

COMP 520 - Compilers

Lecture 8 (Tue Feb 8, 2022)

Abstract Syntax Trees

- **Reading for Feb 10**
 - Finish Chapter 4
 - section 4.6 – Syntactic Analysis in Triangle (pp 124 – 128)
- **Materials on our website**
 - simpleAST
 - illustrates AST construction and evaluation for miniArith
 - AbstractSyntaxTrees.zip
 - `miniJava.AbstractSyntaxTrees` package needed to construct AST's for miniJava in PA2

Topics

- From concrete syntax trees to abstract syntax trees
 - AST “grammar”
 - AST representation choices
 - AST construction
- AST traversal
 - generalize using Visitor
- miniJava AST classes (available on web)
 - use these to construct the AST in PA2
 - you will need to provide some implementation of sourcePosition



Concrete syntax tree

- Concrete syntax is described by an (EBNF) CFG grammar
 - ex: simple arithmetic expressions

$S ::= E \$$

$E ::= E \text{ Op } T \mid T$

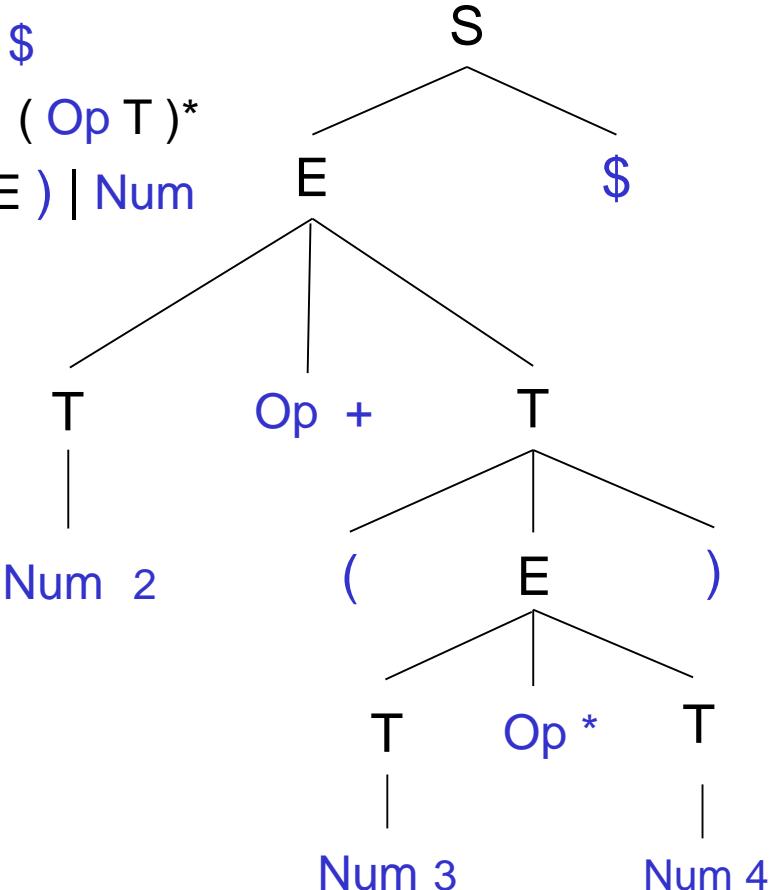
$T ::= (E) \mid \text{Num}$



$S ::= E \$$

$E ::= T (\text{Op } T)^*$

$T ::= (E) \mid \text{Num}$

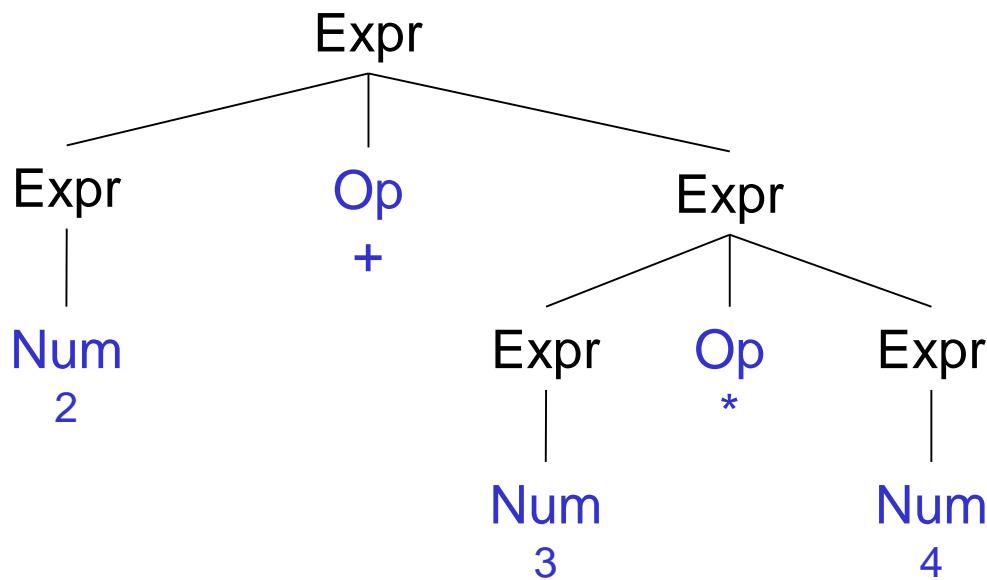


Abstract syntax

- Abstract syntax can be described by a simple CFG grammar
 - ex: simple arithmetic expressions

$$\begin{array}{ll} \text{Expr} ::= \text{Expr } \text{Op } \text{Expr} & (\text{BinExpr}) \\ | \text{ Num} & (\text{NumExpr}) \end{array}$$

- Abstract syntax tree for $2 + (3 * 4)$



AST representation in Java (Strategy 1)

- AST classes define an (n-ary) tree with tagged nodes
 - simple to define, more complex to construct and traverse
 - minimal benefit from Java type checker

```
enum NodeType { UNARYEXPR, BINARYEXPR, NUM }
public abstract class AST {
    public NodeType n;
}

public class Node extends AST {
    public AST[] children;
    public Node(NodeType n, AST[] children) { ... }
}

public class Leaf extends AST {
    public String spelling;
    public Leaf(NodeType n, String spelling) { ... }
}
```



AST representation in Java (Strategy 2)

- AST classes closely follow structure of AST grammar
 - more complex to define AST classes, but simpler to construct AST
 - we gain considerable benefit from the Java type checker
 - more rigorously supports AST traversal using a Visitor interface
- Rules
 - create abstract class AST
 - for each nonterminal in AST grammar
