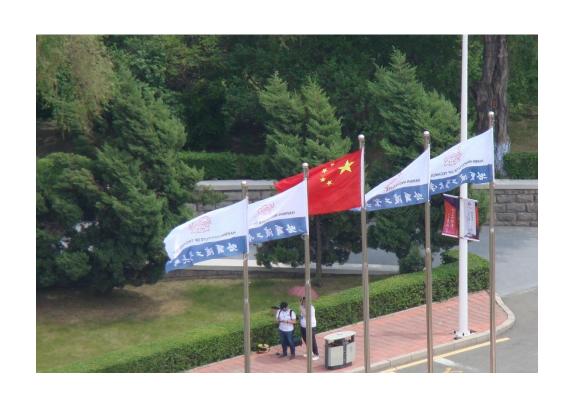
Afternoon

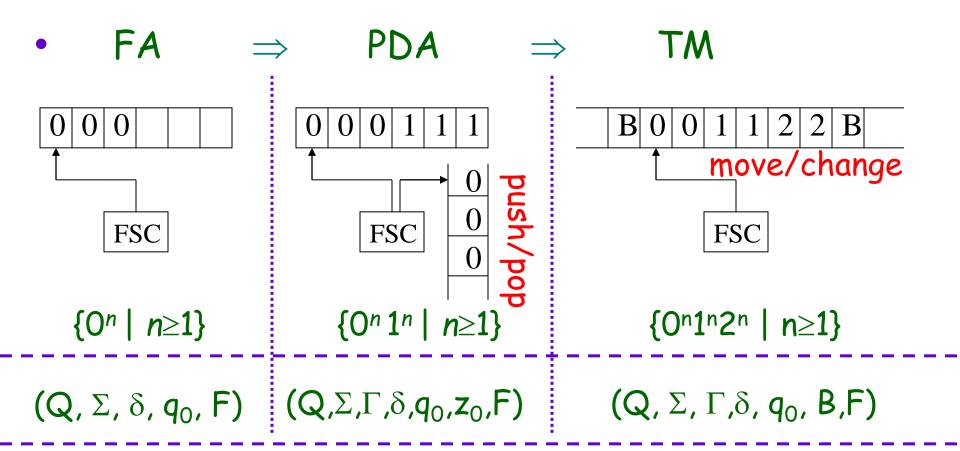


Turing Machine

- Definition
- ◆ Construction
- Language



Different automata

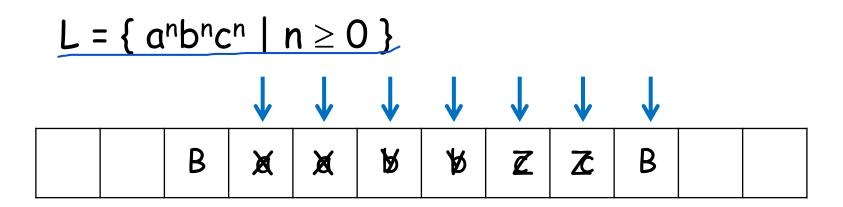


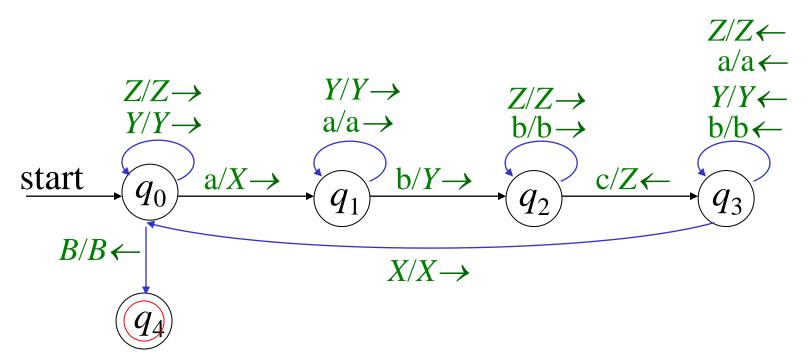
Difinition

TM is a seven-tuple $P = (Q, \Sigma, \Gamma, \delta, q_0, B, F)$

- Q is finite set of states
- Σ is finite set of input symbols
- Γ is finite set of tape symbols
- δ is transition function : $Q \times \Gamma \Rightarrow Q \times \Gamma \times \{R, L\}$ $\delta(q, X) = (p, Y, D)$
- q_0 is start state
- B is blank symbol
- F is finite set of final state

