

T-Axi (declarative)

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Type conversion

$$\frac{\Gamma \text{ ctx}_{\text{nc}}}{\Gamma \vdash \text{Unit} \equiv \text{Unit type}} \quad \frac{\Gamma \text{ ctx}_{\text{nc}}}{\Gamma \vdash \text{Empty} \equiv \text{Empty type}}$$

$$\frac{\Gamma \vdash A_1 \equiv A_2 \text{ type} \quad \Gamma \vdash B_1 \equiv B_2 \text{ type}}{\Gamma \vdash r A_1 \rightarrow B_1 \equiv r A_2 \rightarrow B_2 \text{ type}}$$

$$\frac{\Gamma \vdash A_1 \equiv A_2 \text{ type}}{\Gamma \vdash !_r A_1 \equiv !_r A_2 \text{ type}}$$

$$\frac{\Gamma \vdash A_1 \equiv A_2 \text{ type} \quad \Gamma \vdash B_1 \equiv B_2 \text{ type}}{\Gamma \vdash A_1 \otimes B_1 \equiv A_2 \otimes B_2 \text{ type}}$$

$$\frac{\Gamma \vdash A_1 \equiv A_2 \text{ type} \quad \Gamma \vdash B_1 \equiv B_2 \text{ type}}{\Gamma \vdash A_1 \oplus B_1 \equiv A_2 \oplus B_2 \text{ type}}$$

Type conversion – properties

The rules above are the complete definition of type conversion. For now the context Γ is not used, but soon it will. We don't need to take any closures – type conversion already is an equivalence relation.

- $\Gamma \vdash A \equiv A$ type
- If $\Gamma \vdash A \equiv B$ type then $\Gamma \vdash B \equiv A$ type
- If $\Gamma \vdash A \equiv B$ type and $\Gamma \vdash B \equiv C$ type then $\Gamma \vdash A \equiv C$ type

Proposition conversion

$$\frac{\Gamma \text{ ctx}_{\text{nc}}}{\Gamma \vdash \top \equiv \top \text{ prop}} \quad \frac{\Gamma \text{ ctx}_{\text{nc}}}{\Gamma \vdash \perp \equiv \perp \text{ prop}}$$

$$\frac{\Gamma \vdash P_1 \equiv P_2 \text{ prop} \quad \Gamma \vdash Q_1 \equiv Q_2 \text{ prop}}{\Gamma \vdash P_1 \Rightarrow Q_1 \equiv P_2 \Rightarrow Q_2 \text{ prop}}$$

$$\frac{\Gamma \vdash P_1 \equiv P_2 \text{ prop} \quad \Gamma \vdash Q_1 \equiv Q_2 \text{ prop}}{\Gamma \vdash P_1 \wedge Q_1 \equiv P_2 \wedge Q_2 \text{ prop}}$$

$$\frac{\Gamma \vdash P_1 \equiv P_2 \text{ prop} \quad \Gamma \vdash Q_1 \equiv Q_2 \text{ prop}}{\Gamma \vdash P_1 \vee Q_1 \equiv P_2 \vee Q_2 \text{ prop}}$$

Proposition conversion

$$\frac{\Gamma \vdash A_1 \equiv A_2 \text{ type} \quad \Gamma, x : A_1 \vdash P_1 \equiv P_2 \text{ prop}}{\Gamma \vdash \forall x : A_1. P_1 \equiv \forall x : A_2. P_2 \text{ prop}}$$

$$\frac{\Gamma \vdash A_1 \equiv A_2 \text{ type} \quad \Gamma, x : A_1 \vdash P_1 \equiv P_2 \text{ prop}}{\Gamma \vdash \exists x : A_1. P_1 \equiv \exists x : A_2. P_2 \text{ prop}}$$

$$\frac{\Gamma \vdash A_1 \equiv A_2 \text{ type} \quad \Gamma \vdash e_1 \equiv e_2 : A_1 \quad \Gamma \vdash e'_1 \equiv e'_2 : A_1}{\Gamma \vdash e_1 =_{A_1} e'_1 \equiv e_2 =_{A_2} e'_2 \text{ prop}}$$

Note that we don't worry about α -conversion and assume the variable is the same on both sides.

Term conversion – sanity checks

We'll set up the term conversion judgement so that if
 $\Gamma \vdash a_1 \equiv a_2 : A$ then $\Gamma \vdash_{\text{nc}} a_1 : A$ and $\Gamma \vdash_{\text{nc}} a_2 : A$.

