

## Summary: Lecture 5

Summary for the chapters  $X$  and  $X$ . [6]

### Reduction

#### Examples of NP-problems:

- Travelling Salesman Problem
- SATISFIABLE
- REACHBILITY (in P)
- CIRCUIT VALUE (in P)

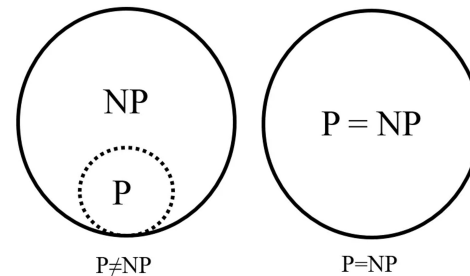


Figure 1: P and NP sets [3]

- reduction: a problem is at least as hard as another
- problem  $A$  is at least as hard as problem  $B$  if  $B$  reduces to  $A$
- $B$  reduces to  $A$  if there is a transformation  $R$ 
  - $R$  produces for every input  $x$  of  $B$  an equivalent input  $R(x)$  of  $A$
  - the answer of input  $x$  on  $B$  and input  $R(x)$  on  $A$  have to be the same
- to solve  $B$  on input  $x$ ,  $A$  can be solved instead with input  $R(x)$

#### Reduction

Problem  $A$  is at least as hard as problem  $B$  if  $B$  reduces to  $A$ .

#### Transformation function:

- transformation function  $R$  should not be too hard to compute  
→  $R$  should be limited
- efficient reduction  $R$ :  $\log n$  space bounded

#### Transformation function

A language  $L_1$  is reducible to  $L_2$  if there is a function  $R$  computable by a deterministic Turing Machine in space  $O(\log n)$  and  $x \in L_1 \Leftrightarrow R(x) \in L_2$ .

$R$  is called a reduction from  $L_1$  to  $L_2$ .

- A Turing Machine  $M$  that computes a reduction  $R$  halts for all inputs  $x$  after a polynomial number of steps.
  - there are  $O(n \cdot c^{\log n})$  possible configurations for  $M$  on an input of length  $n$
  - deterministic: no configuration can be repeated
  - computation of length at most  $O(n^k)$

## Reduction HAMILTONIAN PATH to SATISFIABLE

### Problem: HAMILTON PATH

The Hamiltonian Path problem asks whether there is a route in a directed graph  $G$  from a start node to an ending node, visiting each node exactly once. (Is there a path in  $G$  that visits each node one?) [1]

### Problem: SAT

The SAT (satisfiability) problem is the problem of determining if there exists an interpretation that satisfies a given Boolean formula. [7]

- HAMILTON PATH can be reduced to SAT  
→ demonstrates HAMILTON PATH is not significantly harder than SAT
- construct a boolean expression  $R(G)$  that is satisfiable only if  $G$  has a Hamilton path  
→ write a logical formula that only becomes true when HP is true
- instance: Graph  $G = (V, E)$  with  $n$  nodes  $(1, 2, \dots, n)$
- $R(G)$  has  $n^2$  boolean variables  $x_{i,j}$  then
- node  $j$  is the  $i$ th node in the HAMILTON PATH
- $R(G)$  is in conjunctive normal form (CNF:  $(a \vee b) \wedge (\neg a \vee c)$ )
- conjuncted clauses of  $R(x)$ :
  - each node  $j$  must appear in the path  $x_{1,j} \vee x_{2,j} \vee \dots \vee x_{n,j}$  – for every node  $j$
  - no node  $j$  appears twice in the path:  $\neg x_{i,j} \vee \neg x_{k,j}$  for all  $i, j, k$  with  $i \neq k$
  - every position  $i$  on the path must be occupied –  $x_{i,1} \vee x_{i,2} \vee \dots \vee x_{i,n}$  for each  $i$
  - no two nodes  $j$  and  $k$  occupy the same position in the path –  $\neg x_{i,j} \vee \neg x_{i,k}$  for all  $i, j, k$  with  $j \neq k$
  - nonadjacent nodes  $i$  and  $j$  cannot be adjacent in the path –  $\neg x_{k,i} \vee \neg x_{k+1,j}$  for all  $(i, j) \notin E$  and  $k = 1, 2, \dots, n - 1$

[4]

### Proof idea:

- to show:
  - for any graph  $G$ ,  $R(G)$  has a satisfying truth assignment only if and only if  $G$  has a Hamilton path
  - $R$  can be computed in space  $\log n$
- 

TODO

proof

Questions:

## Boolean Circuit

A Boolean circuit is a mathematical tree model for logic formulas. Boolean circuits are defined in terms of the logic gates they contain. For example, a circuit might contain binary AND and OR gates and unary NOT gates, or be entirely described by binary NAND gates.

