

Recursive Functions of Symbolic Expressions and Their Computation by Machine, Part I

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Introduction

WARNING

- I make mistakes.
 - Sorry if I did!
- Feedback welcome.

John McCarthy?

- Born September 4, 1927 in Boston, Massachusetts.
- Pioneer who contributed to AI, Internet, the web.
- At MIT from 1958 to 1962 (then Stanford university).
- A. M. Turing Award from Association for Computing Machinery, 1971.
- “He flew planes, climbed mountains, and traveled to the Soviet Union and China in the 1960s and 70s at a time when it was a major challenge.”

A bit of Context

- Paper is from 1960.
- The part II of this paper has never been written.
- Lisp and Fortran were written for the IBM 704 (available at MIT).
 - “Only computer which can handle complex math” - Wikipedia
- Second oldest high level programming language (non-assembly language).
 - Only Fortran older than Lisp.
- Today, Lisp is a family of languages
 - Many “dialects”
 - This paper is the foundation of all of these languages... and more (JINGLE!).

IBM 704



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Good Old Languages

- IPL2 (1956) - List processing in Assembly.
- Fortran (1957) - No list processing
- Lisp (for LISt Processing) (1959-1962?)

Why Lisp?

- AI (term coined by McCarthy)
 - They thought they were almost there in the 50s / 60s.
 - McCarthy wrote a “funny paper” later “HUMAN-LEVEL AI IS HARDER THAN IT SEEMED IN 1955”
 - Chess player (beginning of CS game theory).
- “Mathematical language”.
- Programming the Advice Taker (proposed in 1958).

The Paper

Advice Taker

“representing information about the world by sentences in a suitable formal language and a reasoning program that would decide what to do by making logical inferences. Representing sentences by list structure seemed appropriate - it still is - and a list processing language also seemed appropriate for programming the operations involved in deduction - and still is.” - McCarthy, 1979

- From the paper: Lisp handle declarative and imperative sentences using lists.

The notation: S-Expression and M-expressions

S-Expression (paper)

Characters: (.)
(CAR(CONS,x,y))
(CDR(CONS,x,y))

M-Expression (paper, never implemented)

car [cons [x; y]] = x
cdr [cons [x; y]] = y

Actual code (Scheme - dialect of Lisp)

(car (cons x y))
(cdr (cons x y))

Functions

- Functions first class citizen (code as data)
- Precise that it's **not** the usual mathematics term “function”.
- Inspired by Church's lambda-notation (1936).
 - Turing Machine (1936) “too complicated”.
- Use of lambda expression as bedrock for computation.
- Function composition to create new function.

Example (Scheme)

```
(+ 2 3)
(apply + '(1 2))
((lambda (x) (+ x x)) 2)
```

Recursion

- Describe formalism for defining function recursively.
- First programming language with recursion.

Example (Scheme)

```
(define (factorial n)
  (if (= n 0)
      1
      (* n (factorial (- n 1)))))
```

Conditionals

- First use of conditional expressions in a programming language.
 - Predicates: function returning #T or #F.
 - Essential for testing recursive base cases.
- Fortran had IF but it was “very awkward to use” - McCarthy 1979.

Example (Scheme)

```
(cond ((equal? x "#t") #t)
      ((equal? x "#f") #f)
      (else x))
```

Linked List (basic data structure)

Modern Lisp (Scheme)

```
(car (cons x y))  
(cdr (cons x y))
```

Box and Pointer diagram

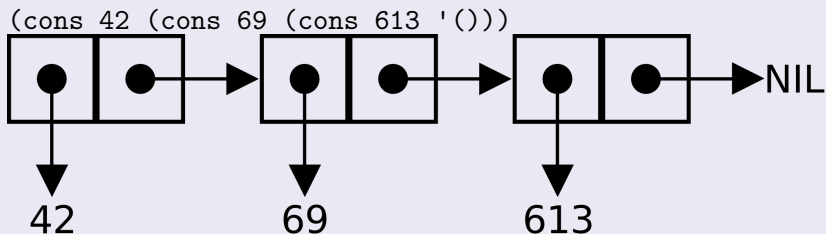


Figure 2: Pointer diagram

???

