

5 Leetcode256房子涂色付钱问题

笔记本: DP Note

创建时间: 10/19/2019 11:54 PM

更新时间: 10/23/2019 10:18 PM

作者: tanziqi1756@outlook.com

256. Paint House

难度 简单 12 收藏 分享 切换为中文 关注

题目描述

评论 (18)

题解 (14) ^{New}

提交记录

There are a row of n houses, each house can be painted with one of the three colors: red, blue or green. The cost of painting each house with a certain color is different. You have to paint all the houses such that no two adjacent houses have the same color.

The cost of painting each house with a certain color is represented by a $n \times 3$ cost matrix. For example, `costs[0][0]` is the cost of painting house 0 with color red; `costs[1][2]` is the cost of painting house 1 with color green, and so on... Find the minimum cost to paint all houses.

Note:

All costs are positive integers.

Example:

Input: `[[17,2,17],[16,16,5],[14,3,19]]`

Output: 10

Explanation: Paint house 0 into blue, paint house 1 into green, paint house 2 into blue.
Minimum cost: $2 + 5 + 3 = 10$.

如果直接套用之前机器人走路思路，可以做到 $O(n^3)$

因为这里只有三种颜色，可以做到 $O(n)$

如果是 N 种颜色，就只能是 $O(n^3)$

256. Paint House

难度 简单 12 收藏 分享 切换为中文 关注

题目描述

评论 (18)

题解 (14) ^{New}

提交记录

执行结果: 通过 显示详情

执行用时: 1 ms, 在所有 java 提交中击败了 100.00% 的用户

内存消耗: 39.9 MB, 在所有 java 提交中击败了 100.00% 的用户

还是用机器人走路的二维dp

这里可以in-place, 即使用原来的costs作为dp矩阵

输入[

[17,2,17],
[16,16,5],
[14,3,19]
]

第一行[17, 2, 17]为dp初始状态

第二行开始进行状态转移方程： $dp[i][j] = \text{Min}(dp[i-1]) + dp[i][j]$

含义：在对第i座房子，涂第j种颜色时，第i-1座房子就不能涂成第j种颜色，

因而，就找第i-1行的最小值（最小花费），
加上当前对第i座房子，涂第j种颜色时的花费。

最终的dp矩阵

17 2 17

18 33 7

21 10 37

遍历最后一行，得到最小值10

```
1  class Solution {
2  public int minCost(int[][] costs) {
3      if( costs.length == 0 ) {
4          return 0;
5      }
6      for( int i = 1; i < costs.length; i++ ) {
7          for( int j = 0; j < costs[0].length; j++ ) {
8              // For every row[i], find the minimum of row[i-1], excluding the element costs[i-1][j]
9              int minCost = findMin(costs[i-1], j);
10             costs[i][j] += minCost;
11         }
12     }
13
14     int ans = Integer.MAX_VALUE;
15     int N = costs.length - 1;
16     int colors = costs[0].length;
17     for( int j = 0; j < colors; j++ ) {
18         if( costs[N][j] < ans ) {
19             ans = costs[N][j];
20         }
21     }
22     return ans;
23 }
24
25 private int findMin(int[] cost, int index) {
26     int minCost = Integer.MAX_VALUE;
27     for( int i = 0; i < cost.length; i++ ) {
28         if( i != index && cost[i] < minCost ) {
29             minCost = cost[i];
30         }
31     }
32     return minCost;
33 }
34 }
35 }
```

Leetcode 265 房子涂颜色升级版+时间优化

每次都要算除了一个数，其他元素中的最小元素，

那我只要直接记录最小元素和次小元素就行了。

```

1  class Solution {
2      public int minCostII(int[][] costs) {
3          if (costs.length == 0) {
4              return 0;
5          } else if (costs[0].length == 1) {
6              return costs[0][0];
7          }
8          int minColour = -1;
9          int minCost = 0;
10         int secondMinCost = 0;
11         for (int[] cost : costs) {
12             int tmpMinColour = -1;
13             int tmpMinCost = Integer.MAX_VALUE;
14             int tmpSecondMinCost = Integer.MAX_VALUE;
15             for (int i = 0; i < cost.length; i++) {
16                 int thisMinCost = cost[i] + (i == minColour ? secondMinCost : minCost);
17                 if (thisMinCost < tmpMinCost) {
18                     tmpSecondMinCost = tmpMinCost;
19                     tmpMinCost = thisMinCost;
20                     tmpMinColour = i;
21                 } else if (thisMinCost < tmpSecondMinCost) {
22                     tmpSecondMinCost = thisMinCost;
23                 }
24             }
25             minCost = tmpMinCost;
26             minColour = tmpMinColour;
27             secondMinCost = tmpSecondMinCost;
28         }
29         return minCost;
30     }
31 }

```