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Snips Documentation & Manual

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1. About SNIPS

The Language

Snips is a C/Java oriented programming language. It is meant to be a simple, easy to use language, that us used best for small programs. The Language does provide a small built in Library, of which you can learn more in a later chapter.

The compiler translates the language into ARM Assembly. The output file contains all functions, imports, direct imports as well as transitive imports. The output file can be seen as a completely independent source file. All that's needed to run the code is already included. This means that no kind of linking has to be done, the assembly only has to be converted into binary.

The language features some basic features, like functions, control structures and loops. More advanced features like Templating, Predicates and Structs are supported as well. The language implements a wide set of operators, from arithmetic to boolean to bitwise and comparison operators.

An import system is also built in. Imports stated by the include directive are automatically processed and imported. This allows for a more distributed programming style and higher reusability.

About this document

This document aims to introduce and familiarize with the language as a whole and give a brief overview over some techniques and the included standard library. Also, a few code examples are provided.

Some words in advance

This project was started and still is for educational purposes. The programming language Snips, the Compiler and all included modules are not following any standards and are built to function well only for this project. Results produced by the compiler and included modules may contain errors and are not thought for any production environment. The project and all its included modules are still under development and are subject to change, both in functionality, as well in language syntax and behaviour.

2. Language Syntax

Type System

Primitives

Snips has a range of built-in primitive types. All of the types are 1 32-bit word large.

Void. The void type acts as a universal primitive type. Using this type requires caution, because of its unique property to match every other type. This is not only true for primitives, but for every possible type. This can be used to create “don’t care” types. But, to improve type security, this should be avoided. Also, since the void type is 1 word large, assigning an expression to this type that is larger than 1 word can cause major problems. Consider **Void Arrays** as a workaround. A Void is declared with the type specifier *void*. For example:

```
void v = 20;
```

Integer. The integer represents the most common numeric type, with a range from $-2^{31} - 1$ to 2^{31} . An Integer is declared with the type specifier *int*. For example:

```
int x = 10;
```

Boolean. The boolean type is mostly used for logic. Its value range is the exact same as the one of the Integer, but when interpreting the value, only the value **0** is interpreted as *false*, all other possible values are interpreted as *true*. A boolean is declared with the type specifier *bool*. For example:

```
bool b = true;
```

Character. The character type can be used to store UTF-8 encoded Characters. The value Range is, because of the 32-bit size, like the one of the Integer, but the usable Range is only from **0** to **255**. The **0** value takes a special role. In a String, a char with the value **0** indicates the end of the String. A character is declared with the type specifier *char*. For example:

```
char c = 'a';
```

Enumerations. The enumeration type is not necessarily a primitive type, but is handled like one on the low level. An enumeration has to be declared, see Enum Typedef. After Declaration, an Enum value can be selected. An Enum is declared with the name of the Enum itself. For Example:

```
State s = State.Normal;
```

Predicates. A predicate type holds a reference to a function. This means that a predicate type can be “called” like a function. You can read more about it in the Section “Predicates”. A predicate type is declared using the *func* type specifier. For example:

```
func pred = my_predicate;
```

In this example, `my_predicate` is the name of a function already declared. By calling this predicate the program will execute the function using the given parameters.

Warning: Using predicates requires caution, especially with anonymous predicates. See the Predicates Section for more information. When handled wrong, predicates can cause the program to crash or the stack to be misaligned.

Composite Types

Arrays. An array holds a fixed amount of values of the same type. Arrays can be created with a fixed size, or by allocating memory on the heap. Arrays can hold any type, both primitive and composite, except the proviso type. An array is declared like this:

```
int [5] arr = {1, 2, 3, 4, 5};
```

In this case, an array of the length 5 with the element type `int` is created. Also, using an array init expression, we assign the array a value. Arrays of arrays are also possible:

```
int [2] [2] mat = {{1, 2}, {3, 4}};
```

In this case, we declare an integer matrix of the size 2 by 2. Again, we initialize the matrix with an array init expression, except the values of this expression are array init expressions themselves. Arrays can be of any dimension. It may be worth considering to flatten the array when using pointers.

Structs. A struct holds a collection of fields. Each of which has its own separate type. Since all fields of the struct have a fixed size, the struct has a fixed size as well. The struct has to be declared to be a type using a Struct Typedef. A struct is initialized like this:

```
MyStruct s = MyStruct::(5, true);
```

In this case, a struct called `MyStruct` is initialized. The Type of `MyStruct` has an integer and boolean field, both of which we have to give an initial value in the Structure Init Expression.

