (j, 0, LOG) t[i][j] = INF;
    FOR (i, 0, SZ(v)) t[i][0] = v[i];
    FOR (j, 1, LOG)
      int len = 1 << (j - 1);
      FOR (i, 0, N - (1 << j))
        t[i][j] = min(t[i][j-1],
        t[i + len][j - 1]);
  int query(int 1, int r)
    int i = lq[r - l + 1];
    return min(t[1][i], t[r - (1 << i) + 1][i]);
```

} st;

```
convex-hull-trick.hpp
                                          06db0c, 74 lines
struct Line
  LL a, b, xLast;
 Line() {}
  Line(LL _a, LL _b): a(_a), b(_b) {}
 bool operator<(const Line& 1) const
    return MP(a, b) < MP(1.a, 1.b);
 bool operator<(int x) const</pre>
    return xLast < x;</pre>
   __int128 getY(__int128 x) const
    return a * x + b;
  LL intersect (const Line& 1) const
    assert(a < 1.a);
    LL dA = 1.a - a, dB = b - 1.b, x = dB / dA;
    if (dB < 0 && dB % dA != 0)
      x--;
    return x;
};
struct ConvexHull: set<Line, less<>>
 bool needErase(iterator it, const Line& 1)
    LL x = it -> xLast;
    if (it->getY(x) > l.getY(x))
      return false;
    if (it == begin())
      return it->a >= 1.a;
    x = prev(it) -> xLast + 1;
    return it->getY(x) < l.getY(x);</pre>
  void add(LL a, LL b)
    Line l(a, b);
    auto it = lower_bound(1);
    if (it != end())
      LL x = it == begin() ? -LINF :
```

```
prev(it)->xLast;
     if ((it == begin()
       | | prev(it) - sqetY(x) >= l.qetY(x) |
        && it->getY(x + 1) >= l.getY(x + 1))
        return;
   while (it != end() && needErase(it, 1))
     it = erase(it);
   while (it != begin()
     && needErase(prev(it), 1))
     erase(prev(it));
   if (it != begin())
     auto itP = prev(it);
     Line lIt = *itP;
     lIt.xLast = itP->intersect(l);
     erase(itP);
     insert(lIt);
   l.xLast = it == end() ? LINF :
       l.intersect(*it);
   insert(1);
 LL getMaxY(LL x)
   return lower_bound(x)->getY(x);
 }
};
```

Graphs (3)

3.1 Decompositions

centroid.hpp

Description: dfsSZ calculates size of subtrees not going to usedc[v] vertices

```
if (sz[to] * 2 > szAll)
{
    v = to;
    break;
}
if (v == -1)
    break;
pr = cent;
cent = v;
}
usedc[cent] = true;

// here calculate f(cent)

for (auto to : g[cent])
{
    if (!usedc[to])
    {
        build(to);
    }
}
```

HLD.hpp

Description: run dfsSZ and dfsHLD to build HLD. Vertex v has index tin[v]. To update on path use process as in get().

```
VI q[N];
int sz[N];
int h[N];
int p[N];
int top[N];
int tin[N];
int tout[N];
int t = 0;
void dfsSZ(int v, int par = -1, int hei = 0)
  sz[v] = 1;
  h[v] = hei;
  p[v] = par;
  for (auto& to : g[v])
    if (to == par)
      continue;
    dfsSZ(to, v, hei + 1);
    sz[v] += sz[to];
    if (g[v][0] == par || sz[g[v][0]] < sz[to])
      swap(g[v][0], to);
```

```
void dfsHLD(int v, int par = -1, int tp = 0)
 tin[v] = t++;
 top[v] = tp;
 FOR (i, 0, SZ(q[v]))
   int to = q[v][i];
   if (to == par)
     continue;
   if (i == 0)
     dfsHLD(to, v, tp);
   else
     dfsHLD(to, v, to);
  tout[v] = t - 1;
LL get(int x, int y)
 LL res = 0;
  while(true)
   int tx = top[x];
   int ty = top[y];
   if (tx == ty)
     int t1 = tin[x];
     int t2 = tin[y];
     if (t1 > t2)
      swap(t1, t2);
     res += query(t1, t2);
     break;
   if (h[tx] < h[ty])
     swap(tx, ty);
     swap(x, y);
   res += query(tin[tx], tin[x]);
   x = p[tx];
  return res;
```

biconnected-components.hpp

2d79e1, 83 lines

```
struct Graph
 vector<PII> edges;
 vector<VI> q;
```

```
VI tin, low;
VI col;
VI par;
VI used;
int t = 1, c = 1;
vector<int> st;
int n, m;
void init(int _n, int _m)
  n = _n;
  m = _m;
  edges.assign(m, {0, 0});
  g.assign(n, {});
  tin.assign(n, 0);
  used.assign(n, 0);
  par.assign(n, -1);
  used.assign(n, 0);
  t = c = 1;
void addEdge(int a, int b, int i)
  assert(0 <= a && a < n);
  assert (0 \le b \& b \le n);
  assert(0 <= i && i < m);
  edges[i] = MP(a, b);
  g[a].PB(i);
  q[b].PB(i);
void dfs (int v, int p = -1)
  used[v] = 1;
  par[v] = p;
  low[v] = tin[v] = t++;
  int cnt = 0;
  for (auto e : g[v])
    int to = edges[e].F;
    if (to == v)
      to = edges[e].S;
    if (p == to) continue;
    if (!used[to])
```

```
cnt++;
        st.PB(e);
        dfs(to, v);
        low[v] = min(low[v], low[to]);
        if ((par[v] == -1 \&\& cnt > 1) | |
        (par[v] != -1 \&\& low[to] >= tin[v])
          while (st.back() != e)
            col[st.back()] = c;
            st.pop_back();
          col[st.back()] = c++;
          st.pop_back();
        }
      else
        low[v] = min(low[v], tin[to]);
        if (tin[to] < tin[v])</pre>
          st.PB(e);
};
```

3.2 Flows

```
kuhn.hpp
```

```
Time: 0.6s for |V| = 10^5, |E| = 2 * 10^5
```

```
39cb20, 93 lines
struct Graph
 int L, R;
 //edges from left to right in 0 indexing
 vector<VI> q;
 VI mt, P, U;
  void init(int 1, int r)
   L = 1, R = r;
   a.resize(L);
   P.resize(L);
    U.resize(L);
    mt.resize(R);
  void addEdge(int from, int to)
```

```
assert(0 <= from && from < L);
  assert (0 \leq to && to \leq R);
  g[from].PB(to);
int iter;
bool kuhn (int v)
  if (U[v] == iter) return false;
  U[v] = iter;
  random_shuffle(ALL(g[v]));
  for(int to : q[v])
    if (mt[to] == -1)
      mt[to] = v;
      P[v] = to;
      return true;
  for(int to : g[v])
    if (kuhn(mt[to]))
      mt[to] = v;
     P[v] = to;
      return true;
    }
  return false;
int doKuhn()
  fill (ALL (mt), -1);
  fill(ALL(P), -1);
  fill(ALL(U), -1);
  int res = 0;
  iter = 0:
  VI order(L);
  iota(ALL(order), 0);
  random shuffle(ALL(order));
  while (true)
    iter++;
    bool ok = false;
    for(int v : order)
```

```
if (P[v] == -1)
          if (kuhn(v))
            ok = true;
            res++;
      if (!ok) break;
    return res;
};
//Minimum number of vertex to cover edges
//equal to size of mathing
//Maximal independent set
//is complement of vertex cover
//To restore cover make directed graph
//matched\ edges\ l <-r,\ not\ matched\ l >> r
//dfs from vl : P[vl] == -1
//cover: used[r] != -1 \&\& mt[r] != -1
// used[l] == -1 \&\& P[l] != -1
dinic.hpp
                                         6afa18, 93 lines
struct Graph
  struct Edge
   int from, to;
   LL cap, flow;
 };
 int _n;
 vector<Edge> edges;
 vector<VI> q;
 VI d, p;
  Graph() : _n(0) {}
  Graph (int n) : _n(n), g(n), d(n), p(n) {}
 void addEdge(int from, int to, LL cap)
    assert(0 <= from && from < _n);
    assert(0 <= to && to < _n);
    assert(0 <= cap);
    g[from].PB(SZ(edges));
    edges.PB({from, to, cap, 0});
```

