

Discrete Mathematics

Chapter 10: Trees

Department of Mathematics
The FPT university

Chapter 10: Introduction

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Topics covered:

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10.1 Introduction to Trees

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10.2 Applications of Trees

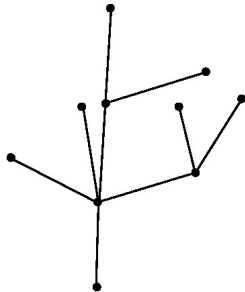
Chapter 10: Introduction

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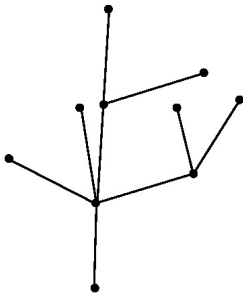
- 10.1 Introduction to Trees
- 10.2 Applications of Trees
- 10.3 Tree Traversal
- 10.4 Spanning Trees
- 10.5 Minimum Spanning Trees

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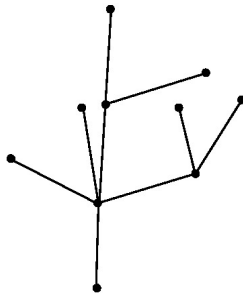
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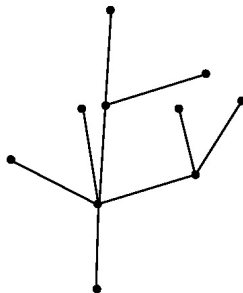


10.1 Introduction to Trees



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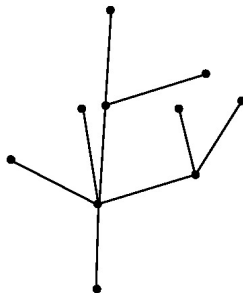
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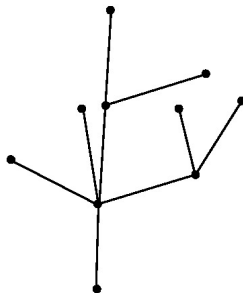


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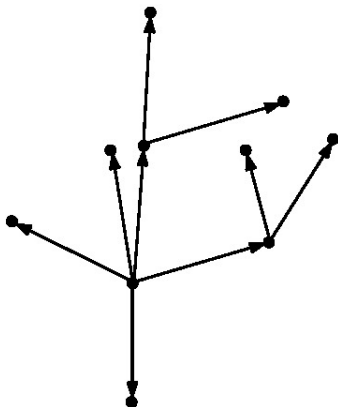
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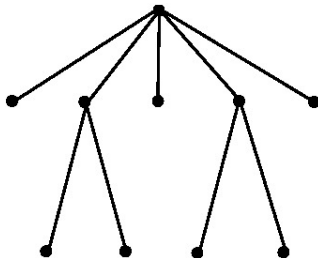
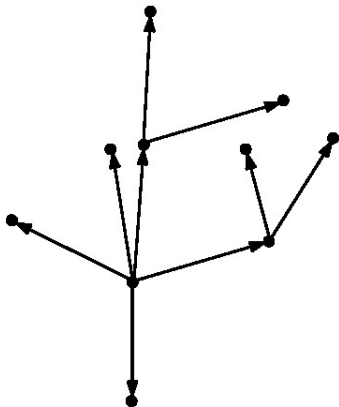
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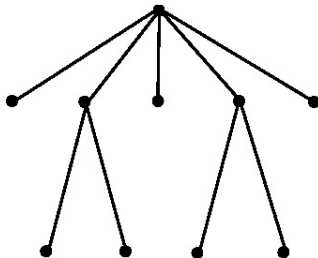
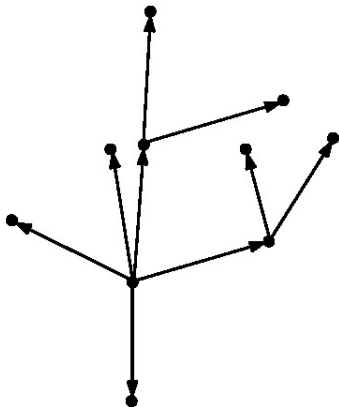
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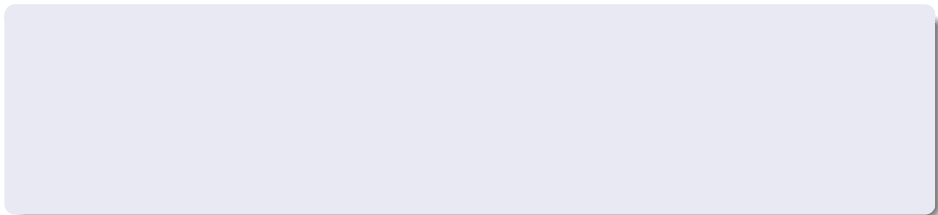
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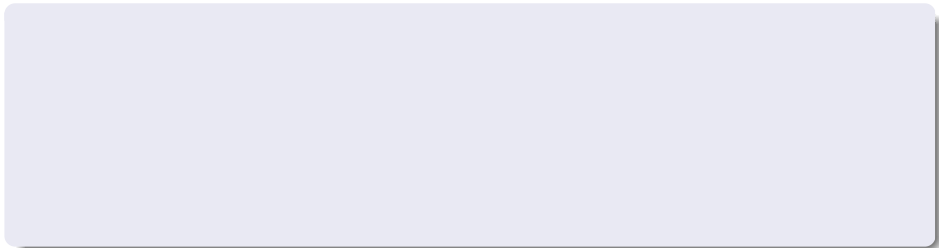
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Example. Does there exist a full m -ary tree with height 4 and 100 leaves?

