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Summary: Lecture 6

Summary for the chapter *8.2.* [1, 4]

Completeness

Let C be a complexity class and let L be a language in C. L is called *C-complete* if any language $L' \in C$ can be reduced to L.

(Every language of a complexity class can be reduced to L.)

- reducitbility is transitive $\rightarrow$ problems are ordered by difficulty
-

TODO
Questions:

Which problems can be reduced to a formal language?

> SAT can be expressed as formal language. [3]
> $\Rightarrow$ SAT can be reduced to a formal language. (?)
>
> Because CIRQUIT SAT can be reduced to SAT: CIRQUIT SAT can be reduced to a formal language. (?)
>
> **Formal language**
>
> Formal languages are abstract languages, which define the syntax of the words that get accepted by that language. It is a set of words that get accepted by the language and has a set of symbols that is called alphabet, which contains all the possible characters of the words. Those characters are called nonterminal symbols. [2, 5]
>
> **Kleene star**
> The Kleene star Σ^* of an alphabet Σ is the set of all words that can be created through concatenation of the symbols of the alphabet Σ. The empty word ϵ is included.
>
> **Formal language**
> A formal language L over an alphabet Σ is a subset of the Kleene star of the alphabet: $L \subseteq \Sigma^*$
>
> Where to set the line between lanuguage decisions and other problems? Can every problem be contrcuted as a formal language?

What does completeness do for us?

- a reduction definition is usefull because the complexity classe are closed under reduction
- examples look helpfull

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- L and R seem to be important:

$$\begin{array}{cc} L' \in P & A \\ L \to L' & R \end{array}$$

- drawing set circle inclusion thing (P and NP)

TODO
Questions:

P-completeness of CIRCUIT VALUE

Problem: Circuit Value

The CIRCUIT VALUE Problem is the problem of computing the output of a given Boolean circuit on a given input.

In terms of time complexity, it can be solved in linear time (topological sort).

- P-complete
- limit of power of reductions
- got a little tired and zoned out

TODO
Questions:

The reduction (?)

Problem: Circuit Sat

The circuit satisfiability problem (CIRCUIT SAT) is the decision problem of determining whether a given Boolean circuit has an assignment of its inputs that makes the output true.

Input: a Boolean circuit C

Question: Is there a truth assignment which makes C output the value true?

CIRCUIT SAT is NP-complete

- circuit decides nondeterministically (?)
- a variable is added n the nondeterministic Turing Machine
- check if one of the variables is tue: use this choice (?)
- problem: can we set thiese variables such that the Turing Machine accepts?
- answer corresponds direct to *is there a choice of nd decisions such that the turing machine accepts?*
- extremely direct reduction
- cooks theorem :)
- SAT is NP-complete

TODO
Questions:

References

[1] Martin Berglund. *Lecture notes in Computational Complexity.*

[2] *Chomsky's Normal Form (CNF).* Website. `https://www.javatpoint.com/automata-chomskys-normal-form`, opened on 26.09.2022.

[3] klaus-joern Lange. "The Boolean Formula Value Problem as Formal Language". In: (Jan. 2012). DOI: `10.1007/978-3-642-31644-9_9`.

[4] Christos H. Papadimitriou. *Computational Complexity.* Addison-Wesley Publishing Company, 1994.

[5] A.J. Kfoury Robert N. Moll Michael A. Arbib. *An Introduction to Formal Language Theory.* Springer-Verlag, 1988.