```
q[to].PB(SZ(edges));
  edges.PB({to, from, 0, 0});
int bfs(int s, int t)
  fill(ALL(d), -1);
  d[s] = 0;
  queue<int> q;
  q.push(s);
  while (!q.empty())
    int v = q.front();
    q.pop();
    for (int e : g[v])
      int to = edges[e].to;
      if (edges[e].flow < edges[e].cap</pre>
        && d[to] == -1)
        d[to] = d[v] + 1;
        q.push(to);
  return d[t];
LL dfs(int v, int t, LL flow)
  if (v == t || flow == 0)
    return flow;
  for (; p[v] < SZ(q[v]); p[v]++)
    int e = g[v][p[v]], to = edges[e].to;
    LL c = edges[e].cap, f = edges[e].flow;
    if (f < c \&\& (to == t || d[to] == d[v] + 1)
      LL push = dfs(to, t, min(flow, c - f));
      if (push > 0)
        edges[e].flow += push;
        edges[e ^ 1].flow -= push;
        return push;
      }
  return 0;
```

```
assert(s != t);
    LL flow = 0:
    while (bfs(s, t) !=-1)
      fill(ALL(p), 0);
      while (true)
        LL f = dfs(s, t, LINF);
        if (f == 0)
          break:
        flow += f;
    return flow;
};
min-cost-flow.hpp
                                         8a8605, 103 lines
struct Graph
  struct Edge
    int from, to;
    int cap, flow;
   LL cost;
  };
  int _n;
  vector<Edge> edges;
  vector<VI> q;
  vector<LL> d;
  VI p, w;
  Graph(): _n(0) {}
  Graph (int n): _n(n), _n(n), _n(n), _n(n), _n(n), _n(n)
  void addEdge(int from,int to, int cap, LL cost)
    assert(0 <= from && from < _n);
    assert(0 <= to && to < _n);
    assert (0 <= cap);
    assert(0 <= cost);
    g[from].PB(SZ(edges));
    edges.PB({from, to, cap, 0, cost});
    g[to].PB(SZ(edges));
```

LL flow(int s, int t)

assert(0 <= s && s < _n);

assert(0 <= t && t < _n);

```
edges.PB({to, from, 0, 0, -cost});
}
pair<int, LL> flow(int s, int t)
  assert(0 <= s && s < _n);
  assert(0 <= t && t < _n);
  assert(s != t);
  int flow = 0;
  LL cost = 0;
  while (true)
    fill(ALL(d), LINF);
    fill(ALL(p), -1);
    fill(ALL(w), 0);
    queue<int> q1, q2;
    w[s] = 1;
    d[s] = 0;
    q2.push(s);
    while (!q1.empty() || !q2.empty())
      int v;
      if (!q1.empty())
        v = q1.front();
        q1.pop();
      else
        v = q2.front();
        q2.pop();
      for (int e : q[v])
        if (edges[e].flow == edges[e].cap)
          continue;
        int to = edges[e].to;
        LL newDist = d[v] + edges[e].cost;
        if (newDist < d[to])</pre>
          d[to] = newDist;
          p[to] = e;
          if (w[to] == 0)
            q2.push(to);
          else if (w[to] == 2)
            q1.push(to);
          w[to] = 1;
      w[v] = 2;
```

```
if (p[t] == -1)
        break;
      int curFlow = INF;
      LL curCost = 0;
      for (int v = t; v != s;)
        int e = p[v];
        curFlow = min(curFlow,
        edges[e].cap - edges[e].flow);
        curCost += edges[e].cost;
        v = edges[e].from;
      for (int v = t; v != s;)
        int e = p[v];
        edges[e].flow += curFlow;
        edges[e ^ 1].flow -= curFlow;
        v = edges[e].from;
      flow += curFlow;
      cost += curCost * curFlow;
    return {flow, cost};
};
```

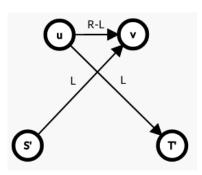
3.3 Flows text

https://atcoder.jp/contests/abc285/
editorial/5535

On the resulting graph, accumulate maximum flow in the following order:

- from S' to T'
- from S' to T
- from S to T'
- from S to T.

An S-T flow that satisfies the minimum capacities exists if and only if, for all outgoing edges from S' and incoming edges to T', the flow and capacity are equal.



$$\sum_{i} a_{i} x_{i} + \sum_{i} b_{i} \overline{x_{i}} + \sum_{i,j} c_{ij} x_{i} \overline{x_{j}} \to \min$$

If $a_i \leq b_i$, add an edge from S to i of capacity $b_i - a_i$ and add a_i to the answer.

Otherwise, add an edge from i to T of capacity $a_i - b_i$ and add b_i to the answer.

Add an edge from i to j of capacity c_{ij} .

Add the S-T minimum cut to the answer.

Specific 3.4

hungarian.hpp

Obaccf, 63 lines

```
LL hungarian (const vector<vector<LL>>& a)
 int n = SZ(a), m = SZ(a[0]);
 assert (n <= m);
  vector<LL> u(n + 1), v(m + 1);
 VI p(m + 1, n), way(m + 1);
 FOR(i, 0, n)
   p[m] = i;
   int j0 = m;
   vector<LL> minv(m + 1, LINF);
   vector<int> used(m + 1);
   while (p[j0] != n)
     used[j0] = true;
     int i0 = p[j0], j1 = -1;
     LL delta = LINF;
     FOR(j, 0, m)
        if (!used[j])
          int cur = a[i0][j] - u[i0] - v[j];
```

```
if (cur < minv[j])</pre>
            minv[j] = cur;
             way[j] = j0;
           if (minv[j] < delta)</pre>
             delta = minv[j];
             j1 = j;
      assert (j1 != -1);
      FOR(j, 0, m + 1)
        if (used[j])
          u[p[j]] += delta;
          v[j] -= delta;
        else
          minv[j] -= delta;
      j0 = j1;
    while (j0 != m)
      int j1 = way[j0];
      p[j0] = p[j1];
      j0 = j1;
 VI ans(n + 1);
 FOR(j, 0, m)
    ans[p[j]] = j;
  LL res = 0;
 FOR(i, 0, n)
    res += a[i][ans[i]];
  assert (res == -v[m]);
  return res;
edmonds-blossom.hpp
                                          76c9ac, 127 lines
struct Graph
 int n;
 vector<VI> q;
 VI label, first, mate;
  Graph() {}
```

```
Graph(int _n): n(_n), g(_n + 1), label(_n + 1),
    first (n + 1), mate (n + 1) {}
void addEdge(int u, int v)
  assert(0 <= u && u < n);
  assert (0 \le v \& v \le n);
  u++;
  v++;
  q[u].PB(v);
  q[v].PB(u);
void augmentPath(int v, int w)
  int t = mate[v];
  mate[v] = w;
  if(mate[t] != v)
    return;
  if(label[v] \le n)
    mate[t] = label[v];
    augmentPath(label[v], t);
    return;
  int x = label[v] / (n + 1);
  int y = label[v] % (n + 1);
  augmentPath(x, y);
  augmentPath(y, x);
int findMaxMatching()
  FOR(i, 0, n + 1)
    assert(mate[i] == 0);
  int mt = 0;
  DSU dsu;
  FOR(u, 1, n + 1)
    if(mate[u] == 0)
      fill(ALL(label), -1);
      iota(ALL(first), 0);
      dsu.init(n + 1);
      label[u] = 0;
      dsu.unite(u, 0);
      queue<int> q;
      q.push(u);
      while(!q.empty())
        int x = q.front();
        q.pop();
```

