Systems and Networking – Unit I

B.Sc. in Applied Computer Science and Artificial Intelligence 2021-2022

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File System API

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File creation, manipulation, protection, etc.

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OS Implementation

OS internal data structures and algorithms

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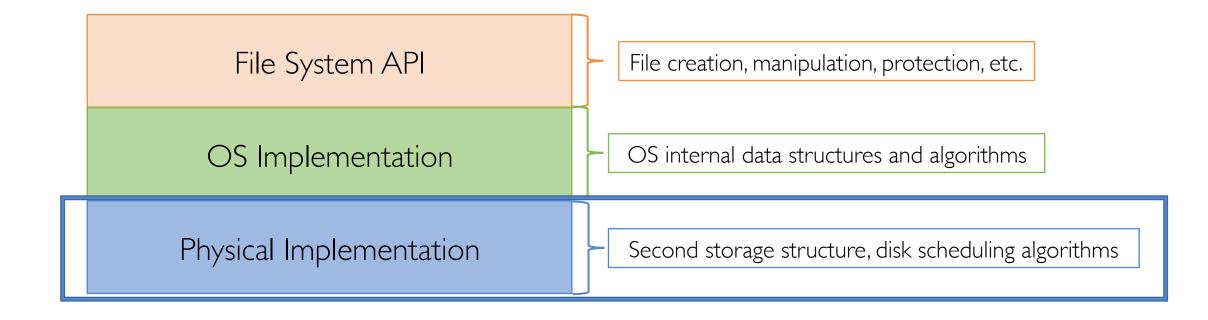
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OS Implementation