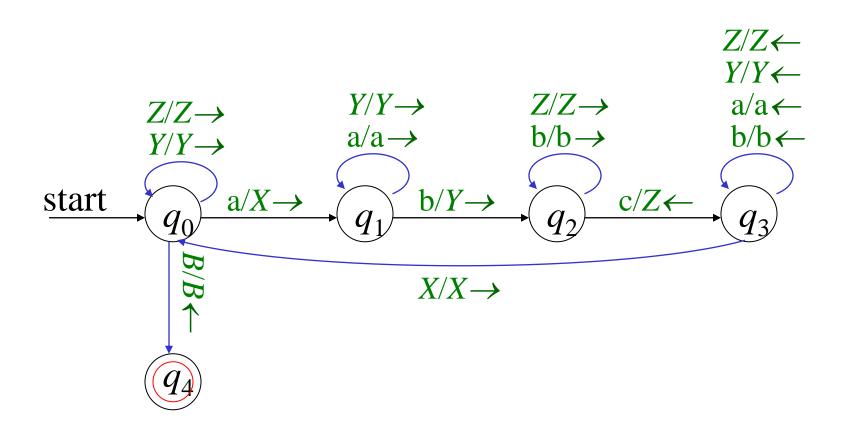
Example 1 TM for





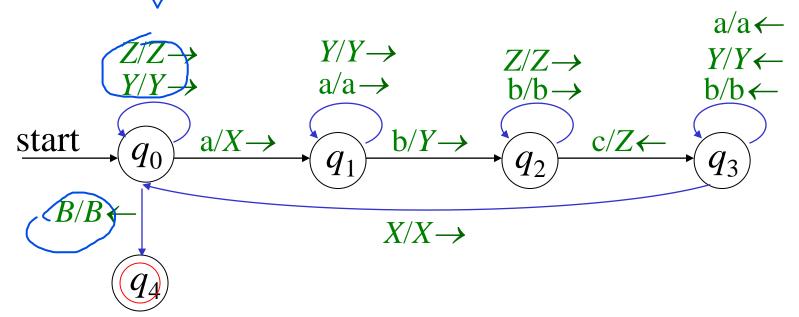
Example 1 TM for 'anbncn'

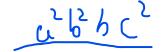
 $M = (\{q_0, q_1, q_2, q_3, q_4\}, \{a,b,c\}, \{a,b,c,B,X,Y,Z\}, \delta, q_0,B, \{q_4\})$



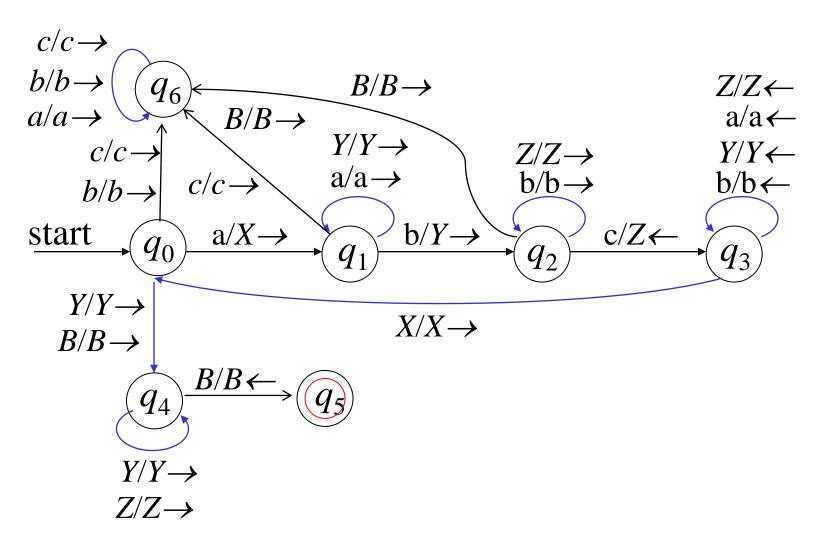
Example 1 TM for 'anbncn'