 - construct an abstract subclass of AST
 - for each choice within a rule
 - construct a concrete subclass of the LHS nonterminal
- We will follow this strategy



AST representation in Java (Strategy 2)

- Example

Expr ::= Expr Oper Expr	(BinExpr)
Num	(NumExpr)

```
abstract public class AST {}

abstract public class Expr extends AST {}

public class BinExpr extends Expr {
    public Terminal oper;
    public Expr left, right;
    public BinExpr(Expr left, Terminal oper, Expr right) { ... }
}

public class NumExpr extends Expr {
    public Terminal num;
    public NumExpr(Terminal num) { ... }
}

public class Terminal extends AST {
    public String spelling;
    public Terminal(String spelling) { ... };
}
```



Building an AST during a concrete syntax parse

- concrete syntax for arithmetic expression grammar

$$\begin{aligned} E &::= T \mid E \ op \ T \\ T &::= (E) \mid num \end{aligned}$$

- transformed and augmented

$$\begin{aligned} S &::= E \$ \\ E &::= T (op \ T)^* \\ T &::= (E) \mid num \end{aligned}$$

- abstract syntax

$$\begin{aligned} Expr &::= Expr \ Op \ Expr \ (BinExpr) \\ &\quad \mid \text{Num} \quad (\text{NumExpr}) \end{aligned}$$

- how to build AST?

- modify parse procedures to return pieces of AST

```
Expr parseS() {
    Expr e = parseE();
    accept(Token.eot);
    return e
}

Expr parseE() {
    Expr e1 = parseT();
    while (curToken.kind == Token.op) {
        Terminal op = new Terminal(curToken);
        acceptIt();
        Expr e2 = parseT();
        e1 = new BinExpr(e1,op,e2);
    }
    return e1;
}

Expr parseT() {
    case (curToken.kind) {
        Token.LPAREN:
        acceptIt();
        Expr e1 = parseE();
        accept(Token.RPAREN);
        return e1;
    }
    Token.num:
        NumExpr e2 = new NumExpr(curToken);
        acceptIt();
        return e2;
    }
}
```



mini-Triangle – Expression ASTs

Concrete Syntax (EBNF)

	Expression	::= primary-Expression	(1.4a)
		Expression Operator primary-Expression	(1.4b)
	primary-Expression	::= Integer-Literal	(1.5a)
		V-name	(1.5b)
		Operator primary-Expression	(1.5c)
		(Expression)	(1.5d)
	V-name	::= Identifier	(1.6)
Tokens	Operator	::= + - * / < > = \	(1.10a-h)
	Identifier	::= Letter Identifier Letter Identifier Digit	(1.11a-c)
	Integer-Literal	::= Digit Integer-Literal Digit	(1.12a-b)

