

## Summary: Lecture 7

Summary for the chapters 9.1 and 9.2. [1]

### NP-Completeness

#### NP

Class of languages decided by nondeterministic Turing machines in polynomial time.  
Most problems are in NP.

#### NP-completeness:

- easiest problems among those we do not know how to solve efficiently
- if  $P \neq NP$  can be proven: exact border of efficient solvability is found
- best bet for proving  $P=NP$ : show that some NP-complete problem is  $P$
- Until then, the NP-complete problems are the least likely ones in NP to be efficiently solved

#### Language $L$

$L = \{x : (x, y) \in R \text{ for some } y\}$   
 $L$  gets an input  $x$  and finds a  $y$  with  $((x, y) \in R$  and the relation  $R \subseteq \Sigma^* \times \Sigma^*$ .

#### Polynomially decidable:

- $R$  is polynomially decidable if there is a deterministic Turing machine deciding the language  $L$  in polynomial time
- then the relation  $R$  (not the language  $L$ ) is polynomially decidable

#### Polynomially balanced:

- $R$  is polynomially balanced if  $(x, y) \in R$  implies  $|y| \leq |x|^k$  for some  $k \geq 1$   
→ length of the second component is bounded by a polynomial in the length of the first
- then the relation  $R$  (not the language  $L$ ) is polynomially balanced

#### NP

The language  $L \subseteq \Sigma^*$  is in NP only if there is a polynomially decidable and polynomially balanced relation  $R$  such that  $L = \{x : (x, y) \in R \text{ for some } y\}$ .

For example: Is there a satisfying assignment ( $y$ ) for a formula ( $x$ )?

Why is  $R \subseteq \Sigma^* \times \Sigma^*$ ? Is the input formula and the truth assignment  $\in \Sigma^*$ ?

TODO

proof

Questions:

## Succinct certificate (for NP-complete problems)

- *yes* instance of  $x$  has a polynomial witness  $y$  (certificate)
- *no* instances don't have such a certificate
- Examples:
  - SAT: certificate is the truth assignment
  - HAMILTONPATH: certificate is the hamilton path of a graph

## Typical problems in NP

- sometimes the optimum needs to be found
- sometimes any object that fits the specification is enough
- constraints can be added to optimization problems

TODO

Questions:

## 3Sat is NP-complete

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TODO

Questions:

## 2Sat in P (graph construction)

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Questions:

## 2Sat in NL

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Questions:

## MaxSat is NP-complete

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Questions:

## NaeSat is NP-complete

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Questions:

## References

- [1] Christos H. Papadimitriou. *Computational Complexity*. Addison-Wesley Publishing Company, 1994.