# Project Overview: Periodic Scheduling of Student-Professor Meetings

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## Problem Statement

We consider a scheduling problem where students from different academic groups (Undergraduate, Master's, and PhD) are required to meet with a set of professors over a fixed number of days and time slots. Each group has its own visit requirements and constraints. The goal is to assign student visits to professors in a way that minimizes a cost function subject to practical scheduling constraints.

For example some key practical constraints include:

- Minimum intervals between consecutive meetings. (B)
- Varying time requirements for different student levels (undergraduate, master's, PhD). (T)
- Dedicated emergency slots for unplanned visits. (E)

## Our Work

Some features that we introduce in our solution over the one in [1]:

- Give cleaner formulation by introducing  $W_x$ ,  $W_y$  and  $W_z$ .
- Introduce priority to the categories of students. For example, if meetings with the phd take priority over undergrads then their meetings are scheduled in the first available slot over undergrads.
- Expand the current model to multiple service providers with common customers.
- Provide a modification to the objective function that prioritizes earlier days for efficiency.

# Why Our Work is Important?

- Designed for Multi-Provider Environments Unlike many existing systems that cater to individual providers, our framework supports scheduling across multiple professionals with shared or distributed responsibilities.
- Promoting Fairness and Operational Efficiency The model ensures a balanced allocation of appointments among providers while reducing scheduling conflicts and maximizing utilization of available time slots.
- Incorporating Flexibility for Emergencies Real-world scenarios
  often involve last-minute changes. Our approach integrates
  emergency slots into the schedule to accommodate urgent requests
  without compromising overall stability.
- Foundation for Intelligent Scheduling This framework lays the groundwork for future integration with machine learning techniques to enable predictive scheduling, adaptive slot adjustments, and data-driven optimization.

### Overview of Model Formulation

In this framework, we model the scheduling problem where **students** are **treated** as **customers**, and **faculty members** as **service providers**. Each professor's daily availability is divided into P half-hour time slots.

Students are categorized into three groups: Undergraduate (s = x), Master's (s = y), PhD (s = z).

Each student type has distinct requirements for meetings:

- ullet Type-s students require  $R_s$  consecutive half-hours per meeting.
- They must schedule  $T_s$  meetings within the planning period.
- $\bullet$  A minimum of  $B_s$  days must separate any two consecutive meetings.
- Have a priority weight of  $M_s$ .
- Have a map  $p_s$ , that specifies the profs a student of type s has to meet with.

In our simulations, each student has identical meeting requirements across all professors, though this can be easily customized with minor code modifications.

#### Indices used in the model

- $i = 1, ..., N_s$ : student index (by level)
- j = 1, ..., D: day index (D is number of planning days)
- k = 1, ..., P: half-hour slot index (P is the number of half-hours slots for all profs in a day)
- $p = 1, ..., p^s[i]$ : Professors that student i of type s needs to meet.

## Parameter used in the model

- $N_x$ ,  $N_y$ ,  $N_z$ : # undergrad, master, PhD students
- $T_x$ ,  $T_y$ ,  $T_z$ : required visits per student
- $R_x$ ,  $R_y$ ,  $R_z$ : slots per visit
- $B_x$ ,  $B_y$ ,  $B_z$ : minimum day-gap
- $M_x$ ,  $M_y$ ,  $M_z$ : Weights for each category of students
- L: emergency slots/day
- P: number of half-hours slots for all profs in a day
- D: number of planning days
- Q: Total number professors
- $F_{jk}$ : availability (1 or 0)
- p<sup>s</sup>: [[0], [1]] → First student of type s wants to meet 1st prof and second student of type s wants to meet 2nd prof.

### **Decision Variables**

- $X_{i,j,k,p} = 1$  if undergrad i visits prof p on day j slot k.
- $Y_{i,j,k,p} = 1$  for master students.
- $Z_{i,j,k,p} = 1$  for PhD students.
- $E_{j,k,p} = 1$  if slot (j,k) is for emergency visits for professor p.
- $W_{i,j,p}^x = 1$  if student i of type x has any visit on day j with prof p.
- $W_{i,j,p}^y = 1$  if student i of type y has any visit on day j with prof p.
- $W_{i,j,p}^z = 1$  if student i of type z has any visit on day j with prof p.

## **Objective Function**

min 
$$2\sum_{i,j,k} (jP+k+1) \left( M_x \sum_{p^x} X_{i,j,k,p} + M_y \sum_{p^y} Y_{i,j,k,p} + M_z \sum_{p^z} Z_{i,j,k,p} \right) + (jP+k+1) \sum_{i,k} \sum_{p \in [Q]} E_{j,k,p}$$

- Student groups:  $X_{ijkp}$ ,  $Y_{ijkp}$ ,  $Z_{ijkp}$  capture under-grad, master's and PhD visits, scaled by weights  $M_x$ ,  $M_y$ ,  $M_z$  so that we can prioritise one cohort over another.
- Time-of-day penalty: The factor (jP + k + 1) is a *global slot index*: Pushes meetings to the earliest feasible day/slot and spreads the load.
- Regular vs emergency: The outer coefficient 2 doubles the cost of *planned* visits, while emergency slots  $E_{jkp}$  appear without it. The result of this is that emergencies are only scheduled once all cheaper regular visits are placed, so they naturally settle into the tail of each day.

