Spring 2022, Homework 3 (20 points in total)

Q1. (2pts) Multiplication with shift instructions

MOV R0, #100 MUL R2, R1, R0

The above code computes R2 = R1 * 100. Carry out the same computation, (i.e., a multiplication by 100) with another code, using only LSL and ADD instructions. **Don't use any loops, (i.e., no use of branches).** The result should be placed into R2 like the original code. You may use additional registers such as R3 and R4.

Note that you cannot run the provided code in VisUAL, because MUL is not a supported instruction in VisUAL.

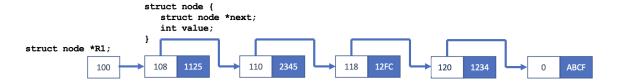
Q2. (6 pts) Memory map

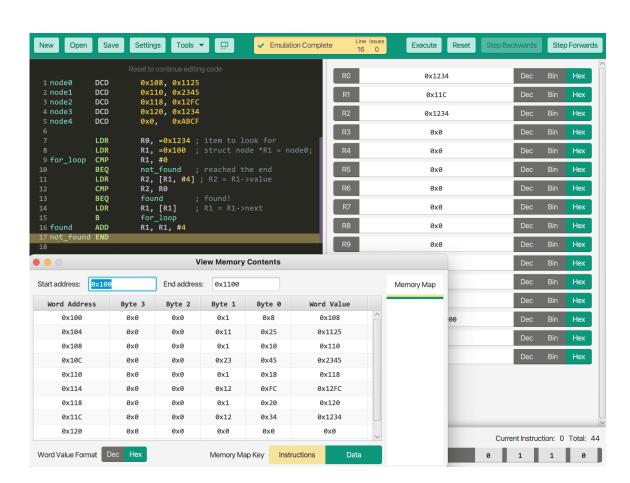
Let's assume that you're using Keil uVersion. Fill out the blanks of the memory map (address 0x00000004 to address 0x00000014) when running the following assembly program.

| StackSize | | THUMB EQU | 0x00000100 |
|---------------|---|--|---|
| MyStackMem | | AREA SPACE | STACK, NOINIT, READWRITE, ALIGN=3 StackSize |
| Vectors | | AREA EXPORT | RESET, READONLYVectors |
| | | DCD DCD | MyStackMem + StackSize Reset_Handler |
| dst | | AREA SPACE | MYDATA, DATA, READWRITE 8 |
| src0 src1 | | AREA DCB DCB | MYDCODE, CODE, READONLY "UWB", 0 "CSS", 0 |
| Reset Handler | | ALIGN ENTRY EXPORT | Reset_Handler |
| | nandier | LDR LDR LDR | R0, =src0 R1, =src1 R2, =dst |
| loop1 | | LDRB CBZ STRB B | R3, [R0], #1 R3, next R3, [R2], #1 |
| next | | MOV | R3, '.' R3, [R2], #1 |
| | | STRR | R3. [R2]. #1 |
| loop2 | | STRB LDRB CBZ STRB | R3, [R1], #1 R3, end_prog |
| loop2 | og | LDRB CBZ | R3, [R1], #1 |
| | og | LDRB CBZ STRB B | R3, [R1], #1 R3, end_prog R3, [R2], #1 loop2 |
| | | LDRB CBZ STRB B | R3, [R1], #1 R3, end_prog R3, [R2], #1 loop2 end_prog |
| | Pg Address | LDRB CBZ STRB B | R3, [R1], #1 R3, end_prog R3, [R2], #1 loop2 end_prog |
| | | LDRB CBZ STRB B | R3, [R1], #1 R3, end_prog R3, [R2], #1 loop2 end_prog Contents DRAM |
| | Address | LDRB CBZ STRB B B END | R3, [R1], #1 R3, end_prog R3, [R2], #1 loop2 end_prog Contents DRAM Peripherals |
| | Address 0x60000 | LDRB CBZ STRB B B END | R3, [R1], #1 R3, end_prog R3, [R2], #1 loop2 end_prog Contents DRAM |
| | Address | LDRB CBZ STRB B B END | R3, [R1], #1 R3, end_prog R3, [R2], #1 loop2 end_prog Contents DRAM Peripherals |
| | Address 0x60000 0x40000 0x20000 | LDRB CBZ STRB B B END | R3, [R1], #1 R3, end_prog R3, [R2], #1 loop2 end_prog Contents DRAM Peripherals |
| | Address 0x60000 0x40000 0x20000 0x20000 | LDRB CBZ STRB B B END | R3, [R1], #1 R3, end_prog R3, [R2], #1 loop2 end_prog Contents DRAM Peripherals |
| | Address 0x60000 0x40000 0x20000 0x00000 | LDRB CBZ STRB B B END 00000 00000 00000 00000 00000 | R3, [R1], #1 R3, end_prog R3, [R2], #1 loop2 end_prog Contents DRAM Peripherals |
| | Address 0x60000 0x40000 0x20000 0x20000 0x00000 | LDRB CBZ STRB B B END 0000 0000 0000 0008 00000 0014 0012 | R3, [R1], #1 R3, end_prog R3, [R2], #1 loop2 end_prog Contents DRAM Peripherals |
| | Address 0x60000 0x40000 0x20000 0x20000 0x00000 0x00000 | LDRB CBZ STRB B B END 0000 0000 0000 0008 00000 0014 0012 0010 | R3, [R1], #1 R3, end_prog R3, [R2], #1 loop2 end_prog Contents DRAM Peripherals |
| | Address 0x60000 0x40000 0x20000 0x20000 0x0000 0x00000 | LDRB CBZ STRB B B END 0000 0000 0008 0000 0014 0012 0010 000C | R3, [R1], #1 R3, end_prog R3, [R2], #1 loop2 end_prog Contents DRAM Peripherals |
| | Address 0x60000 0x40000 0x20000 0x0000 0x0000 0x0000 0x0000 | LDRB CBZ STRB B B END 0000 0000 0000 0008 00000 0014 0012 00010 0000C 00008 | R3, [R1], #1 R3, end_prog R3, [R2], #1 loop2 end_prog Contents DRAM Peripherals |
| | Address 0x60000 0x40000 0x20000 0x20000 0x0000 0x00000 | LDRB CBZ STRB B B END 0000 0000 0000 0000 00014 0012 0010 0000C 00008 00004 | R3, [R1], #1 R3, end_prog R3, [R2], #1 loop2 end_prog Contents DRAM Peripherals |

