

Machine Architecture - Lecture 6

```
# Get n from user and save
                               $v0, 5
                                          # read integer
               syscall
                               $t0, $v0 # syscall result
               move
# Initialize registers
               1i
                               $t1, 0 # initialize counter i
               li
                               $t2, 0 # initialize sum
               # Main loop body
                               $t1, $t1, 1
loop:
                               $t2, $t2, $t1
               add
                               $t0, $t1, exit
                                                  # break from loop
               beq
exit:
               # Print sum
                               $v0, 1
                                           # print string
                               $a0, $t2
               move
               syscall
               # Exit
               li.
                               $v0, 10
                                          # exit
               syscall
```

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MIPS – programming

Slide material acknowledgements: Magnus Bordewich

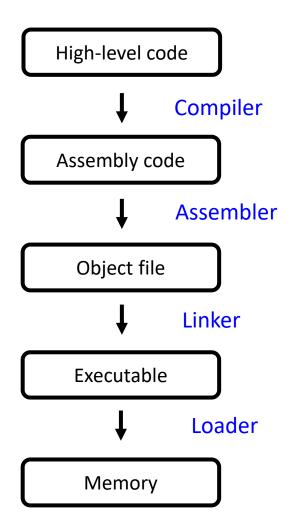


Compiling, assembling, and loading

Translate a program from a high-level language into machine language and start executing it



Translating and starting a program



The compiler translates the high-level code into assembly language.

The assembler turns the assembly code into machine language code (object file).

The linker combines the object file with other machine language code (e.g. from already compiled and assembled libraries).

The loader puts the executable into the memory.



Assembler

main:	addi sw addi sw addi sw	\$ra, \$a0, \$a0, \$a0, \$sw,	\$sp, -4 0(\$sp) \$0, 2 f \$0, 3 \$a1, g
sum:	jal sw lw addi jr add jr	<pre>\$v0, \$ra, \$sp, \$ra</pre>	y 0(\$sp) \$sp, 4 \$a0, \$a1

The assembler makes two passes through the assembly code and turns it into the object file.

On the first pass, the assembler assigns instruction addresses and finds all symbols, such as labels and global variable names, and makes a symbol table.

In this example, symbols are shown in bold red.



Assembler

0x00400000	main:	addi	\$sp,	\$sp, -4
0x00400004		SW	\$ra,	0(\$sp)
0x00400008		addi	\$a0,	\$0 , 2
0x0040000C		SW	\$a0,	f
0x00400010		addi	\$a0,	\$0, 3
0x00400014		SW	\$sw,	\$a1, g
0x00400018		jal	sum	
0x0040001C		SW	\$v0,	У
0x00400020		lw	\$ra,	0(\$sp)
0x00400024		addi	\$sp,	\$sp, 4
0x00400028		jr	\$ra	
0x0040002C	sum:	add	\$v0,	\$a0, \$a1
0x00400030		jr	\$ra	

symbol	address
f	0x10000000
þ	0x10000004
У	0x10000008
main	0x00400000
sum	0x0040002C

After the first pass of the assembler, instruction addresses have been assigned and the symbol table has been created.



Assembler, linker and loader

0x80000000	Reserved		
0 x 10010000	Dynamic Data		
0X10010000	:	← \$gb=0x10008000	
0x10000008	variable y		
0x10000004	variable g		
0x10000000	variable f		
	:		
	0x03E00008		
	0x00851020		
	0x03E00008		
	0x23BD0004		
	0x8FBF0000		
	0xAF828008		
	0x0C10000B		
	0xAF858004		
	0x20050003		
	0xAF848000		
0×00400008	0x20040002		
0×00400004	0xAFBF0000		
0×00400000	0x23BDFFFC	\leftarrow PC=0x00400004	
0x00000000	Reserved		

The second pass of the assembler generates the object file.

The linker creates the executable by combining the object file with other machine language code, e.g. code corresponding to libraries.

The loader puts the executable into the memory and its execution can start.



MIPS programming

To be read in conjunction with the material in the practical of Week 9



Functions

Functions are pieces of code which can be accessed from other parts of the program.

Using functions makes code more modular and readable.

MIPS assembly functions have inputs, called arguments, and an output called return value.