- Expression ASTs (E):

IntegerExpression

|
IntegerLiteral

(1.16a)

VnameExpression

|
V

(1.16b)

UnaryExpression

|
Operator E
|
spelling

(1.16c)

BinaryExpression

|
E₁ Operator E₂
|
spelling

(1.16d)

- V-name ASTs (V):

SimpleVname

|
Identifier

(1.17)

|
spelling

Abstract Syntax (ASTs)

mini-Triangle – Command ASTs

Concrete Syntax

Program	::=	single-Command	(1.1)
Command	::=	single-Command	(1.2a)
		Command ; single-Command	(1.2b)
single-Command	::=	V-name := Expression	(1.3a)
		Identifier (Expression)	(1.3b)
		if Expression then single-Command else single-Command	(1.3c)
		while Expression do single-Command	(1.3d)
		let Declaration in single-Command	(1.3e)
		begin Command end	(1.3f)

ASTs



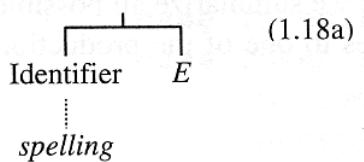
mini-Triangle – Declaration ASTs

Concrete Syntax

Declaration	::=	single-Declaration	(1.7a)
		Declaration ; single-Declaration	(1.7b)
single-Declaration	::=	const Identifier ~ Expression	(1.8a)
		var Identifier : Type-denoter	(1.8b)
Type-denoter	::=	Identifier	(1.9)

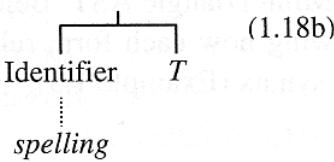
- Declaration ASTs (D):

ConstDeclaration



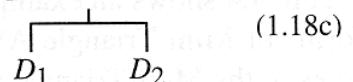
(1.18a)

VarDeclaration



(1.18b)

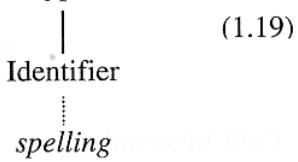
SequentialDeclaration



(1.18c)

- Type-denoter ASTs (T):

