

Product of As Expanded

Theorem

$$\begin{aligned} & \bar{A}_\ell^{[N]} \bar{A}_{\ell-1}^{[N]} \dots \bar{A}_3^{[N]} \bar{A}_2^{[N]} \\ &= \text{pow} \left(\overbrace{\left[\sum_{S \subseteq \{\ell, \ell-1, \dots, 3, 2\}} \exp(|S| B_\Delta^{[1]}[e]) \exp(W_\Delta^{[E]}[:, e] \cdot \sum_{i \in S} x_i^{[E]}) \right]}^{[1]}, \bar{A}[e]^{[N]} \right) \end{aligned}$$

Context

We have

$$\begin{aligned} \Delta_{\ell,e}^{[1]} &= \text{softplus}(x_\ell^{[E]} \cdot W_\Delta^{[E]}[:, e] + B_\Delta^{[1]}[e]) \\ \bar{A}_{\ell,e}^{[N]} &= \exp(\Delta_{\ell,e}^{[1]} \bar{A}[e]^{[N]}) \\ \bar{B}_\ell^{[N]} &= W_B^{[N,E]} x_\ell^{[E]} \\ \bar{B}_{\ell,e}^{[N]} &= \Delta_{\ell,e}^{[1]} \bar{B}_\ell^{[N]} \\ h_{\ell,e}^{[N]} &= \bar{A}_{\ell,e}^{[N]} h_{\ell-1,e}^{[N]} + \bar{B}_{\ell,e}^{[N]} x_{\ell,e}^{[1]} \\ \bar{C}_\ell^{[N]} &= W_C^{[N,E]} x_\ell^{[E]} \\ y_{\ell,e}^{[1]} &= h_{\ell,e}^{[N]} \cdot \bar{C}_\ell^{[N]} \\ h_{\ell,e}^{[N]} &= \exp(\Delta_{\ell,e}^{[1]} \bar{A}[e]^{[N]}) h_{\ell-1,e}^{[N]} + \Delta_{\ell,e}^{[1]} W_B^{[N,E]} x_{\ell,e}^{[E]} x_{\ell,e}^{[1]} \end{aligned}$$

Lets expand the recurrence out a bit(temporarily omitting the e and b index for readability)

$$\begin{aligned} h_1 &= \bar{A}_1^{[N]} \mathbf{0}^{[N]} + \bar{B}_1^{[N]} x_{1,e}^{[1]} = \bar{B}_1^{[N]} x_{1,e}^{[1]} \\ h_2 &= \bar{A}_2^{[N]} h_1 + \bar{B}_2^{[N]} x_{2,e}^{[1]} = \bar{A}_2^{[N]} \bar{B}_1^{[N]} x_{1,e}^{[1]} + \bar{B}_2^{[N]} x_{2,e}^{[1]} \\ h_3 &= \bar{A}_3^{[N]} h_2 + \bar{B}_3^{[N]} x_{3,e}^{[1]} = \bar{A}_3^{[N]} \bar{A}_2^{[N]} \bar{B}_1^{[N]} x_{1,e}^{[1]} + \bar{A}_3^{[N]} \bar{B}_2^{[N]} x_{2,e}^{[1]} + \bar{B}_3^{[N]} x_{3,e}^{[1]} \\ h_4 &= \bar{A}_4^{[N]} h_3 + \bar{B}_4^{[N]} x_{4,e}^{[1]} = \bar{A}_4^{[N]} \bar{A}_3^{[N]} \bar{A}_2^{[N]} \bar{B}_1^{[N]} x_{1,e}^{[1]} + \bar{A}_4^{[N]} \bar{A}_3^{[N]} \bar{B}_2^{[N]} x_{2,e}^{[1]} + \bar{A}_4^{[N]} \bar{B}_3^{[N]} x_{3,e}^{[1]} + \bar{B}_4^{[N]} x_{4,e}^{[1]} \end{aligned}$$

$$\begin{aligned}
& \dots \\
& h_\ell = \\
& \begin{aligned} & \overset{[N]}{A}_\ell \overset{[N]}{A}_{\ell-1} \dots \overset{[N]}{A}_3 \overset{[N]}{A}_2 \overset{[N]}{B}_1 \overset{[1]}{x}_{1,e} + \\ & \overset{[N]}{A}_\ell \overset{[N]}{A}_{\ell-1} \dots \overset{[N]}{A}_4 \overset{[N]}{A}_3 \overset{[N]}{B}_2 \overset{[1]}{x}_{2,e} + \\ & \vdots \\ & \overset{[N]}{A}_\ell \overset{[N]}{A}_{\ell-1} \overset{[N]}{A}_{\ell-2} \overset{[N]}{B}_{\ell-3} \overset{[1]}{x}_{\ell-3,e} + \\ & \overset{[N]}{A}_\ell \overset{[N]}{A}_{\ell-1} \overset{[N]}{B}_{\ell-2} \overset{[1]}{x}_{\ell-2,e} + \\ & \overset{[N]}{A}_\ell \overset{[N]}{B}_{\ell-1} \overset{[1]}{x}_{\ell-1,e} + \\ & \overset{[N]}{B}_\ell \overset{[1]}{x}_{\ell,e} \end{aligned}
\end{aligned}$$