Need an agreement on how to:

call and return from a function

access the input arguments and the return value



Call and return

```
# The caller function "main"

0x00400200 main: jal simple

0x00400204 ....

# The callee function "simple"

0x00401020 simple: jr $ra
```

The function main calls the function simple using the jal instruction.

jal simple (jump and link) jumps to the address 0x00401020 (as j would do) but also stores in register \$ra the address where the program should return after simple has been executed (here 0x00400204).

jr \$ra (jump register) jumps to the address stored in a register (here \$ra). Notice, that jr is an R-type not a J-type instruction.



Arguments and return value

```
main:
        1i
           $a0, 10
                                # argument 0 gets the value 10
        li $a1, 5
                                # argument 1 gets the value 5
        li $a2, 20
                                # argument 2 has value 20
        li $a3, 10
                                # argument 3 has value 10
                                # call the function
        jal diffofsums
        move $s0, $v0
                                # put return value to $s0
diffofsums:
        add $t0, $a0, $a1 # sum of the first two arguments
        add $t1, $a2, $a3 # sum of the other two arguments
        sub $t2, $t0, $t1
                                # difference of the two sums
        move $v0, $t2
                                # put return value at $v0
                                # return to caller
          $ra
        jr
```

move \$s0, \$v0 is another pseudo-instruction like li. It copies the value of a register into another register. It is implemented as:

```
add $s0, $v0, $0
```



Arguments and return value

According to MIPS conventions on the behaviour of caller and callee:

the caller places the arguments into the registers \$a0 - \$a3

the return value is placed into the registers \$v0 - \$v1

the saved registers \$s0 - \$s7 are not modified by the callee

This convention can be quite restrictive, especially if the callee is going to call another function (or even call itself recursively).

Instead of directly conforming to this convention, the callee can first save all important registers in a stack (a data structure covered in the ADS module) and restore them before returning to the caller.

Machine Architecture (Ioannis Ivrissimtzis)



Loops

Exercise: Compute the n-th triangular number.

The program should take the input n and output T(n) where:

$$T(n) = 1 + 2 + ... + n$$



Loops

The code fragment shows register initialisation and the main loop.

```
# Initialize registers
      li
            $t0, 10 # load the value of N
      li $t1, 0
                         # initialize the counter (i)
      li $t2, 0
                       # initialize sum
      # Main loop body
     addi $t1, $t1, 1 #i = i + 1 (increment the counter)
loop:
      add $t2, $t2, $t1 # sum = sum + i
      beq $t0, $t1, exit # if i = N, break from the loop
      j
            loop
exit:
```



Loops

The main loop is implemented through an unconditional jump instruction jand a branch on equal instruction beq, which branches out of the loop when the values in the registers \$t0 and \$t1 become equal.

```
# Main loop body
loop: addi  $t1, $t1, 1  # i = i + 1 (increment the counter)
    add  $t2, $t2, $t1  # sum = sum + i
    beq  $t0, $t1, exit  # if i = N, break from the loop
    j  loop
exit: ...
```



Input / Output

Can we get the value of n from the user, through the keyboard, instead of hard encoding it?

The syscall instruction (system call) suspends the execution of the program to provide an operating-system-like service, such as input, output, termination.



Input

The type of the syscall service is specified by a code, which should be stored in \$v0.

The code 5 used here corresponds to reading an integer from keyboard.



Examples of syscall services

service	syscall code	arguments	result
print integer	1	\$a0 = integer	-
print string	4	\$a0 = string	-
read integer	5	-	integer (in \$v0)
exit	10	-	-



Output and exit

After exiting the main loop, we print the output and stop the execution.

It is important, always, to declare the end of the program.

Otherwise, the computer will fetch the word stored immediately after the last instruction and try to execute it with unpredictable behaviour.



Putting it all together ... so far

```
# Get n from user and save
               $v0, 5 # read integer (syscall code is 5)
       1 i
       syscall
               $t0, $v0 # syscall result (returned in $v0) move to $t0
       move
# Initialize registers
       1i $t1, 0
                               # initialize the counter (i)
             $t2, 0
       li
                                # initialize sum
       # Main loop body
       addi $t1, $t1, 1 # i = i + 1  (increment the counter)
loop:
       add $t2, $t2, $t1 # sum = sum + i
       beg $t0, $t1, exit # if i = N, break from the loop
       j loop
exit:
       # Print sum
                              # print string syscall code = 1
       1i $v0, 1
       move $a0, $t2
       syscall
       # Exit
               $v0, 10
       li
                              # exit
       syscall
```



Text output

To make the program more user friendly, we want to add some text output describing the format of the requested input and also what the output is.

See the program triangularNumber.asm uploaded on DUO (you will be able to run it on the MARS MIPS simulator after the practical of Week 9).