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$$\log(n!) = \log 1 + \log 2 + \cdots + \log n = \Theta(n \log n)$$

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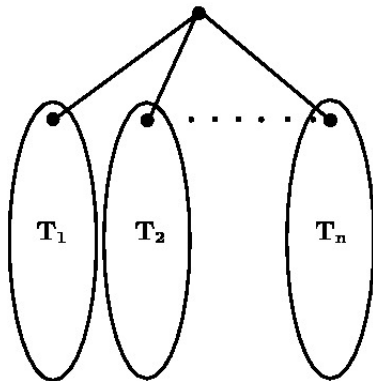
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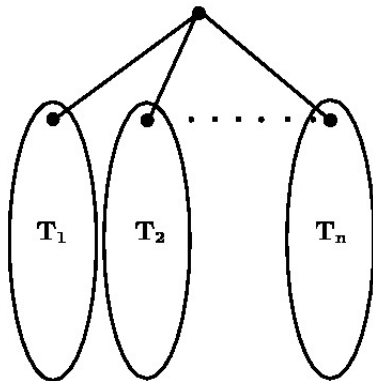
Read textbook!

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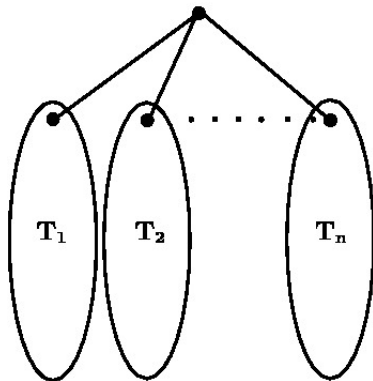


10.3 Tree Traversal



Preorder traversal

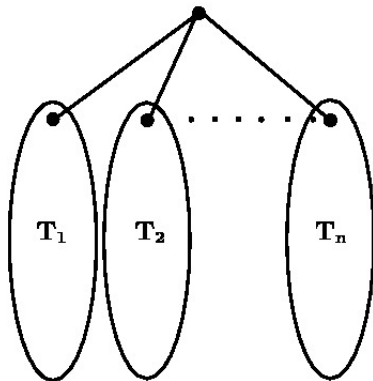
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Preorder traversal

- Visit the root

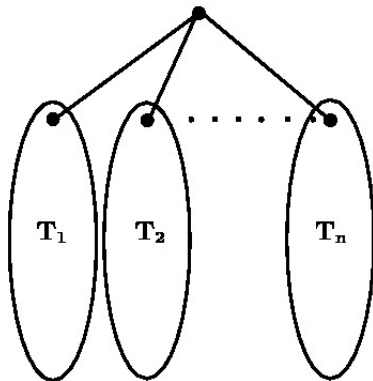
10.3 Tree Traversal



Preorder traversal

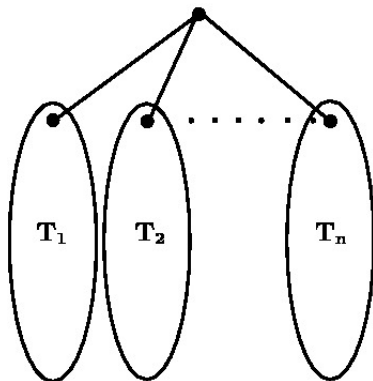
- Visit the root
- Traverse each of T_1, T_2, \dots, T_n in preorder.