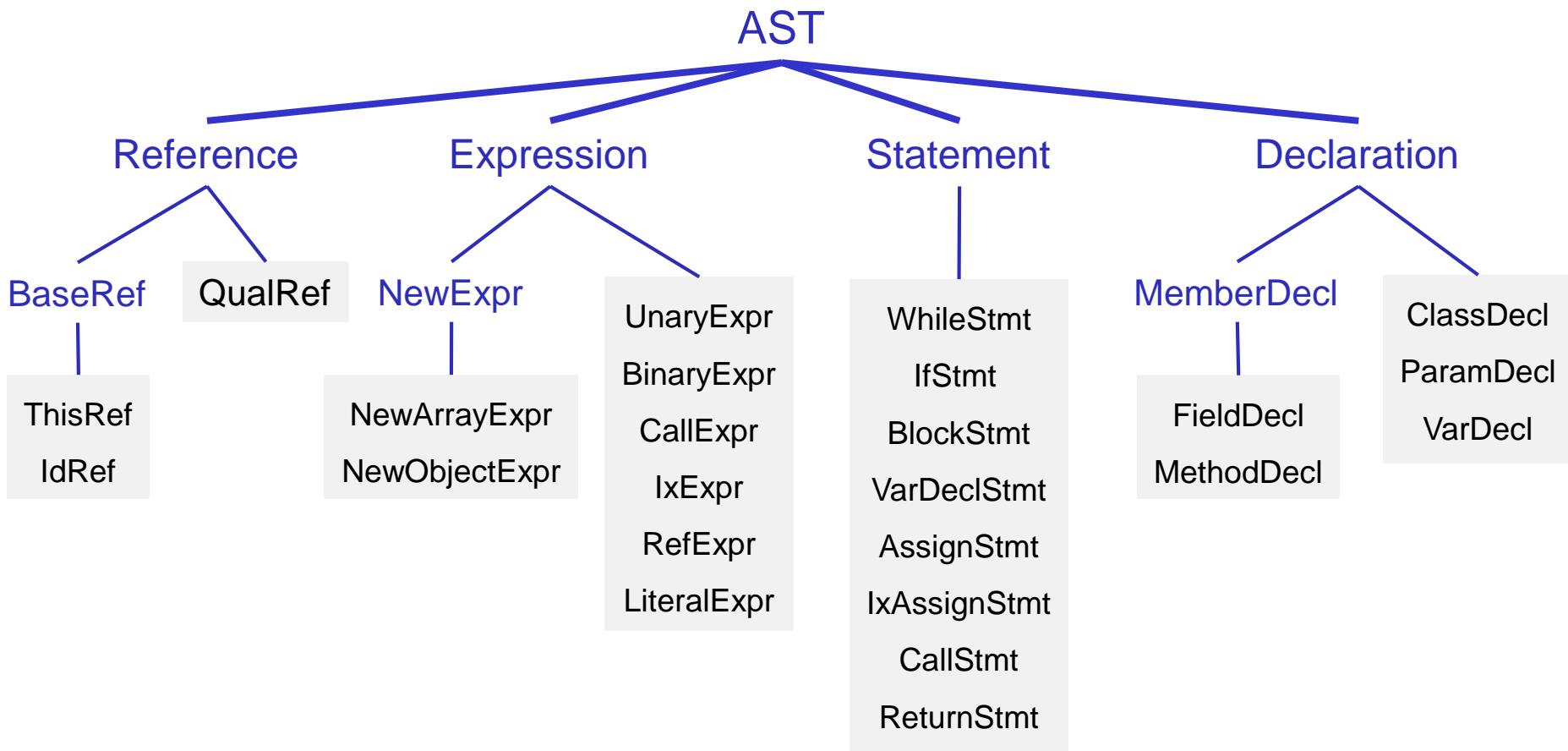
SimpleTypeDenoter



(1.19)

ASTs

The miniJava AST classes



Additional constructors for lists of class instances, with an iterator for traversal of the list:

`ClassDeclList`, `FieldDeclList`, `MethodDeclList`, `ParamDeclList`, `StatementList`, `ExprList`





miniJava AST types

- ✓ **C^A AST**
- ✓ **C^A Declaration**
 - C ClassDecl
 - ✓ **C^A LocalDecl**
 - C ParameterDecl
 - C VarDecl
 - ✓ **C^A MemberDecl**
 - C FieldDecl
 - C MethodDecl
- ✓ **C^A Expression**
 - C BinaryExpr
 - C CallExpr
 - C IxExpr
 - C LiteralExpr
 - ✓ **C^A NewExpr**
 - C NewArrayExpr
 - C NewObjectExpr
 - C RefExpr
 - C UnaryExpr
 - C Package

- ✓ **C^A Reference**
 - ✓ **C^A BaseRef**
 - C IdRef
 - C ThisRef
 - C QualRef
- ✓ **C^A Statement**
 - C AssignStmt
 - C BlockStmt
 - C CallStmt
 - C IfStmt
 - C IxAssignStmt
 - C ReturnStmt
 - C VarDeclStmt
 - C WhileStmt
- ✓ **C^A Terminal**
 - C BooleanLiteral
 - C Identifier
 - C IntLiteral
 - C Operator
- ✓ **C^A TypeDenoter**
 - C ArrayType
 - C BaseType
 - C ClassType
- C ASTDisplay

Implementing a general AST traversal

- **Strategy 1**
 - add methods to each AST class for each kind of traversal
 - Example
 - display methods for AST display
 - eval methods for AST evaluation
 - drawback
 - Classes become very large when traversals get complex
 - Contextual analysis
 - Code generation
 - Code for each kind of traversal is scattered over many classes



Implementing a general AST traversal

- Strategy 2
 - use a *visitor* pattern
 - A visitor Interface
 - specifies a visitX method for each concrete class X in AST
 - a visit method in each AST class
 - public Object visit(Visitor v, Object arg)
 - each separate traversal implements the Visitor interface
 - sample Visitor instances
 - AST display
 - Identification
 - Type checking
 - Code generation
 - Code for each type of visitor is collected in one class
 - but a bit cumbersome to write out



Simple set of AST classes

```
abstract class AST {}

abstract class Expr extends AST {}

class BinExpr extends Expr {
    public Token oper;
    public Expr left;
    public Expr right;

    public BinExpr(Expr left, Token oper, Expr right) { . . . }
}

class NumExpr extends Expr {
    public int val;

    public NumExpr(int val) { . . . }
}
```



The Visitor interface

- Visitor interface requires a **visitX** method for every (non-abstract) AST class X

```
public interface Visitor {  
    visitBinaryExpr(BinaryExpr e);  
    visitNumExpr(NumExpr e);  
}
```

- Each AST class is augmented with a single **visit** method

```
class NumExpr extends Expr {  
    public int val;  
    public NumExpr(int v) { . . . }  
  
    public void visit(Visitor v) { v.visitNumExpr(this); }  
}
```