OS internal data structures and algorithms

Physical Implementation

Second storage structure, disk scheduling algorithms

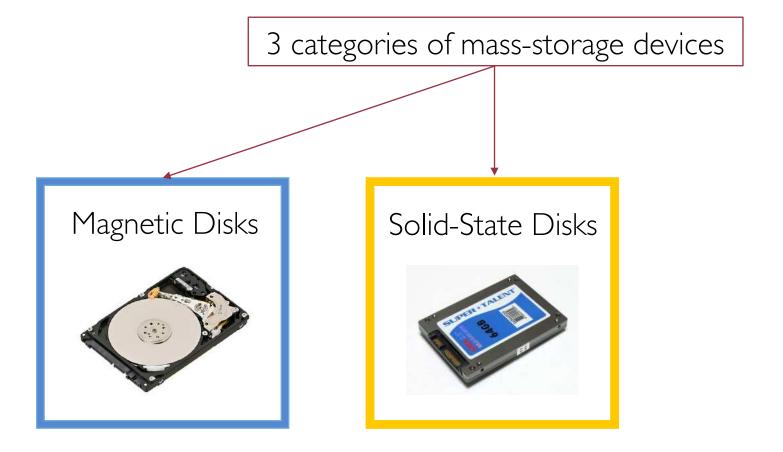


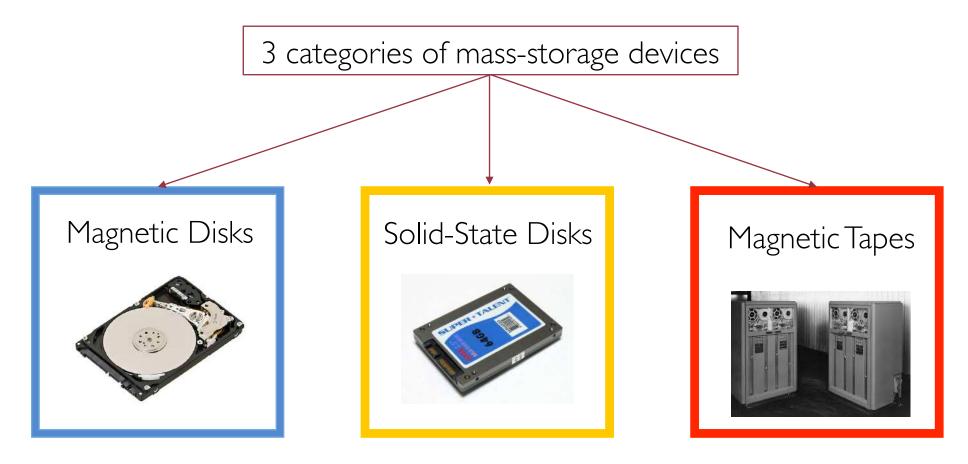
Part V: Storage Management

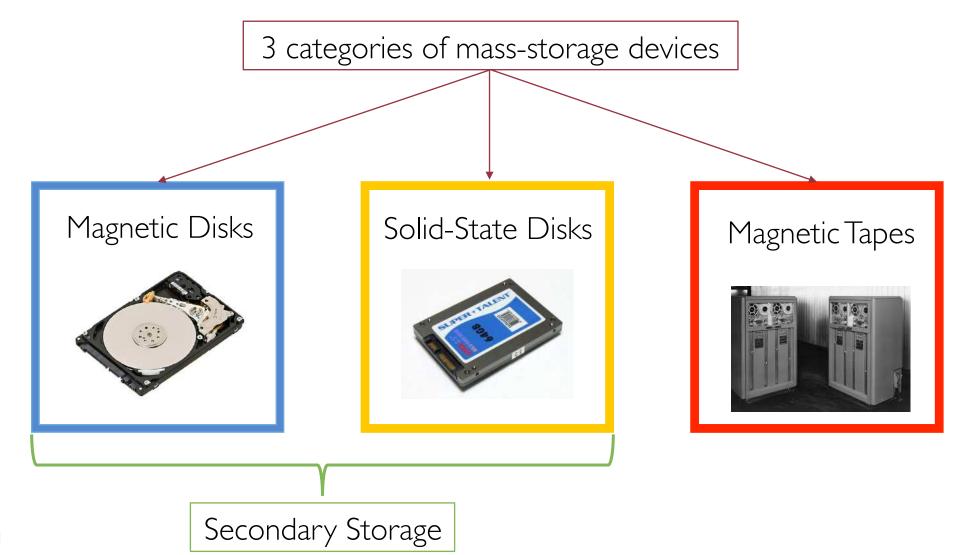
3 categories of mass-storage devices

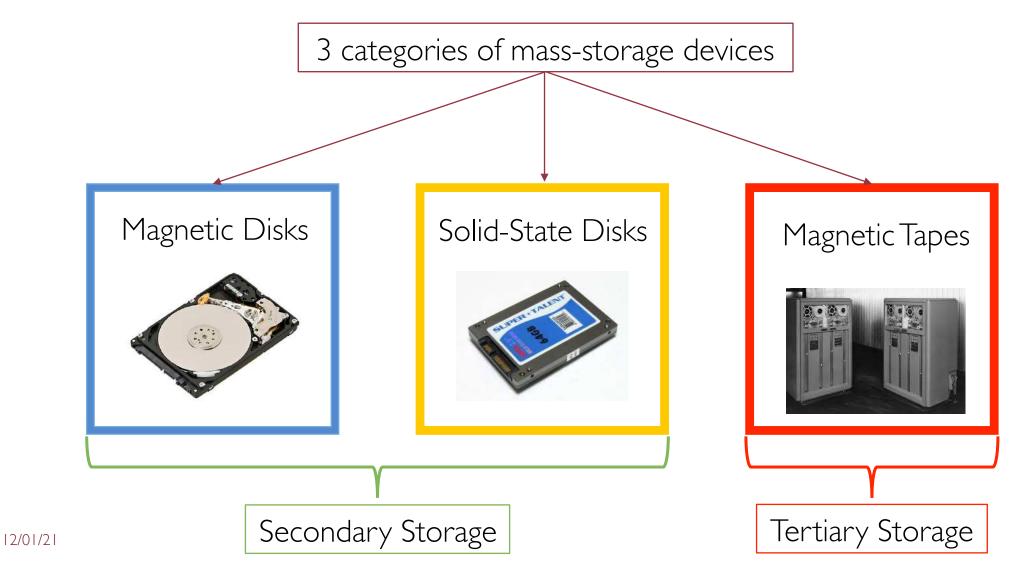
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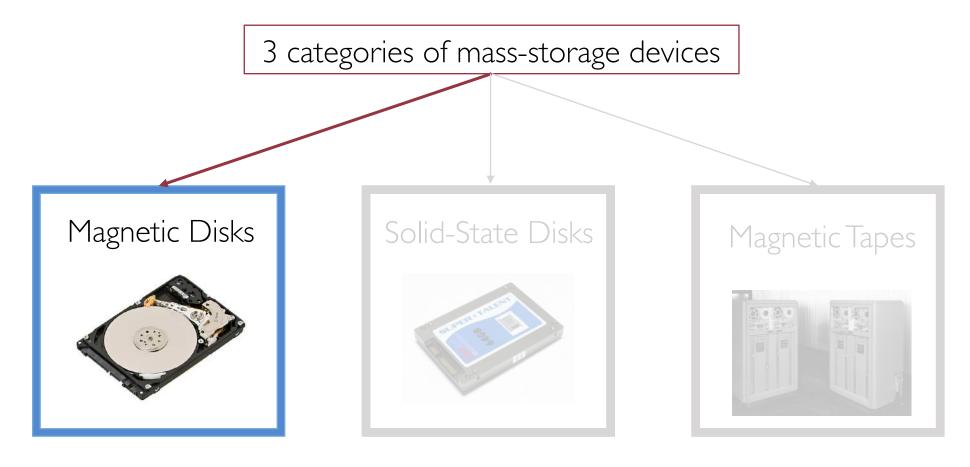