Q3. (12 pts) Pointer operations

On slide deck 6.ARM-InstrMem, we studied how to traverse a linked list using pre-indexed/register offset addressing. The following code intends to travers a linked list in search for a given value in R0 and returns the address of this value into R1 (but not the address of the node). If the value was not found, it returns 0 in R1, (i.e., a null address).

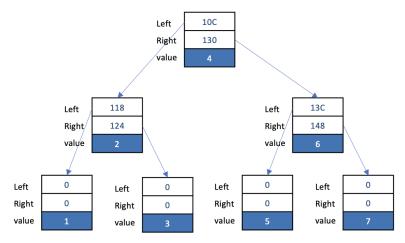




How about traversing a binary search tree?

R0 maintains a value to search. R1 first points to the tree root and is used to access each tree node. You can access its left pointer, right pointer, and value with

```
struct node {
    struct node *left; // R1
    struct node *right;// R1 + 4
    int value; // R1 + 8
}
```

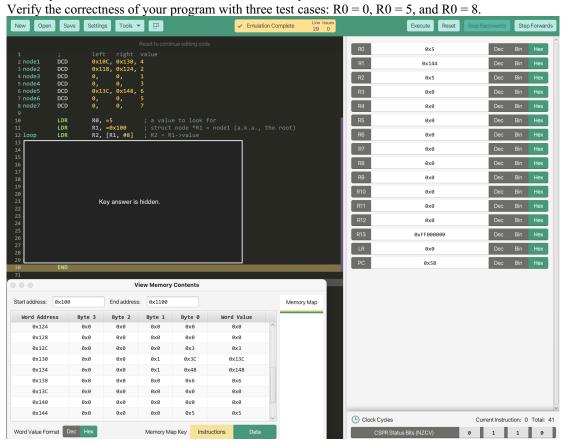


Using VisUAL, write a binary-tree search program. Your program searches for a value given in R0 and returns the address of this value into R1 (but not the address of the node). If the value was not found, it returns 0 in R1 (i.e., a null address).

Initialize a tree with the following code:

| | ; | left right value |
|-------|-----|------------------|
| node1 | DCD | 0x10C, 0x130, 4 |
| node2 | DCD | 0x118, 0x124, 2 |
| node3 | DCD | 0, 0, 1 |
| node4 | DCD | 0, 0, 3 |
| node5 | DCD | 0x13C, 0x148, 6 |
| node6 | DCD | 0, 0, 5 |
| node7 | DCD | 0, 0, 7 |

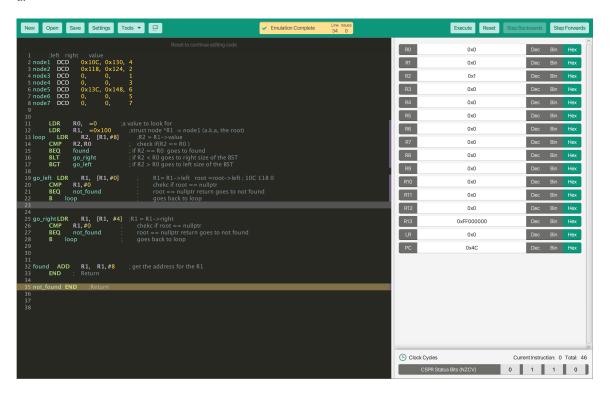
You may assume that the tree root is located at memory address 0x100.



