Modelling for Combinatorial Optimisation (Course 1DL451 = Part 1 of Course 1DL441) Uppsala University – Autumn 2018 Project Report

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27th November 2018

A Alternative Model

As the discussion in the presentation session on 20 September and the meeting on 17 October 2018, the bin_packing_load constraint could be used as an alternative model. The bin_packing_load(array[int] of var int: load, array[int] of var int: bin, array[int] of int: w) constraint requires that each item i with weight w[i] be put into bin bin[i] such that the sum of the weights of items in each bin b is equal to load[b]. In this problem, with the view point of video serving network, capacity load[i] must be no greater than given capacity X of each cache server. The weights of each item, w[i], corresponds to the videos size reqVid[i]. Each cache server is corresponding to one bin, so C cache servers corresponds to C bins. While the videos that are not requested or exceed the capacity of cache servers will be stored in the data center. In the section Task of the problem description at page 1 requires to decide which videos to put in which cache server, and at page 4 of File format section saying that "it is not necessary to describe all cache servers: ...", so, it's not necessary to consider data center as a cache server in the scrope of the Streaming Videos problem.

A.1 Description

The bin_packing_load model includes constraints that consider the caches as bins, with maximum capacity and loading capacity. The bin_packing_load MiniZinc model could be found in Listing 1. The videos that are stored in data center are implicitly captured by parameter reqVid.

but this unquoted

Listing 1: A bin_packing_load MiniZinc model for the Streaming Videos problem include "globals.mzn";

```
3 int: V;% the number of videos
4 int: E;% the number of endpoints
5 int: R;% the number of request descriptions
6 int: C;% the number of cache servers
7 int: X;% the capacity of each cache server in MB

9 set of int: VID = 1..V; % index range for videos
10 set of int: ENDPOINT = 1..E; % index range for endpoint S
```

neu

```
> not used anywhere?
                                  \stackrel{/}{=} 1...C; % index range for cache serverS
        11 set of int: CACHE
                                 = 0..C; % index range for load of bin
        12 set of int: (CACHEO)
                                 = 0..X; % index range for capacity of caches
        13 set of int: CAPO
                                                 > Value
                                                                    Lodds
           server
                                 = 0..1000; % index range for videos
        14 set of int: SIZEO
                                                                sizes ot
        16 % videoSize[v] = the size of video v
        17 array[VID] of SIZE∅: videoSize;
        19 % Ld: data center/latency, K: number of connected caches
        20 enum LDK = \{Ld, K\};
        22 % endpoint[e, {Ld, K}] = the latency Ld, number of connected caches
                                                      > max (4000, 1000), to be precise
        23 array[ENDPOINT, LDK] of 0.,4000: endpoint;
etter Pradabilit
        % Rv: requested video, Re: coming from endpoint
        26 % Rn: number of requests
        27 enum RVEN = \{Rv, Re, Rn\};
                                                                              your data tiles use of instead of the latency
        29 % request[r, {Rv,Re,Rn}] = requested video Rv, coming from endpoint
        30 % number of requests Rn of request r
        31 array[RREQ, RVEN] of int: request;
        33 % number of distinct requestS
        34 int: realReq;
                                same concept?
                                                                                             emors: Ex
        _{36} % index range for real requests
        37 set of int: RREQ = 1..realReq;
                                                                                convention it
        39 % eConCache[e, c] = latency of endpoint e and cache c
                                                                               e is not connected
        40 array[ENDPOINT, CACHE] of 0. (1000; eConCache;
                                                                       > do we always have
        42 % total number of requests
        43 int: nReq = sum(re in RREQ) (request[ re, Rn ]);
        44 number of
                                          > so are they all valid?
        45 % valid requested videos
       1 46 int: (nReqVid;
        48 % reqVid[v] = size of valid requested videqs
        49 array[VID] of SIZEO: reqVid;
        51 % videoInCap[v] = size of videos that less than or equal to max
           cache capacity
                                                                           > cadable?
       52 array[VID] of SIZEO: videoInCap;
                    > so are they all cachable?
        54 % bin[v] = cacheli-th that stores video v
        55 array[VID] of var CACHE: bin;
                     there seems to be an assumption that non-cachable from requested predicts videos have size of 2 same for videos only requested from enapoints not connected to caches?
```

```
57 % load[c] = load of cache cath
  array[CACHE] of var CAPO: load;
59
  % total saving time
60
  var int: savingTime = sum(req in RREQ where request[req, Rn] > 0)(
     let { int: rv = request[req, Rv];
            int: re = request[req, Re];
            int: rn = request[req, Rn];
64
            int: ld = endpoint[re, Ld];
65
            var int: lc = eConCache[re, bin[request[req, Rv]]];
66
67
         (1d - 1c) was video In Cap in v3: probably the source for ts discrepancy in v1

(bool2int (reqvid[rv] > 0))

objective value with v1
69
70
71
         rn
72
  ));
73
  % find out which videos to assign to which bins very procedural explanation constraint bin_packing load/load bin
74
77
  solve :: int_search(
78
                               S$ in CACHE?
           bin ++
79
           load
80
81
           first_fail, indomain_min, complete) maximize savingTime;
82
83
  output [show((savingTime/nReq)*1000) ++ "\n" ++
84
           show(sum(vi in VID)(bool2int(load(vi)] > 0)))] ++
85
           [ if vi = 1 then \sqrt[n]{b} " else "" endif ++
             if (fix(feqVid[vi]) > 0 / fix(bin[vi]) = b)
                             To how could it not be fixed it's a parameter, no.
            then "\(vi) "
             else "" endif
89
             | b in CACHE, vi in VID];
90
91
    for test server
var int: nReq; nReq = sum(re in RREQ)(request[ re, Rn ]);
95 % output ["\nbin = " ++ show(bin) ++ "\nload = " ++ show(load)];
```