- ϵ ?
 $a^2b^2c^2a^3b^3c^3$??
 $a^2b^2c^2a^3b^3c^3$??
- $a^2b^3c^4$, ac, bbc???





Example 1 TM for 'anbncn'



Instantaneous

- how to describe the configuration of TM
 - > sequence of symbols in tape
 - > state of TM
 - > read/wtite head of TM
 - $> X_1 \dots X_{i-1} q X_i X_{i+1} \dots X_n$

Instantaneous

• ID of the above TM for $w=aabbcc \in \{a^nb^nc^n \mid n \ge 1\}$

 q_0 aabbcc $\vdash Xq_1$ abbcc $\vdash X$ a q_1 bbcc $\vdash XaYq_2$ bcc $XaYbq_2$ cc $\vdash X$ a Yq_3 bZc $\vdash X$ a q_3 YbZc $\vdash Xq_3$ aYbZc $\vdash Xq_3$ aYbZc $\vdash Xq_0$ aYbZc $\vdash XXq_1$ YbZc $\vdash XXYq_1$ bZc $\vdash XXYYq_2$ Zc $\vdash XXYYZq_2$ c $\vdash XXYYQ_3$ ZZ $\vdash ...$

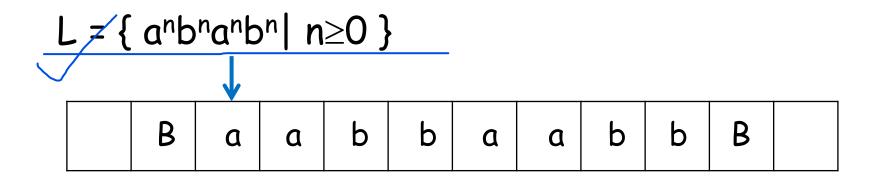
Language of TM

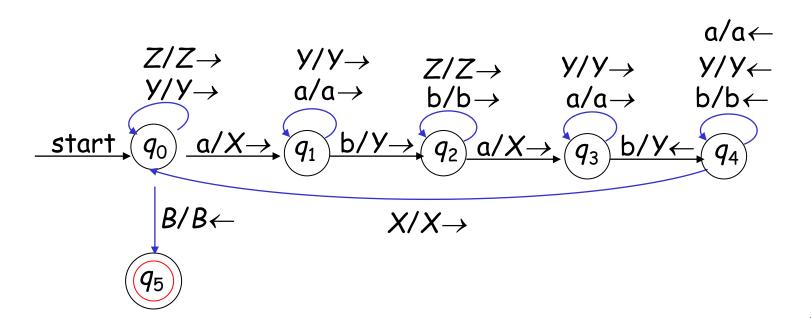
language accepted by a TM

$$\{w \mid q_0 w \not \succeq \alpha p \beta, p \in F, \alpha, \beta \in \Gamma^*\}$$

 The language accepted by TM is called recursively enumerable(RE) language.

