HOCHSCHULE FÜR TECHNIK STUTTGART

SOMMERSEMESTER 2017

PRÜFUNGSLEISTUNG im		STUDIENGANG	S oftware T echnology			
MODUL:	CPL	FACH:	Concepts	Concepts of Programming Langauges 1) Lisp, C, Assembler		
DATUM:	10.07.2017	NAME:		3) Assembler, C, C++ 4) Python, Prolog, Lisp,		
ZEIT:	08:30-10:30	SEMESTER:	ST1			
PRÜFER:	Prof. Dr. Ulrike Padó, Prof. Dr. Peter Heusc	h 				
MODUL-TEILLEISTUNGEN		TEILPUNKTE	0	SESAMTPUNKTE	NOTE	
FACH 1:	CPL		_			
HILFSMITTEL: Orignal Slides + Notes + Seminar Papers						
ANLAGEN:						

Please read the following general rules before starting the exam:

- Put your name on every sheet that you use for noting your exam, but do not put your name on otherwise empty sheets
- Do not write final answers on concept paper, do not write final answers with pencils, do not use your own paper
- Write legibly, leave at least 5 mm of empty space between your lines and at least 40 mm of empty space at the right margin
- Put the solutions for questions 5 and 6 on a separate sheet of paper. Failing to do so will result in 50% deduction of points for these two questions.
- The points for every question reflect the amount of minutes you need to answer it.

1. Question (30P)

Answer the following questions concerning the languages from the seminar. Every question has 3 points. Some questions contain a limitation of results after the question. Failing to comply (and giving more names) results in fewer points! Also denote in one sentence why the language of your choice belongs into that list.



- 1) Which of the languages are macro languages or contain macro features (e.g. C or C++ are no macro languages, but contain macro elements). Give three names.
- 2) Contrary to most other programming languages FORTH does not allow you to use the name of the word being defined inside the definition to call the word recursively, but you must use the predefined word recurse. What does this tell you about the compilation process in FORTH?
- 3) Which languages from the seminar have been used to implement operating systems? Give three names.
- 4) Which languages would you use to solve problems of artificial intelligence? Give three names.
- 5) Which languages have a large number of keywords, which languages do only have few? Name two languages for every case.
- 6) Which language is the oldest, which is the newest one of all the languages of the seminar? Name one language for every case.
- 7) Describe typical features of interpreted and compiled languages. Give two names for every category.
- 8) Which languages are specialized on web environments? Give three names.

- 9) Which languages were designed for commercial programs? Give three names.
- 10)Describe in 80 to 100 words which language you liked most and why.

2. Question (20P)

A lot of modern programming languages support type inference. E.g. Java 8 allows you to write

IntStream.range(0,100).forEach((x) -> {System.out.printf("%d%n", x&5); }) while it does not allow you to write

IntStream.range(0,100).asDoubleStream().forEach((x) ->

{System.out.printf("%d%n", x&5); }),

because the binary operator & is not defined for values of type double. However, this practice has become mainstream only during the last ten years. Describe the difference between this way of type inference and weakly typed languages.

A language is stringly typed if type errors are aways detected.

It is weakly typed if types are only loosely checked. Java ist almost stongly typed with a few exceptions.

The example above does not work because the type of x is checked at compile time and Look at the following grammar:

The type x is not specificly typed but javas type inference automaticly assings a data types

The second rule contains a left recusrive production. Therefore it is not LL(1) parsable.

By reformulating the grammar as a tail recusring we can make it LL(1) parsable.

based on the context.

 $S'-> aS'b \mid bS'a \mid S'S' \mid \varepsilon$

Describe the words that this grammar produces and show that this grammar is not LL(1)parsable. Define an equivalent grammar that is LL(1)-parseable.

4. Question (15P)

3. Question (20P)

S -> S'\$

Proof by creating first set h FIRST(S') -> ... FIRST(S') = FIRST(S') = FIRST(S')

Look at the following context free grammar:

 $E \rightarrow E + T \mid T$ S' -> aS'bS" | bS'aS" S" -> S' | epsilon $T \rightarrow T * F | F$

alpha = null $F \rightarrow '(' E ')' | < num >$ beta = aS'b and bS'a

Construct the abstract and the concrete syntax tree for the input (2 + 3) * 4

5. Question (15P)

- a) What is a recursive function, and why is tail recursion important? (3P)
- b) Explain how we can convert recursive functions into tail-recursive functions. Use the following Scheme code in your explanation (you do not need to produce Scheme code of your own unless you wish to do so) (6P):
 - a. (define sum 1. (lambda (n) (if (= n 1) 1)ii. (+ n (sum (- n 1)))))
- c) We have seen functional elements in many languages that are not functional in basic design. Why do you think this is? Please explain your opinion. (6P)



ab ba abba baab

6. Question (15)

- a) Prolog and Scheme are very close to two mathematical concepts. Why is this, and why were those concepts chosen? (8P)
- b) You are looking for a pet. It should be small, harmless and furry. Give the search tree that Prolog follows when asked to find all possible suggestions for a good pet from its knowledge base. Be sure to show the goal stacks and unification steps as well as success or failure of a search branch. (7P)

Knowledge Base:

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pet(X):- small(X), harmless(X), furry(X).
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small(fly).

small(tarantula).

small(guineaPig).

harmless(fly).

harmless(guineaPig).

furry(tarantula).
furry(guineaPig).

- Uses expressions (evaluate to a value) instead of statements (alter machine state)

- Uses recursion instead of iterative loops

- Built around the concept of functional

application (from λ calculus)

????

pet(X)