Search Annotations. In the alternative bin_packing_load model for Streaming Videos problem, the ::int_search annotation is used to compute the final score with array of variables, which concatenate the bin array and load array. The next argument first_fail specifies that the variables are chosen in the order that appear. To those chosen variables, the assignment annotation indomain_min will assign the smallest video size in the bin and load domain.

Ultimately, the strategy annotation complete is specified.

load is not a variable, but an array of variables:

Cand array reg Vid

A.2 Implementation

In this model, the load of each cache is computed and guaranteed that its load can not exceed the cache capacity. Ultimately, search strategy is included to find the maximal score. An array of decision variables array [CACHE] of var CAPO: load; is declared to capture the capacity of the cache servers.

The variables bin's index is ranged over the VID to indicate that which videos are stored in which bin. Similarly, the variable load's index is also ranged over CACHE to capture the capacity of all cache. The parameters reqVid and videoInCap can be used alternatively without affecting the final results, but fewer variables usually means faster solving. Since the videoInCap captures all the videos that is no greater than the given capacity of cache servers. The unrequested videos are also stored in the videoInCap in this case. While the reqVid is stricter since it only stores the requested videos that are not greater than the given capacity of cache servers. Those parameters can be alternatively used because just the valid requested videos are used to compute the final scores. In other word, unrequested videos in videoInCap doesn't affect to final scores.

The prepocessing has been done in the following ways. First, the duplicated requested videos are aggregated and corresponding real requests are captured by parameter realReq. Second, the valid requested videos are preprocessed and stored in array videoInCap.

of the same endpoint?

A.3 Bin Packing Model Evaluation

We have chosen the backends for Gecode, Chuffed, Gurobi, OscaR.cbls, and Lingeling. Table 2 gives the results for various instances at Table 1 on the Streaming Videos model[1, 2]. The time-out was 900000 milliseconds. To the instance me_at_the_zoo and warm_up, the backend fzn-oscar-cbls, picat-sat, Gurobi, Chuffed, and Gecode give the results in details at table 2. When testing with much bigger instances such as trending_today, and video_worth_spreading, all backends couldn't produce the final results after 900000 milliseconds. The instance kittens is the biggest and toughest instance that defeats all the backends, and ends up with the time-out.

Name	Videos	Cache Servers	Endpoints	Requests
warm_up	5	3	2	4
me_at_the_zoo	100	10	10	81
video_worth_spreading	10000	100	100	40317
trending_today	10000	100	100	95180
kittens	10000	500	1000	197987

Table 1: Instances of Streaming Videos model.

the Marken unstated assumption at the Wine Hor of page 3

This stems and/or water comment on top of page 3

Technology LCG CPMIP **CBLS** SAT Solver Gecode Chuffed Gurobi OscaR.cbls Lingeling Backend Chuffed Gecode Gurobi fzn-oscar-cbls Picat-sat instance score time score time score time score time score time 562500 562500 562500 warm_up 443 2428 562500 677 562500 734 me_at_the_zoo 607330 607330 607330 518 607330 2356 736 video_worth_spreading

Table 2: Results for the alternative bin packing model with MiniZinc 2.2/1

see v3: no explanation is provided here

Upon examining the produced solution, _______
it terms out that some videos are cached at caches not accessible to the endpoints where their reguests originale from:
the results can be super-optimal.

trending_today

kittens

these runtimes are basically as good (and sometimes) better!?) than on the much tinier warm-up: this is suspicious

5

B Conclusion

In this project, the disadvantage of those backends is the division computation such as / and div, which can be avoided by putting the division computation in the output phase.

References

- [1] Google. Streaming videos, 2017. Available from https://hashcode.withgoogle.com/2017/tasks/hashcode2017_qualification_task.pdf.
- [2] Google. Streaming videos data, 2017. Available from https://hashcode.withgoogle.com/2017/tasks/qualification_round_2017.in.zip.