Example 2 TM for

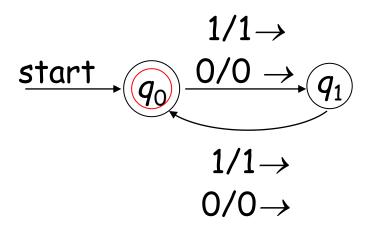




Example 3 TM for

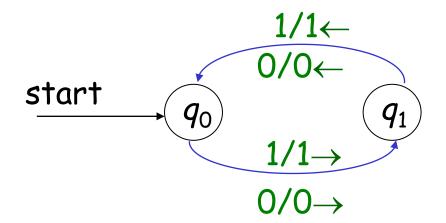
$$L = \{ w \mid w \in \{0,1\}^* \text{ and } |w| \text{ is even } \}$$

$$B \mid 0 \mid 0 \mid 1 \mid 1 \mid 0 \mid 1 \mid B$$



Halting

We say a TM halts if it enters a state q, scanning a tape symbol X, and there is no move in this situation.

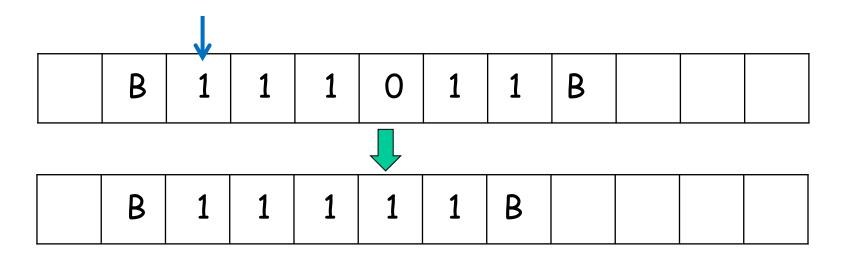


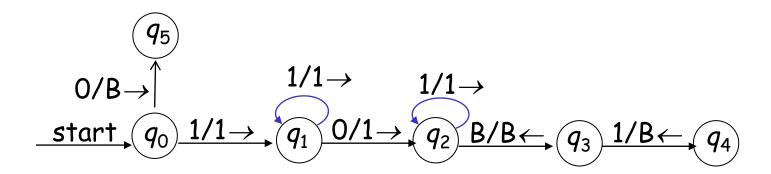
Given two positive integers x and y, design a TM to compute x + y.

Notation for x and y

$$x \Rightarrow w(x) \in \{1\}^+ \text{ and } |w(x)| = x$$

 $x + y \Rightarrow w(x + y) \in \{1\}^+ \text{ and } |w(x + y)| = x + y$
 $3 - 111, 2 - 11, 3 + 2 - 11111$





Example 5 褐龙

Compute the function nomus $(m, n) = \max(m-n, 0)$

• put 1^mO1^n into tape as input



- delete a 1 from 1^m and a 1 from 1^n
- three cases:

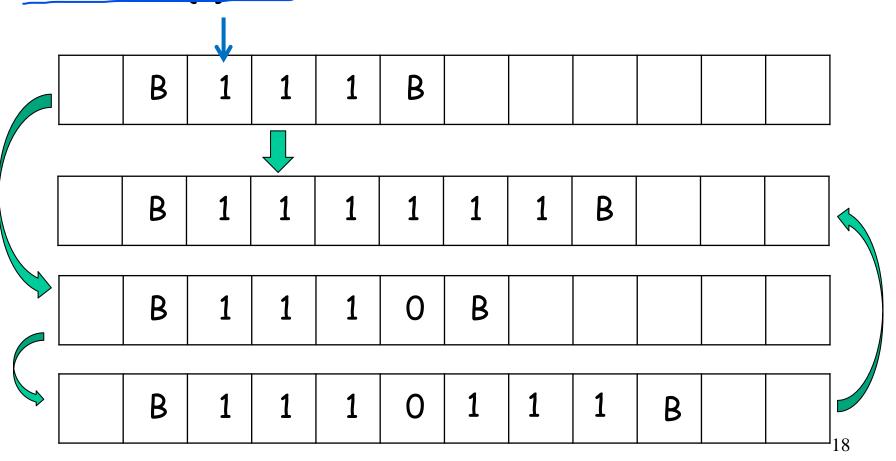
$$m>n \Leftrightarrow 1^{m-n-1}$$
 at the left of 1