Pointers. A pointer holds a reference to another object. When dereferencing a pointer, we get the value a pointer points to. A pointer can be created manually, or by retrieving the address of an object:

```
int* p = &var;
```

```
int* p = (int*) 0;
```

In the first example, we get the address of an object with an address-of expression. This will return a pointer to the object. The type of the pointer is determined through what type it points to. The * after the type signals that this is in fact a pointer. In the second example we create a Null-Pointer by casting the value 0 to an integer pointer. This can be useful in data structures to signal an invalid pointer or Null-Value. Pointers can point to a pointer as well:

```
int** p0 = &p;
```

In this example, p is an integer pointer. By getting the address of this pointer, we now have a pointer of the second grade. This means this pointer is now pointing to a pointer that points to the target value. We can retrieve the pointer by dereferencing this pointer twice.

Expressions

Operators

Snips brings a wide range of operators to the table. You can see all down below. The higher up, the higher precedence has the operator.

Operator	Description
Atom, Enum Selection: <i>a, 10, Enum. v</i>	Basic building blocks. Atoms can be immediate or variables. Enum Selection does count as a new Enum Immediate value.
Array Selection: <i>a [0], m [0][1]</i>	Select an element from an array.
Struct Select: <i>st. v, st-> v</i>	Select a field from a struct.
Increment, Decrement: <i>i ++, i --</i>	Increment or decrement a primitive value.
Unary Minus: <i>-val</i>	Negates value.
Not: <i>!b, ~i</i>	Boolean negation via <code>!</code> , bitwise complement with <code>~</code> .
Type Cast: <i>(int) b, (void) k</i>	Cast value to type.
Dereference: <i>*val</i>	Treat value as address, load value at this address.
Address Of: <i>&v</i>	Get the address of a value.
Size of: <i>sizeof(v), sizeof(Struct)</i>	Get the word size of a value or of a type.
Multiplication, Division: <i>a * b, a / b</i>	Multiply, Divide.
Addition, Subtraction: <i>a + b, a - b</i>	Add, Subtract.
Shift: <i>a << b, a >> b</i>	Shift logical.
Comparison: <i>a < b, <=, ==, !=, >, >=</i>	Compare two values based on comparator.
Bitwise And, xor and or: <i>a & b, a ^ b, a b</i>	Perform bitwise and, xor and or operation.
Logical And, Or: <i>a && b, a b</i>	Perform boolean and, or operation
Ternary: <i>(a) ? x : y</i>	Select one of two values based on compared value.
Array Initialization: <i>{a, b, ..., c}</i>	Create a new array of values of the same type.
Struct Initialization: <i>Struct::(a, b, ..., c)</i>	Create a new Struct instance by providing values for the structs fields.

3. ARM Assembly & Output

Compiler Assembly Conventions

Register Usage

The Snips Compiler maps certain functionality to the target machine registers, the mapping can be seen below:

Register	Functionalities
R0	Holds the result of an arithmetic operation, can hold a parameter in a function call, holds the return value of a function call if the return type word size is equal to 1.
R1, R2	Used as operands for arithmetic operations, can hold a parameter in a function call.
R3-R9	Part of the Register Stack, used to hold variables with word size 1.
R10	Holds PC backup during syscalls.
R11	Acts as the frame pointer.
R12	Holds the exception code when a exception is thrown.
R13	Acts as the stack pointer.
R14	Acts as the link register.
R15	Acts as the program counter.

Heap Management

The heap is managed by a custom linked list implementation that is optimized for a minimal memory footprint, and easy allocation and de-allocation. A single node consists of a size entry and the data block.

The size entry contains the word size of the memory block plus the size entry itself. The memory block simply contains the payload. Starting from the heap start, memory blocks can be created. When creating a new memory block, `resv` iterates over the heap. It starts at the heap start, loads the first size entry field n . If $n = 0$, no entry is here, and the memory location is free. If $n > 0$, a memory block of the size n lies ahead. The routine can then jump n words further and try again. If $n < 0$, a free memory section of the size $-n$ lies ahead. If the requested size s is less than n , the memory section ahead is large enough. Write s to the current cell, write $n + s$ to the current cell plus s cells. By doing this we split the free section, use the whole section or a part, and keep the heap structure intact. A pointer that points to the heap always points to the first word of the payload. This is especially true for pointers created by the `resv` routine. This means that $p - 1$ would point to the size entry, where p is a pointer to the heap.

The `free` routine now only has to negate the size entry of a memory section to mark it as unused. The routine can locate this entry through $p - 1$, where p is the passed pointer.