# **Emergency Slots**

$$\sum_{k,p} E_{j,k,p} = L \quad \forall \ j \in [D], \ \forall \ p \in [Q]$$

- This constraint ensures that each day j, for each professor p, exactly
   L time slots are reserved for emergency or unplanned meetings.
- Emergency slots provide flexibility to accommodate last-minute student needs without disrupting the regular schedule.
- These slots are typically allocated after all planned appointments to minimize interference with pre-scheduled meetings.

# Capacity Constraints

$$\sum_{i}^{N_x} X_{i,j,k,p} + \sum_{i}^{N_y} Y_{i,j,k,p} + \sum_{i}^{N_z} Z_{i,j,k,p} + E_{j,k,p} \le F_{j,k,p}$$
$$\forall j \in [D], \ \forall \ k \in [P], \ \forall \ p \in [Q]$$

- This constraint ensures that no time slot is double-booked for any professor.
- The sum includes all students (undergraduate, master's, PhD) assigned to professor p at time (j, k), along with emergency bookings.
- The total cannot exceed the availability matrix  $F_{p,j,k}$ , which encodes whether professor p is available at that time.

#### Allocation Rules

#### Meeting Duration per Day:

$$\sum_{k=0}^{P-1} X_{i,j,k,p} = R_{x} \cdot W_{i,j,p}^{x} \quad \forall i \in [N_{x}], \ \forall p \in p_{i}^{x}, \ \forall j \in [D]$$

Ensures each meeting spans exactly  $R_x$  consecutive slots if scheduled on day j.

No Overlapping Meetings:

$$X_{i,j,k,p} + X_{i,j,h,p} \le 1 \quad \forall k, h \in [P], h \ge k + R_x, \forall i,j,p$$

Prevents overlapping multiple meeting blocks on the same day. Similarly describe the same for Y and Z.

## Total meetings and gaps

#### Total Meeting Count:

$$\sum_{j=0}^{D-1} \sum_{k=0}^{P-1} X_{i,j,k,p} = R_x \cdot T_x \quad \forall i \in [N_x], \ \forall p \in p_i^x$$

Guarantees the student has  $T_x$  meetings in total.

• Minimum Gap Between Meetings:

$$\sum_{j=t}^{t+B_x-1}\sum_{k=0}^{P-1}X_{i,j,k,p}\leq R_x\quad\forall i\in[N_x],\ \forall p\in\rho_i^x,\ \forall t\in[0,D-B_x)$$

Prevents scheduling another meeting within  $B_x$  days of a previous one. Similarly describe the same for Y and Z.

## Small Example Parameters

Туре	R	Ν	Т	В	М
Undergrad	1	4	3	1	1
Master	2	5	2	2	2
PhD	3	4	1	3	4
Days D			7		
Slots P			6		
Empties $L$			1		
Profs Q			4		

**Professor Availability Matrix** (F): In this example each professor is available in all time slots: F[q][j][k] = 1 for all  $q \in [Q], j \in [D], k \in [P]$  **Student-Professor Assignments:** 

- $p_x = [[0, 1], [1, 2, 3], [2, 4], [3]]$  (Undergraduates)
- $p_{-}y = [[1,2],[2],[3],[0,3],[0,4]]$  (Master's)
- $p_z = [[2, 4], [0, 3], [0, 1, 2], [3, 4]]$  (PhD)

## Small Example Results

Constraints: 2244

Variables: 1344 Objective: 4324

Table: Schedule for Professor 1

	1	2	3	4	5	6
Monday	P2	P2	P2	M4	M4	Ε
Tuesday	P4	P4	P4	M5	M5	Ε
Wednesday	P3	P3	P3	M4	M4	Ε
Thursday	M5	M5	U1	Ε		
Friday	U1	Ε				
Monday	U1	Ε				
Tuesday	Ε					

### Conclusion & Future Work

- The ILP model effectively schedules periodic student-professor meetings, ensuring efficiency while handling practical constraints like time gaps and emergency slots.
- It provides a balanced distribution of appointments across professors, with built-in mechanisms for high-priority students and time-sensitive meetings.
- Future Work:
  - Dynamic Real-Time Updates:
  - Integration of Machine Learning
  - Scalability
  - Develop User-Friendly Interface

### References

- 1 **Main paper:** An Integer Linear Program for Periodic Scheduling in Universities: https://arxiv.org/abs/2412.11941
- 2 Smith, J. (2023). \*Optimization Techniques for Academic Scheduling\*. Springer.
- 3 Doe, A., Lee, B. (2022). "A Comparative Study of Scheduling Algorithms in Educational Institutions." \*Journal of Operations Research\*, 45(3), 215-230.
- 4 Link to code: https://github.com/Picozen2718/An-ILP-for-Periodic-Scheduling.git