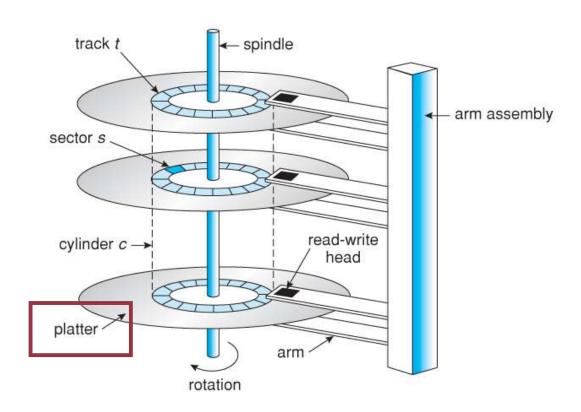




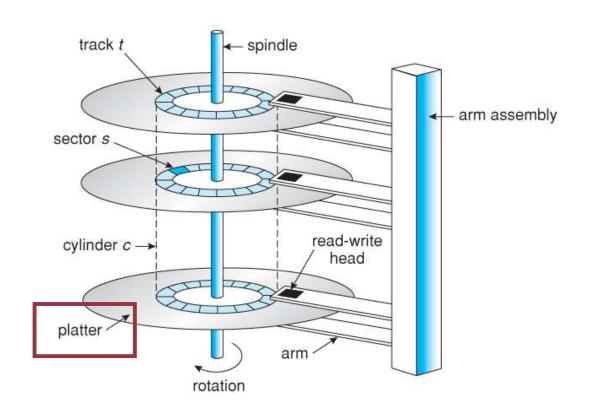


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14

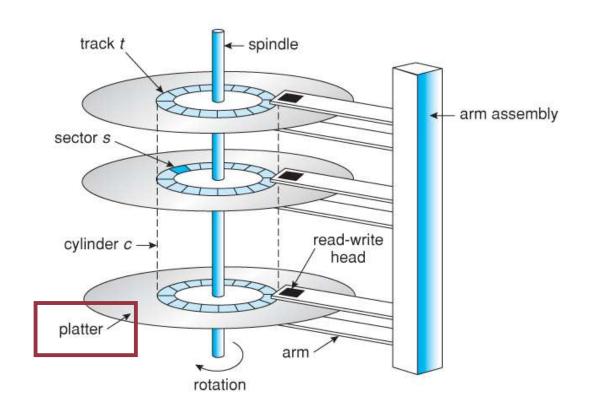


One or more platters covered with magnetic media



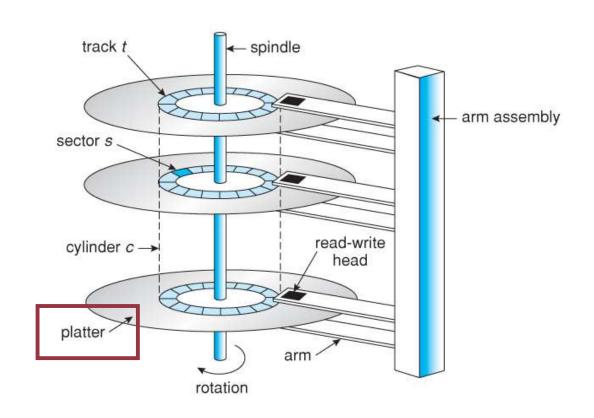
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Hard disk rigid metal



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Hard disk rigid metal Floppy disk flexible plastic