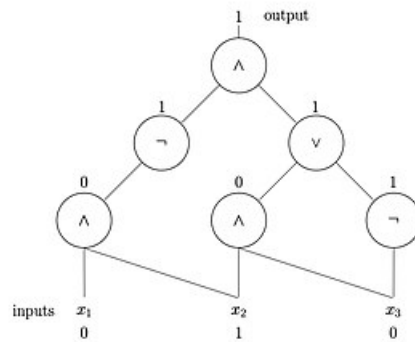


Figure 2: Boolean circuit example [2]

## Reduction REACHABILITY PATH to CIRCUIT VALUE

### Problem: GRAPH REACHABILITY

Given a graph  $G$  and two nodes  $n_1, n_2 \in V$ , is there path from  $n_1$  to  $n_2$ ?  
A graph  $G = (V, E)$  is a finite set  $V$  of nodes and a set  $E$  of edges as node pairs.

REACHABILITY can be nondeterministically solved in space  $\log n$ .

### Problem: CIRCUIT VALUE

The CIRCUIT VALUE Problem is the problem of computing the output of a given Boolean circuit on a given input.

In terms of time complexity, it can be solved in linear time (topological sort).

The problem is closely related to the SAT (Boolean Satisfiability) problem which is complete for NP and its complement, which is complete for co-NP.

- REACHABILITY can be reduced to CIRCUIT VALUE  
→ demonstrates REACHABILITY is not significantly harder than CIRCUIT VALUE
- construct a variable-free circuit  $R(G)$  that has *true* as output only if  $G$  has a path from the start node to node  $n$
- $R(x)$  uses no  $\neg$  gates (monotone circuit)
- (both problems are in P)
- idea: use the Floyd-Warshall algorithm (dynamic programming)
- instance: Graph  $G = (V, E)$  with  $n$  nodes  $(1, 2, \dots, n)$
- the gates:

- $g_{i,j,k}$  with  $1 \leq i, j \leq n$  and  $0 \leq k \leq n$ 
    - \* there is a path from node  $i$  to node  $j$  without passing through a node bigger than  $k$
    - \*  $g_{i,j,0}$  is true if and only if  $i = j$  or  $i$  and  $j$  are neighbours
  - $h_{i,j,k}$  with  $1 \leq i, j, k \leq n$ 
    - \* there is a path from node  $i$  to node  $j$  passing through  $k$  but not any node bigger than  $k$
  - $h_{i,j,k}$  is an *and* gate with predecessors  $g_{i,k,k-1}$  and  $g_{k,j,k-1}$  where  $k = 1, 2, \dots, n$
  - $g_{i,j,k}$  is an *or* gate with predecessors  $g_{i,j,k-1}$  and  $h_{i,j,k}$  where  $k = 1, 2, \dots, n$
  - $g_{1,n,n}$  is the output gate
- [5]

TODO

proof

Questions:

## Reduction CIRCUIT SAT to SAT

### Problem: CIRCUIT SAT

The circuit satisfiability problem (CIRCUIT SAT) is the decision problem of determining whether a given Boolean circuit has an assignment of its inputs that makes the output true.

- CIRCUIT SAT can be reduced to SAT
  - demonstrates CIRCUIT SAT is not significantly harder than SAT
- given a circuit  $C$ , construct a boolean expression  $R(C)$  such that  $R(C)$  is satisfiable if and only if  $C$  is satisfiable
- $R(C)$  will turn out to be a CNF.
- the variables of  $R(C)$  are those of  $C$  plus  $g$  for each gate  $g$  of  $C$
- clauses of  $R(C)$ :
  - $g$  is a variable gate  $x$ :
    - Add clauses  $(\neg g \vee x)$  and  $(g \vee \neg x) \equiv g \Leftrightarrow x$
  - $g$  is a *true* gate:
    - Add clause  $(g) \rightarrow g$  must be true to make  $R(C)$  true
  - $g$  is a *false* gate:
    - Add clause  $(\neg g) \rightarrow g$  must be false to make  $R(C)$  true
  - $g$  is a  $\neg$  gate with predecessor gate  $h$ :
    - Add clauses  $(\neg g \vee \neg h)$  and  $(g \vee h) \equiv g \Leftrightarrow \neg h$
  - $g$  is a  $\vee$  gate with predecessor gates  $h$  and  $h'$ :
    - Add clauses  $(\neg h \vee g)$ ,  $(\neg h' \vee g)$  and  $(h \vee h' \vee \neg g)$ , meaning:  $g \Leftrightarrow (h \vee h')$
  - $g$  is a  $\wedge$  gate with predecessor gates  $h$  and  $h'$ :
    - Add clauses  $(\neg g \vee h)$ ,  $(\neg g \vee h')$ , and  $(\neg h \vee \neg h' \vee g)$ , meaning:  $g \Leftrightarrow (h \wedge h')$
  - $g$  is the output gate:
    - Add clause  $(g)$ , meaning:  $g$  must be true to make  $R(C)$  true

[5]

TODO  
example  
proof idea  
Questions:

### **Further examples**

TODO  
Questions:

### **Closedness under Composition**

TODO  
Questions:

## References

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