SHARED MEMORY PROGRAMMING WITH OPENMP (2)

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Reference material for part of this presentation :

An Introduction to Parallel Programming by Peter Pacheco

SCHEDULING LOOPS

We want to parallelize this loop.

```
sum = 0.0;
for (i = 0; i <= n; i++)
   sum += f(i);
```

Thread	Iterations
0	$0,\ n/t,\ 2n/t,\ \ldots$
1	$1,\ n/t+1,\ 2n/t+1,\ \ldots$
$\vdots$	$\vdots$
$t-1$	$t-1,\ n/t+t-1,\ 2n/t+t-1,\ \ldots$

Assignment of work
using cyclic partitioning.

```c
double f(int i) {
   int j, start = i*(i+1)/2, finish = start + i;
   double return_val = 0.0;

   for (j = start; j <= finish; j++) {
      return_val += sin(j);
   }
   return return_val;
}  /* f */
```

Our definition of function *f*.

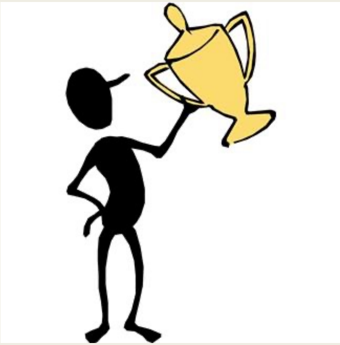
Results

- f(i) calls the sin function *i* times.
- Assume the time to execute f(2i) requires approximately twice as much time as the time to execute f(i).

- n = 10,000
  - *one thread*
  - *run-time = 3.67 seconds.*

Results

- n = 10,000
  - *two threads*
  - *default assignment*
  - *run-time = 2.76 seconds*
  - *speedup = 1.33*
- n = 10,000
  - *two threads*
  - *cyclic assignment*
  - *run-time = 1.84 seconds*
  - *speedup = 1.99*



The Schedule Clause

- Default schedule:

```
        sum = 0.0;
#       pragma omp parallel for num_threads(thread_count) \
           reduction(+:sum)
        for (i = 0; i <= n; i++)
           sum += f(i);
```

- Cyclic schedule:

```
        sum = 0.0;
#       pragma omp parallel for num_threads(thread_count) \
           reduction(+:sum) schedule(static,1)
        for (i = 0; i <= n; i++)
           sum += f(i);
```

schedule (type , chunksize)

- Type can be:
  - *static: the iterations can be assigned to the threads before the loop is executed.*
  - *dynamic or guided: the iterations are assigned to the threads while the loop is executing.*
  - *auto: the compiler and/or the run-time system determine the schedule.*
  - *runtime: the schedule is determined at run-time.*
- The chunksize is a positive integer.

The Static Schedule Type

twelve iterations, 0, 1, . . . , 11, and three threads

```
schedule(static,1)
```

$$\begin{array}{ll} \text{Thread } 0: & 0,3,6,9 \\ \text{Thread } 1: & 1,4,7,10 \\ \text{Thread } 2: & 2,5,8,11 \end{array}$$

The Static Schedule Type

twelve iterations, 0, 1, . . . , 11, and three threads

```
schedule(static,2)
```

Thread 0 :	0, 1, 6, 7
Thread 1 :	2, 3, 8, 9
Thread 2 :	4, 5, 10, 11

The Static Schedule Type

twelve iterations, 0, 1, . . . , 11, and three threads

```
schedule(static,4)
```

Thread 0 :	0, 1, 2, 3
Thread 1 :	4, 5, 6, 7
Thread 2 :	8, 9, 10, 11

The Dynamic Schedule Type

- The iterations are also broken up into chunks of chunksize consecutive iterations.
- Each thread executes a chunk, and when a thread finishes a chunk, it requests another one from the run-time system.
- This continues until all the iterations are completed.
- The chunksize can be omitted. When it is omitted, a chunksize of 1 is used.

The Guided Schedule Type

- Each thread also executes a chunk, and when a thread finishes a chunk, it requests another one.
- However, in a guided schedule, as chunks are completed the size of the new chunks decreases.
- If no chunksize is specified, the size of the chunks decreases down to 1.
- If chunksize is specified, it decreases down to chunksize, with the exception that the very last chunk can be smaller than chunksize.

Thread	Chunk	Size of Chunk	Remaining Iterations
0	1 – 5000	5000	4999
1	5001 – 7500	2500	2499
1	7501 – 8750	1250	1249
1	8751 – 9375	625	624
0	9376 – 9687	312	312
1	9688 – 9843	156	156
0	9844 – 9921	78	78
1	9922 – 9960	39	39
1	9961 – 9980	20	19
1	9981 – 9990	10	9
1	9991 – 9995	5	4
0	9996 – 9997	2	2
1	9998 – 9998	1	1
0	9999 – 9999	1	0

Assignment of trapezoidal rule iterations 1–9999 using a guided schedule with two threads.