$$m=n \Leftrightarrow \text{no } 1$$

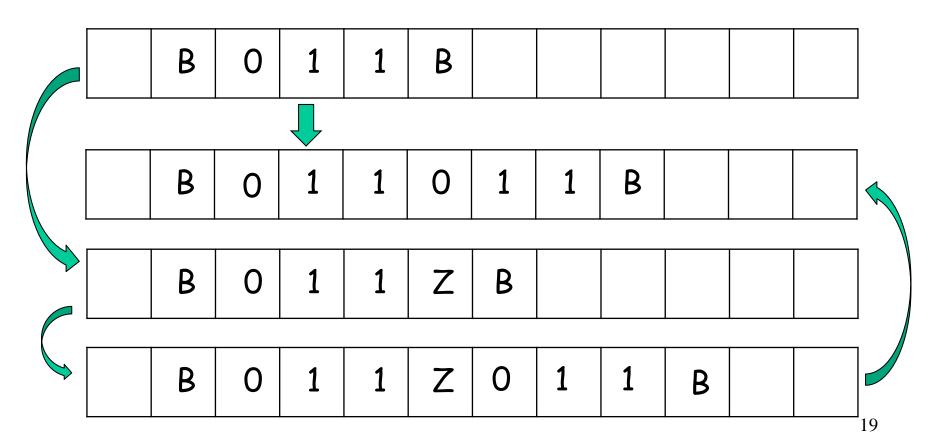
 $m < n \Leftrightarrow 1^{n-m}$ at the right of 1

Example 6 河流

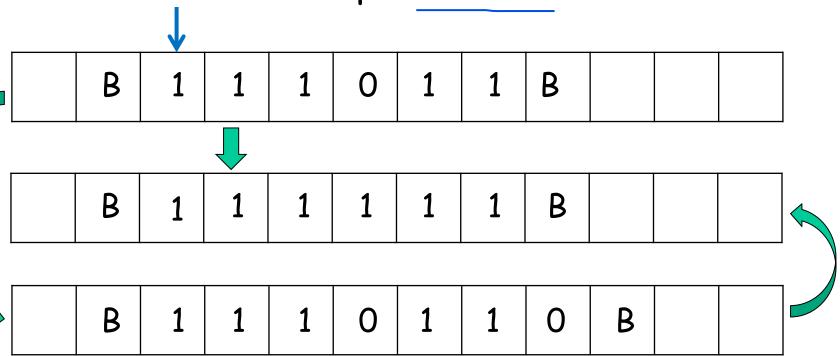
Construct a TM to compute the function f(w)=wwwhere $w \in \{1\}^+$.



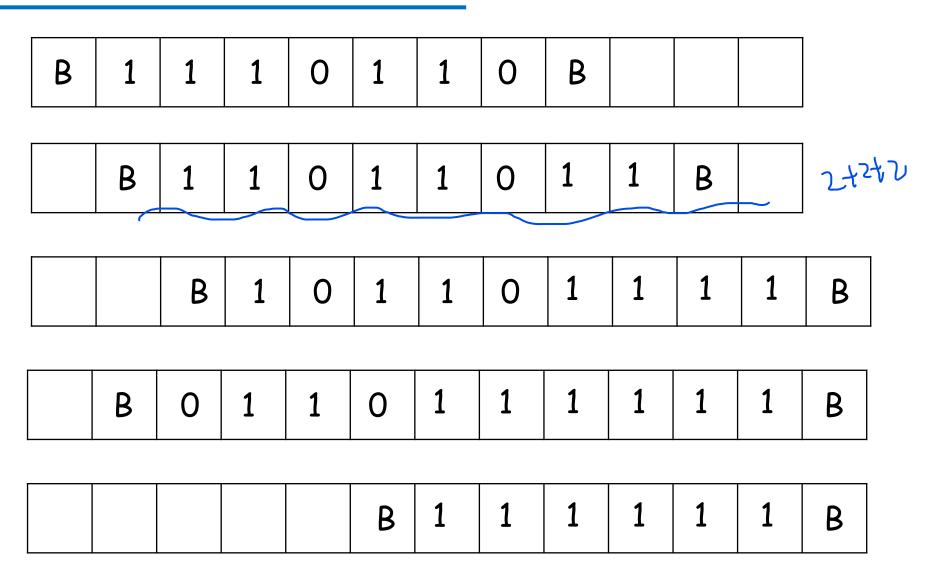
Construct a TM to compute the function $\underline{f(w)=ww}$ where $w \in \{0, 1\}^+$.

