

# Machine Architecture - Lecture 8



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## Programming in AVR assembly

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# Programming in AVR assembly

Through a simple example we will study some basic aspects of AVR assembly programming :

- directives

- I/O

- the status register

- arithmetic and logical operations

- branching

Our program will wait for a signal (e.g. the press of a button) and then a LED light will blink three times. The code was uploaded in Ultra. It is a simplified version of the example in

<https://akuzechie.blogspot.com/2021/10/assembly-programming-using-arduino-uno.html>

# The C++ shell

An easy way to program in AVR assembly through the Arduino IDE, is to create a C++ shell file, and a separate linked file for the assembly code.

In the `setup()` function we will define the program's inputs and outputs.

The `loop()` function repeats for ever, which usually is the desirable behaviour in an embedded system.

```
//-----  
// C Code: LED blinking  
//-----  
extern "C"  
{  
    void start();  
    void btnLED();  
}  
//-----  
void setup()  
{  
    start();  
}  
//-----  
void loop()  
{  
    btnLED();  
}
```

# Directives

The assembly code contains some statements called **directives**.

Directives do not correspond to assembly instructions, but the assembler will take them into account when compiling the assembly code to machine language.

The use of directives, for example for defining a constant, makes easier to develop and maintain assembly code.

In AVR assembly a directive is precented by a dot (.).

# Directives

The `equ` directive assigns a value to a label, and the label can then be used in later expressions.

A label assigned a value by the `equ` directive is a constant, it cannot be redefined. We define a constant named `delayVal` with value 10000 by:

```
.equ delayVal, 10000
```

The directives `.global start` and `.global btnLED` declare the functions `start` and `btnLED` as global.

This information will be used by the linker when it puts various pieces of code from several source files together.

# Declaring input output pins

start:

```
SBI  DDRB, 4      ;set PB4   (declare pin PB4 as output pin (o/p) - LED)
CBI  DDRD, 2      ;clear PD2 (declare pin PD2 as input pin (i/p) - e.g. a button)
RET              ;return to setup()
```

First, we choose a pin for output (to light the LED light), and one pin for input (to receive the signal for the light to start blinking). Here we chose pins PB4 and PD2, respectively.

The same example was used in lecture 7.

# Reading values from the input pin

btnLED:

SBIC PIND, 2	;skip next statement if button not pressed
RJMP blink	;jump to label blink
RJMP btnLED	;return to label btnLED

The status of the input pin PD4 is recorded on register **PIND** (lecture 7).

The SBIC instruction (**s**kip if **b**it in I/O register **i**s **c**leared) checks the corresponding bit in PIND and if it is 0 (cleared) skips one instruction.

So, if the bit (PIND,2) has value 0 (no incoming electric signal from pin PD2), we skip one instruction and go to the jump with **RJMP btnLED**, that is, to the beginning of the function to check again the value of that bit.

Otherwise, we jump with **RJMP btnLED** to the function that will blink the LED three times.

# Controlling the output pin

LDI initialises the counter R21 to 3, so it would blink three times.

Pin PB4 has already been declared as output pin (bit **DDRB,4** set). Its behaviour is controlled by the bit 4 of the corresponding I/O register **PORTB**.

When the value of bit **PORTB,4** is 1 (set), the microcontroller sends electric signal through pin PB4 and the LED is on. When it is 0 (cleared) no electric signal is sent, and the light is off.

```
blink:
    LDI  R21, 3           ;initialise register R21

blinkOnce:
    SBI  PORTB, 4         ;turn ON LED
    RCALL myDelay         ;call subroutine myDelay
    CBI  PORTB, 4         ;turn OFF LED
    RCALL myDelay         ;call subroutine myDelay
    SUBI R21, 1           ;decrement counter by 1
    BRNE blinkOnce       ;loop if counter not zero
    RJMP btnLED          ;return to label btnLED
```



# Controlling the output pin

The value of the bit `PORTB,4` is controlled (set or cleared) with the instructions `SBI` and `CBI`, which we have seen already.

In between, the program calls the function `myDelay`, which does some computations (countdown) just to waste some time. Otherwise, the blinking would be so rapid that our eyes wouldn't be able to notice it.

```
blink:
    LDI  R21, 3           ;initialise register R21

blinkOnce:
    SBI  PORTB, 4         ;turn ON LED
    RCALL myDelay         ;call subroutine myDelay
    CBI  PORTB, 4         ;turn OFF LED
    RCALL myDelay         ;call subroutine myDelay
    SUBI R21, 1           ;decrement counter by 1
    BRNE blinkOnce        ;loop if counter not zero
    RJMP btnLED           ;return to label btnLED
```

# Decrementing the counter

The counter is decremented with the SUBI instruction (subtract immediate).