Which can be written as

$$\begin{aligned}
& h_\ell = \\
& \left(\overset{[N]}{A}_\ell \overset{[N]}{A}_{\ell-1} \dots \overset{[N]}{A}_3 \overset{[N]}{A}_2, \overset{[N]}{A}_\ell \overset{[N]}{A}_{\ell-1} \dots \overset{[N]}{A}_4 \overset{[N]}{A}_3, \dots \overset{[N]}{A}_\ell \overset{[N]}{A}_{\ell-1} \overset{[N]}{A}_{\ell-2}, \overset{[N]}{A}_\ell \overset{[N]}{A}_{\ell-1}, \overset{[N]}{A}_\ell, 1 \right) \\
& \left(\overset{[N]}{B}_1, \overset{[N]}{B}_2, \dots, \overset{[N]}{B}_{\ell-3}, \overset{[N]}{B}_{\ell-2}, \overset{[N]}{B}_{\ell-1}, \overset{[N]}{B}_\ell \right) \\
& \left(\overset{[1]}{x}_{1,e}, \overset{[1]}{x}_{2,e}, \dots, \overset{[1]}{x}_{\ell-3,e}, \overset{[1]}{x}_{\ell-2,e}, \overset{[1]}{x}_{\ell-1,e}, \overset{[1]}{x}_{\ell,e} \right)
\end{aligned}$$

(doing element-wise multiplication of these three terms, than summing all the elements together.

Now consider the first term in

$$\left(\overset{[N]}{A}_\ell \overset{[N]}{A}_{\ell-1} \dots \overset{[N]}{A}_3 \overset{[N]}{A}_2, \overset{[N]}{A}_\ell \overset{[N]}{A}_{\ell-1} \dots \overset{[N]}{A}_4 \overset{[N]}{A}_3, \dots \overset{[N]}{A}_\ell \overset{[N]}{A}_{\ell-1} \overset{[N]}{A}_{\ell-2}, \overset{[N]}{A}_\ell \overset{[N]}{A}_{\ell-1}, \overset{[N]}{A}_\ell, 1 \right)$$

specifically this term:

$$\overset{[N]}{A}_\ell \overset{[N]}{A}_{\ell-1} \dots \overset{[N]}{A}_3 \overset{[N]}{A}_2$$

Proof

lets expand it out

$$= \exp(\overset{[1]}{\Delta}_{\ell,e} \overset{[N]}{A}[e]) \exp(\overset{[1]}{\Delta}_{\ell-1,e} \overset{[N]}{A}[e]) \dots \exp(\overset{[1]}{\Delta}_{3,e} \overset{[N]}{A}[e]) \exp(\overset{[1]}{\Delta}_{2,e} \overset{[N]}{A}[e])$$

$$\begin{aligned}
&= \exp(\Delta_{\ell,e}^{[1]} A[e]^{[N]} + \Delta_{\ell-1,e}^{[1]} A[e]^{[N]} + \cdots + \Delta_{3,e}^{[1]} A[e]^{[N]} + \Delta_{2,e}^{[1]} A[e]^{[N]}) \\
&= \exp((\Delta_{\ell,e}^{[1]} + \Delta_{\ell-1,e}^{[1]} + \cdots + \Delta_{3,e}^{[1]} + \Delta_{2,e}^{[1]}) A[e]^{[N]})
\end{aligned}$$

What is this inner delta term? Well

$$\Delta_{\ell,e}^{[1]} = \text{softplus}(x_{\ell}^{[E]} \cdot W_{\Delta}[:, e] + B_{\Delta}^{[1]}[e])$$

Lets call $\delta_{\ell,e}^{[1]} = x_{\ell}^{[E]} \cdot W_{\Delta}[:, e] + B_{\Delta}^{[1]}[e]$

Then, consider this simple case first:

$$\begin{aligned}
&\Delta_{1,e}^{[1]} + \Delta_{2,e}^{[1]} \\
&= \text{softplus}(\delta_{1,e}^{[1]}) + \text{softplus}(\delta_{2,e}^{[1]})
\end{aligned}$$