- All AST traversals use the same “**visit**” method in each node type
 - the visit method “connects” a specific visitor v to **this** specific node



Inorder traversal of the AST

- The `InorderWalk` AST traversal implements the `Visitor` interface

```
public class InorderWalk implements Visitor {
    public void visitBinExpr(BinExpr be) {
        be.left.visit(this);
        System.out.println(be.oper.spelling);
        be.right.visit(this);
    }

    public void visitNumExpr(NumExpr ne) {
        System.out.println(ne.val);
    }

    // print nodes of AST inorder
    public void walk(AST a) {
        a.visit(this);
    }
}
```



Adding arguments to the traversal (Book method)

- methods implemented by a Visitor have an Object arg and Object result

```
public interface Visitor {  
    Object visitBinExpr(BinExpr e, Object arg);  
    Object visitNumExpr(NumExpr e, Object arg);  
}
```

- AST class visit method has an Object arg and Object result

```
class NumExpr extends Expr {  
    public int val;  
  
    public NumExpr(int val) { . . . }  
  
    public Object visit(Visitor v, Object arg) {  
        return v.visitNumExpr(this, arg);  
    }  
}
```



Example use of the Visitor

- Individual traversals implement Visitor with custom logic for each node type

```
class Checker implements Visitor {  
  
    public Object visitAssignment(Assignment s, Object arg) {  
        Type t1 = (Type) s.ref.visit(this, arg);  
        Type t2 = (Type) s.exp.visit(this, arg);  
        return (t1.equals(t2) ? t1 : Type.ERROR)  
    }  
}
```

- Good solution?
 - (+) appropriately OO
 - (+) compiler insures visitor defined for all node types
 - (+) specific definitions gathered together in a single class
 - (-) Object parameters and results will require a lot of explicit casting



Adding arguments to the traversal (parameterized types)

```
public interface Visitor<ArgType, ResultType> {
    public ResultType visitBinExpr(BinExpr expr, ArgType arg);
    public ResultType visitNumExpr(NumExpr expr, ArgType arg);
}
```

```
class NumExpr extends Expr {
    public int val;
    public NumExpr(int val) { . . . }

    public <ArgType, ResType> ResType
        visit(Visitor<ArgType, ResType> v, ArgType arg) {
            return v.visitNumExpr(this, arg);
    }
}
```



Example use of the Visitor

- Individual traversals implement Visitor with custom logic for each node type

```
class Checker implements Visitor<Type, Type> {  
    ...  
  
    public Type visitAssignment(Assignment s, Type arg) {  
        Type t1 = s.ref.visit(this, arg);  
        Type t2 = s.exp.visit(this, arg);  
        if (!t1.equals(t2))  
            typeError("incompatible types in assignment", s);  
        return Types_STMT_TYPE;  
    }  
}
```

- Good solution?
 - Improves type checking in the visitors
 - Improves readability



PA2 Submission Details

1. Your compiler project must follow these requirements with respect to the package names and structure
 2. Your `Compiler.java` mainclass must have the functionality called out on the following page, but you may vary your implementation.



```

package miniJava;
import java.io.FileInputStream;
import java.io.FileNotFoundException;
import java.io.InputStream;

import miniJava.SyntacticAnalyzer.Parser;
import miniJava.SyntacticAnalyzer.Scanner;
import miniJava.AbstractSyntaxTrees.*;

/** 
 * Recognize whether input file named in args[0] contains a syntactically valid
 * miniJava program, and, if valid, display corresponding AST.
 */
public class Compiler {

    public static void main(String[] args) {
        InputStream inputStream = null;
        try {
            inputStream = new FileInputStream(args[0]);
        }
        catch (FileNotFoundException e) {
            System.out.println("Input file " + args[0] + " not found");
            System.exit(3); ←----- use this
        }
        ErrorReporter errorReporter = new ErrorReporter();
        Scanner scanner = new Scanner(inputStream, errorReporter);
        Parser parser = new Parser(scanner, errorReporter);

        System.out.println("Syntactic analysis ... ");
        AST ast = parser.parseProgram();
        System.out.print("Syntactic analysis complete: ");
        if (errorReporter.hasErrors()) {
            System.out.println("Invalid miniJava program");
            System.exit(4); ←----- use these
        }
        else {
            System.out.println("Valid miniJava program:");
            new ASTDisplay().showTree(ast);
            System.exit(0); ←----- termination codes for
        }
    }
}

```

create AST → use this
termination code if
unable to open
input file

output AST using
ASTDisplay(),
for accepted
miniJava
programs

use these
termination
codes for
invalid/valid miniJava
programs,
respectively