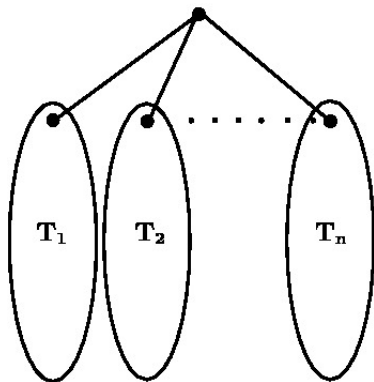
10.3 Tree Traversal



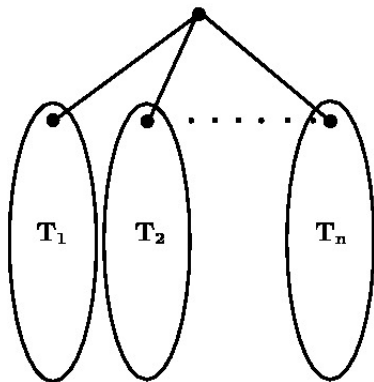
Preorder traversal

- Visit the root
- Traverse each of T_1, T_2, \dots, T_n in preorder.



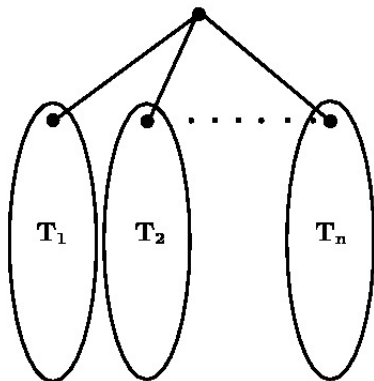


Inorder traversal



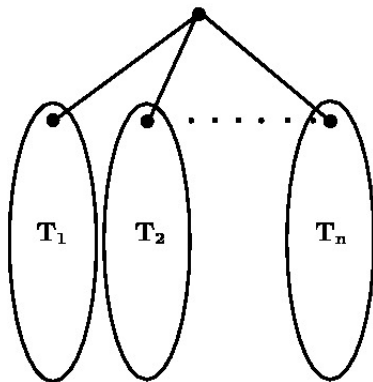
Inorder traversal

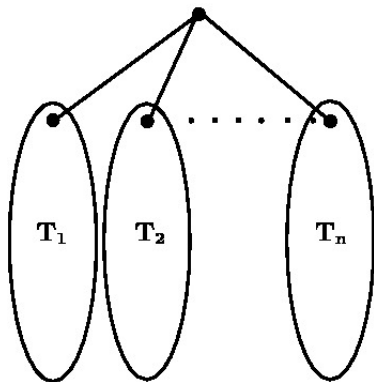
- Traverse T_1 in inorder.



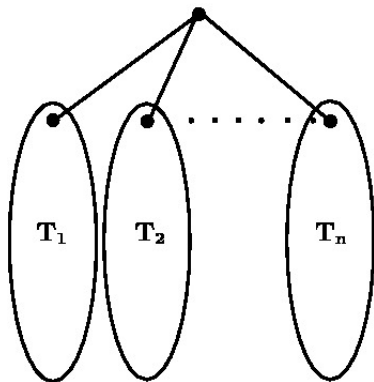
Inorder traversal

- Traverse T_1 in inorder.
- Visit the root.
- Traverse each of T_2, \dots, T_n in inorder.



