

# Problem Solving

## Lab #12 (Chapter 11)

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## *Distinct Subsequences*

**PC/UVa IDs: 111102/10069, Popularity: B, Success rate: average Level: 3**

- A subsequence of a given sequence  $S$  consists of  $S$  with zero or more elements deleted. Formally, a sequence  $Z = z_1 z_2 \dots z_k$  is a subsequence of  $X = x_1 x_2 \dots x_m$  if there exists a strictly increasing sequence  $\langle i_1, i_2, \dots, i_k \rangle$  of indices of  $X$  such that for all  $j = 1, 2, \dots, k$ , we have  $x_{i_j} = z_j$ . For example,  $Z = bcdb$  is a subsequence of  $X = abcb dab$  with corresponding index sequence  $\langle 2, 3, 5, 7 \rangle$ .
- Your job is to write a program that counts the number of occurrences of  $Z$  in  $X$  as a subsequence such that each has a distinct index sequence.

### **Input**

- The first line of the input contains an integer  $N$  indicating the number of test cases to follow. The first line of each test case contains a string  $X$ , composed entirely of lowercase alphabetic characters and having length no greater than 10,000. The second line contains another string  $Z$  having length no greater than 100 and also composed of only lowercase alphabetic characters. Be assured that neither  $Z$  nor any prefix or suffix of  $Z$  will have more than  $10^{100}$  distinct occurrences in  $X$  as a subsequence.

## Output

For each test case, output the number of distinct occurrences of Z in X as a subsequence.

Output for each input set must be on a separate line.

### Sample Input

2  
babgbag  
bag  
rabbbit  
rabbit

### Sample Output

5  
3

# What you have to do

## 1. Analysis of the problem

- Understand what the problem is.

## 2. Make test cases

- Make test cases more than 5 for the problem
- A test case is usually a single step, or occasionally a sequence of steps, to test the correct behavior and functionalities.
- Test cases can be used to find some error in your partner program

## 3. Design

- Design a solution strategy for solving the problem and **write its pseudo code**

## 4. Implement the program

- Implement the program according to the pseudo code

## 5. Test

- Test implemented program by using the test cases

**\* Submit the pseudo code and the implemented program to the assignment submission board**

# SOLUTION TIPS

# UNDERSTANDING PROBLEM

## *Examples: Distinct Subsequences*

- Example 1:  $X = \text{babgbag}$ ,  $Z = \text{bag}$ 
  - **b**abgbag
  - **b**abgbag
  - **b**abgb**a**g
  - ba**b**gb**a**g
  - babg**b**a**g**
  - So, there are 5 distinct subsequences.
- Example 2:  $X = \text{rabbbit}$ ,  $z = \text{rabbit}$ 
  - **r**ab**b**bit
  - **r**ab**b**bit
  - **r**ab**b**it
  - So, there are 3 distinct subsequences

DESIGN



# *Design Sketch*

- This is an simple instance of dynamic programming problems.
  - What is the recurrence relation in this problem?
- This is also an instance of counting problems and the maximum number of occurrences is  $10^{100}$ .
  - We need to use a Big Integer representation.
  - For today's lab, design and implement class `BigInt` for representing big integers.

# Dynamic Programming

- Let  $X$  be the sequence, and  $Y$  be the subsequence you're looking for.
- Let  $a[i][j]$  be the number of times the sequence  $(Y_1 \dots Y_i)$  appears inside the string  $(X_1 \dots X_j)$ .
- Initially,  $a[0][j] = 1$  for all  $j$ , as the null string appears once in any string. Also  $a[i][0] = 0$  for all  $i > 0$  as a non-empty substring of the subsequence cannot appear in an empty string.
- Iterate through  $a[][]$  row-by-row with the following recurrence:
  - If  $Y_i == X_j$ ,  $a[i][j] = a[i-1][j-1] + a[i][j-1]$   
(You have 1 new match for every occurrence of  $Y_1 \dots Y_i$  in  $X_1 \dots X_j$ )
  - If  $Y_i != X_j$ ,  $a[i][j] = a[i][j-1]$   
(No new occurrence, so copy previous value)
- At the end, output the last cell of the array
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# Pseudo Code

- Initialization

```
for j = 0 to X.length: a[0][j] = 1;
for i = 1 to Z.length: a[i][0] = 0;
```

- Recurrence Steps

```
for i = 1 to Z.length
  for j = 1 to X.length
    if X[j] == Z[i]:
      a[i][j] = a[i][j-1] + a[i-1][j-1];
    else
      a[i][j] = a[i][j-1];
```

- Goal: we get the result at index  $a[Z.length][X.length]$ ;

# Example

$\begin{matrix} \diagdown \\ = \end{matrix}$	X =	b	a	b	g	b	a	g
b	1	1	1	1	1	1	1	1
a	0							
g	0							

Initial Step of matrix a

$\begin{matrix} \diagdown \\ = \end{matrix}$	X =	b	a	b	g	b	a	g
b	1	1	1	1	1	1	1	1
a	0	1	1	2	2	3	3	3
g	0							

After 1<sup>st</sup> element scan

if  $X(j) == Y(i)$ :  
otherwise:

$a[i][j] = a[i][j-1] + a[i-1][j-1];$   
 $a[i][j] = a[i][j-1];$

# Example cont...

$Y \backslash X =$	b	a	b	g	b	a	g
b	1	1	1	1	1	1	1
a	0	1	1	2	2	3	3
g	0	0	1	1	1	1	4

Arrows indicate the calculation of the 'g' row from the 'a' row and the 'b' row. For example, the value 1 in the 'g' row under 'a' is derived from the 0 in the 'a' row under 'a' and the 1 in the 'b' row under 'a'.