Term conversion – Empty

$$\frac{\Gamma \vdash_{\text{nc}} e_1 : \text{Empty} \quad \Gamma \vdash_{\text{nc}} e_2 : \text{Empty}}{\Gamma \vdash e_1 \equiv e_2 : \text{Empty}}$$

$$\frac{|\Gamma| \vdash A \text{ type} \quad \Gamma \vdash e_1 \equiv e_2 : \text{Empty}}{\Gamma \vdash \text{Empty-elim } e_1 \equiv \text{Empty-elim } e_2 : A}$$

Term conversion – Unit

$$\frac{\Gamma \vdash_{\text{nc}} a : A}{\Gamma \vdash \text{let unit} = \text{unit} \text{ in } a \equiv a : A}$$

$$\frac{\Gamma \text{ ctx}_{\text{nc}} \quad \Gamma \sqsubseteq 0\Gamma}{\Gamma \vdash \text{unit} \equiv \text{unit} : \text{Unit}}$$

$$\frac{\Gamma \vdash_{\text{nc}} u_1 : \text{Unit} \quad \Gamma \vdash_{\text{nc}} u_2 : \text{Unit}}{\Gamma \vdash u_1 \equiv u_2 : \text{Unit}}$$

$$\frac{\Gamma \sqsubseteq \Gamma_1 + \Gamma_2 \quad \Gamma_1 \vdash u_1 \equiv u_2 : \text{Unit} \quad \Gamma_2 \vdash a_1 \equiv a_2 : A}{\Gamma \vdash \text{let unit} = u_1 \text{ in } a_1 \equiv \text{let unit} = u_2 \text{ in } a_2 : A}$$

Term conversion – functions

$$\frac{\Gamma \sqsubseteq \Gamma_1 + r\Gamma_2 \quad \Gamma_1, r_A^i x : A \vdash_{\text{nc}} b : B \quad \Gamma_2 \vdash_{\text{nc}} a : A}{\Gamma \vdash (\lambda_r x : A. b) \ a \equiv b[x := a] : B}$$

$$\frac{\Gamma, r_A^i x : A \vdash f \ x \equiv g \ x : B}{\Gamma \vdash f \equiv g : r A \rightarrow B}$$

$$\frac{\Gamma \sqsubseteq \Gamma_1 + r\Gamma_2 \quad \Gamma_1 \vdash f_1 \equiv f_2 : r A \rightarrow B \quad \Gamma_2 \vdash a_1 \equiv a_2 : A}{\Gamma \vdash f_1 \ a_1 \equiv f_2 \ a_2 : B}$$

Term conversion – box

$$\frac{\Gamma \sqsubseteq r \Gamma_1 + \Gamma_2 \quad \Gamma_1 \vdash_{\text{nc}} a : A \quad \Gamma_2, r_A^i x : A \vdash_{\text{nc}} b : B}{\Gamma \vdash \text{let box } x = \text{box}_r a \text{ in } b \equiv b[x := a] : B}$$

$$\frac{\Gamma \sqsubseteq r \Gamma' \quad \Gamma' \vdash a_1 \equiv a_2 : A}{\Gamma \vdash \text{box } a_1 \equiv \text{box } a_2 : !_r A}$$

$$\frac{\Gamma \sqsubseteq \Gamma_1 + \Gamma_2 \quad \Gamma_1 \vdash a_1 \equiv a_2 : !_r A \quad \Gamma_2, r_{A_1}^i x : A_1 \vdash b_1 \equiv b_2 : B}{\Gamma \vdash \text{let box } x = a_1 \text{ in } b_1 \equiv \text{let box } x = a_2 \text{ in } b_2 : B}$$

Term conversion – product

$$\frac{\Gamma \sqsubseteq \Gamma_1 + \Gamma_2 \quad \Gamma_1 \vdash a_1 \equiv a_2 : A \quad \Gamma_2 \vdash b_1 \equiv b_2 : B}{\Gamma \vdash (a_1, b_1) \equiv (a_2, b_2) : A \otimes B}$$

$$\frac{\Gamma \sqsubseteq \Gamma_1 + \Gamma_2 \quad \Gamma_1 \vdash e_1 \equiv e_2 : A \otimes B \quad \Gamma_2, 1_A^i x : A, 1_B^i y : B \vdash c_1 \equiv c_2 : C}{\Gamma \vdash \text{let } (x, y) = e_1 \text{ in } c_1 \equiv \text{let } (x, y) = e_2 \text{ in } c_2 : C}$$

Term conversion – sum

$$\frac{\Gamma \vdash a_1 \equiv a_2 : A \quad |\Gamma| \vdash B \text{ type}}{\Gamma \vdash \text{inl } a_1 \equiv \text{inl } a_2 : A \oplus B}$$

$$\frac{|\Gamma| \vdash A \text{ type} \quad \Gamma \vdash b_1 \equiv b_2 : B}{\Gamma \vdash \text{inr } b_1 \equiv \text{inr } b_2 : A \oplus B}$$

$$\frac{\Gamma \sqsubseteq \Gamma_1 + \Gamma_2 \quad \Gamma_1 \vdash e_1 \equiv e_2 : A \oplus B \quad \begin{array}{c} \Gamma_2 \vdash f_1 \equiv f_2 : 1 \, A \rightarrow C \\ \Gamma_2 \vdash g_1 \equiv g_2 : 1 \, B \rightarrow C \end{array}}{\Gamma \vdash \text{case } e_1 \text{ of } (f_1, g_1) \equiv \text{case } e_2 \text{ of } (f_2, g_2) : C}$$