```
for(int y: g[x])
  if (mate[y] == 0 && y != u)
   mate[y] = x;
    augmentPath(x, y);
    while(!q.empty())
     q.pop();
   mt++;
   break;
  if(label[y] < 0)
    int v = mate[v];
    if(label[v] < 0)
     label[v] = x;
     dsu.unite(v, y);
     q.push(v);
  }
  else
    int r = first[dsu.find(x)], s =
       first[dsu.find(y)];
    if(r != s)
      int edgeLabel = (n + 1) * x + y;
      label[r] = label[s] = -edgeLabel;
      int join;
      while(true)
        if(s != 0)
          swap(r, s);
        r = first[dsu.find(label[mate[r
           ]])];
        if(label[r] == -edgeLabel)
          join = r;
          break;
        label[r] = -edgeLabel;
      for(int z: {x, y})
        for(int v = first[dsu.find(z)];
            v != join;
          v = first[dsu.find(label[mate
              [v]]))
```

```
label[v] = edgeLabel;
                    if (dsu.unite(v, join))
                       first[dsu.find(join)] =
                           join;
                    q.push(v);
            }
   return mt;
  int getMate(int v)
   assert (0 \leq v && v \leq n);
   v++;
   int u = mate[v];
   assert(u == 0 || mate[u] == v);
   u--;
   return u;
 }
};
```

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Strings (4)

```
Aho-Corasick.hpp
```

```
b80725, 70 lines
```

```
const int AL = 26;
struct Node
  int p;
  int c;
  int g[AL];
  int nxt[AL];
  int link;
 void init()
   c = -1;
   p = -1;
   fill(q, q + AL, -1);
   fill(nxt, nxt + AL, -1);
   link = -1;
 }
};
struct AC
 Node A[MAX];
  int sz;
  void init()
   A[0].init();
   sz = 1;
  int addStr(const string& s)
   int x = 0;
   FOR (i, 0, SZ(s))
      int c = s[i] - 'A'; // change to [0 AL)
      if (A[x].nxt[c] == -1)
       A[x].nxt[c] = sz;
       A[sz].init();
       A[sz].c = c;
       A[sz].p = x;
        sz++;
      x = A[x].nxt[c];
    return x;
```

```
int go(int x, int c)
    if (A[x].q[c] != -1)
      return A[x].g[c];
    if (A[x].nxt[c] != -1)
     A[x].q[c] = A[x].nxt[c];
    else if (x != 0)
     A[x].q[c] = qo(qetLink(x), c);
    else
     A[x].q[c] = 0;
    return A[x].q[c];
  int getLink(int x)
    if (A[x].link != -1)
      return A[x].link;
   if (x == 0 | | A[x].p == 0)
      return 0;
    return A[x].link=go(getLink(A[x].p), A[x].c);
} A;
automaton.hpp
                                         2a29a0, 60 lines
const int AL = 26;
struct Node
 int g[AL];
 int link;
 int len;
 int cnt;
 void init()
   fill(g, g + AL, -1);
   link = -1;
    len = -1;
};
struct Automaton
 Node A[MAX * 2];
 int sz;
  int head;
 void init()
   sz = 1;
    head = 0;
   A[0].init();
```

```
void add(char c)
    int ch = c - 'a'; // change to [0 AL)
    int nhead = sz++;
    A[nhead].init();
    A[nhead].len = A[head].len + 1;
    int cur = head;
    head = nhead;
    while (cur != -1 && A[cur].g[ch] == -1)
      A[cur].q[ch] = head;
      cur = A[cur].link;
    if (cur == -1)
      A[head].link = 0;
      return;
    int p = A[cur].g[ch];
    if (A[p].len == A[cur].len + 1)
      A[head].link = p;
      return;
    int q = sz++;
    A[q] = A[p];
    A[q].len = A[cur].len + 1;
    A[p].link = A[head].link = q;
    while (cur != -1 && A[cur].g[ch] == p)
     A[cur].q[ch] = q;
      cur = A[cur].link;
};
suffix-array.hpp
                                         d69c2d, 60 lines
void countSort(VI& p, const VI& c)
  int n = SZ(p);
  VI cnt(n);
  FOR(i, 0, n)
   cnt[c[i]]++;
  VI pos(n);
  FOR (i, 1, n)
   pos[i] = pos[i - 1] + cnt[i - 1];
  VI p2(n);
  for (auto x : p)
```

```
int i = c[x];
   p2[pos[i]++] = x;
 p = p2;
VI suffixArray(const string& t)
 string s = t + "$";
 int n = SZ(s);
 VI p(n), c(n);
 FOR (i, 0, n) p[i] = i;
  sort(ALL(p), [&](int i, int j)
   return s[i] < s[j];
 });
 int x = 0;
 c[p[0]] = 0;
 FOR (i, 1, n)
   if (s[p[i]] != s[p[i-1]])
    x++;
   c[p[i]] = x;
  int k = 0;
  while ((1 << k) < n)
   FOR (i, 0, n)
     p[i] = (p[i] - (1 << k) + n) % n;
   countSort(p, c);
   VI c2(n);
   PII pr = {c[p[0]], c[(p[0] + (1 << k)) % n]};
   FOR (i, 1, n)
     PII nx = \{c[p[i]], c[(p[i] + (1 << k)) % n]\};
     c2[p[i]] = c2[p[i - 1]];
     if (pr != nx)
       c2[p[i]]++;
     pr = nx;
   c = c2;
   k++;
  p.erase(p.begin());
  return p;
```

```
lcp.hpp
                                          72ff1e, 24 lines
VI lcpArray(const string& s, const VI& sa)
   int n = SZ(s);
   VI rnk(n);
    FOR (i, 0, n)
        rnk[sa[i]] = i;
   VI lcp(n-1);
   int h = 0;
    FOR (i, 0, n)
        if (h > 0)
      h--;
        if (rnk[i] == 0)
      continue;
        int j = sa[rnk[i] - 1];
        for (; j + h < n && i + h < n; h++)</pre>
            if (s[i + h] != s[i + h])
        break;
        lcp[rnk[i] - 1] = h;
    return lcp;
z.hpp
                                          e27ac7, 23 lines
VI zFunction(const string& s)
 int n = SZ(s);
 VI z(n);
  int 1 = 0;
  int r = 0;
  FOR (i, 1, n)
   z[i] = 0;
    if (i <= r)
     z[i] = min(r - i + 1, z[i - 1]);
    while(i + z[i] < n \&\& s[i + z[i]] == s[z[i]])
      z[i]++;
    if(i + z[i] - 1 > r)
      r = i + z[i] - 1;
     1 = i;
  return z;
```

```
prefix.hpp
                                         500608, 16 lines
VI prefixFunction(const string& s)
  int n = SZ(s);
  VI p(n);
  p[0] = 0;
  FOR (i, 1, n)
    int j = p[i - 1];
    while(j != 0 && s[i] != s[j])
     j = p[j - 1];
    if (s[i] == s[j]) j++;
    p[i] = j;
  return p;
manacher.hpp
                                         9f7ea5, 39 lines
int d1[MAX], d2[MAX];
void manacher(const string& s)
    int n = SZ(s);
    int 1 = -1;
    int r = -1;
    FOR (i, 0, n)
        if (i <= r)
            d1[i] = min(r - i + 1,
              d1[1 + (r - i)];
        while (i + d1[i] < n \&\& i - d1[i] >= 0
          && s[i + d1[i]] == s[i - d1[i]])
            d1[i]++;
        if (i + d1[i] - 1 > r)
            r = i + d1[i] - 1;
            1 = i - (d1[i] - 1);
        }
    1 = -1;
    r = -1;
    FOR (i, 0, n)
        if (i <= r)
            d2[i] = min(r - i + 1,
              d2[1 + (r - i) + 1]);
```

