# - Open Question -Does the uncloneable bit exist?

## Towards Unconditional Uncloneable Encryption [6]

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## 1 Goal

Have a secure cryptographic scheme against cloning attacks.

#### 2 Idea

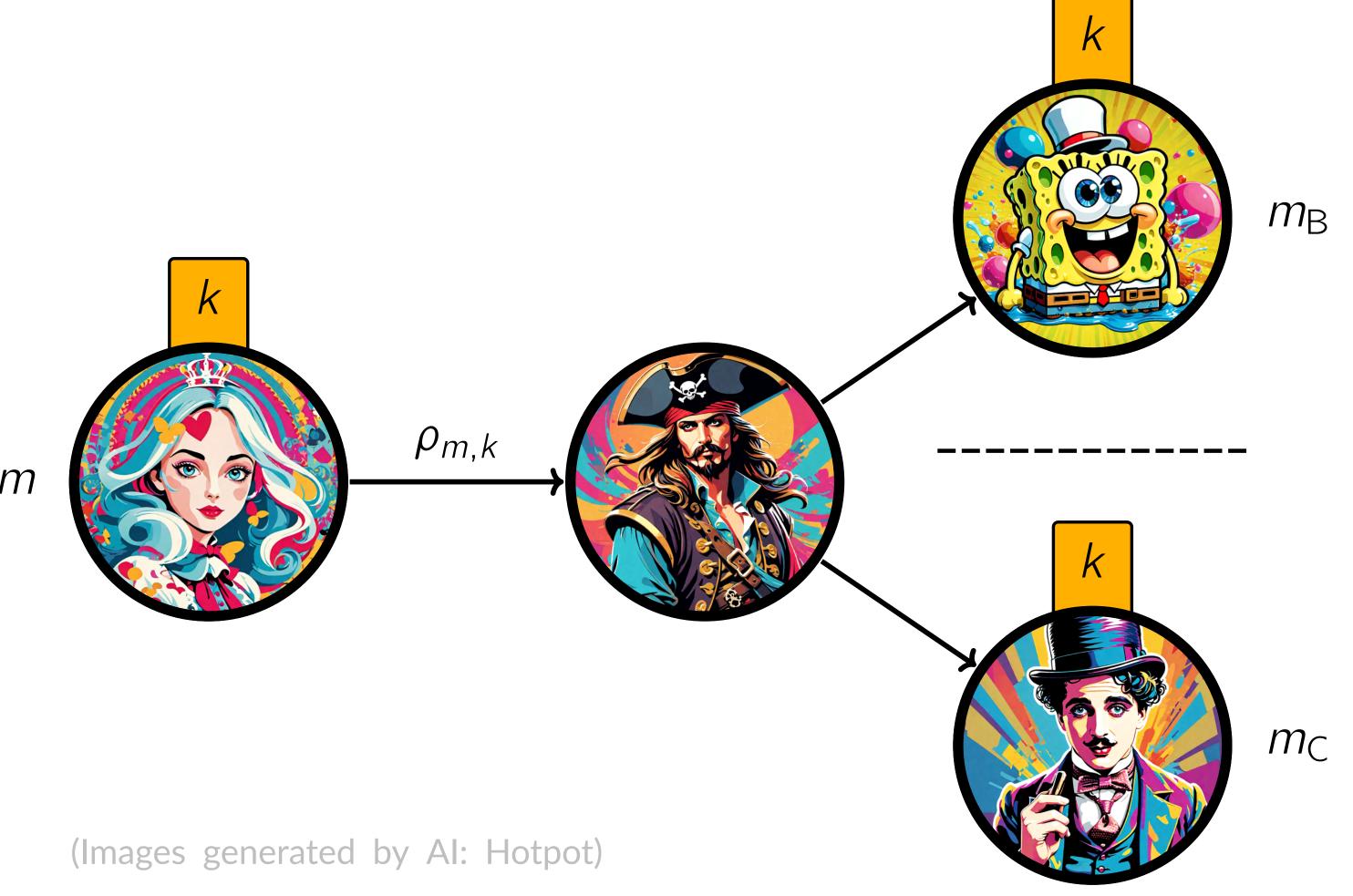
Laverage the quantum no-cloning theorem.

## **3** Consequences

**Applications** private-key quantum money [8], preventing [8], attacks quantum storage functional encryption [13], quantum copy-protection [2], uncloneable decryption [11, 15, 12], and quantum position verification [10].

## 1. Uncloneable Bit

**4 No-Cloning Game.** Alice encrypts a message  $m \in \{0, 1\}$  with a key  $k \in \{1, ..., K\}$ into a quantum state  $\rho_{m,k}$ . She sends it via a quantum channel but a pirate (P) intercepts it. Without knowing the key k, the pirate tries to share the information with two non-communicating parties, Bob (B) and Charlie (C), so that both of them may retrieve Alice's message m. We say that the adversary team (P, B, C) wins the game if both Bob's and Charlie's guesses are correct, i.e. if  $m_B = m_C = m$ .



## **5** Definition (Correctness)

The encryption protocol  $(m, k) \mapsto \rho_{m,k}$  is said to be correct if there exists a way to retrieve m from  $\rho_{m,k}$  and k.

#### **6** Definition (Security)

A protocol  $(m, k) \mapsto \rho_{m,k}$  is said to be uncloneable-indistinguishable secure if:

$$\mathbb{P}\Big((\mathsf{P},\mathsf{B},\mathsf{C}) \text{ win the game}\Big) \leqslant \frac{1}{2} + f(\lambda),$$

where  $f(\lambda) \to 0$  as  $\lambda \to \infty$ , and where  $\lambda$  is the security parameter. Additionally, this security is said to be strong if moreover  $f(\lambda) = \text{negl}(\lambda)$ .

#### **8** Former Work

have focused **Efforts** on achievability under various models and definitions, including:

- in the quantum random oracle model (QROM) [8, 3, 4],
- in an interactive version of the scenario [7],
- in a device-independent variant with variable keys [12],
- assuming the existence of specific types of obfuscation [1, 9],
- in a variant with quantum keys [5], and a "succinct" variant [14].

## 7 Open Question (Uncloneable Bit) [8]

Is there an encryption scheme  $(m,k) \mapsto \rho_{m,k}$  that is both correct and uncloneable-indistinguishable secure?

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## 2. Candidate Scheme and Conjecture

## **9** Candidate Scheme: Clifford Algebra

Let  $\Gamma_1, \ldots, \Gamma_K$  be Hermitian unitaries that anti-commute. Consider the following encryption:

$$\rho_{m,k} := \frac{2}{d} \frac{I_d + (-1)^m \Gamma_k}{2},$$

where  $m \in \{0, 1\}$  and  $k \in \{1, ..., K\}$ . This encryption protocol is correct since one can retrieve m from measuring  $\rho_{m,k}$  in the eigenbasis of  $\Gamma_k$ . It remains to show the security.

**10** Example. It is possible to produce such  $\Gamma_k$ 's using pairwise anti-commuting Pauli strings. Indeed, if K = 2n, consider:

$$\Gamma_k:=X^{\otimes (k-1)}\otimes Y\otimes I^{\otimes (n-k)} \ \ ext{and} \ \ \Gamma_{n+k}:=X^{\otimes (k-1)}\otimes Z\otimes I^{\otimes (n-k)} \ ,$$

for  $k \in \{1, ..., n\}$ . Otherwise, if K = 2n + 1, consider the same operators and add  $\Gamma_{2n+1} := X^{\otimes n}$ .

## 11 Conjecture

This scheme is uncloneable-indistinguishable secure:

$$\mathbb{P}\Big((\mathsf{P},\mathsf{B},\mathsf{C}) \text{ win the game}\Big) \leqslant \frac{1}{2} + \frac{1}{2\sqrt{K}}.$$

**Remark.** Here  $K \sim 2\lambda$ , but ideally  $K \sim 2^{\lambda}$  (strong security).

## 3. Results

## 12 Proposition (Sufficient Formula)

To achieve the security of the Conjecture, it is sufficient to prove the following upper bound for all Hermitian unitaries  $\{U_k\}$ :

$$\left\| \underbrace{\sum_{k=1}^{K} \left( \Gamma_k \otimes U_k \otimes I + \Gamma_k \otimes I \otimes U_k + I \otimes U_k \otimes U_k \right)}_{=:W_K} \right\|_{\text{op}} \leqslant K + 2\sqrt{K}.$$

**13** Remarks. The value  $K + 2\sqrt{K}$  is achieved when considering  $U_k = I$  for all k. Moreover, the formula trivially holds if we assume that the operators  $U_k$  commute.

## Theorem 1

The Conjecture is valid for  $K \leq 7$ .

*Proof.* When  $K \leq 7$ , we find the following sum-of-squares (SoS) decomposition:

$$(K + 2\sqrt{K})I - W_K = \sum_{k=1}^{K} \alpha_k A_k^2$$

for some explicit coefficients  $\alpha_k \geqslant 0$  and operators  $A_k$ . Hence  $(K+2\sqrt{K})I-W_K \geq 0$  and therefore  $K+2\sqrt{K} \geq ||W_K||_{op}$ .

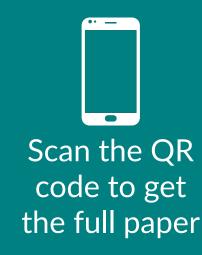
## **15** Numerical Results

The Conjecture is numerically confirmed for  $K \leq 17$  (NPA level-2) algorithm) and  $K \leq 18$  (Seesaw algorithm).

## Theorem 2

Asymptotically, the winning probability of the no-cloning game for our candidate scheme is upper-bounded by 5/8.





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