Best/Worst/Always and Asymptotic Notation

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  - $n(n-1)/2$ in the worst case
  - $\geq n - 1$ and $\leq n(n-1)/2$ on every input

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- The number of comparisons that InsertionSort makes on inputs of size n is:
  - $\Theta(n)$ in the best case
  - $\Theta(n^2)$ in the worst case
  - $\Omega(n)$ and $O(n^2)$ on every input

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- Typically we just say: InsertionSort makes $O(n^2)$ comparisons (suggesting implicitly that there is no better bound for the worst case).