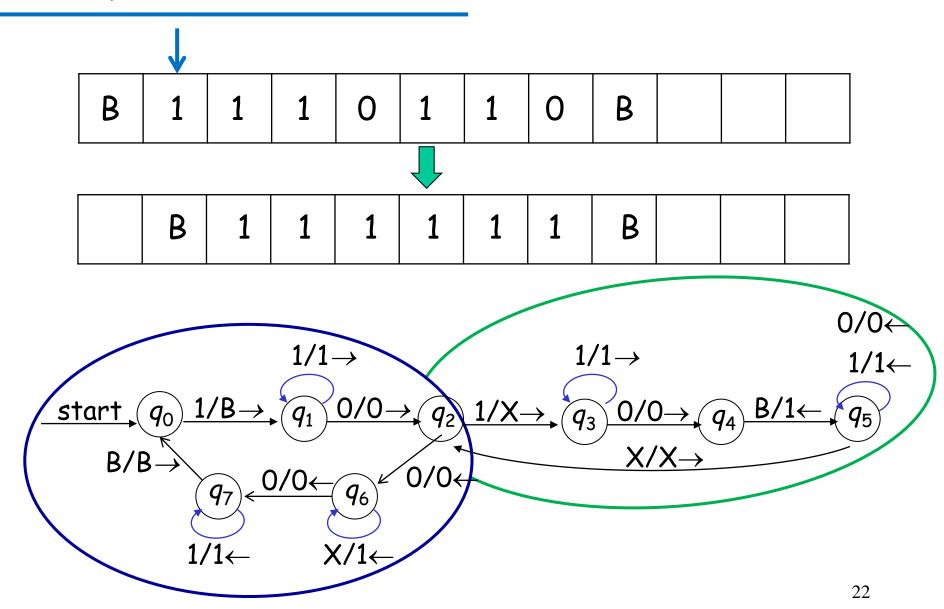


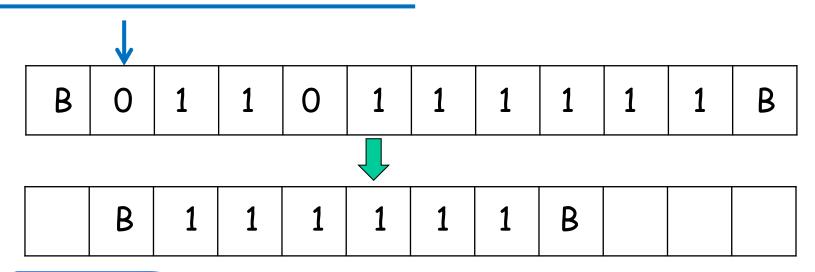
Construct a TM to compute $m \times n$.