Apart from subtracting the immediate 1 from register R21, SUBI will also check the result of this operation and update certain bits of the status register, see lecture 7.

In particular, if the result of the subtraction is 0, the value of the Z bit of the status register will become 1 (set). Otherwise, it will become 0 (cleared).

```
blink:
    LDI  R21, 3           ;initialise register R21

blinkOnce:
    SBI  PORTB, 4         ;turn ON LED
    RCALL myDelay         ;call subroutine myDelay
    CBI  PORTB, 4         ;turn OFF LED
    RCALL myDelay         ;call subroutine myDelay
    SUBI R21, 1           ;decrement counter by 1
    BRNE blinkOnce       ;loop if counter not zero
    RJMP btnLED          ;return to label btnLED
```

# SUBI instruction

Subtract Immediate

Operation:  $Rd \leftarrow Rd - K$

Syntax: SUBI Rd, K

Operands:  $16 \leq d \leq 31$ ,  $0 \leq K \leq 255$

Machine language code:

0101	KKKK	dddd	KKKK
------	------	------	------

The Z bit of the status register is set (value 1) when the result  $Rd - K$  is equal to 0; it is cleared (value 0) otherwise.

I	T	H	S	V	N	Z	C	bit name
								value

# Conditional branching

The **BRNE** instruction (Branch if Not Equal), branches to **blinkOnce** at the top of the loop, if and only if the Z bit of the status counter is 1.

Despite its name, BRNE does not compare numbers. It only checks the Z bit in the status register. The idea is that the value of the Z bit represents the result of a comparison done by the previous instruction, here SUBI.

Finally, the **RJMP**, at the exit of the blink loop, branches back to the **btnLED**, that is, the microcontroller waits again for an input signal.

```
blink:
    LDI  R21, 3           ;initialise register R21

blinkOnce:
    SBI  PORTB, 4         ;turn ON LED
    RCALL myDelay         ;call subroutine myDelay
    CBI  PORTB, 4         ;turn OFF LED
    RCALL myDelay         ;call subroutine myDelay
    SUBI R21, 1           ;decrement counter by 1
    BRNE blinkOnce        ;loop if counter not zero
    RJMP btnLED           ;return to label btnLED
```

# The myDelay function

```
.equ delayVal, 10000      ;equate delayVal with initial count value

myDelay:
    LDI R20, 100          ;initial count value for outer loop
outerLoop:
    LDI R30, lo8(delayVal) ;low byte of delayVal in R30
    LDI R31, hi8(delayVal) ;high byte of delayVal in R31
innerLoop:
    SBIW R30, 1            ;subtract 1 from 16-bit value in R31, R30
    BRNE innerLoop        ;jump if countVal not equal to 0

    SUBI R20, 1            ;subtract 1 from R20
    BRNE outerLoop        ;jump if R20 not equal to 0
    RET
```

Two nested loops counting down from 10000 and 100, respectively.

The **SUBI** instruction and the similar **SBIW** (subtract immediate from word) will be called  $10,000 \times 100 = 1,000,000$  times in total.

# The myDelay function

```
.equ delayVal, 10000      ;equate delayVal with initial count value

myDelay:
    LDI R20, 100          ;initial count value for outer loop
outerLoop:
    LDI R30, lo8(delayVal) ;low byte of delayVal in R30
    LDI R31, hi8(delayVal) ;high byte of delayVal in R31
innerLoop:
    SBIW R30, 1           ;subtract 1 from 16-bit value in R31, R30
    BRNE innerLoop        ;jump if countVal not equal to 0
    SUBI R20, 1           ;subtract 1 from R20
    BRNE outerLoop        ;jump if R20 not equal to 0
    RET
```

Notice the split of the iterations as  $10000 \times 100$  rather than  $1000^2$ .

This way the **outerLoop** counter is stored in a word (pair of registers R30 and R31), while the counter of **innerLoop** in a byte, in register R20.

Assembly programming affords this type of optimisations.

# Exercise

Can you compute (at least approximately) the duration of the delay introduced by the `myDelay` function?

**Hint:** First, find how many times each instruction will be executed.

Then, find how many clock cycles will be required using the following table:

<code>LDI</code>	1 cycle
<code>SUBI</code>	1 cycle
<code>SBIW</code>	2 cycles
<code>BRNE</code>	1 cycle if condition is false 2 cycles if condition is true
<code>RET</code>	4 cycles

Finally, compute the timing knowing that ATmega328 runs at 16MHz.

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Next ...

... addressing modes  
and the  
instruction set