Message Passing Interface

Week 10

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Welcome to MPI Message Passing Interface

- First developed in 1994.
- A genereral framework for parallel computing
 - Distributed systems—co-ordinates across nodes
 - Shared-memory systems within each node
- Defines a standard interface, with various implementations
 - We use OpenMPI
- Provides a convenient abstraction for parallelism
 - Assumes that you must explicitly send and receive data from one process to another
 - Within a single process, threads can share data.

MPI Basics Hello, World

- Every MPI program has
 - At least one communicator to pass mesages among processes
 - A rank for each process within a communicator.
- The MPI program is sandwiched between Init and Finalize calls.
- Compile with MPI headers and libraries
 - For convenience, use the OpenMPI mpic++ wrapper compiler.

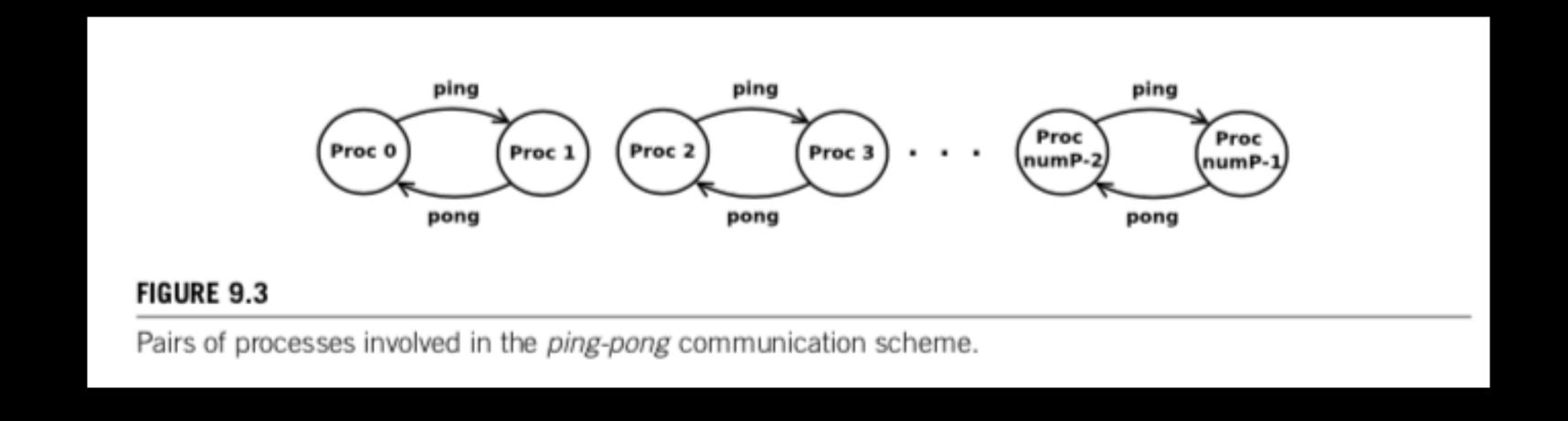
ExecutionDistributed System Management

- You run the MPI program with mpirun.
- This causes your master node to send copies of the executable to the other nodes in your distributed system.
- It launches execution on the other nodes.
- The output is routed back to the master node.

Data Transfer Send and Receive

- Data must be explicitly transferred from process to process.
- void Send(const void * data, int count, const MPI::Datatype& data_type, int destination_id, int tag);
- void Recv(const void * data, int count, const MPI::Datatype& data_type, int source_id, int tag);
- Note that data are assumed to uniform in type!
 - Scalar or array.
- The receiver must "know" how many objects to expect and whence to expect them.
- The **tag**s must agree.
- This kind of communication is blocking.
 - The sender has to wait until its buffer is free for re-use.
 - The receiver has to wait to receive the message from the sender.