```
while (i + d2[i] < n
          && i - (d2[i] + 1) >= 0
          && s[i + d2[i]] == s[i - (d2[i] + 1)])
            d2[i]++;
        if (i + d2[i] > r)
            r = i + d2[i] - 1;
            1 = i - d2[i];
palindrome-tree.hpp
                                         a5d480, 64 lines
struct Node
  int to[ALP];
  int link;
  int len;
  void clear()
   fill(to, to + ALP, -1);
   link = -1;
   len = -1;
};
struct PalTree
  string s;
  vector<Node> A;
  int sz;
  int last;
  void init(string t)
   A.resize(2 * SZ(t));
    A[0].clear();
    A[1].clear();
    A[1].len = 0;
    A[1].link = 0;
    sz = 2;
    last = 1;
    s = t;
  void add(int idx)
    int cur = last;
    while (cur !=-1)
      int pos = idx - A[cur].len - 1;
      if (pos >= 0 && s[pos] == s[idx])
```

```
break;
     cur = A[cur].link;
   assert (cur !=-1);
   if (A[cur].to[s[idx] - 'a'] == -1)
     A[cur].to[s[idx] - 'a'] = sz;
     A[sz].clear();
     A[sz].len = A[cur].len + 2;
     int link = A[cur].link;
     while (link !=-1)
       int pos = idx - A[link].len - 1;
       if (pos >= 0 && s[pos] == s[idx])
         break;
       link = A[link].link;
     if (link == -1)
       link = 1;
     else
       link = A[link].to[s[idx] - 'a'];
     A[sz].link = link;
     sz++;
   last = A[cur].to[s[idx] - 'a'];
} pt;
```

Geometry (5)

In general, try to build programs that are resistant to the oddities of floating-point numbers. Imagine that some evil demon is slightly modifying every result you compute in the way that is most likely to make your program fail. And try to write clean code that is clearly correct at first glance. If you need long explanations to justify why your program will not fail, then it is more likely that your program will in fact fail.

Victor Lecomte, Handbook of geometry for competitive programmers

geometry.hpp

f709e0, 386 lines

```
struct Pt
  db x, y;
  Pt operator+(const Pt& p) const
   return {x + p.x, y + p.y};
  Pt operator-(const Pt& p) const
   return {x - p.x, y - p.y};
  Pt operator* (db d) const
   return {x * d, y * d};
  Pt operator/(db d) const
   return {x / d, y / d};
// Returns the squared absolute value
db sq(const Pt& p)
 return p.x * p.x + p.y * p.y;
// Returns the absolute value
db abs(const Pt& p)
 return sqrt(sq(p));
// Returns -1 for negative numbers, 0 for zero,
// and 1 for positive numbers
int sqn(db x)
```

```
return (EPS < x) - (x < -EPS);
// Returns 'p' rotated counter-clockwise by 'a'
Pt rot(const Pt& p, db a)
  db co = cos(a), si = sin(a);
  return {p.x * co - p.y * si,
    p.x * si + p.y * co;
// Returns 'p' rotated counter-clockwise by 90
Pt perp(const Pt& p)
  return {-p.y, p.x};
// Returns the dot product of 'p' and 'q'
db dot (const Pt& p, const Pt& q)
 return p.x * q.x + p.y * q.y;
// Returns the angle between 'p' and 'q'
db angle (const Pt& p, const Pt& q)
 return acos(clamp(dot(p, q) / abs(p) /
    abs(q), (db)-1.0, (db)1.0);
// Returns the cross product of 'p' and 'q'
db cross(const Pt& p, const Pt& q)
 return p.x * q.y - p.y * q.x;
// Positive if R is on the left side of PQ,
// negative on the right side,
// and zero if R is on the line containing PQ
db orient(const Pt& p, const Pt& q, const Pt& r)
  return cross (q - p, r - p) / abs <math>(q - p);
// Checks if a polygon 'v' is convex
bool isConvex(const vector<Pt>& v)
 bool hasPos = false, hasNeg = false;
 int n = SZ(v);
 FOR(i, 0, n)
    int o = sgn(orient(v[i], v[(i + 1) % n],
     v((i + 2) % n));
    hasPos |= o > 0;
    hasNeq |= o < 0;
```