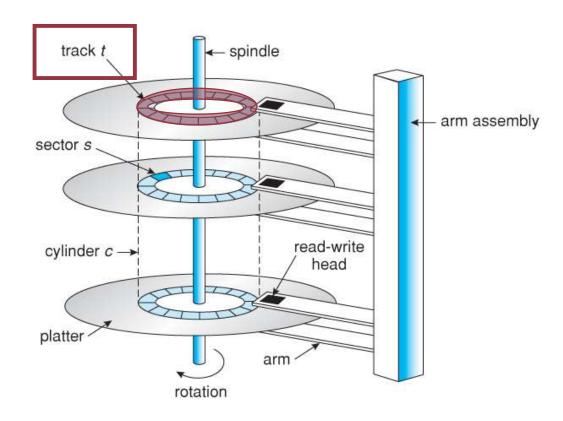


One or more platters covered with magnetic media

Hard disk rigid metal Floppy disk flexible plastic

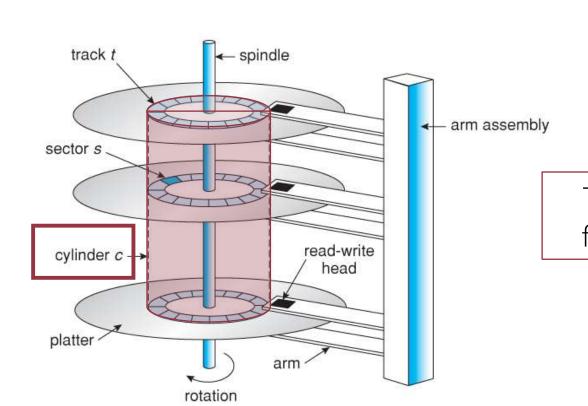
Each platter has 2 working surfaces

Magnetic Disks: Tracks and Cylinders



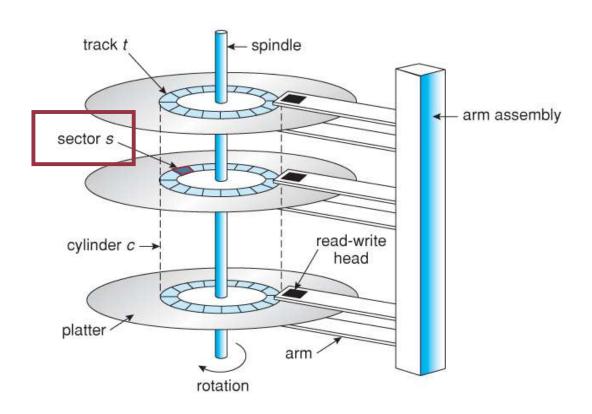
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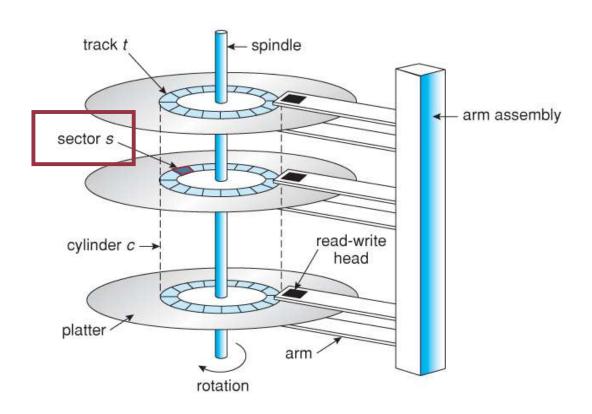


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The set of all tracks that are the same distance from the edge of the platter is called a cylinder

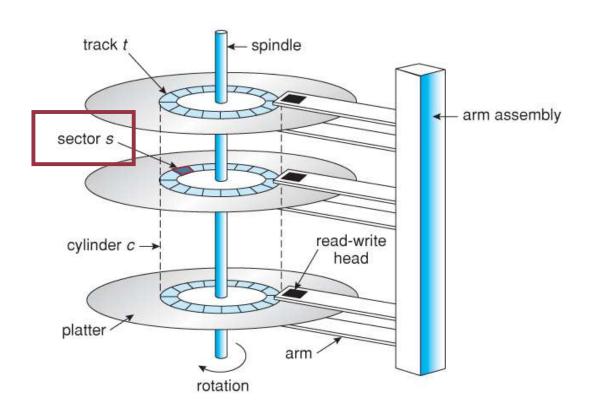


Each track is further divided into sectors



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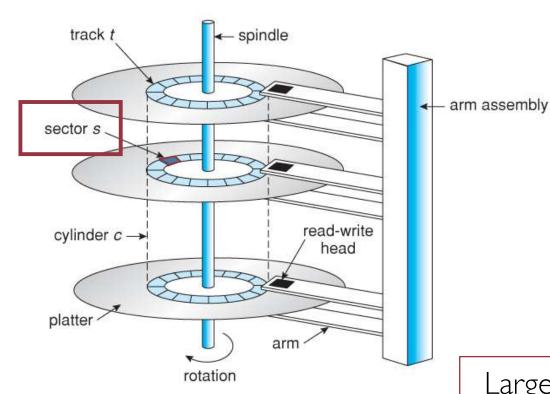
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Sectors also include a header and a trailer, and checksum information

