

Graph coloring

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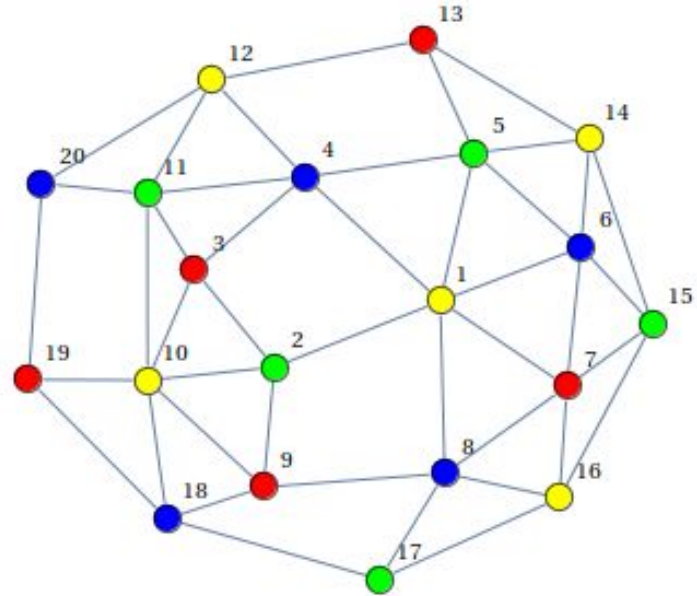
Graph coloring

- **Graph coloring** is one of the basic problems of Graph theory, widely used in informatics, to solve practical problems related to zoning, grouping, and maps.



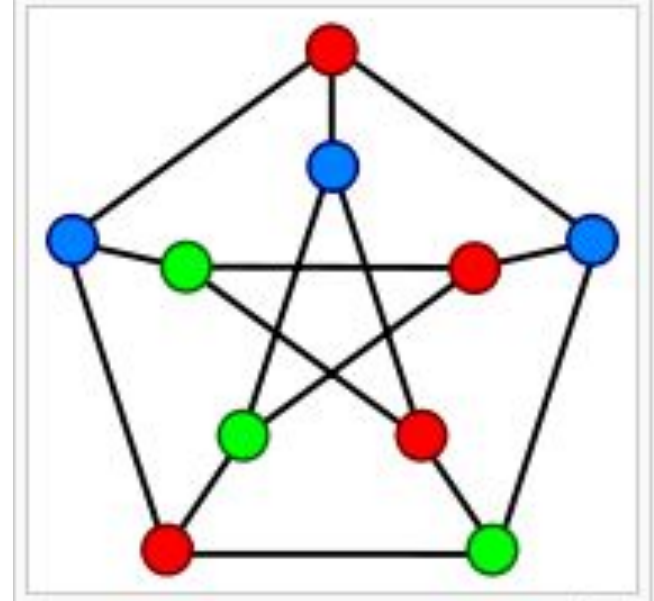
Definition

- In graph theory, graph coloring is a way of coloring the vertices of a graph such that no two adjacent vertices share the same color.



Chromatic number - Definition

- The **chromatic number** of a graph is the minimum number of colors one can use to color the vertices of the graph so that no two adjacent vertices are the same color.
- **Chromatic number** of the graph G is denoted by $\chi(G)$.
- $1 \leq \chi(G) \leq n$

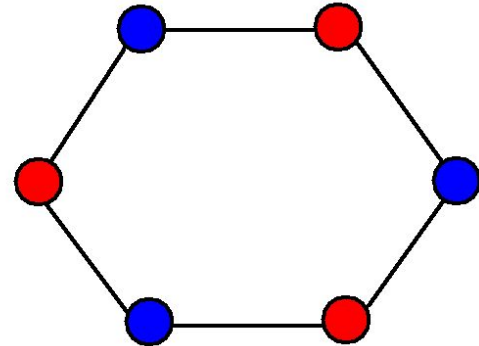


Note: n are vertices of graph

Chromatic number - Types of graphs

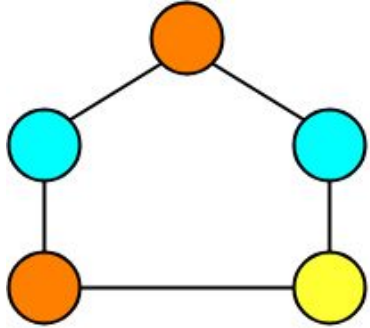
Cycle graph

- A simple graph of ' n ' vertices ($n \geq 3$) and ' n ' edges forming a cycle of length ' n ' is called as a cycle graph.
- In a cycle graph, all the vertices are of degree 2.