The hsize routine can load the size entry using the method described earlier. Note that the function will load a random value if the pointer does not point to the heap but to the stack. The resv and free routine also implement some heap defragmentation mechanisms that are automatically executed. These mechanisms make sure to defragment multiple free memory sections after another into one big free section. This way the search time for a free section is shorter and the utilization is denser.

4. Built in Libraries

Memory & Heap Routines

resv.sn

Full Path: -
 Package Type: Base Package
 Namespacing: None
 Description:

Contains the memory reserve routine that is responsible for allocating a requested block size in the heap. This file is included dynamically when the resv function is called.

Function Header	Description
void* resv(int size)	Reserves a memory block of given size + 1. The additional word is used for heap management. Returns a pointer to the memory location. The returned pointer points to the second data word of the block. The first word contains the heap management data. $O(n)$ where n is the number of heap elements.

free.sn

Full Path: -
 Package Type: Base Package
 Namespacing: None
 Description:

Contains the memory free routine that is responsible for freeing memory in the heap. This file is included dynamically when the free function is called.

Function Header	Description
void free(void* p)	Frees the given pointer from the heap. This is done by negating the heap management data on the block head. The data will remain in the heap, but can from now on be overwritten. $O(n)$, where n is the amount of heap elements.

hsize.sn

Full Path: -
Package Type: Base Package
Namespacing: None
Description:

Contains the memory hsize routine that can determine the size of a heap memory block. This file is included dynamically when the hsize function is called.

Function Header	Description
Int hsize(void* p)	Returns the size of the memory section. The pointer should point to a memory block in the heap. The size is determined by reading the heap management data and subtracting 1. $O(1)$

System Operators

__op_div.sn

Full Path: -

Package Type: Base Package

Namespacing: None

Description:

Contains a routine to divide two integers. This file is included dynamically when the division operator is used.

Function Header	Description
int __op_div(int a, int b)	Calculates $\frac{a}{b}$ and returns the result.

__op_mod.sn

Full Path: -

Package Type: Base Package

Namespacing: None

Description:

Contains a routine to compute the rest of a integer division. This file is included dynamically when the modulo operator is used.

Function Header	Description
int __op_mod(int a, int b)	Calculates $a \% b$ and returns the result.

Types

boolean.sn

Full Path: lib/std/type/boolean.sn
 Package Type: Includes: string.sn
 Namespacing: The entire package is namespaced in 'Boolean'.
 Description:
 Contains utility around the boolean type.

Function Header	Description
bool parseBool(char* str)	Parses a boolean value from given String. The result will be <i>true</i> if the String is equal to "true". In any other case, the result will be false.
char* toString(bool b)	Converts given boolean into String representation. The result will be, depending on given value, "true" or "false". Returns a pointer to the created String.

integer.sn

Full Path: lib/std/type/integer.sn
 Package Type: Base Package
 Namespacing: The entire package is namespaced in 'Integer'.
 Description:
 Contains utility around the integer type.

Function Header	Description
int parseInt(char* str)	Parses a integer from given String. The String has to match the pattern: $-?([0-9])^*$. The result is the parsed int.
char* toString(int num)	Converts given int to String representation. The resulting string will contain a sign if the int was negative. Returns a pointer to the created String.

Namespace Array:

Function Header	Description
void sort(int* arr, int size, func (int a, int b) -> bool pred)	Sorts given integer array. The sorting is done using a bubble sort algorithm, so the runtime complexity equals $O(n^2)$. The sorting is determined by given predicate. Writes back the result in the array at the given pointer.

string.sn

Full Path: lib/std/type/string.sn
 Package Type: Includes: linked_list.sn
 Namespacing: The entire package is namespaced in 'String'.
 Description:
 Contains utility for String operations.

Function Header	Description
bool equals(char* str0, char* str1)	Checks if the two given Strings are equal on a char level. All contained chars must have the same value, and both Strings have to have the same case.
char* substring(char* str, int begin, int end)	Cuts out a part of the String specified by the bounds. The begin index marks the first char to be included in the new String, the end index marks the last char to be included. The 0 char is inserted at the end of the resulting String. The begin index should be less than the end index. In case the indexes are out of the bounds of the String, the entire String is copied. Returns a pointer to the new String on the heap.
int length(char* str)	Returns the length of the String. This includes the 0 char. Returns the length.
char* concat(char* str0, char* str1)	Concatenates the two given Strings. The resulting String will contain the entire first String followed by the entire second String. The 0 char of the first String is discarded. Returns a pointer to the new String.

Data Structures

linked_list.sn

Full Path: lib/std/data/linked_list.sn

Package Type: Base Package

Namespacing: The entire package is namespaced in 'List'.