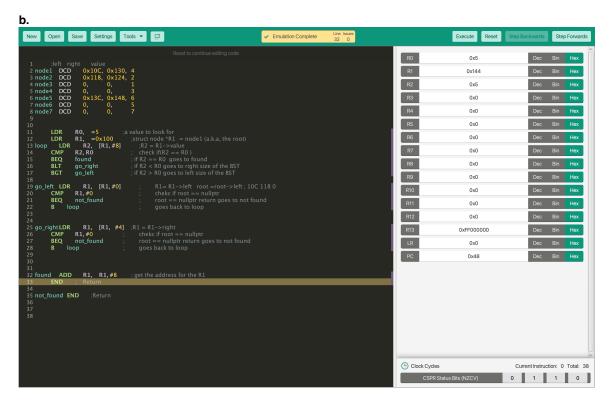
What to submit: source code, screen shorts, and short explanations

- 1. (6 pts) Your source code named hw3q3.s: You need to add comment to each line of your code, otherwise, you get 0 for the coding part!
- 2. (6 pts) In the same file recording your answers to Q1 and Q2, add screenshots and explanations for the following three test cases
 - a. Test case 1 (where R1 = 0)'s screenshot of registers (R1 R13, LR, and PC) and a short explanation: 2pt
 - b. Test case 2 (where R1 = 5)'s screenshot of registers (R1 R13, LR, and PC) and a short explanation: 2pt
 - c. Test case 3 (where R1 = 8)'s screenshot of registers (R1 R13, LR, and PC) and a short explanation: 2pt
 - d. Copy your source code to the file after the test case screenshots and explanations.

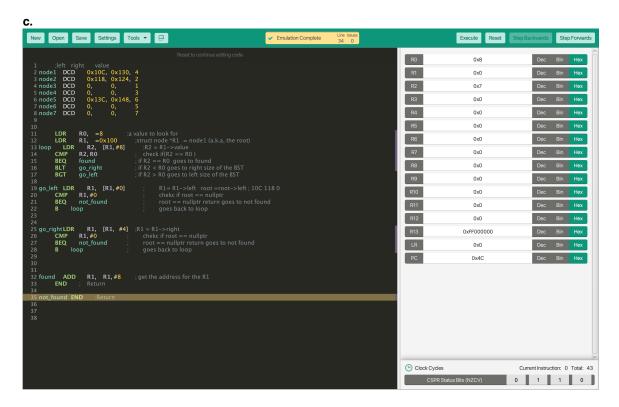
2. a.



Case 1: When R0 = 0. The program will set R0 = 0, then get the value for what R1 point to, which is 0x4, and place it in R2. After that, the program will compare R2 and R0. Since 0 is less than 4, it will go to the left side of the binary search tree. Then move root point to root->left, and go back to the loop to compare R2 and R0. At this time, R2 changes to 0x2, which is still greater than R0. It will keep moving left. Then R2 will become 0x1, and the root will reach the end, which means the program does not find 0 in the BST. Out of the loop.



Case 1: When R0 = 5. The program will set R0 = 0, then get the value for what R1 point to, which is 0x4, and place it in R2. After that, the program will compare R2 and R0. Since 5 is greater than 4, it will go to the right side of the binary search tree. Move the root to current = root->right, and check if it reaches the end. If not, go back to the loop, reset the value for the current node, and compare it with R0 since R6 is greater than R0. The node will go to the left side of the current node and check if it reaches the end. Then goes back to the loop to modify the current node's value. Which is 0x5, and it is equal to R0, and it will go to found and return the address of R1.



Case 1: When R0 = 8. The program will set R0 = 0, then get the value for what R1 point to, which is 0x4, and place it in R2. After that, the program will compare R2 and R0. Since 8 is greater than 4, it will go to the right side of the binary search tree. Then move root point to root->right, and go back to the loop to compare R2 and R0. At this time, R2 changes to 0x6, which is still less than R0. It will keep moving right. Then R2 will become 0x7, and the root will reach the end, which means the program does not find 0 in the BST. Out of the loop.

```
d.
                :left
                       right
                                       value
node1
       DCD
                       0x10C,
                               0x130,
node2
                       0x118, 0x124, 2
       DCD
node3
       DCD
                       0,
                                               0,
node4
       DCD
                                                                      3
                       0x13C, 0x148, 6
node5
       DCD
node6
                                               0,
                                                                      5
       DCD
node7
       DCD
                               R0,
               LDR
                                       =8
                                                               ;a value to look for
                               R1,
                                       =0x100
                                                               struct node *R1
               LDR
                                                                                      = node1
(a.k.a, the root)
loop
                                       [R1, #8]
                                                                       R2 = R1->value
               LDR
                               R2, R0
                                                                              check if(R2 ==
               CMP
R0 )
                                                                      ; if R2 == R0 goes to
               BEQ
                               found
found
                               go_right
               BLT
                                                                              ; if R2 < R0
goes to right size of the BST
```

left size of the BST go_left ; if R2 > R0 goes to

go_left LDR R1, root =root->left ; 10C 118 0 [R1, #0] R1= R1->left

R1, #0 CMP

chekc if root == nullptr

BEQ not_found root == nullptr return goes to not found

loop goes back to loop

R1, R1, #0 go_right [R1, #4] ;R1 = R1->right LDR

chekc if root CMP

== nullptr root == nullptr BEQ not_found

return goes to not found loop

В goes back to loop

;getting the address for R1 R1, R1, #8 found ADD

END return ;

not_found return **END**