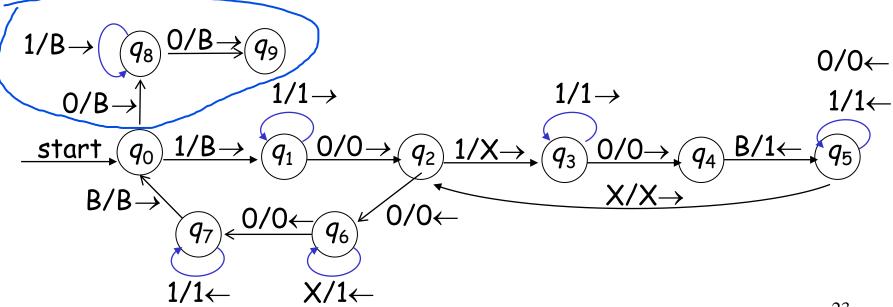


$$3 \times 2 \Rightarrow 2 + 2 + 2$$

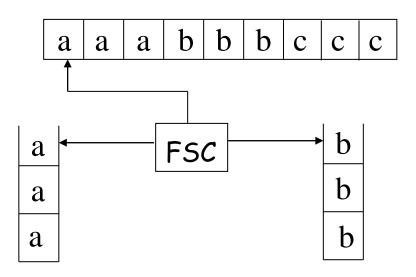






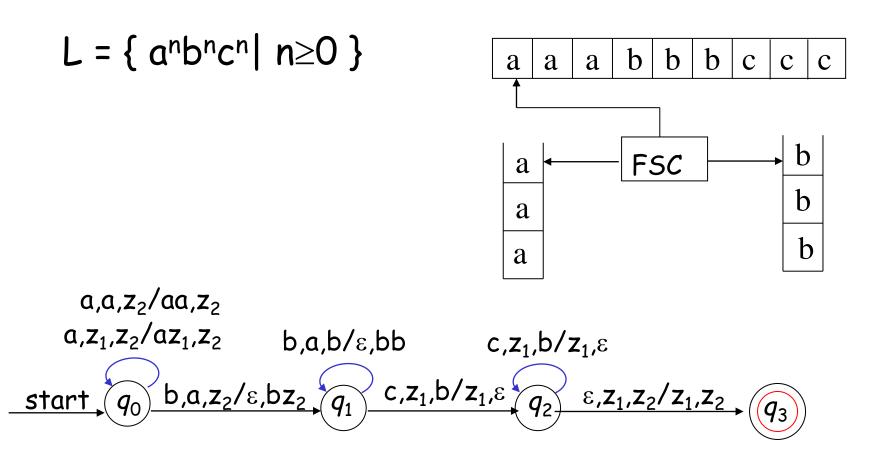


Two Stack Machine



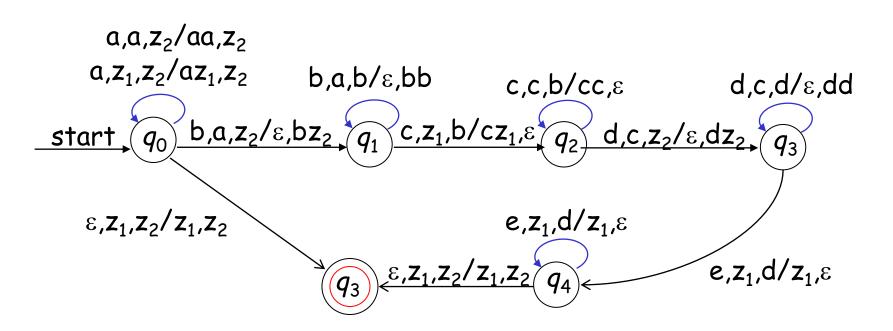
$$\delta(q, a, X, Y) = (p, \alpha, \beta)$$

Construct a two stack machine for



Construct a two stack machine for

$$L = \{ a^n b^n c^n d^n e^n | n \ge 0 \}$$



Good good still day day up

Never confuse education with intelligence.

Intelligence isn't ability to remember and repeat like they teach you in school.

Intelligence is the ability to learn from experience, solve problems, and use our knowledge to adapt to new situation.

—— Richard Feynman