```
return ! (hasPos && hasNeg);
// Checks if argument of 'p' is in [-pi, 0]
bool half(const Pt& p)
  assert (sqn(p.x) != 0 || sqn(p.y) != 0);
  return sqn(p.y) == -1 \mid \mid
    (sgn(p.v) == 0 \&\& sgn(p.x) == -1);
// Polar sort of vectors in 'v' around 'o'
void polarSortAround(const Pt& o, vector<Pt>& v)
  sort (ALL(v), [o] (const Pt& p, const Pt& q)
    bool hp = half (p - o), hq = half (q - o);
    if (hp != hq)
      return hp < hq;
    int s = sgn(cross(p, q));
    if (s != 0)
      return s == 1;
    return sq(p - o) < sq(q - o);
  });
// Returns the distance of the closest points
db closestPair(vector<Pt> v)
  sort (ALL(v), [] (const Pt& p, const Pt& q)
    return sgn(p.x - q.x) < 0;
  });
  set<pair<db, db>> s;
  int n = SZ(v), ptr = 0;
  db h = 1e18;
  FOR(i, 0, n)
    for (auto it = s.lower bound(
      MP(v[i].y - h, v[i].x)); it != s.end()
      && sgn(it->F - (v[i].y + h)) <= 0; it++)
      Pt q = \{it->S, it->F\};
      h = min(h, abs(v[i] - q));
    for (; sgn(v[ptr].x - (v[i].x - h)) <= 0;</pre>
      s.erase({v[ptr].y, v[ptr].x});
    s.insert(\{v[i].y, v[i].x\});
  return h;
// Example:
```

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```
// \ cout << a + b << " " << a - b << " ";
ostream& operator<<(ostream& os, const Pt& p)
 return os << "(" << p.x << "," << p.y << ")";
struct Line
  // Equation of the line is dot(n, p) + c = 0
 Pt n;
 db c;
 // The line containing two points 'p' and 'q'
 Line (const Pt& p, const Pt& q):
   n(perp(q - p)), c(-dot(n, p)) {}
 // The "positive side": dot(n, p) + c > 0
 // The "negative side": dot(n, p) + c < 0
  db side (const Pt& p) const
   return dot(n, p) + c;
  // Returns the distance from 'p'
  db dist(const Pt& p) const
   return abs(side(p)) / abs(n);
  // Returns the squared distance from 'p'
  db sqDist(const Pt& p) const
   return side(p) * side(p) / (db)sq(n);
  // Returns the perpendicular line through 'p'
 Line perpThrough(const Pt& p) const
   return {p, p + n};
  // Compares 'p' and 'q' by their projection
 bool cmpProj(const Pt& p, const Pt& q) const
   return sgn(cross(p, n) - cross(q, n)) < 0;</pre>
 // Returns the orthogonal projection of 'p'
 Pt proj(const Pt& p) const
   return p - n * side(p) / sq(n);
  // Returns the reflection of 'p' by the line
 Pt refl(const Pt& p) const
   return p - n * 2 * side(p) / sq(n);
};
```

```
// Checks if 'll' and 'l2' are parallel
bool parallel (const Line& 11, const Line& 12)
 return sqn(cross(11.n, 12.n)) == 0;
// Returns the intersection point
Pt inter(const Line& 11, const Line& 12)
 db d = cross(11.n, 12.n);
 assert (sgn(d) != 0);
 return perp(12.n * 11.c - 11.n * 12.c) / d;
// Checks if 'p' is in the disk of diameter [ab]
bool inDisk (const Pt& a, const Pt& b,
 const Pt& p)
 return sgn(dot(a - p, b - p)) <= 0;
// Checks if 'p' lies on segment [ab]
bool onSegment (const Pt& a, const Pt& b,
  const Pt& p)
 return sgn(orient(a, b, p)) == 0
    && inDisk(a, b, p);
// Checks if the segments [ab] and [cd] intersect
// properly (their intersection is one point
// which is not an endpoint of either segment)
bool properInter(const Pt& a, const Pt& b,
 const Pt& c, const Pt& d)
 db oa = orient(c, d, a);
 db ob = orient(c, d, b);
 db oc = orient(a, b, c);
 db od = orient(a, b, d);
 return sqn(oa) * sqn(ob) == -1
    && sgn(oc) * sgn(od) == -1;
// Returns the distance between [ab] and 'p'
db seqPt(const Pt& a, const Pt& b, const Pt& p)
 Line l(a, b);
  assert(sgn(sq(1.n)) != 0);
 if (l.cmpProj(a, p) && l.cmpProj(p, b))
   return l.dist(p);
 return min(abs(p - a), abs(p - b));
// Returns the distance between [ab] and [cd]
db segSeg(const Pt& a, const Pt& b, const Pt& c,
 const Pt& d)
```

```
if (properInter(a, b, c, d))
    return 0;
  return min({segPt(a, b, c), segPt(a, b, d),
      segPt(c, d, a), segPt(c, d, b)});
// Returns the area of triangle abc
db areaTriangle (const Pt& a, const Pt& b,
  const Pt& c)
  return abs (cross (b - a, c - a)) / 2.0;
// Returns the area of polygon 'v'
db areaPolygon(const vector<Pt>& v)
  db area = 0.0;
  int n = SZ(v);
  FOR(i, 0, n)
    area += cross(v[i], v[(i + 1) % n]);
  return abs(area) / 2.0;
// Returns true if 'p' is at least as high as 'a'
bool above (const Pt& a, const Pt& p)
  return sqn(p.y - a.y) >= 0;
// Checks if [pq] crosses the ray from 'a'
bool crossesRay(const Pt& a, const Pt& p,
  const Pt& q)
  return sqn((above(a, q) - above(a, p))
    * orient(a, p, q)) == 1;
// Checks if point 'a' is inside a polygon
// If 'strict', false when 'a' is on the boundary
bool inPolygon (const vector < Pt > & v, const Pt & a,
  bool strict = true)
  int numCrossings = 0;
  int n = SZ(v);
  FOR(i, 0, n)
  {
    if (onSegment(v[i], v[(i + 1) % n], a))
      return !strict;
    numCrossings +=
      crossesRay(a, v[i], v[(i + 1) % n]);
  return numCrossings & 1;
// Returns the counter-clockwise convex hull
```

```
vector<Pt> convexHull(vector<Pt> v)
  sort (ALL(v), [] (const Pt& p, const Pt& q)
   int dx = sgn(p.x - q.x);
   if (dx != 0)
     return dx < 0;
   return sgn(p.y - q.y) < 0;
 });
  vector<Pt> lower, upper;
  for (const Pt& p : v)
   while (SZ(lower) > 1
      && sgn(orient(lower[SZ(lower) - 2],
     lower.back(), p)) < 0)
     lower.pop_back();
   while (SZ(upper) > 1
      && sgn(orient(upper[SZ(upper) - 2],
     upper.back(), p)) > 0)
     upper.pop_back();
   lower.PB(p);
   upper.PB(p);
  reverse (ALL (upper));
 lower.insert(lower.end(), upper.begin() + 1,
   prev(upper.end()));
 return lower;
// Returns the circumcenter of triangle abc
Pt circumCenter(const Pt& a, Pt b, Pt c)
 b = b - a;
 c = c - a;
 assert (sgn(cross(b, c)) != 0);
 return a + perp(b * sq(c) - c * sq(b))
   / cross(b, c) / 2;
// Returns circle-line intersection points
vector<Pt> circleLine(const Pt& o, db r,
  const Line& 1)
 db h2 = r * r - l.sqDist(o);
 if (sgn(h2) == -1)
   return {};
 Pt p = 1.proj(0);
 if (sqn(h2) == 0)
   return {p};
 Pt h = perp(l.n) * sqrt(h2) / abs(l.n);
 return {p - h, p + h};
```

```
// Returns circle-circle intersection points
vector<Pt> circleCircle(const Pt& o1, db r1,
  const Pt& o2, db r2)
 Pt d = 02 - 01;
 db d2 = sq(d);
 if (sqn(d2) == 0)
   assert (sgn(r2 - r1) != 0);
   return {};
 db pd = (d2 + r1 * r1 - r2 * r2) / 2;
 db h2 = r1 * r1 - pd * pd / d2;
 if (sqn(h2) == -1)
   return {};
 Pt p = o1 + d * pd / d2;
 if (sqn(h2) == 0)
   return {p};
 Pt h = perp(d) * sqrt(h2 / d2);
 return \{p - h, p + h\};
// Finds common tangents (outer or inner)
// If there are 2 tangents, returns the pairs of
// tangency points on each circle (p1, p2)
// If there is 1 tangent, the circles are tangent
// to each other at some point p, res contains p
// 4 times, and the tangent line can be found as
// line (o1, p). perp Through (p)
// The same code can be used to find the tangent
// to a circle through a point by setting r2 to 0
// (in which case 'inner' doesn't matter)
vector<pair<Pt, Pt>> tangents (const Pt& o1,
 db r1, const Pt& o2, db r2, bool inner)
 if (inner)
   r2 = -r2;
 Pt d = 02 - 01;
 db dr = r1 - r2, d2 = sq(d),
   h2 = d2 - dr * dr;
 if (sqn(d2) == 0 | | sqn(h2) < 0)
   assert (sgn(h2) != 0);
   return {};
 vector<pair<Pt, Pt>> res;
 for (db sign : {-1, 1})
   Pt v = (d * dr + perp(d) * sqrt(h2)
     * sign) / d2;
   res.PB(\{01 + v * r1, 02 + v * r2\});
```

}
return res;
}

Math (6)

```
gcd.hpp
                                            e001bc, 16 lines
int gcd(int a, int b, int& x, int& y)
  x = 1, y = 0;
  int x2 = 0, y2 = 1;
  while (b)
    int k = a / b;
    x -= k * x2;
    y -= k * y2;
    a %= b;
    swap(a, b);
    swap(x, x2);
    swap(y, y2);
  return a;
fast-chinese.hpp
Description: x \% p_i = m_i, lcm(p) < 10^{18}, p < 10^9.
no solution -> return -1
Time: \mathcal{O}(nloq)
                                            9b392a, 25 lines
LL FastChinese(VI m, VI p)
  assert(SZ(m) == SZ(p));
  LL aa = p[0];
  LL bb = m[0];
  FOR(i, 1, SZ(m))
    int b = (m[i] - bb % p[i] + p[i]) % p[i];
    int a = aa % p[i];
    int c = p[i];
    int x, y;
    int d = gcd(a, c, x, y);
    if(b % d != 0)
      return -1;
    a /= d;
    b /= d;
    c /= d;
    b = b * (LL)x % c;
    bb = aa * b + bb;
    aa = aa * c;
  return bb;
```

```
chinese.hpp
Description: result < 10^{18}, p < 10^9.
Time: \mathcal{O}\left(n^2\right)
                                               c692b5, 36 lines
LL Chinese (VI m, VI p)
  int n = SZ(m);
  FOR(i, 1, n)
    LL a = 1;
    LL b = 0;
    RFOR(j, i, 0)
      b = (b * p[j] + m[j]) % p[i];
      a = a * p[j] % p[i];
    b = (m[i] - b + p[i]) % p[i];
    int c = p[i];
    int x, y;
    int d = gcd(a, c, x, y);
    if(b % d != 0)
       return -1;
    a /= d;
    b /= d;
    c /= d;
    b = b * x % c;
    m[i] = b;
    p[i] = c;
  LL res = m[n - 1];
  RFOR(i, n - 1, 0)
    res *= p[i];
    res += m[i];
  return res;
gauss.hpp
Description: a[i].back() is right side element
Time: \mathcal{O}\left(m^2*n\right)
                                               e404cd, 50 lines
VI Gauss (vector<VI> a)
 int n = SZ(a);
 if(n == 0)
    return {};
```