After 2<sup>nd</sup> element scan

$Y \backslash X =$	b	a	b	g	b	a	g
b	1	1	1	1	1	1	1
a	0	1	1	2	2	3	3
g	0	0	1	1	1	4	5

Arrows indicate the calculation of the 'g' row from the 'a' row and the 'b' row. The final value 5 in the 'g' row under 'g' is circled in red, indicating the final output.

After 3<sup>rd</sup> element scan

if  $X(j) == Y(i)$ :  
otherwise:

$a[i][j] = a[i][j-1] + a[i-1][j-1];$   
 $a[i][j] = a[i][j-1];$

Finally we find output 5

## *C Strings Indexing Problem*

- Indices of C strings range from 0 to string length – 1.
- So, we need to be careful for using indices when comparing strings. We have to refer to  $X[j-1]$  for  $j$ th character of string  $X$  if  $X$  is represented as C string.

```
for i = 1 to Z.length
  for j = 1 to X.length
    if  $X[j-1] == Z[i-1]$ :
       $a[i][j] = a[i][j-1] + a[i-1][j-1]$ ;
    else
       $a[i][j] = a[i][j-1]$ ;
```

## *Another Solution: C Strings Indexing Problem*

- Another approach to solve this problem is to use different interpretation for matrix  $a$ .
  - Old: Let  $a[i][j]$  be the number of times the sequence  $(Y_1 \dots Y_i)$  appears inside the string  $(X_1 \dots X_j)$ .
    - Assumption: The strings are stored from index 1.
  - New: Let  $a[i][j]$  be the number of times the sequence  $(Y_0 \dots Y_i)$  appears inside the string  $(X_0 \dots X_j)$ .
    - Assumption: The strings are stored from index 0.

- The new definition requires to change the initialization.
- According to the new definition,  $a[0][j]$  should be interpreted as # of times the sequence  $(Y_0)$  appears in inside the string  $(X_0 \dots X_j)$ . So, we have

```
a[0][0] = (X[0] == Z[0]) ? 1:0;
for (int j = 1; j < X.length; j++)
    a[0][j] = a[0][j-1] + ((X[j] == Z[0]) ? 1:0);
```

- Similarly,  $a[i][0]$  represents # of times the sequence  $(Y_0 \dots Y_i)$  appears in inside the string  $(X_0)$ . It is 0 when  $i > 0$ . So, we have
- ```
for (int i = 1; i < Z.length; i++) a[i][0] = 0;
```



## Optimization: Required Memory for Dynamic Programming

- This solution requires a dynamic programming matrix whose size is  $Z.length * X.length$ . Can we reduce its size?
- It is an interesting observation that  $a[i][j]$  only refers to the previous row,  $a[i-1][*]$ . At any time, we are using only two rows: the previous and the current.

|   | X = | b | a | b | g | b | a | g |
|---|-----|---|---|---|---|---|---|---|
| 1 | 1   | 1 | 1 | 1 | 1 | 1 | 1 | 1 |
| b | 0   | 1 | 1 | 2 | 2 | 3 | 3 | 3 |
| a | 0   | 0 | 1 | 1 | 1 | 1 | 4 | 4 |
| g | 0   |   |   |   |   |   |   |   |

After 2<sup>nd</sup> element scan

- So we can do dynamic programming with only two rows. We declare the matrix as `a[2][X.length]`.
- When processing row `i`, `a[i%2]` is used as the current row while the other the previous.
- Of course, we need reset the elements of the array representing the current row at the beginning of each row process.

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# BIG INTEGER NUMBER & SUM

- Suppose one two big integer number

A = 1234858548957489574895784750894758

B = 3843248439048392493849038493489034

$$\begin{array}{r}
 1234858548957489574895784750894758 \\
 3843248439048392493849038493489034 \\
 \hline
 \text{Sum} = 50780 \dots\dots\dots 83792
 \end{array}$$



At first, 2 big numbers are placed two arrays, respectively. Then add their values based on indexing while adding carry to the next indexing value.

## Optimization: Big Integer Representation

- Suppose  $\text{length}(X) = 10,000$  and  $\text{length}(Y) = 100$  then we need a matrix of size  $10,000 * 100 = 1,000,000$ .
- If each element of the matrix has array of size 100, then we need 100,000,000, a big number. It is also very time consuming to perform a huge number of add operations. You can get TLE.
- So we need to modified BIG INTEGER NUMBER & SUM calculation technique.
- Instead of storing one digit in one element of the array used for storing big numbers, we can store more digits. Suppose we store 4 digits. Then we need an array of 25 integers to represent  $10^{100}$  and computation time is reduced by factor 4.

$$\begin{array}{r}
 1234858548957489574895784750894758 \\
 3843248439048392493849038493489034 \\
 \hline
 \text{Sum} = 50780 \dots\dots\dots 83792
 \end{array}$$

1  
 Carry =  $13792 / 10000$   
 sum =  $13792 \% 10000$