Useless fun fact

Note,

$$\begin{aligned}
\text{softplus}(x) &= \log(1 + e^x) \\
&= \log\left(\frac{e^{-x} + 1}{e^{-x}}\right) \\
&= \log(e^{-x} + 1) - \log(e^{-x}) \\
&= \log(1 + e^{-x}) + x
\end{aligned}$$

Means that

$$\begin{aligned}
&\Delta_{1,e}^{[1]} + \Delta_{2,e}^{[1]} \\
&= \text{softplus}(\delta_{1,e}^{[1]}) + \text{softplus}(\delta_{2,e}^{[1]}) \\
&= \log\left((1 + \exp(-\delta_{1,e}^{[1]}))(1 + \exp(-\delta_{2,e}^{[1]}))\right) + \delta_{1,e}^{[1]} + \delta_{2,e}^{[1]} \\
&= \log\left(\exp(-\delta_{1,e}^{[1]} - \delta_{2,e}^{[1]}) + \exp(-\delta_{1,e}^{[1]}) + \exp(-\delta_{2,e}^{[1]}) + 1\right) + \delta_{1,e}^{[1]} + \delta_{2,e}^{[1]}
\end{aligned}$$

vs

Back to the proof

$$\text{softplus}(x) = \log(1 + e^x)$$

$$\begin{aligned}
& \Delta_{1,e}^{[1]} + \Delta_{2,e}^{[1]} \\
&= \text{softplus}(\delta_{1,e}^{[1]}) + \text{softplus}(\delta_{2,e}^{[1]}) \\
&= \log \left((1 + \exp(\delta_{1,e}^{[1]})) (1 + \exp(\delta_{2,e}^{[1]})) \right) \\
&= \log \left(\exp(\delta_{1,e}^{[1]} + \delta_{2,e}^{[1]}) + \exp(\delta_{1,e}^{[1]}) + \exp(\delta_{2,e}^{[1]}) + 1 \right)
\end{aligned}$$

I think I prefer this way of writing it because either way you gotta do all the subsets. Speaking of the subsets, we can rewrite this:

$$\begin{aligned}
& \log \left(\exp(\delta_{1,e}^{[1]} + \delta_{2,e}^{[1]}) + \exp(\delta_{1,e}^{[1]}) + \exp(\delta_{2,e}^{[1]}) + 1 \right) \\
&= \log \left(\sum_{S \subseteq \{1,2\}} \exp \left(\sum_{i \in S} \delta_{i,e} \right) \right)
\end{aligned}$$

Where $S \subseteq \{1, 2\}$ means loop through all subsets (including $\{1, 2\}$ and the empty set, where empty set gives $\exp(0) = 1$)

Lets expand this out more: since $\delta_{\ell,e}^{[1]} = x_{\ell}^{[E]} \cdot W_{\Delta}^{[E]}[:, e] + B_{\Delta}^{[1]}[e]$

$$\begin{aligned}
& \log \left(\sum_{S \subseteq \{1,2\}} \exp \left(\sum_{i \in S} \delta_{i,e} \right) \right) \\
&= \log \left(\sum_{S \subseteq \{1,2\}} \exp \left(\sum_{j \in S} x_j^{[E]} \cdot W_{\Delta}^{[E]}[:, e] + B_{\Delta}^{[1]}[e] \right) \right) \\
&= \log \left(\sum_{S \subseteq \{1,2\}} \exp \left(|S| B_{\Delta}^{[1]}[e] + W_{\Delta}^{[E]}[:, e] \cdot \sum_{i \in S} (x_i^{[E]}) \right) \right) \\
&= \log \left(\sum_{S \subseteq \{1,2\}} \exp \left(|S| B_{\Delta}^{[1]}[e] \right) \exp \left(W_{\Delta}^{[E]}[:, e] \cdot \sum_{i \in S} (x_i^{[E]}) \right) \right)
\end{aligned}$$

Thus

$$\begin{aligned}
& \Delta_{1,e}^{[1]} + \Delta_{2,e}^{[1]} \\
&= \log \left(\sum_{S \subseteq \{1,2\}} \exp \left(|S| B_{\Delta}^{[1]}[e] \right) \exp \left(W_{\Delta}^{[E]}[:, e] \cdot \sum_{i \in S} (x_i^{[E]}) \right) \right)
\end{aligned}$$