Description:

Provides a data structure that can hold a variable number of items. Each element is capsuled in a own node. Using pointers, these nodes are chained together. The nodes are stored on the heap.

Struct Name	Struct Fields	Description
ListNode<T>	ListNode<T>* next T value	Capsules a single data element. Also contains a pointer to the next list node in the list. Is set to 0 if this is the last element in the list.
LinkedList<T>	ListNode<T>* head ListNode<T>* tail T defValue	Capsules two list node pointers, one pointing to the first, and one to the last element. Both pointers are initialized to <i>null</i> . Also contains a default value, that is returned when attempting to get a value out of bounds.

Function Header	Description
LinkedList<T>* create<T>()	Initializes a new linked list struct. Sets the pointers to 0. $O(1)$
void destroy(LinkedList<void>* l)	Destroys the given list and all of its contained elements by freeing every node and the list itself from the heap. $O(n)$
void add<T>(LinkedList<T>* lp, T x)	Adds given x to the list by encapsulating it in a new list node and appending the element at the end of the list. This uses the tail pointer of the list, so the complexity is $O(1)$.
T get<T>(LinkedList<T>* lp, int i)	Returns the value of the i-th element in the given list or the default value when out of bounds. $O(n)$
ListNode<T>* getNode<T>(LinkedList<T>* lp, int i)	Returns a pointer to the list node in the list at the i-th position in the list. $O(n)$
bool contains<T>(LinkedList<T>* lp, T x)	Checks whether given value is the value of one of the contained list nodes. Returns true if the value is found. $O(n)$
ListNode<T>* find<T>(LinkedList<T>* lp, T x)	Attempts to find the first list node in order in the list that has the given value x. Returns a pointer to this list node. $O(n)$
int size(LinkedList<void>* lp)	Returns the number of nodes in the list. $O(n)$
void remove(LinkedList<void>* lp, int i)	Returns the list node at the i-th position. Updates the pointers of the neighbours to fill the gap. Also frees the removed node. $O(n)$

binary_tree.sn

Full Path: lib/std/data/binary_tree.sn

Package Type: Base Package

Namespacing: The entire package is namespaced in 'Tree'.

Description:

Provides a data structure that can hold a variable number of elements and organizes the elements to minimize search time.

Struct Name	Struct Fields	Description
TreeNode<T>	TreeNode<T>* left TreeNode<T>* right T value	Capsules pointers to left and right child nodes, and a value.

Function Header	Description
TreeNode<T>* create<T>(T value)	Initializes a new Tree. Initializes pointers to 0. Sets given value to node value. Returns pointer to newly created node. $O(1)$
TreeNode<T>* insert<T>(TreeNode<T>* root, T value)	Capsules given value in a new Node and inserts it into given Tree. The insertion location is determined by numerically comparing the value. Returns a pointer to the newly created node. $O(\log n)$

queue.sn

Full Path: lib/std/data/queue.sn

Package Type: Base Package

Namespacing: The entire package is namespaced in 'Queue'.

Description:

Provides a data structure that can hold a fixed number of elements and acts like a bypassing FIFO. The queue is implemented with a circular array.

Struct Name	Struct Fields	Description
CyclicQueue<T>	T* storage T defValue int size int head int tail bool isEmpty	Capsules a pointer to the storage array that holds the contained elements. The field defValue holds the default value that is returned f.E. when a dequeue is attempted while the queue is empty. Also capsules a field that holds the max number of elements. Holds two indices that act as pointers in the array, used to determine where the head and tail of the queue in the array is. Finally, it holds a boolean that stores wether the queue is full or not in the case that the head is equal to the tail.

Function Header	Description
CyclicQueue<T>* create<T>(int size)	Initializes a new Queue with given size. Creates a storage array on the heap and initializes the CyclicQueue struct. Returns a pointer to the created queue. $O(1)$
void destroy(CyclicQueue<void>* queue)	Frees the given queue object and frees the storage array of the queue in the heap. $O(1)$.
bool isEmpty(CyclicQueue<void>* queue)	Returns whether the queue contains any elements or not. $O(1)$
bool isFull(CyclicQueue<void>* queue)	Returns whether the number of contained elements is equal to the maximum of elements. $O(1)$
void enqueue<T>(CyclicQueue<T>* queue, T value)	Adds a new element to the tail of the queue. If the queue is full, the element won't be added. $O(1)$
T dequeue<T>(CyclicQueue<T>* queue)	Removes the head of the queue and returns it. If the queue is empty, the function will return 0. Note that this may cause problems with types larger than 1 word. $O(1)$
void clear(CyclicQueue<void>* queue)	Removes all elements from the queue. This is done by setting the pointers, rather than clearing the actual values. $O(1)$
int size(CyclicQueue<void>* queue)	Returns the number of currently stored elements in the queue. $O(1)$

stack.sn

Full Path: lib/std/data/stack.sn

Package Type: Includes: linked_list.sn

Namespacing: The entire package is namespaced in 'Stack'.