Ping Pong Pairwise Blocking Communication



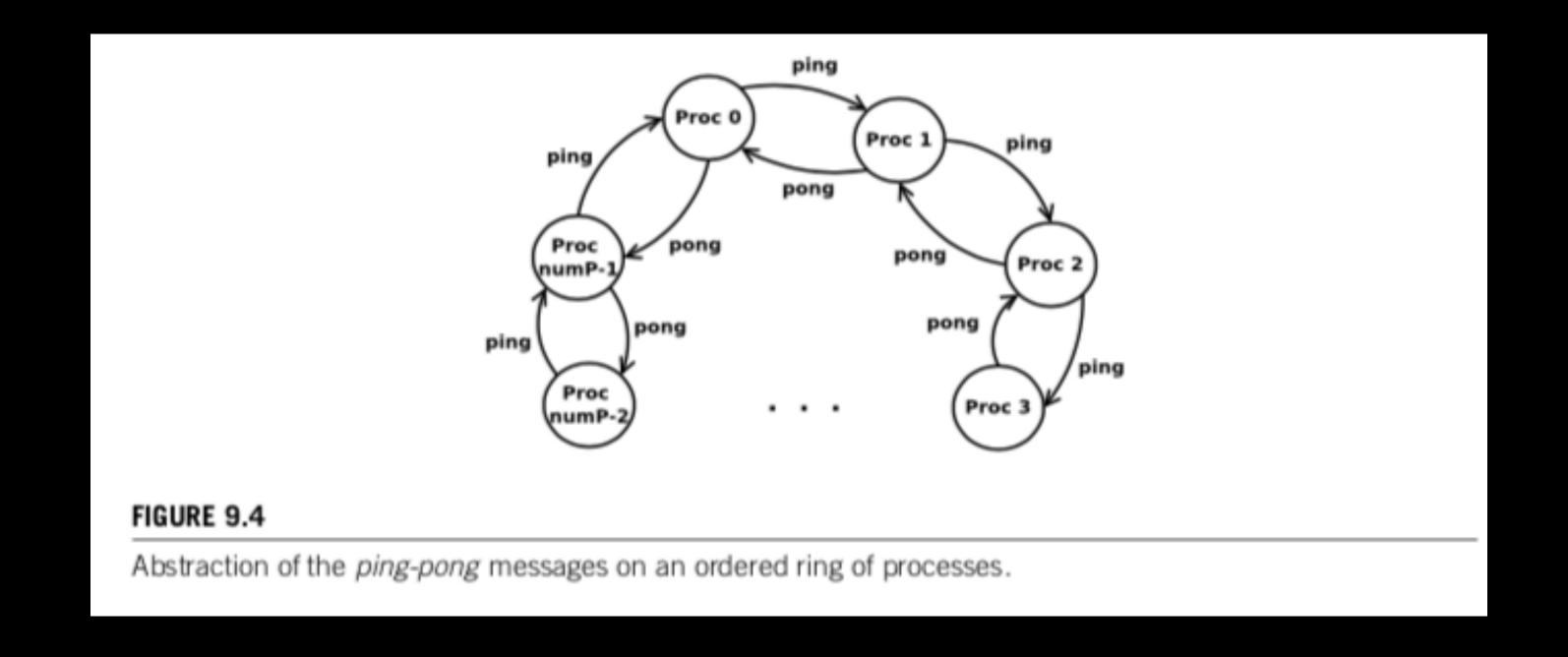
What If You Ping in a Ring?

The Problem
With Blocking
Communication

```
void ping_pong(int num_ping_pongs, int id) {
  int ping_pong_count = 0;
  int next_id = id + 1, prev_id = id - 1;
  if(next_id >= num_p) {
    next_id = 0;
  if(prev_id < 0) {
    prev_id = num_p - 1;
  while(ping_pong_count < num_ping_pongs) {
    ping_pong_count++;
    // Send the ping.
    MPI::COMM_WORLD.Send(&ping_pong_count, 1, MPI::INT, next_id, 0);
    // Wait and receive the ping.
    MPI::COMM_WORLD.Recv(&ping_pong_count, 1, MPI::INT, prev_id, 0);
    // Send the pong
    MPI::COMM_WORLD.Send(&ping_pong_count, 1, MPI::INT, prev_id, 0);
    // Wait and receive the pong
     MPI::COMM_WORLD.Recv(&ping_pong_count, 1, MPI::INT, next_id, 0);
```

What If You Ping in a Ring?

The Problem of Blocking Communication



DeadlockThe Need for Non-Blocking Communication

- Each thread sends to id + 1 and then receives from id 1.
- Thread id's send may not return until thread id + 1's receive completes.
- Each thread never gets to receive because it is waiting for send to return.
- No send can return because each receive is waiting.
- This only actually happens when the data load is big enough.

Non-Blocking Communication

Safe Against Deadlock

- MPI::Request Isend(const void * data, int count, const MPI::Datatype& data_type, int destination_id, int tag);
- MPI::Request Irecv(const void * data, int count, const MPI::Datatype& data_type, int source_id, int tag);
- These functions return immediately.
- To synchronize, i.e., wait until data are available, use Wait or Test.

Deadlock-Safe Ping Pong

Using Non-Blocking Communication

```
MPI::Request rq_receive;
while (ping_pong_count < num_ping_pongs) {
  ping_pong_count++;
  // Send the ping.
  MPI::COMM_WORLD.Isend(&ping_pong_count, 1, MPI::INT,next_id, 0);
  // Wait and receive the ping.
  rq_receive = MPI::COMM_WORLD.Irecv(&ping_pong_count, 1, MPI::INT, prev_id, 0);
  rq_receive.Wait();
  // Send the pong
  MPI::COMM_WORLD.Isend(&ping_pong_count, 1, MPI::INT, prev_id, 0);
  // Wait and receive the pong
  rq_receive = MPI::COMM_WORLD.Irecv(&ping_pong_count, 1, MPI::INT, next_id, 0);
  rq_receive.Wait();
```