int m = SZ(a[0]) - 1; //number of variables

```
assert (n >= m);
  int vars = m;
  FOR(i, 0, m)
    if (a[i][i] == 0)
      //for double find row
      //with max abs value
      int row = -1;
      FOR(k, i + 1, n)
        if (a[k][i] != 0)
          row = k;
      if(row == -1)
        //variable i can be any
        vars--;
        continue;
      swap(a[i], a[row]);
    int d = inv(a[i][i]);
    FOR(k, i + 1, n)
      int c = mult(a[k][i], d);
      FOR(j, 0, m + 1)
        updSub(a[k][j], mult(c, a[i][j]));
    }
  FOR(i, vars, n)
    if(a[i].back() != 0)
      cout << "No solution\n";</pre>
  VI \times (m);
  RFOR(i, m, 0)
    x[i] = a[i].back();
    FOR(j, i + 1, m)
      updSub(x[i], mult(a[i][j], x[j]));
    x[i] = mult(x[i], inv(a[i][i]));
  return x;
miller-rabin.hpp
// to speed up change candidates to at least 4
    random values rng() \% (a - 3) + 2;
// use __int128 in mult
```

```
29, 31, 47};
bool MillerRabin(LL a)
  if (a == 1)
    return false;
  if (a == 2 || a == 3)
    return true;
  LL d = a - 1;
  int s = __builtin_ctzll(d);
  d >>= s;
  for (LL b : candidates)
    if (b >= a)
      break;
    b = binpow(b, d, a);
    if (b == 1)
      continue;
    bool ok = false;
    FOR (i, 0, s)
      if (b + 1 == a)
        ok = true;
        break;
      b = mult(b, b, a);
    if (!ok)
      return false;
  return true;
pollard.hpp
// uses Miller-Rabin test
// rho finds divisor of n
// use __int128 in mult
// works in O(n^{(1/4)} * log n)
LL f(LL x, LL c, LL n)
  return add(mult(x, x, n), c, n);
LL rho(LL n)
  const int iter = 47 * sqrt(sqrt(n));
  while (true)
```

VI candidates = $\{2, 3, 5, 7, 11, 13, 17, 19, 23,$

```
LL \times 0 = rng() % n;
    LL c = rng() % n;
    LL x = x0;
    LL y = x0;
    LL q = 1;
    FOR (i, 0, iter)
      x = f(x, c, n);
      y = f(y, c, n);
      y = f(y, c, n);
      g = gcd(abs(x - y), n);
      if (g != 1)
        break;
    if (q > 1 \&\& g < n)
      return g;
VI primes = \{2, 3, 5, 7, 11, 13, 17, 19, 23\};
vector<LL> factorize(LL n)
 vector<LL> ans;
  for (auto p : primes)
    while (n % p == 0)
      ans.PB(p);
      n /= p;
    }
  queue<LL> q;
  q.push(n);
  while (!q.empty())
    LL x = q.front();
    q.pop();
    if (x == 1)
      continue;
    if (MillerRabin(x))
      ans.PB(x);
    else
      LL y = rho(x);
      q.push(y);
      q.push(x / y);
```

```
return ans;
simplex.hpp
Description: c^T x \to \max, Ax < b, x > 0.
                                          3805fb, 142 lines
struct Simplex
private:
  int m, n;
  VI nonBasic, basic;
  vector<vector<db>> a;
  vector<db> b;
 vector<db> c;
  db v;
public:
  void pivot(int 1, int e)
    assert(0 <= 1 \&\& 1 < m);
    assert (0 \le e \&\& e \le n);
    assert(abs(a[l][e]) > EPS);
    b[1] /= a[1][e];
    FOR(j, 0, n)
     if (j != e)
        a[1][j] /= a[1][e];
    a[1][e] = 1 / a[1][e];
    FOR(i, 0, m)
      if (i != 1)
        b[i] -= a[i][e] * b[l];
        FOR(j, 0, n)
          if (j != e)
            a[i][j] -= a[i][e] * a[l][j];
        a[i][e] *= -a[l][e];
      }
    v += c[e] * b[1];
    FOR(j, 0, n)
      if († != e)
        c[j] -= c[e] * a[l][j];
    c[e] *= -a[l][e];
    swap(nonBasic[e], basic[l]);
  void findOptimal()
    vector<db> delta(m);
    while (true)
```

```
int e = -1;
    FOR(j, 0, n)
      if (c[j] > EPS \&\& (e == -1 || nonBasic[j])
           < nonBasic[e]))
        e = i;
    if (e == -1)
      break:
    FOR(i, 0, m)
      delta[i] = a[i][e] > EPS ? b[i] / a[i][e]
          : LINF;
    int 1 = min_element(ALL(delta)) - delta.
       begin();
    if (delta[1] == LINF)
      // unbounded
      assert (false);
    pivot(l, e);
void initializeSimplex(const vector<vector<db</pre>
   >>& _a, const vector<db>& _b, const vector<
   db>& c)
 m = SZ(b);
 n = SZ(c);
 nonBasic.resize(n);
 iota(ALL(nonBasic), 0);
 basic.resize(m);
 iota(ALL(basic), n);
 a = _a;
 b = _b;
 c = _c;
 v = 0;
  int k = min_element(ALL(b)) - b.begin();
 if (b[k] > -EPS)
    return;
 nonBasic.PB(n);
  iota(ALL(basic), n + 1);
 FOR(i, 0, m)
   a[i].PB(-1);
 c.assign(n, 0);
  c.PB(-1);
 n++;
  pivot(k, n - 1);
  findOptimal();
  if (v < -EPS)
    // infeasible
    assert (false);
```

```
int l = find(ALL(basic), n - 1) - basic.begin
      ();
  if (1 != m)
    int e = -1;
    while (abs(a[l][e]) < EPS)</pre>
     e++;
   pivot(l, e);
  int p = find(ALL(nonBasic), n) - nonBasic.
     begin();
  assert (p < n + 1);
  nonBasic.erase(nonBasic.begin() + p);
  FOR(i, 0, m)
    a[i].erase(a[i].begin() + p);
  c.assign(n, 0);
  FOR(j, 0, n)
    if (nonBasic[i] < n)</pre>
      c[j] = \_c[nonBasic[j]];
    else
      nonBasic[j]--;
  FOR(i, 0, m)
    if (basic[i] < n)
     v += _c[basic[i]] * b[i];
     FOR(j, 0, n)
        c[j] = c[basic[i]] * a[i][j];
    else
      basic[i]--;
pair<vector<db>, db> simplex(const vector<
   vector<db>& _a, const vector<db>& _b,
   const vector<db>& _c)
  initializeSimplex(_a, _b, _c);
  assert(SZ(a) == m);
  FOR(i, 0, m)
    assert(SZ(a[i]) == n);
  assert(SZ(b) == m);
  assert(SZ(c) == n);
  assert(SZ(nonBasic) == n);
  assert(SZ(basic) == m);
  findOptimal();
```

```
vector<db> x(n);
FOR(i, 0, m)
   if (basic[i] < n)
      x[basic[i]] = b[i];
return {x, v};
}
};</pre>
```