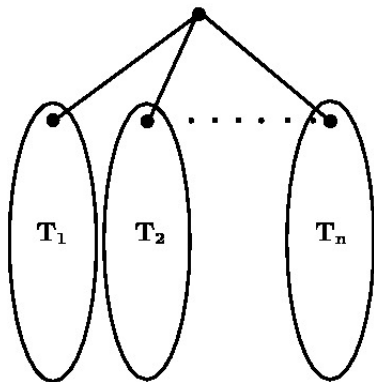


Postorder traversal



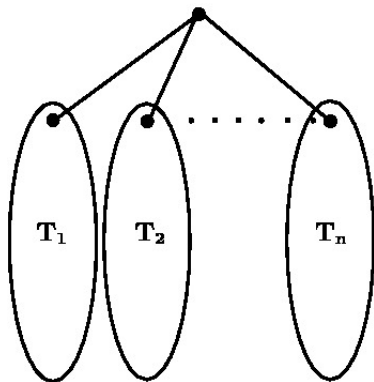
Postorder traversal

- Traverse each of T_1, T_2, \dots, T_n in postorder.
- Visit the root.



Postorder traversal

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Postorder traversal

- Traverse each of T_1, T_2, \dots, T_n in postorder.
- Visit the root.

Infix, Prefix, and Postfix Notation

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Consider the expression $(x + y)^2 + (x - 4)/3$.

Infix, Prefix, and Postfix Notation

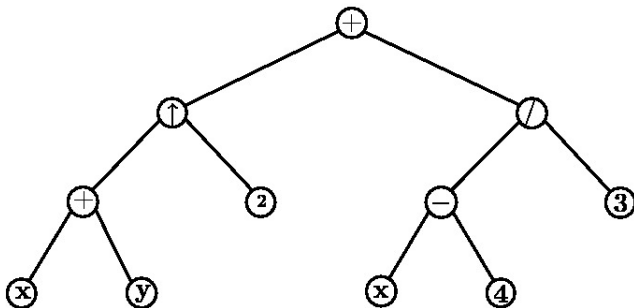
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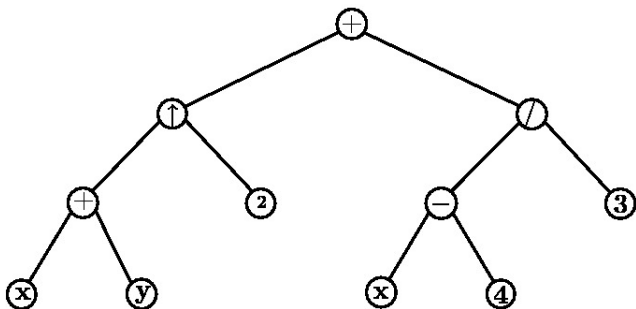
We can use a binary tree to represent this expression, where the internal vertices represent operations, and the leaves represent numbers or variables.

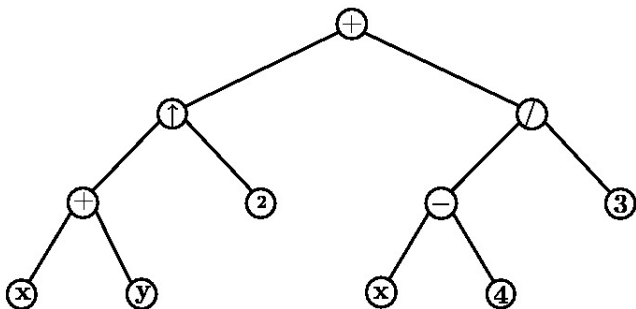
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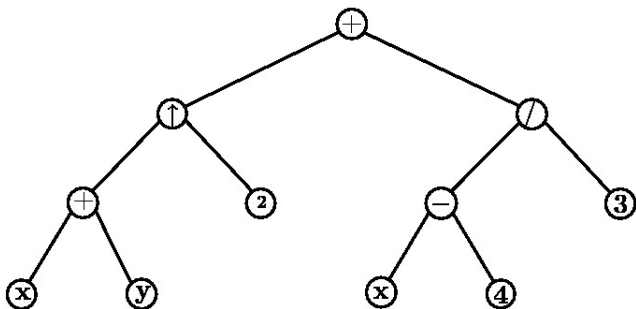
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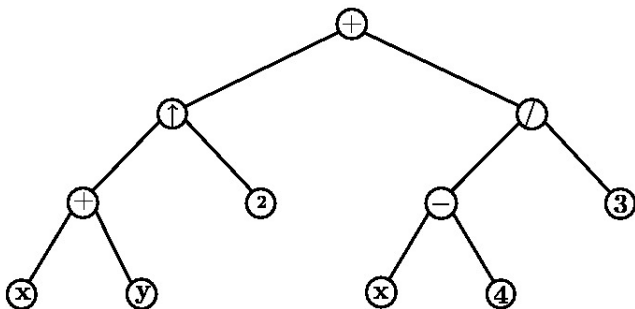