Description:

Provides a data structure that can hold a variable amount of elements and acts like a LIFO buffer.

Struct Name	Struct Fields	Description
StackedList<T>	LinkedList<T>* list	Capsules only a linked list. This list will hold the elements.

Function Header	Description
StackedList<T>* create<T>(T defValue)	Initializes a new Stacked List and the capsuled linked list. Returns a pointer to the newly created stack. $O(1)$
void destroy(StackedList<void>* stack)	Frees the stack and the capsuled linked list and all of the lists contained elements. $O(n)$

<code>void push<T>(StackedList<T>* stack, T value)</code>	Adds given element to the top of the stack. $O(1)$
<code>bool isEmpty(StackedList<void>* stack)</code>	Returns whether no elements are stored in the stack. $O(n)$
<code>int size(StackedList<void>* stack)</code>	Returns the number of contained elements in the stack. $O(n)$
<code>T peek<T>(StackedList<T>* stack)</code>	Returns the top element of the stack, but does not remove it. Returns the default value if the stack is empty. $O(n)$
<code>T pop<T>(StackedList<T>* stack)</code>	Returns the top element of the stack and removes it. Returns the default value if the stack is empty. $O(n)$

Math

math.sn

Full Path: lib/std/math/math.sn
 Package Type: Base Package
 Namespacing: The entire package is namespaced in 'Math'.
 Description:
 Contains math functions and operations.

Function Headers	Description
int pow(int x, int n)	Returns x to the n-th power.
int abs(int x)	Returns the absolute value of x.
int fac(int n)	Returns the n-th faculty.

matrix.sn

Full Path: lib/std/math/matrix.sn
 Package Type: Base Package
 Namespacing: The entire package is namespaced in 'Matrix'.
 Description:
 Contains matrix utility. Currently only matrix multiplication.

Function Headers	Description
int mult(int* m, int* n, int dim0, int dim1)	Multiplies the two given matrices. Stores the result in an integer array on the heap. Returns a pointer to the result. The first matrix should have dim0 columns, the second matrix should have dim1 columns. The resulting matrix will have the dimensions <i>dim0</i> × <i>dim1</i> . Does not check matrix size compatibility.

vector.sn

Full Path: lib/std/math/vector.sn
 Package Type: Base Package
 Namespacing: The entire package is namespaced in 'Vector'.
 Description:
 Contains utility functions for vector operations, like the scalar product.

Function Headers	Description
int scalarProd(int* v, int* w, int l)	Calculates the scalar product of the two given vectors. Both vectors have to have the given length. Returns the result.

Utility

bits.sn

Full Path: lib/util/bits.sn

Package Type: Base Package

Namespacing: The entire package is namespaced in 'Bits'.

Description:

This library can be used to create something like bit-banks. Meaning a single Integer can hold 32 States for an automaton, flag-set etc. These bits can easily be modified with this library.

Function Headers	Description
bool isBitSet(int x, int i)	Checks whether the bit at the i-th place in given word x is a 1. The lowest bit has the index 0.
int setBit(int target, int i, bool val)	Sets the bit in the target word at the i-th place to given val. Returns the resulting data word.
int toggleBit(int target, int i)	Sets the bit in given target word at the i-th place to its complement. Returns the resulting word.
int clearBit(int target, int i)	Sets the bit in given target word at the i-th place to 0. Returns the resulting word.

color.sn

Full Path: lib/util/color.sn

Package Type: Base Package

Namespacing: The entire package is namespaced in 'Colors'.

Description:

This library contains utility related to colours, specifically rgba-coloring.

Struct Name	Struct Fields	Description
Color	int rgba	Encodes the R, G, B and A color channels, by creating a single integer where each byte contains the value in range of 0-255.

Function Headers	Description
Color* create(int r, int g, int b, int a)	Encodes the four given color values and creates a new Color Struct on the heap. Returns a pointer to the struct.
int getRed(Color* c)	Returns the value of the red channel.
int getGreen(Color* c)	Returns the value of the green channel.
int getBlue(Color* c)	Returns the value of the blue channel.
int getAlpha(Color* c)	Returns the value of the alpha channel.