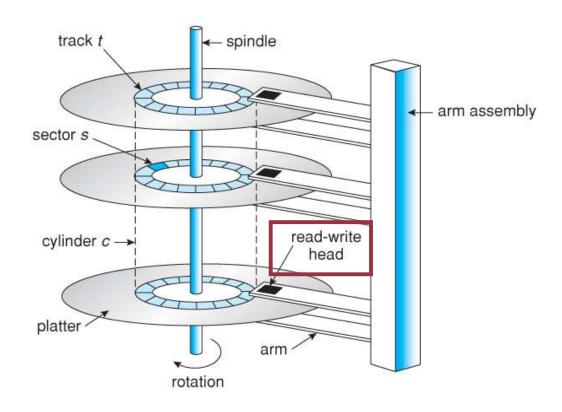


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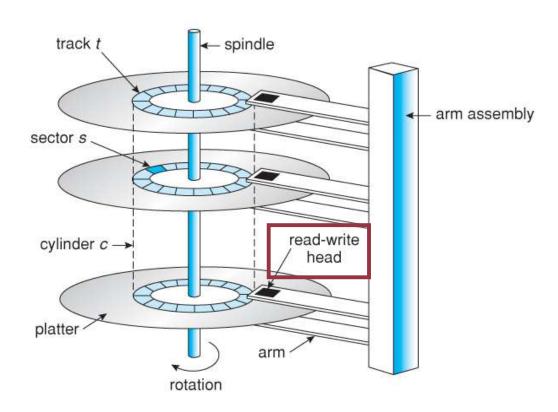
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Larger sector sizes reduce the space wasted by headers and trailers, but increase internal fragmentation

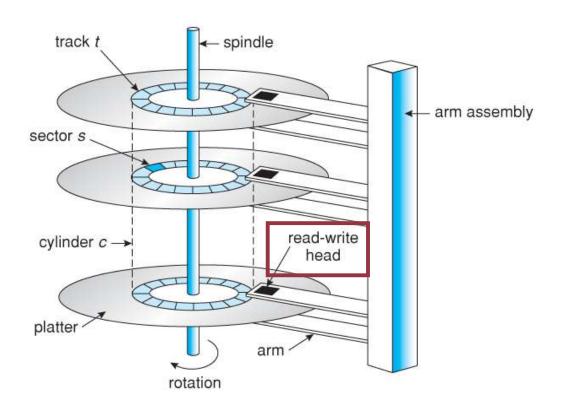


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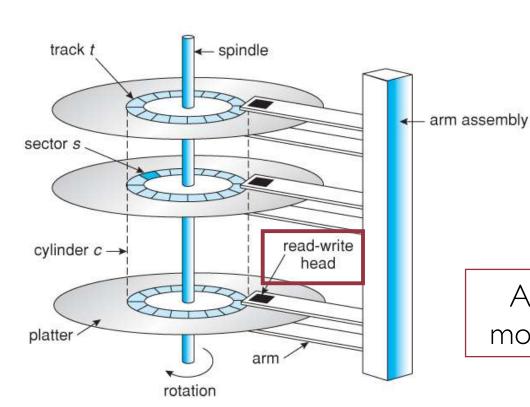
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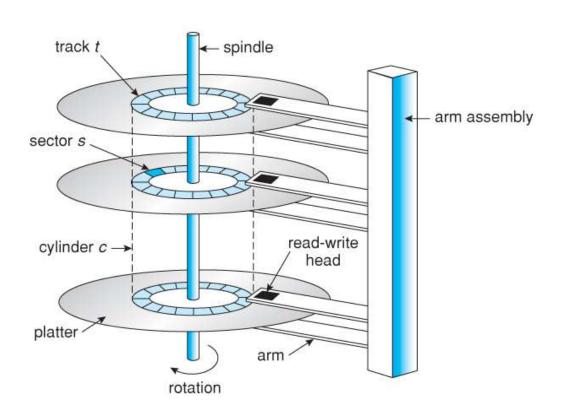


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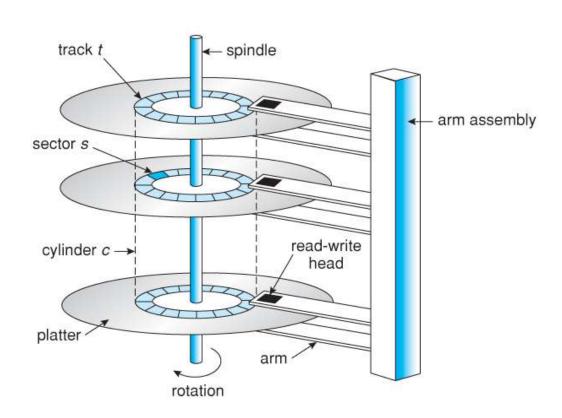
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Arms are controlled by a common arm assembly moving simultaneously from one cylinder to another