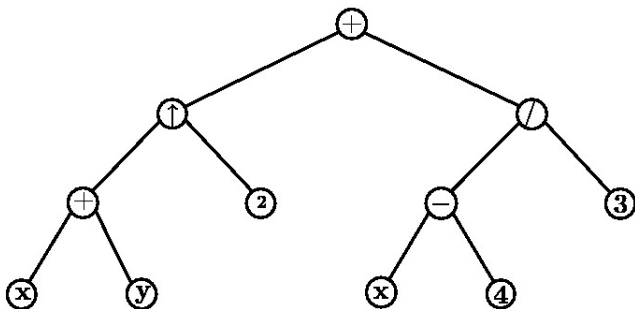
- Prefix notation (Polish notation):



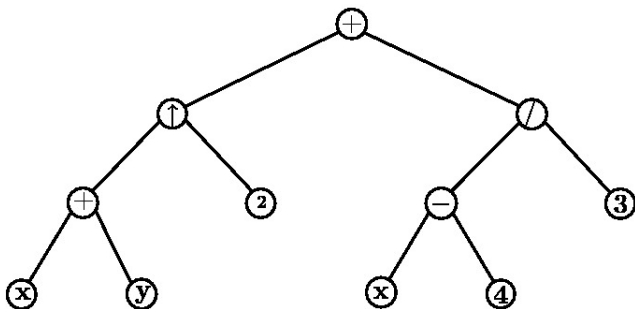
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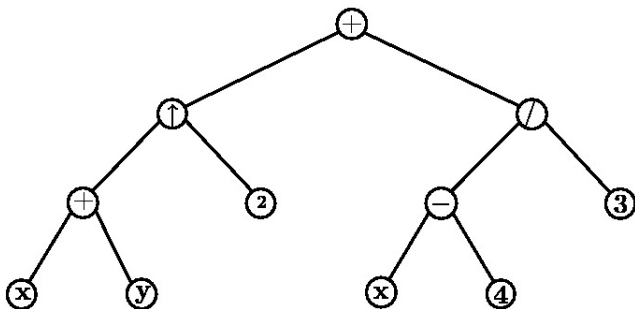
- Prefix notation (Polish notation):
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- Postfix notation (reverse Polish notation):



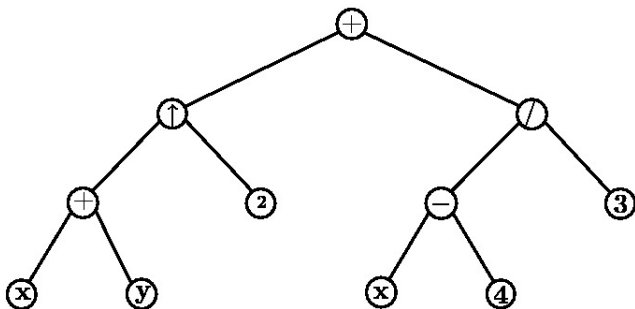
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 $x y + 2 \uparrow x 4 - 3 / +$



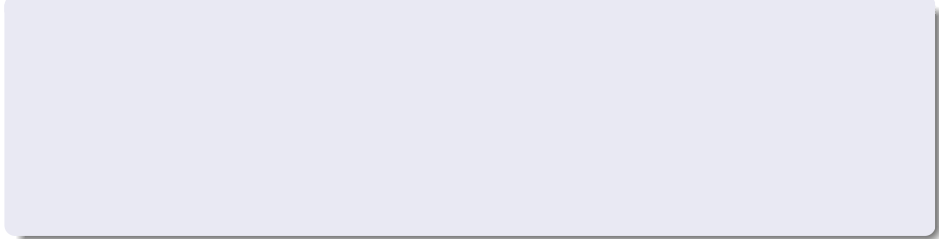
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