- car? cdr? WAT???
- Nice heritage from IBM 704 memory access.

caddr[x] for car[cdr[cdr[x]]]
(caddr x) for (car (cdr (cdr x)))

- About functions which have not one letter
" ... *is convenient to allow English words and phrases to stand for atomic entities for mnemonic reasons*" - John McCarthy (paper)
- ... what about other programmers?

Garbage collection

"... formerly pointed cannot be reached by a car — cdr chain from any base register. Such a register may be considered abandoned by the program because its contents can no longer be found by any possible program; hence its contents are no longer of interest, and so we would like to have it back on the free-storage list. This comes about in the following way. Nothing happens until the program runs out of free storage. When a free register is wanted, and there is none left on the free-storage list, a reclamation cycle starts.

Apply and Eval

- “Universal S-Function apply”
 - Universal turing machine
 - `apply[f ; args] = eval[cons[f ; appq[args]]; NIL]`,
“Writing eval required inventing a notation representing LISP functions as LISP data, and such a notation was devised for the purposes of the paper with no thought that it would be used to express LISP programs in practice” - McCarthy, 1979

Apply (Lisp Programmer's Manual 1.5)

```
apply[fn;x;a] =  
  [atom[fn] → [eq[fn;CAR] → caar[x];  
    eq[fn;CDR] → cdar[x];  
    eq[fn;CONS] → cons[car[x];cadr[x]];  
    eq[fn;ATOM] → atom[car[x]];  
    eq[fn;EQ] → eq[car[x];cadr[x]];  
    T → apply[eval[fn;a];x;a]];  
  eq[car[fn];LAMBDA] → eval[caddr[fn];pairlis[cadr[fn];x;a]];  
  eq[car[fn];LABEL] → apply[caddr[fn];x;cons[cons[cadr[fn];  
    caddr[fn]];a]]]
```

Figure 3: Apply

Eval (Lisp Programmer's Manual 1.5)

```
eval[e;a] = [atom[e] → cdr[assoc[e;a]];  
             atom[car[e]] →  
               [eq[car[e],QUOTE] → cadr[e];  
                eq[car[e],COND] → evcon[cdr[e];a];  
                T → apply[car[e];evlis[cdr[e];a];a];  
             T → apply[car[e];evlis[cdr[e];a];a]]
```

Figure 4: Eval

Eval (Scheme)

```
(define (eval exp env)
  (cond ((self-evaluating? exp) exp)
        ((variable? exp) (lookup-variable-value exp env))
        ((quoted? exp) (text-of-quotation exp))
        ((assignment? exp) (eval-assignment exp env))
        ((definition? exp) (eval-definition exp env))
        ((if? exp) (eval-if exp env))
        ((lambda? exp)
         (make-procedure (lambda-parameters exp)
                          (lambda-body exp)
                          env))
        ((begin? exp)
         (eval-sequence (begin-actions exp) env))
        ((cond? exp) (eval (cond->if exp) env)))
```

Apply (Scheme)

```
(define (apply procedure arguments)
  (cond ((primitive-procedure? procedure)
        (apply-primitive-procedure procedure arguments))
        ((compound-procedure? procedure)
         (eval-sequence
          (procedure-body procedure)
          (extend-environment
           (procedure-parameters procedure)
           arguments
           (procedure-environment procedure))))
        (else
         (error
          "Unknown procedure type -- APPLY" procedure))))
```

Legacy

The most important

- Everything we saw before (recursion, conditionals, lambdas...).
 - Lambdas implemented in mainstream programming languages (Java, C#, C++ ...).
- First functional programming language.
- Smalltalk (70s) was very influenced by Lisp (today: Pharo).
 - One of the first OOP language.
 - One rule “almost everything is function” was appealing to Alan Kay.
 - In smalltalk, everything is an object.
 - Garbage collection.

Lisp today

- Clojure
- Common Lisp
- Emacs Lisp
- Scheme

References

- Slides - <https://github.com/Phantas0s/presentations/>
- HUMAN-LEVEL AI IS HARDER THAN IT SEEMED IN 1955
- History of Lisp
- Page on John McCarthy (all papers)
- xkcd