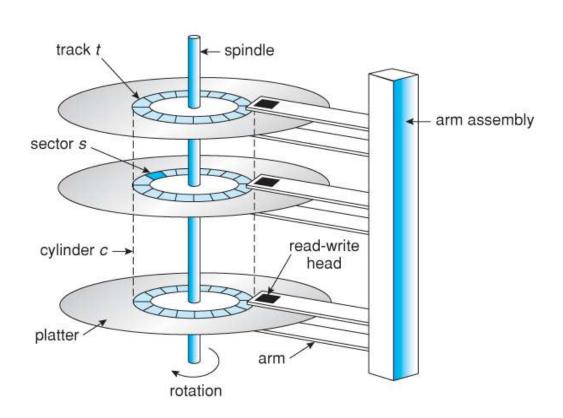


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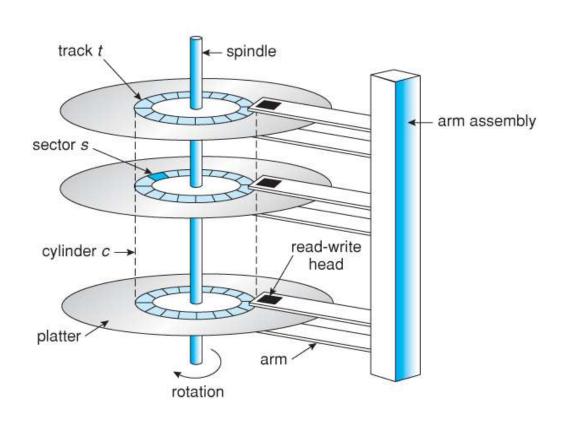
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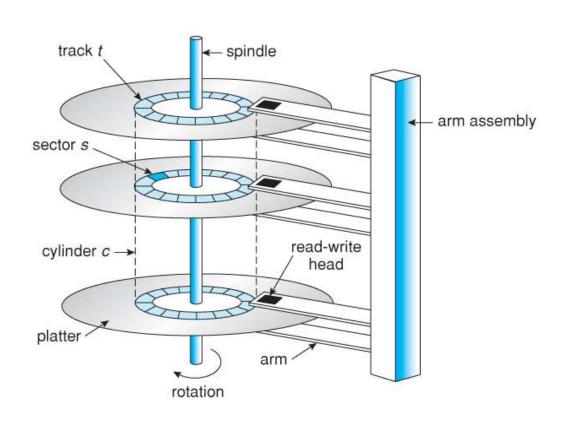


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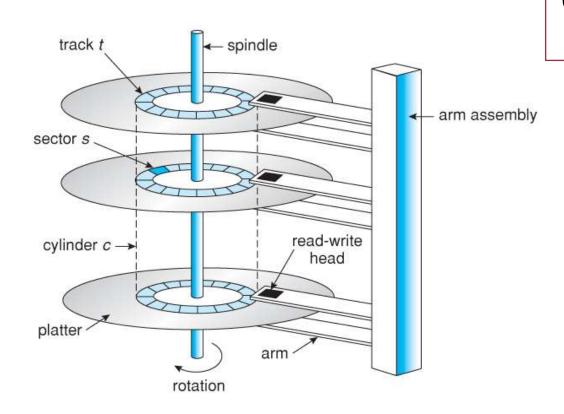
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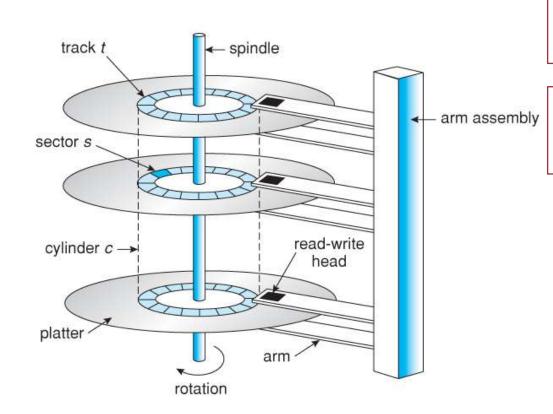
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$$C = H * T * S * B$$