The Runtime Schedule Type

- The system uses the environment variable OMP_SCHEDULE to determine at run-time how to schedule the loop.
- The OMP_SCHEDULE environment variable can take on any of the values that can be used for a static, dynamic, or guided schedule.

PRODUCERS AND CONSUMERS

Queues

- Can be viewed as an abstraction of a line of customers waiting to pay for their groceries in a supermarket.
- A natural data structure to use in many multithreaded applications.
- For example, suppose we have several “producer” threads and several “consumer” threads.
  - *Producer threads might “produce” requests for data.*
  - *Consumer threads might “consume” the request by finding or generating the requested data.*

Message-Passing

- Each thread could have a shared message queue, and when one thread wants to “send a message” to another thread, it could enqueue the message in the destination thread’s queue.
- A thread could receive a message by dequeuing the message at the head of its message queue.

Message-Passing

```
for (sent_msgs = 0; sent_msgs < send_max; sent_msgs++) {
   Send_msg();
   Try_receive();
}

while (!Done())
   Try_receive();
```

Sending Messages

```
    mesg = random();
    dest = random() % thread_count;
#   pragma omp critical
    Enqueue(queue, dest, my_rank, mesg);
```

Receiving Messages

```
if (queue_size == 0) return;
else if (queue_size == 1)
#    pragma omp critical
        Dequeue(queue, &src, &mesg);
else
        Dequeue(queue, &src, &mesg);
Print_message(src, mesg);
```

Termination Detection

```
queue_size = enqueued - dequeued;
if (queue_size == 0 && done_sending == thread_count)
   return TRUE;
else
   return FALSE;
```

each thread increments this after completing its for loop

Startup (1)

- When the program begins execution, a single thread, the master thread, will get command line arguments and allocate an array of message queues: one for each thread.
- This array needs to be shared among the threads, since any thread can send to any other thread, and hence any thread can enqueue a message in any of the queues.

Startup (2)

- One or more threads may finish allocating their queues before some other threads.
- We need an explicit barrier so that when a thread encounters the barrier, it blocks until all the threads in the team have reached the barrier.
- After all the threads have reached the barrier all the threads in the team can proceed.

```
# pragma omp barrier
```

The Atomic Directive (1)

- Unlike the critical directive, it can only protect critical sections that consist of a single C assignment statement.

```
# pragma omp atomic
```

- Further, the statement must have one of the following forms:

```
x <op>= <expression>;
x++;
++x;
x--;
--x;
```

The Atomic Directive (2)

- Here <op> can be one of the binary operators

  `+, *, -, /, &, ^, |, <<, or >>`

- Many processors provide a special load-modify-store instruction.
- A critical section that only does a load-modify-store can be protected much more efficiently by using this special instruction rather than the constructs that are used to protect more general critical sections.

Critical Sections

- OpenMP provides the option of adding a name to a critical directive:

```
# pragma omp critical(name)
```

- When we do this, two blocks protected with critical directives with different names can be executed simultaneously.
- However, the names are set during compilation, and we want a different critical section for each thread's queue.

Locks

- A lock consists of a data structure and functions that allow the programmer to explicitly enforce mutual exclusion in a critical section.

Locks

```
/* Executed by one thread */
Initialize the lock data structure;
. . .
/* Executed by multiple threads */
Attempt to lock or set the lock data structure;
Critical section;
Unlock or unset the lock data structure;
. . .
/* Executed by one thread */
Destroy the lock data structure;
```

Using Locks in the Message-Passing Program

```
#   pragma omp critical
    /* q_p = msg_queues[dest] */
    Enqueue(q_p, my_rank, mesg);
```

```
/* q_p = msg_queues[dest] */
omp_set_lock(&q_p->lock);
Enqueue(q_p, my_rank, mesg);
omp_unset_lock(&q_p->lock);
```

Using Locks in the Message-Passing Program

```
~~# pragma omp critical~~
~~/* q_p = msg_queues[my_rank] */~~
~~Dequeue(q_p, &src, &mesg);~~
```

```
/* q_p = msg_queues[my_rank] */
omp_set_lock(&q_p->lock);
Dequeue(q_p, &src, &mesg);
omp_unset_lock(&q_p->lock);
```