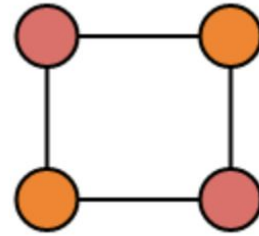
Chromatic number - Types of graphs

Cycle graph



Chromatic Number = 3

Any odd length cycle will have a chromaticity of 3 : $\chi(C_{2n+1}) = 3$



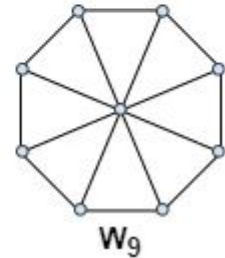
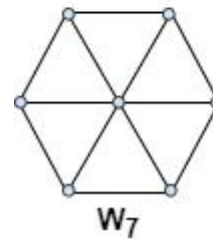
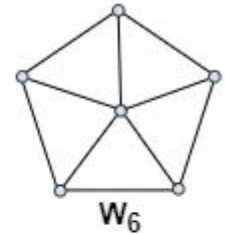
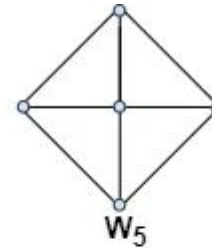
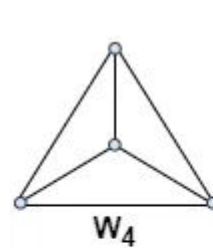
Chromatic Number = 2

Any even length cycle will have a chromaticity of 2 : $\chi(C_{2n}) = 2$

Chromatic number - Types of graphs

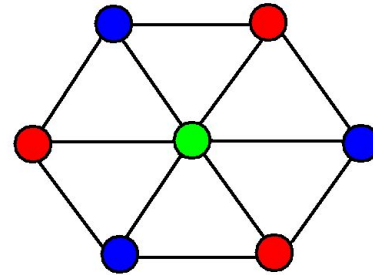
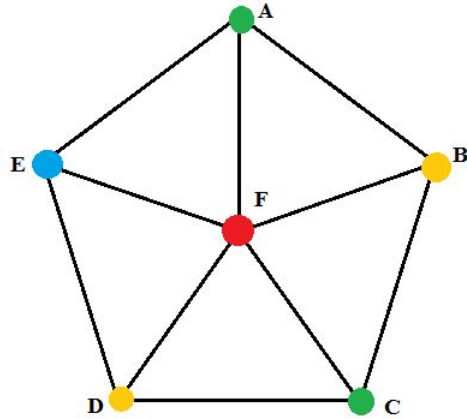
Wheel graph

- A wheel graph is obtained by connecting a vertex to all the vertices of a cycle graph.
- It is denoted by W_n , for $n > 3$ where n is the number of vertices in the graph.



Chromatic number - Types of graphs

Wheel graph



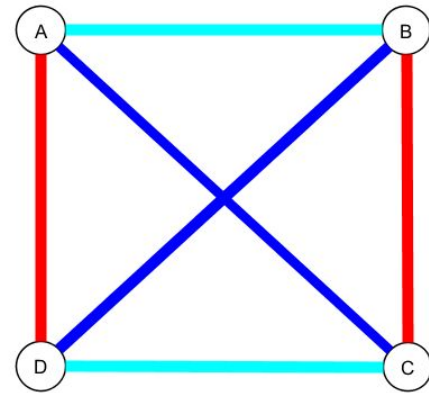
$\chi(W_n) = 4$ if n is even

$\chi(W_n) = 3$ if n is odd

Chromatic number - Types of graphs

Complete Graph

- A complete graph is a graph in which every **two distinct vertices** are joined by **exactly one edge**.
- In a complete graph, each vertex is **connected** with every other vertex.

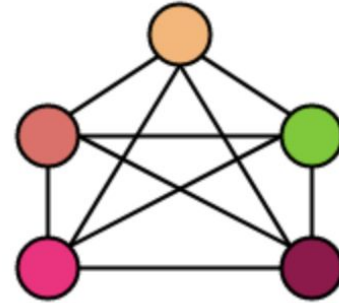


Chromatic number - Types of graphs

Complete Graph



Chromatic Number = 4



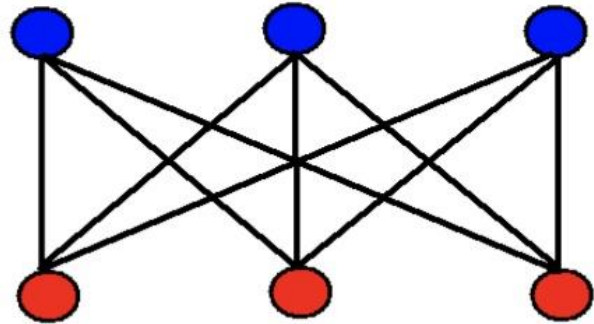
Chromatic Number = 5

- Chromatic Number of any Complete Graph = Number of vertices in that Complete Graph ($\chi(K_n) = n$)

