Flow Control

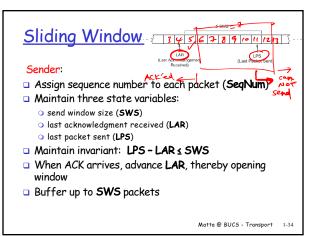
- Goal: control the flow of packets on the link so that receiver always has sufficient buffers to accept them until they can be processed
- □ Sliding Window:
 - Imposes a limit on the number of outstanding unACKed packets, i.e. length of retransmission list, called send window
 - For stop-and-wait, send window = 1 → poor link utilization
 - The size of the send window is chosen to achieve both high link utilization and flow control
 - Send Window Size =

 MIN(delay x bandwidth, available buffer space at receiver)

 Min (delay x bandwidth, available buffer space at receiver)

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Sliding Window (cont'd) Receiver: Maintain three state variables: receive window size (RWS) last packet acceptable (LPA) next packet expected (NPE) Maintain invariant: LPA - NPE + 1 ≤ RWS Packet SeqNum arrives: if NPE ≤ SeqNum ≤ LPA → accept if SeqNum < NPE or SeqNum > LPA → discarded Send ACK/NAK

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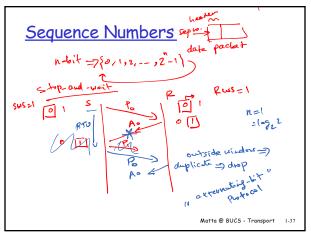
Sliding Window (cont'd)

- □ With Go-Back-N, RWS = 1
- □ With Selective Repeat, RWS = SWS.

 Receiver can then maintain sequence numbers of packets that the sender can send, and so can detect whether a received packet is new or duplicate

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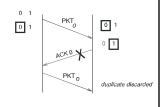
36



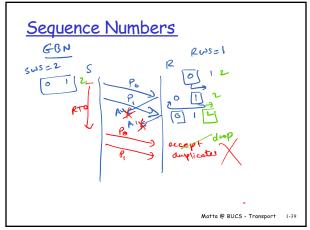
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Sequence Numbers

- SeqNum field is finite; sequence numbers wrap around
- The size of the sequence number space must be larger than the number of outstanding packets
- Stop-and-Wait: sequence numbers {0, 1}



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Sequence Numbers (cont'd)

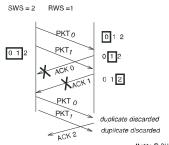
□ Go-Back-N: sequence numbers {0, 1, ..., ??}

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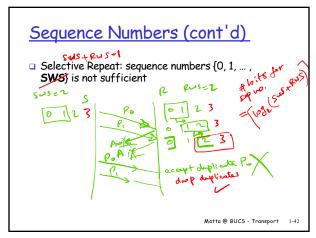
40

Sequence Numbers (cont'd)

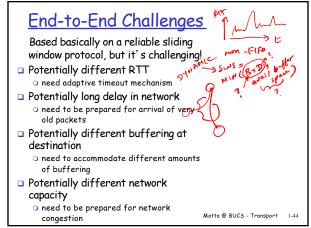
□ Go-Back-N: sequence numbers {0, 1, ..., SWS}

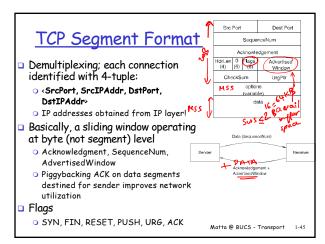


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Selective Repeat: sequence numbers {0, 1, ... , SWS} is not sufficient Size of sequence number space must be at least SWS+RWS = 2 SWS Intuitively, SeqNum ``slides'' between two halves of sequence number space number space Matta @ BUCS - Transport 1-43





TCP Reliability & Flow Control Like SR with Explicit Rtx/cumulative ACKs: storing out-of-order bytes using one timer for all unacked bytes using duplicate ACK to fast retransmit On retransmission, only one segment retransmitted ACK (SBO) ACK (SBO) ACK (SBO) ACK (SBO) Matta @ BUCS - Transport 1-46

TCP Reliability & Flow Control Like SR with Explicit Rtx/cumulative ACKs: storing out-of-order bytes using one timer for all unacked bytes using duplicate ACK to fast retransmit On retransmission, only one segment retransmitted A new version, with SACK option, is more like GBN with selective repeats! At sender: LastByteSent - LastByteAcked & AdvertizedWindow If zero, sender keeps sending 1-byte data segments