Some Caveats

1. You shouldn’t mix the different types of mutual exclusion for a single critical section.
2. There is no guarantee of fairness in mutual exclusion constructs.
3. It can be dangerous to “nest” mutual exclusion constructs.

```
#    pragma omp atomic
   x += f(y);
#    pragma omp critical
   x = g(x);
```

```
   while(1) {
      …
#        pragma omp critical
      x = g(my_rank);
      …
   }
```

```
#    pragma omp critical
   y = f(x);
   …
   double f( double x) {
#        pragma omp critical
      z = g(x);    /* z is shared */
      …
   }
```

Matrix-vector multiplication

$$y_i = a_{i0}x_0 + a_{i1}x_1 + \cdots + a_{i,n-1}x_{n-1}$$

$$\begin{array}{|c|c|c|c|}
\hline
a_{00} & a_{01} & \cdots & a_{0,n-1} \\ \hline
a_{10} & a_{11} & \cdots & a_{1,n-1} \\ \hline
\vdots & \vdots & & \vdots \\ \hline
a_{i0} & a_{i1} & \cdots & a_{i,n-1} \\ \hline
\vdots & \vdots & & \vdots \\ \hline
a_{m-1,0} & a_{m-1,1} & \cdots & a_{m-1,n-1} \\ \hline
\end{array}
\begin{array}{|c|}
\hline
x_0 \\ \hline
x_1 \\ \hline
\vdots \\ \hline
x_{n-1} \\ \hline
\end{array}
=
\begin{array}{|c|}
\hline
y_0 \\ \hline
y_1 \\ \hline
\vdots \\ \hline
y_i = a_{i0}x_0 + a_{i1}x_1 + \cdots a_{i,n-1}x_{n-1} \\ \hline
\vdots \\ \hline
y_{m-1} \\ \hline
\end{array}$$

```
for (i = 0; i < m; i++) {
   y[i] = 0.0;
   for (j = 0; j < n; j++)
      y[i] += A[i][j]*x[j];
}
```

Matrix-vector multiplication

```
#  pragma omp parallel for num_threads(thread_count) \
      default(none) private(i, j) shared(A, x, y, m, n)
   for (i = 0; i < m; i++) {
      y[i] = 0.0;
      for (j = 0; j < n; j++)
         y[i] += A[i][j]*x[j];
   }
```

Run-times and efficiencies of matrix-vector multiplication (times are in seconds)

Threads	Matrix Dimension					
	$8,000,000 \times 8$		8000×8000		$8 \times 8,000,000$	
	Time	Eff.	Time	Eff.	Time	Eff.
1	0.322	1.000	0.264	1.000	0.333	1.000
2	0.219	0.735	0.189	0.698	0.300	0.555
4	0.141	0.571	0.119	0.555	0.303	0.275

Why does efficiency decrease in certain configurations more than others?

Tokenization

Thread-Safety

```
void Tokenize(
      char*   lines[]        /* in/out */,
      int     line_count     /* in     */,
      int     thread_count   /* in     */) {
   int my_rank, i, j;
   char *my_token;

#  pragma omp parallel num_threads(thread_count) \
      default(none) private(my_rank, i, j, my_token) \
      shared(lines, line_count)
   {
      my_rank = omp_get_thread_num();
#     pragma omp for schedule(static, 1)
      for (i = 0; i < line_count; i++) {
         printf("Thread %d > line %d = %s", my_rank, i, lines[i]);
         j = 0;
         my_token = strtok(lines[i], " \t\n");
         while ( my_token != NULL ) {
            printf("Thread %d > token %d = %s\n", my_rank, j, my_token);
            my_token = strtok(NULL, " \t\n");
            j++;
         }
      } /* for i */
   }  /* omp parallel */

}  /* Tokenize */
```

Concluding Remarks (1)

- OpenMP is a standard for programming shared-memory systems.
- OpenMP uses both special functions and preprocessor directives called pragmas.
- OpenMP programs start multiple threads rather than multiple processes.
- Many OpenMP directives can be modified by clauses.

Concluding Remarks (2)

- A major problem in the development of shared memory programs is the possibility of race conditions.
- OpenMP provides several mechanisms for insuring mutual exclusion in critical sections.
  - *Critical directives*
  - *Named critical directives*
  - *Atomic directives*
  - *Simple locks*

Concluding Remarks (3)

- By default most systems use a block-partitioning of the iterations in a parallelized for loop.
- OpenMP offers a variety of scheduling options.
- In OpenMP the scope of a variable is the collection of threads to which the variable is accessible.

Concluding Remarks (4)

- A reduction is a computation that repeatedly applies the same reduction operator to a sequence of operands in order to get a single result.