Convolutions (7)

```
fft.hpp

const int mod = 998244353;
```

```
44d94c, 87 lines
int add(int a, int b)
  return (a + b < mod) ? (a + b) : (a + b - mod);
int sub(int a, int b)
  return (a - b \ge 0) ? (a - b) : (a - b + mod);
int mult(int a, int b)
  return a * (LL) b % mod;
int binpow(int a, int n)
  int res = 1;
  while (n)
    if(n & 1)
      res = mult(res, a);
    a = mult(a, a);
    n /= 2;
  return res;
const int LEN = 1 << 23;</pre>
const int GEN = 31;
const int IGEN = binpow(GEN, mod - 2);
void fft(VI& a, bool inv)
 int lq = 0;
  while((1 << lg) < SZ(a)) lg++;</pre>
  FOR(i, 0, SZ(a))
    int x = 0;
```

```
FOR(j, 0, lq)
      x = ((i >> j) \& 1) << (lq - j - 1);
   if(i < x)
      swap(a[i], a[x]);
  for (int len = 2; len <= SZ(a); len *= 2)
   int ml = binpow(inv ? IGEN : GEN, LEN / len);
   for(int i = 0; i < SZ(a); i += len)</pre>
      int pw = 1;
      FOR(j, 0, len / 2)
        int v = a[i + j];
        int u = mult(a[i + j + len / 2], pw);
        a[i + j] = add(v, u);
        a[i + j + len / 2] = sub(v, u);
        pw = mult(pw, ml);
   }
  if(inv)
   int m = binpow(SZ(a), mod - 2);
   FOR(i, 0, SZ(a))
     a[i] = mult(a[i], m);
 }
VI mult (VI a, VI b)
 int sz = 0;
  int sum = SZ(a) + SZ(b) - 1;
 while((1 << sz) < sum) sz++;
  a.resize(1 << sz);
 b.resize(1 << sz);
  fft(a, 0);
  fft(b, 0);
 FOR(i, 0, SZ(a))
   a[i] = mult(a[i], b[i]);
  fft(a, 1);
  a.resize(sum);
  return a;
```

```
inverse.hpp
                                         a4673f, 32 lines
VI inverse (const VI& a, int k)
 assert(SZ(a) == k \&\& a[0] != 0);
 if(k == 1)
    return {binpow(a[0], mod - 2)};
 VI ra = a;
 FOR(i, 0, SZ(ra))
   if(i & 1)
      ra[i] = sub(0, ra[i]);
  int nk = (k + 1) / 2;
  VI t = mult(a, ra);
  t.resize(k);
 FOR(i, 0, nk)
   t[i] = t[2 * i];
 t.resize(nk);
 t = inverse(t, nk);
 t.resize(k);
 RFOR(i, nk, 1)
   t[2 * i] = t[i];
   t[i] = 0;
 VI res = mult(ra, t);
 res.resize(k);
 return res;
exp-log.hpp
                                         5549eb, 52 lines
VI deriv(const VI& a, int k)
 VI res(k);
 FOR(i, 0, k)
   if(i + 1 < SZ(a))
     res[i] = mult(a[i + 1], i + 1);
 return res;
VI integr(const VI& a, int k)
 VI res(k);
 RFOR(i, k, 1)
   res[i] = mult(a[i - 1], inv[i]);
```

```
VI log(const VI& a, int k)
  assert (a[0] == 1);
 VI ml = mult(deriv(a, k), inverse(a, k));
 return integr(ml, k);
VI exp(VI a, int k)
  assert(a[0] == 0);
  VI Qk = \{1\};
  int pw = 1;
  while(pw <= k)</pre>
    pw *= 2;
    Qk.resize(pw);
    VI lnQ = log(Qk, pw);
    FOR(i, 0, SZ(lnQ))
      if(i < SZ(a))
        lnQ[i] = sub(a[i], lnQ[i]);
      else
        lnQ[i] = sub(0, lnQ[i]);
    updAdd(lnQ[0], 1);
    Qk = mult(Qk, lnQ);
  Ok.resize(k);
  return Ok;
modulo.hpp
                                         8b6a95, 34 lines
void removeLeadingZeros(VI& a)
 while (SZ(a) > 0 \&\& a.back() == 0)
    a.pop_back();
pair<VI, VI> modulo(VI a, VI b)
  //assert(a.back() != 0 \& b.back() != 0);
  int n = SZ(a), m = SZ(b);
```

res[0] = 0;

return res;

```
if(m > n)
  return MP(VI{}, a);
reverse(ALL(a));
reverse (ALL(b));
VI d = b;
d.resize(n - m + 1);
d = mult(a, inverse(d, n - m + 1));
d.resize(n - m + 1);
reverse (ALL(a));
reverse (ALL(b));
reverse (ALL(d));
VI res = mult(b, d);
res.resize(SZ(a));
FOR(i, 0, SZ(a))
  res[i] = sub(a[i], res[i]);
removeLeadingZeros(d);
removeLeadingZeros(res);
return MP (d, res);
```

multipoint-eval.hpp

```
8f6f41, 33 lines
int x[LEN];
VI P[2 * LEN];
void build(int v, int tl, int tr)
 if(t1 + 1 == tr)
   P[v] = {sub(0, x[t1]), 1};
    return;
  int tm = (tl + tr) / 2;
 build(2 * v + 1, tl, tm);
 build(2 * v + 2, tm, tr);
 P[v] = mult(P[2 * v + 1], P[2 * v + 2]);
int ans[LEN];
void solve(int v, int tl, int tr, const VI& Q)
//Q != Q \% P[0] \implies wa
 if(SZ(O) == 0)
    return;
 if(t1 + 1 == tr)
```

```
ans[t1] = Q[0];
    return;
 int tm = (tl + tr) / 2;
  solve(2 * v + 1, t1, tm,
 modulo(Q, P[2 * v + 1]).S);
 solve (2 * v + 2, tm, tr,
 modulo(0, P[2 * v + 2]).S);
newton.hpp
                                         9ffaac, 50 lines
VI newton(VI a, int k)
 //c_n = a_n + sum(i = 0, n - 1) c_i * c_i(n-1-i)
 //Q = A + x * Q * Q
 //F(Q) = Q - x * Q * Q - A
 //F'(Q) = 1 - 2 * x * Q
 VI Qk = {a[0]};
 int pw = 1;
  while (pw <= k)
    assert(SZ(Qk) == pw);
    pw *= 2;
    VI F1(pw);
    F1[0] = 1;
    FOR(i, 0, pw / 2)
     F1[i + 1] = sub(0, mult(2, Qk[i]));
    //F' = 1 - 2 * x * Q
    VI F = mult(Qk, Qk);
    F.resize(pw);
    RFOR(i, pw, 1)
     F[i] = sub(0, F[i - 1]);
    F[0] = 0; // F = -x * Q*Q
    FOR(i, 0, pw / 2)
     F[i] = add(F[i], Qk[i]);
     //F = Q - x * Q * Q
    FOR(i, 0, min(pw, SZ(a)))
     F[i] = sub(F[i], a[i]);
     //F = Q - x * Q * Q - A
    F = mult(F, inverse(F1, pw));
    F.resize(pw);
    FOR(i, 0, pw)
     F[i] = sub(0, F[i]); //-F/F'
    FOR(i, 0, pw / 2)
```

```
F[i] = add(F[i], Qk[i]); //Q - F/F'
  //new Qk = Qk - F(Qk) / F'(Qk) mod(x ^ pw)
  Qk = F;
Qk.resize(k);
return Qk;
```

berlekamp-massey.hpp

866c28, 36 lines

20

```
VI berlekampMassey(const VI& a)
 VI c = \{1\}, bp = \{1\};
 int 1 = 0, b = 1, x = 1;
  FOR(i, 0, SZ(a))
    assert(SZ(c) == 1 + 1);
    int d = a[j];
   FOR(i, 1, 1 + 1)
    updAdd(d, mult(c[i], a[j - i]));
    if (d == 0)
    {
     x++;
      continue;
    VI t = c;
    int coef = mult(d, binPow(b, mod - 2));
    if (SZ(bp) + x > SZ(c))
     c.resize(SZ(bp) + x);
    FOR(i, 0, SZ(bp))
     updSub(c[i + x], mult(coef, bp[i]));
    if (2 * 1 > j)
    {
     x++;
      continue;
    1 = j + 1 - 1;
    bp = t;
   b = d;
    x = 1;
  }
  c.erase(c.begin());
  for (int& ci : c)
    ci = mult(ci, mod - 1);
 return c;
```

```
botsan-mori.hpp
// c - coefficients c[1], ..., c[k] but 0-index
// a - initial values <math>a[0], a[1], \ldots, a[k-1]
int botsanMori(VI c, VI a, LL n) {
  int k = SZ(c);
  assert(SZ(a) == k);
  VIq(k+1);
  q[0] = 1;
  FOR(i, 0, k)
    q[i + 1] = sub(0, c[i]);
  VI p = mult(a, q);
  p.resize(k);
  while (n) {
   VI qMinus = q;
    for (int i = 1; i <= k; i += 2)
      qMinus[i] = sub(0, qMinus[i]);
    VI newP = mult(p, qMinus);
    VI newQ = mult(q, qMinus);
    FOR(i, 0, k)
      p[i] = newP[2 * i + (n & 1)];
    FOR(i, 0, k + 1)
      q[i] = newQ[2 * i];
    n >>= 1;
  return mult(p[0], binPow(q[0], mod - 2));
conv-xor.hpp
                                         8ce066, 12 lines
void convXor(VI& a, int k)
  FOR(i, 0, k)
    FOR(j, 0, 1 << k)
      if((j & (1 << i)) == 0)
        int u = a[j];
        int v = a[j + (1 << i)];
        a[j] = u + v;
        a[j + (1 << i)] = u - v;
conv-and.hpp
void convAnd(VI& a, int k, bool inverse)
```

```
FOR(i, 0, k)
    FOR(j, 0, 1 \ll k)
      if((j \& (1 << i)) == 0)
        if(inverse)
          a[j] -= a[j + (1 << i)];
        else
          a[j] += a[j + (1 << i)];
conv-or.hpp
                                          6ffcf0, 12 lines
void convOr(VI& a, int k, bool inverse)
 FOR(i, 0, k)
    FOR (j, 0, 1 << k)
      if((j & (1 << i)) == 0)
        if(inverse)
          a[j + (1 << i)] -= a[j];
        else
          a[j + (1 << i)] += a[j];
```