OVERALL CAPACITY

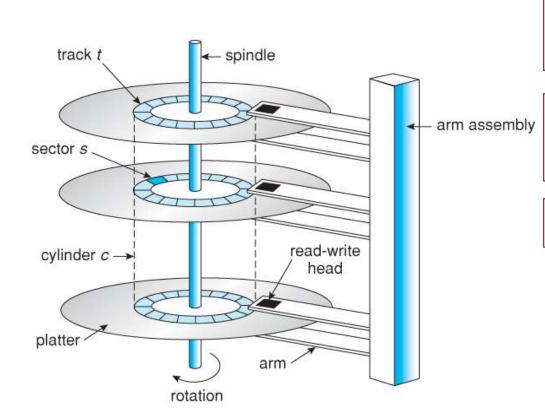


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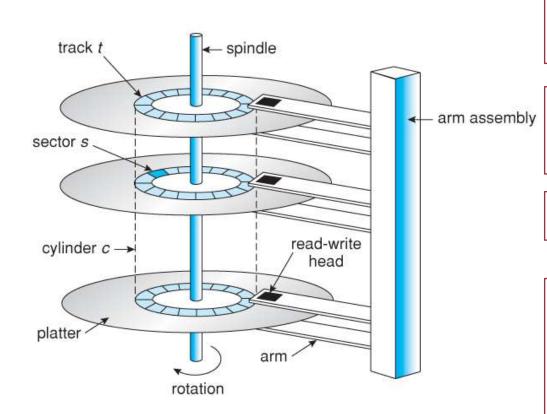
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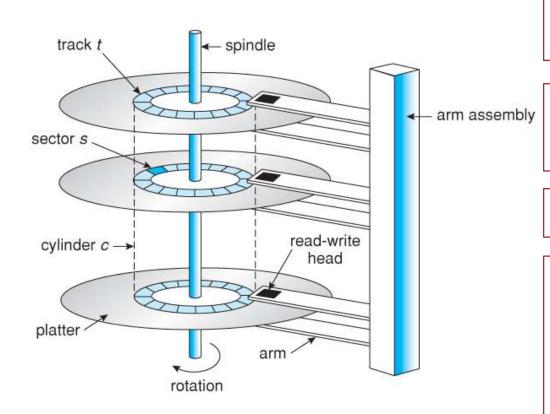
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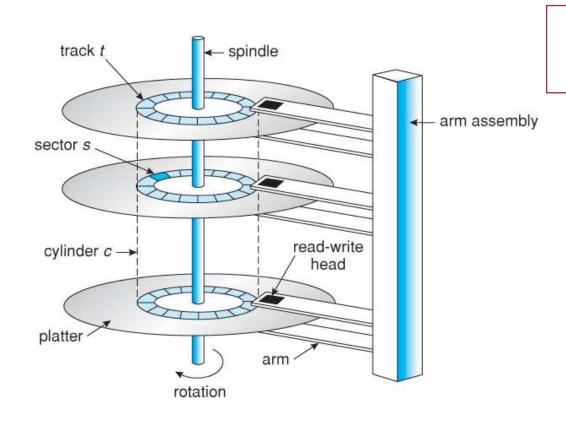
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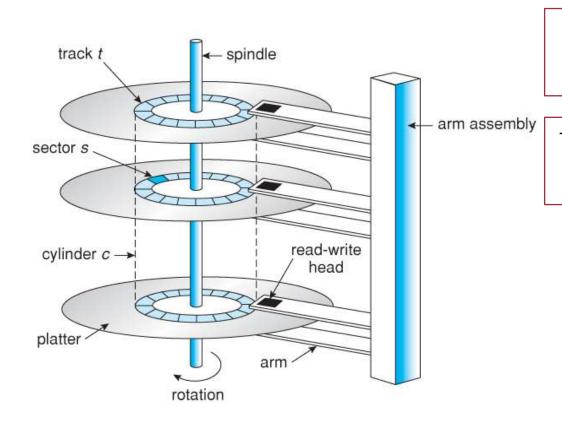
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- Different frequencies and timing from innermost to outermost tracks