Chromatic number - Types of graphs

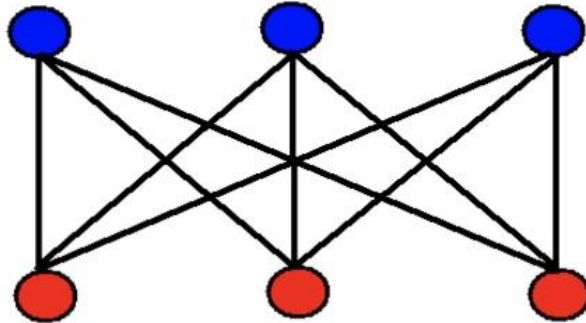
Bipartite Graphs

- A Bipartite Graph consists of **two sets of vertices** X and Y.
- The edges **only join vertices** in X to vertices in Y, not vertices within a set.



Chromatic number - Types of graphs

Bipartite Graphs

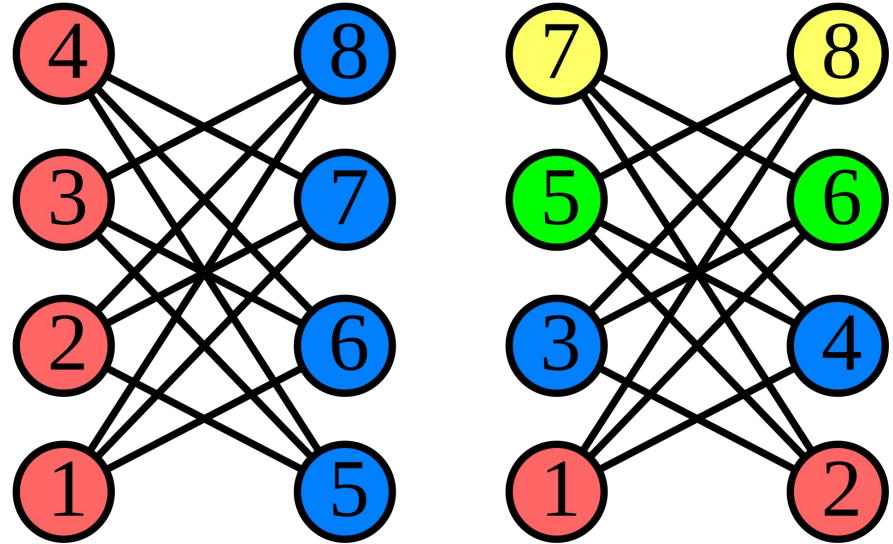


Chromatic Number of any Bipartite Graph = 2

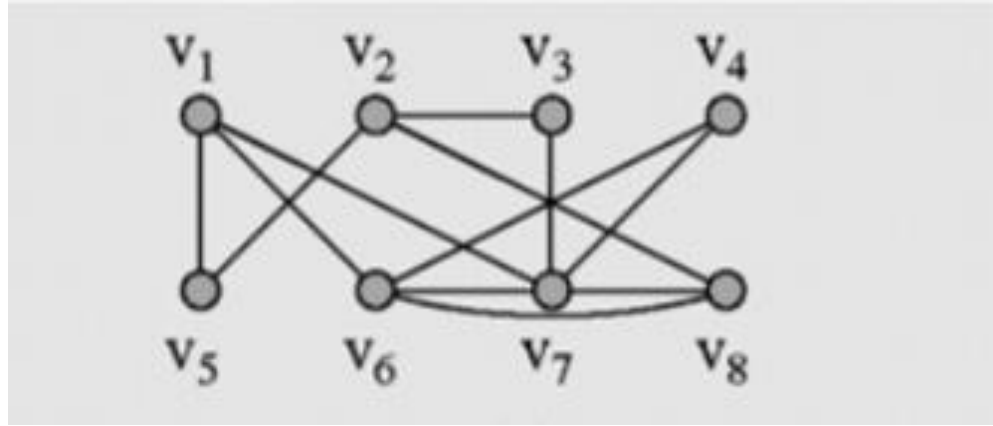
$$(\chi(G) \leq 2)$$

Graph coloring - Sequential coloring

- **Sequential coloring** establishes the **sequence of vertices** before coloring them, and then color the next vertex with the lowest number possible