Various (8)

```
mobius.hpp
```

fba6c5, 19 lines

```
void mobius()
  fill(pr, pr + N, 1);
  fill(mu, mu + N, 1);
  pr[1] = false;
  FOR (i, 2, N)
    if (!pr[i])
      continue;
    mu[i] = mod - 1;
    for (int j = 2 * i; j < N; j += i)
      pr[j] = false;
      if (j % (i * i) == 0)
        mu[j] = 0;
      mu[j] = mult(mu[j], mod - 1);
triangles.hpp
                                          a22d8b, 30 lines
int triangles(int n, int m)
```

```
FOR (i, 0, m)
  auto [u, v] = edges[i];
      if (MP(deg[u], u) < MP(deg[v], v))
          g[u].PB(v);
      else
    q[v].PB(u);
  int cnt = 0;
  FOR (v, 0, n)
      for (auto u : q[v])
    used[u] = 1;
      for (auto u : q[v])
          for (auto w : q[u])
      if (used[w])
        cnt++;
```

```
for (auto u : q[v])
      used[u] = 0;
   return cnt;
ternary.hpp
                                         3db54c, 29 lines
const db phi = (3. - sqrt(5.0)) / 2.;
db get (db L, db R)
 db M1, M2, v1, v2;
 M1 = L + (R - L) * phi;
 M2 = R - (R - L) * phi;
 v1 = f(M1);
 v2 = f(M2);
 FOR (i, 0, 74)
      if (v1 > v2) // for minimum
      L = M1;
      M1 = M2;
      v1 = v2;
      M2 = R - (R - L) * phi;
      v2 = f(M2);
   else
      R = M2;
      M2 = M1;
      v2 = v1;
     M1 = L + (R - L) * phi;
      v1 = f(M1);
 return L; // or f(L);
```

lader-nim.txt

Players have stone piles of size a0, a1, ..., an. In one move player can take $0 < x \le ai$ stones

i-th pile and move them to (i-1)-th pile.

In this game you can forget about even piles. Take stones from odd is equal to remove in NIM and from even equal to add in NIM. Adding in NIM useless.

Formulas (9)

9.1Trigonometry

$$\sin(v+w) = \sin v \cos w + \cos v \sin w$$

$$\cos(v+w) = \cos v \cos w - \sin v \sin w$$

$$\tan(v+w) = \frac{\tan v + \tan w}{1 - \tan v \tan w}$$

$$\sin v + \sin w = 2\sin \frac{v+w}{2}\cos \frac{v-w}{2}$$

$$\cos v + \cos w = 2\cos \frac{v+w}{2}\cos \frac{v-w}{2}$$

9.2Triangles

Area:
$$A = \sqrt{p(p-a)(p-b)(p-c)}$$

 $A = \frac{abc}{4R} = pr$
Length of median: $m_a = \frac{1}{2}\sqrt{2b^2 + 2c^2 - a^2}$

Derivatives/Integrals

$$\frac{d}{dx}\arcsin x = \frac{1}{\sqrt{1-x^2}} \quad \frac{d}{dx}\arccos x = -\frac{1}{\sqrt{1-x^2}}$$

$$\frac{d}{dx}\tan x = 1 + \tan^2 x \quad \frac{d}{dx}\arctan x = \frac{1}{1+x^2}$$

$$\int \tan ax = -\frac{\ln|\cos ax|}{a} \quad \int x\sin ax = \frac{\sin ax - ax\cos ax}{a^2}$$

$$\int e^{-x^2} = \frac{\sqrt{\pi}}{2}\operatorname{erf}(x) \quad \int xe^{ax}dx = \frac{e^{ax}}{a^2}(ax-1)$$

Integration by parts:

$$\int_{a}^{b} f'(x)g(x)dx = [f(x)g(x)]_{a}^{b} - \int_{a}^{b} f(x)g'(x)dx$$

Series 9.4

$$f(x) = \sum_{n=0}^{\inf} \frac{f^{(n)}(a) * (x-a)^n}{n!}$$

$$e^x = \sum_{n=0}^{\infty} \frac{x^n}{n!}$$

$$\ln(1-x) = -\sum_{n=0}^{\infty} \frac{x^n}{n!}, (-1 \le x < 1)$$

9.5 Number Theory

9.5.1 Mobius

$$forn \ge 1g(n) = \sum_{d|n} f(d)$$

$$then f(n) = \sum_{d|n} \mu(d) g(n/d)$$

$$M(n) = \sum_{k=1}^{n} \mu(k), \sum_{n=1}^{n} x M([x/n]) = 1$$

9.5.2 Catalan

$$C_n = \sum_{k=0}^{n-1} C_k C_{n-1-k}$$

$$C_n = \frac{1}{n+1}C_{2n}^n$$

$$C_n = C_{2n}^n - C_{2n}n - 1$$

9.5.3 Binomials

$$\sum_{k=0}^{n} C_n^k = 2^n$$

$$\sum_{m=0}^{n} C_m^k = C_{n+1}^{k+1}$$

$$\sum_{k=0}^{m} C_{n+k}^{k} = C_{n+m+1}^{m}$$

$$\sum_{k=0}^{n} (C_n^k)^2 = C_{2n}^n$$

$$\sum_{j=0}^{k} C_{m}^{j} C_{n-m}^{k-j} = C_{n}^{k}$$

$$\sum_{j=0}^{m} C_m^j C_{n-m}^{k-j} = C_{n+1}^{k+1}$$

$$\sum_{k=0}^{n} C_{n-k}^{k} = F_{n+1}$$

9.6 Runge-Kutta 4th Order Method for Ordinary Differential Equations

$$\frac{dy}{dx} = f(x, y), y(0) = y_0$$

$$x_{i+1} - x_i = h$$

$$y_{i+1} = y_i + \frac{1}{6}(k_1 + 2k_2 + 2k_3 + k_4)h$$

$$k_1 = f(x_i, y_i)$$

$$k_2 = f(x_i + \frac{1}{2}h, y_i + \frac{1}{2}k_1h)$$

$$k_3 = f(x_i + \frac{1}{2}h, y_i + \frac{1}{2}k_2h)$$

$$k_4 = f(x_i + h, y_i + k_3 h)$$

9.7 Burnside's lemma

Let G be a finite group that acts on a set X.

The *orbit* of an element x in X is the set of elements in X to which x can be moved by the elements of G. The orbit of x is denoted by $G \cdot x$:

$$G \cdot x = \{ g \cdot x \mid g \in G \}.$$

For each g in G, let X^g denote the set of elements in X that are fixed by g (also said to be left invariant by g), that is, $X^g = \{x \in X \mid g \cdot x = x\}$. Burnside's lemma asserts the following formula for the number of orbits, denoted |X/G|:

$$|X/G| = \frac{1}{|G|} \sum_{g \in G} |X^g|.$$