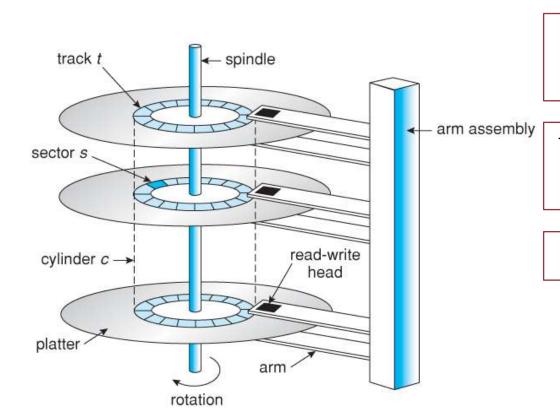


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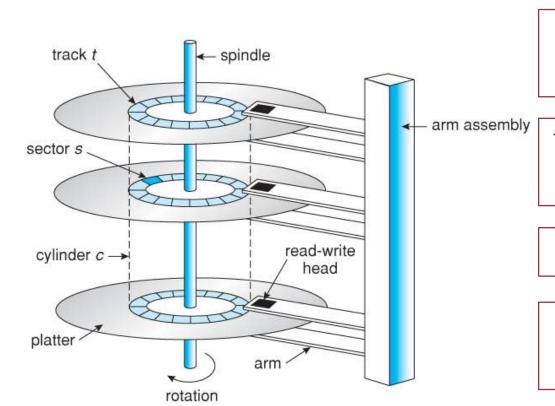
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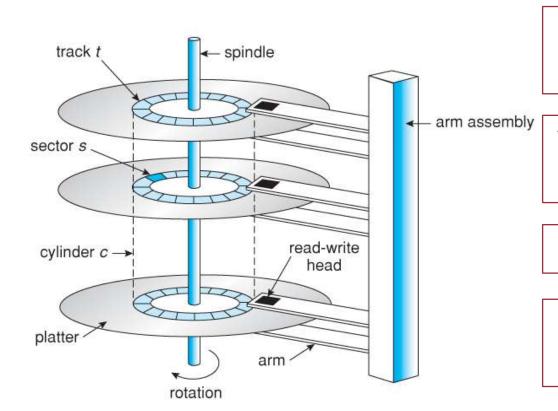


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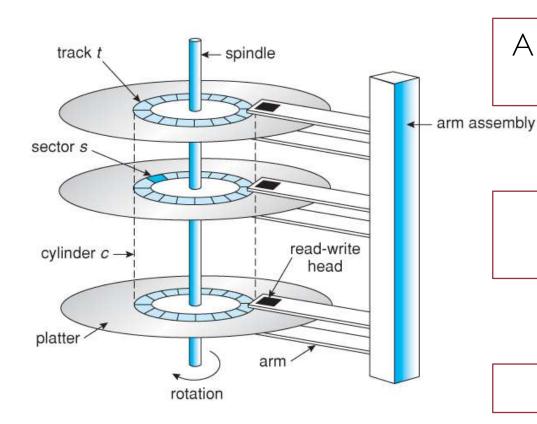
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Zone Bit Recording (ZBR)

Magnetic Disks: (Logical) Referencing



A physical block of data is specified by the (head, cylinder, sector) number

Disk blocks are numbered starting at the outermost cylinder, identified by 0

Note that cylinder coincides with track

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12/01/21 45

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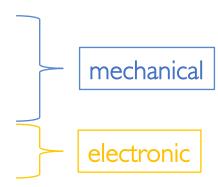
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mechanical

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Magnetic Disks: Positioning (Seek) Time

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The second highest bottleneck

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Data Transfer Time = Seek Time + Rotational Delay + Transfer Time

Sometimes the term transfer rate is used to refer to the overall data transfer time

12/01/21 59

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- The array of blocks is mapped onto disk sectors sequentially
- Sector 0 is the first sector of the first track of the outermost cylinder
 - The mapping proceeds in order through that track
 - Then through the rest of tracks in the same cylinder
 - Then through other cylinders (from the outermost to innermost)

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- Parking heads means to move them off the disk or to an area where no data is stored

Magnetic Disks: Interfaces

- Hard drives may be removable as floppy disks, and some are even hotswappable
 - they can be removed while the computer is running
- Disk drives are connected to the computer via the I/O bus
- Some of the common interface formats include:
 - Enhanced Integrated Drive Electronics (EIDE);
 - Advanced Technology Attachment (ATA) and Serial ATA (SATA);
 - Universal Serial Bus (USB);
 - Fiber Channel (FC);
 - Small Computer Systems Interface (SCSI)

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- Finally, data is transferred from that cache to the host controller and the motherboard memory at electronic speeds

74

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12/01/21 75

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Hardware Optimization

- How can the OS help minimize data transfer time?
- Schedule disk operations so as to minimize head movement
- Lay out data on disk so that related data are located on close tracks
- Place commonly-used data on a specific portion of the disk
- Pick carefully the block size contained on each sector:
 - Too small → more seeks are needed to transfer the same amount of data
 - Too large → more internal fragmentation and space wasted

Summary

- Disks are slow devices compared to CPUs (and main memory)
- Manage those device efficiently is crucial
- Minimize seek and rotational delay on magnetic disks
- HW optimizations are limited > OS needs to take the lead here!