Going back to

$$\begin{aligned}
& \bar{A}_\ell^{[N]} \bar{A}_{\ell-1}^{[N]} \dots \bar{A}_3^{[N]} \bar{A}_2^{[N]} \\
&= \exp(\Delta_{\ell,e}^{[1]} A[e]^{[N]}) \exp(\Delta_{\ell-1,e}^{[1]} A[e]^{[N]}) \dots \exp(\Delta_{3,e}^{[1]} A[e]^{[N]}) \exp(\Delta_{2,e}^{[1]} A[e]^{[N]}) \\
&= \exp(\Delta_{\ell,e}^{[1]} A[e]^{[N]} + \Delta_{\ell-1,e}^{[1]} A[e]^{[N]} + \dots + \Delta_{3,e}^{[1]} A[e]^{[N]} + \Delta_{2,e}^{[1]} A[e]^{[N]}) \\
&= \exp((\Delta_{\ell,e}^{[1]} + \Delta_{\ell-1,e}^{[1]} + \dots + \Delta_{3,e}^{[1]} + \Delta_{2,e}^{[1]}) A[e]^{[N]})
\end{aligned}$$

From above, which

$$\begin{aligned}
& \Delta_{\ell,e}^{[1]} + \Delta_{\ell-1,e}^{[1]} + \dots + \Delta_{3,e}^{[1]} + \Delta_{2,e}^{[1]} \\
&= \text{softplus}(\delta_{\ell,e}^{[1]}) + \text{softplus}(\delta_{\ell-1,e}^{[1]}) + \dots + \text{softplus}(\delta_{3,e}^{[1]}) + \text{softplus}(\delta_{2,e}^{[1]}) \\
&= \log \left(\sum_{S \subseteq \{\ell, \ell-1, \dots, 3, 2\}} \exp \left(\sum_{i \in S} \delta_{i,e}^{[1]} \right) \right)
\end{aligned}$$

And since $\delta_{\ell,e}^{[1]} = x_\ell^{[E]} \cdot W_\Delta[:, e] + B_\Delta[e]$

$$\begin{aligned}
& \Delta_{\ell,e}^{[1]} + \Delta_{\ell-1,e}^{[1]} + \dots + \Delta_{3,e}^{[1]} + \Delta_{2,e}^{[1]} \\
&= \log \left(\sum_{S \subseteq \{\ell, \ell-1, \dots, 3, 2\}} \exp \left(\sum_{i \in S} x_i^{[E]} \cdot W_\Delta[:, e] + B_\Delta[e] \right) \right) \\
&= \log \left(\sum_{S \subseteq \{\ell, \ell-1, \dots, 3, 2\}} \exp \left(|S| B_\Delta[e] + W_\Delta[:, e] \cdot \sum_{i \in S} x_i^{[E]} \right) \right) \\
&= \log \left(\sum_{S \subseteq \{\ell, \ell-1, \dots, 3, 2\}} \exp \left(|S| B_\Delta[e] \right) \exp \left(W_\Delta[:, e] \cdot \sum_{i \in S} x_i^{[E]} \right) \right)
\end{aligned}$$

Thus,

$$\begin{aligned}
& \bar{A}_\ell^{[N]} \bar{A}_{\ell-1}^{[N]} \dots \bar{A}_3^{[N]} \bar{A}_2^{[N]} \\
&= \exp(\log \left(\sum_{S \subseteq \{\ell, \ell-1, \dots, 3, 2\}} \exp \left(|S| B_\Delta[e] \right) \exp \left(W_\Delta[:, e] \cdot \sum_{i \in S} x_i^{[E]} \right) \right) A[e]^{[N]})
\end{aligned}$$

Now, note that $\exp(ab) = \exp(a)^b$ (assuming a and b are real). Thus

$$\exp(\log(a)b) = \exp(\log(a))^b = a^b$$

Giving us

$$\overline{A}^{[N]}_{\ell} \overline{A}^{[N]}_{\ell-1} \dots \overline{A}^{[N]}_3 \overline{A}^{[N]}_2$$

$$= \text{pow}\Big(\overbrace{\Big[\sum_{S \subseteq \{\ell, \ell-1, \dots, 3, 2\}} \exp\big(|S| B^{[1]}_{\Delta}[e]\big) \exp\big(W^{[E]}_{\Delta}[:, e] \cdot \sum_{i \in S} x^{[E]}_i\big)\Big]}^{[1]}, A^{[N]}[e]\Big)$$

Note that, as labeled, the big summation term results in a single float. So this is

$$r = \text{pow}(\overbrace{a}^{[1]}, \overbrace{b}^{[N]})$$

where

$$r_i = \text{pow}(a, b_i) = a^{b_i}$$