

# Laboration 2

## Uppaal

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## 1 Getting started

For the first query  $E \langle \rangle P.s3$  which means that there exists a possible future where  $P.s3$  holds. The result is that the property is satisfied which is true since there is a possible sequence of state transitions leading from the initial state to  $s3$ . An example sequence :  $S0 \rightarrow S1 \rightarrow S3$

For the second query  $A \langle \rangle P.s3$  which means that all possible sequence of states that eventually leads to  $P.s3$ . The property is not satisfied since there is no way to ensure that there is such a sequence. This is due to the fact that there is nothing forcing the states to change. An example sequence:  $S0$  The state machine will stay in  $S0$  for infinity.

## 2 Fisher 1

For this assignment we try to predict the verification time for fisher with  $n = 12$ . We used a android phone to measure time since uppaal doesn't report it.

- $n : 8$  1s
- $n : 9$  2s
- $n : 10$  6s
- $n : 11$  20s
- $n : 12$  120s (Our guess)

It seems to be a factorial increment of time.

### 3 Fisher 2

If  $m < k$  then the mutex requirement will not be satisfied. An example of this is the sequence:

$(-, -, -) \rightarrow (\text{req}, -, -) \rightarrow (\text{req}, \text{req}, -) \rightarrow (\text{wait}, \text{req}, -) \rightarrow (\text{cs}, \text{req}, -) \rightarrow (\text{cs}, \text{wait}, -) \rightarrow (\text{cs}, \text{cs}, -)$

If  $m \geq k$  then mutex requirement will be satisfied.

### 3.1 Traffic Light Controller

The Traffic light timed automata to have three templates. Two for the traffic lights and one controller. The reason that we chose to have two templates for the traffic light instead of one is that there is a slight difference between north and south compared to east and west.

We have an integer called transfer that behaves like a four+bit number. We use this to define seven states of the system, north, south, east, west, north-south, east-west or none. To do this each traffic direction is initialized with an id, 1,2,4,8; these are then added and subtracted from transfer when a light goes to and from green respectively. This enables us to define states by checking the value of transfer, if it's less than 4, for example, the state is 1&2. Less than 2 the state is 1 and so on.

### 3.2 Traffic Light Template

The template has three states, *idle* (red light), *wait* (there are cars waiting) and *Green* (green light). To force the system to transit state the edge from *wait* to *Green* is triggered by a channel from the controller while *Green* has a timed invariant.

### 3.3 Controller

The controller has two states that it switches between on a timed interval. The edges from NS to WE triggers one channel for the WE traffic light and vice versa for the edge from WE to NS but with another channel.

## 4 3.5 Alternating Bit Protocol

The ABP (Alternating Bit protocol) system have three components, sender, receiver and a channel.

### 4.1 Sender

The sender creates a message that it tries to send through the channel to the receiver. It has four nodes, *mode zero*, *new message*, *mode one* and a *transition node* between mode one and zero. From *mode zero* the sender can move to *new message*. In *new message* the sender waits for a acknowledgment from the receiver, this can be either one or zero. If zero the sender transitions to *mode one*. Otherwise it re-transmits the message with state 0. In *mode one* the sender awaits another acknowledgment from the receiver. If one the sender goes to *mode zero*. Otherwise it re-transmits the message with state 1.

In *new message* and *mode one* (*the transition node*) the sender can time out. If so it re-transmits the message with 0 and 1 respectively.

### 4.2 Receiver

The receiver receives a message from the channel. It has four nodes, *mode zero*, *mode one* and *two transition nodes*. Form *mode zero* the receiver can move to *mode one* tough a transition node or back to itself through a transition node. In this case, it re-transmits an acknowledgment with 1. In *mode one* the receiver can move to *mode zero* through a transition node or back to itself through a transition node.

In *mode zero* and *mode one* the receives can time out. If so it re-transmits the message with 1 and 0 respectively.

When the receiver moves to *mode one* it transmits an acknowledgment 0 to the channel and if it moves to *mode zero* it transmits an acknowledgment 1 to the channel.

### 4.3 Channel

The channel has to nodes. From the first it has four edges to move to the second. They are sensitive to the receivers acknowledgment and the sender's message. There are also four edges form the second node to the first. They activate the acknowledgment in the sender and the state change in the receiver.