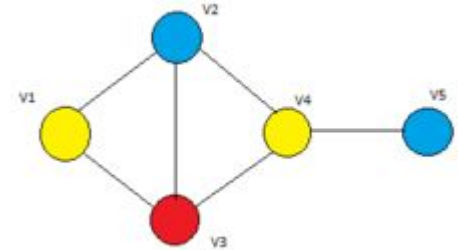
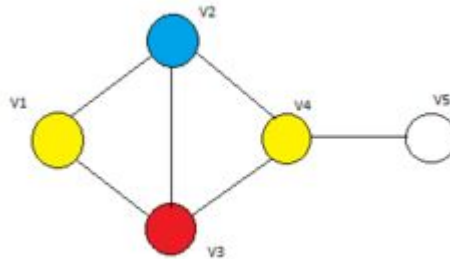
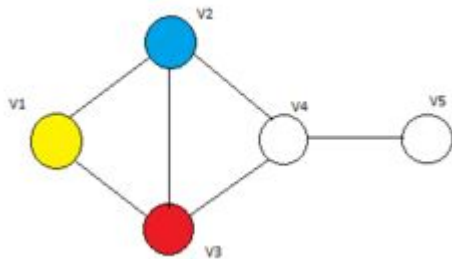
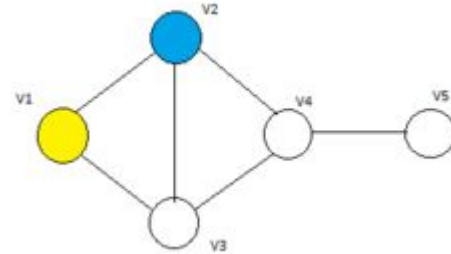
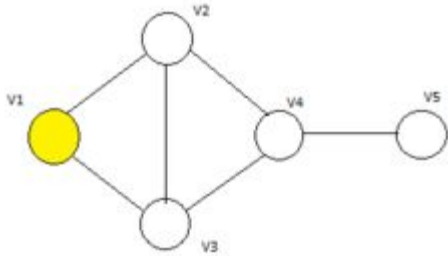


Graph coloring - Sequential coloring

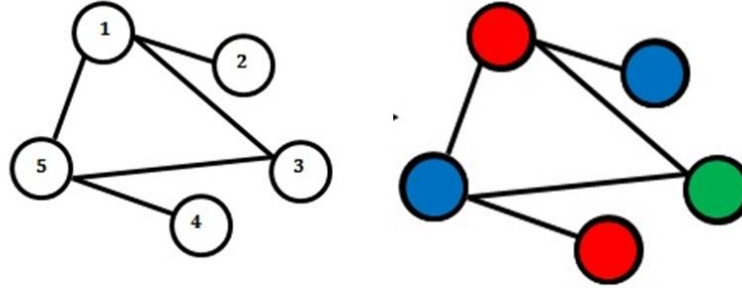


- **Step 1:** Color first vertex with the first color.
- **Step 2:** consider the remaining $(V-1)$ vertices one by one:
 - + Color the currently picked vertex with the **lowest numbered color** if it has not been used to color any of its adjacent vertices.
 - + If it has been used, then choose the **next least numbered color**. If all the previously used colors have been used, then assign a **new color** to the currently picked vertex.

Sequential coloring - EXAMPLE



Graph coloring - Largest first sequence



- **Step 1:** We have a graph with 5 vertices numbered 1, 2, 3, 4, 5 with the ranks corresponding to each vertex in the order of 3, 1, 2, 1, 3. Hence V' initially The order is [1, 5, 3, 2, 4]. Assign $i = 1$ (number of shaded colors).
- **Step 2:** Color 1 (red) for vertex 1. In turn, browse the remaining vertices in V' . We have: Point 5 is adjacent to vertex 1 (vertex 1 is colored 1 - red), so there is no color for vertex 5. Similarly, vertices 3, 2 are adjacent to vertex 1, so vertices 3, 2 have not been colored. Vertex 4 is not adjacent to vertex 1, so color 1 for vertex 4. Vertex 4 has color 1 - red.
- **Step 3:** Check that there are still unpainted vertices in V , so go to step 4.
- **Step 4:** Remove the colored vertices 1, 4 from V' , rearrange V' in descending order, we get $V' = [5, 3, 2]$. We have $i = 2$. Repeat step 2: