

## Summary: Lecture 4

Summary for the chapters  $X$  and  $X$ . [1]

### Asymmetry of non-deterministic revisited

to check if formula is non satisfiable  
question whether  $NP = coNP$

similar issues in non-deterministic space:  
write down  $n$  nodes on a ... non-det. will succeed if it is possible  
to show that it is impossible to ... in SPACE was unsolved until the 80s  
 $h$  for function (yes or no or  $h$ ?)  
 $h$  is yes  
everything that fails stops in state with no

### Immerman-Szelepcsenyi

how many distinct nodes can be reached in a graph if you start from a graph  $x$   
 $s(0)$  will contain node 1 and  $s(1)$  will contain all neighbours of 1  
we will have actual names  
4 nested for loops and algorithm happens in the middle  
outer for loop:  
computes number of nodes reachable from initial node (for  $k$  steps with  $k$  as the iterative thingy  
in the for loop)  
in each step we override the previous set with the next one because we only have limited space  
second loop:  
we get how far we got in the previous steps and sum up how far we can get (because we can get  
previous set size?)  
third loop:  
the actual magic happens here: checking something  
Aux sounds like a port for headphones  
return no when all guesses were correct? we remember solution that we were supposed to reach  
beforehand

## References

- [1] Christos H. Papadimitriou. *Computational Complexity*. Addison-Wesley Publishing Company, 1994.