

COMP1071 - Digital Electronics

Karnaugh Maps

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Slide acknowledgements: Eleni Akrida and Farshad Arvin

Overview of today's lecture

Karnaugh Maps: a more systematic way to simplify Boolean formulae

Example: 7-segment display driver



Input A	Input B	Input C	Input D	Output F(A,B,C,D)
0	0	0	0	1
0	0	0	1	0
0	0	1	0	1
\bigcirc 0	1	0	0	
0	0	1	1	0
O	1	0	1	1
O	1	1	0	1
0	1	1	1	1
≤ 1	0	0	0	1
\bigcirc 1	0	0	1	1
≤ 1	0	1	0	1
1	1	0	0	0
1	0	1	1	0
$\bigcirc 1$	1	0	1	1
1	1	1	0	0
1	1	1	1	0



$$F = \overline{A} \cdot \overline{B} \cdot \overline{C} \cdot \overline{D} + \overline{A} \cdot \overline{B} \cdot C \cdot \overline{D} + \overline{A} \cdot B \cdot \overline{C} \cdot \overline{D} + \overline{A} \cdot B \cdot \overline{C} \cdot D + \overline{A} \cdot B \cdot C \cdot \overline{D} + \overline{A} \cdot B \cdot \overline{C} \cdot D + \overline{A} \cdot B \cdot \overline{C} \cdot D + \overline{A} \cdot B \cdot \overline{C} \cdot D + \overline{A} \cdot \overline{B} \cdot \overline{C} \cdot \overline{D} + \overline{A} \cdot \overline{D} \cdot \overline{D} + \overline{A} \cdot \overline{D} \cdot \overline{D} + \overline{A} \cdot \overline{D} \cdot \overline{D} \cdot \overline{D} + \overline{A} \cdot \overline{D} \cdot \overline{D} \cdot \overline{D} + \overline{A} \cdot \overline{D} \cdot \overline{D} + \overline{A} \cdot \overline{D} \cdot \overline{D} + \overline{A} \cdot \overline{D} \cdot \overline{D} \cdot \overline{D} + \overline{A} \cdot \overline{D} \cdot \overline{D} \cdot \overline{D} + \overline{A} \cdot \overline{D} \cdot \overline{D} \cdot \overline{D} + \overline{D} \cdot \overline{D} \cdot \overline{D} \cdot \overline{D} \cdot \overline{D} + \overline{D} \cdot \overline{D} \cdot \overline{D} \cdot \overline{D} \cdot \overline{D} \cdot \overline{D} + \overline{D} \cdot \overline{D}$$

Simplifying Boolean expressions

Key to simplifying is spotting terms of the form $PA + P\overline{A}$ (since this = P).

Karnaugh Maps are a graphical way of representing equations to make spotting these terms easier.

Α	В	C	Y	Y Al	D.				Y	_			
0	0	0 1	1	C	00	01	11	10	C^{A}	00 B	01	11	10
0	1	0	0	0	1	0	0	0	0	ABC	_ ABC	ABC	ABC
1	0	0	0	-									
1 1	0	1 0	0	1	1	0	0	0	1	ABC	_ ABC	ABC	ABC
1	1	1	0	L									
(a)				(b)					(c)				

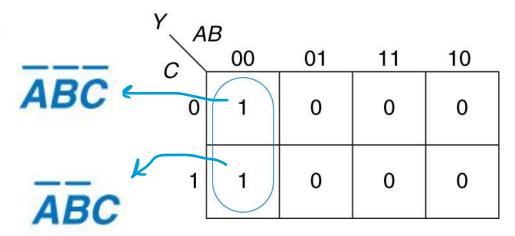
Each cell represents a minterm, and has a zero or one depending on the value of Y corresponding to that minterm. SoP form is given by adding the minterms corresponding to 1s.

The order of minterms is such that each cell differs in the negation of exactly one variable from its neighbours (including wrap around).



Simplifying Boolean expressions

Terms of the form $PA + P\overline{A}$ are neighbouring 1s in the K-map.



Rather than writing out full SoP by taking every 1 as a term, we circle neighbouring 1s, and use the single reduced implicant for both 1s in the circle.

Full SoP: $Y = \overline{A}\overline{B}\overline{C} + \overline{A}\overline{B}C$

Reduced: $Y = \overline{AB}$



Karnaugh Maps

- Create the map so that neighbouring terms differ in the negation of one variable (including wrap around).
- 2. Circle *exactly* all **ones** in the map using as **few circles** as possible, and making each circle as **large** as possible.
- 3. Each circle must span a rectangular block that is a power of 2 in each dimension (i.e. 1,2,4).
- 4. Read off the implicants that were circled.

If a Boolean expression is minimal then it is the sum of **prime implicants**: implicants that cannot be combined with each other.

Each circle represents an implicant. Largest possible circles represent prime implicants.

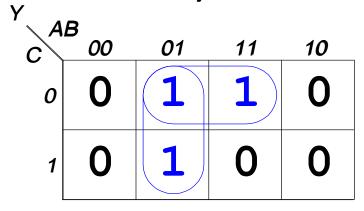


Example

Truth Table

A_	В	C	Y
0	0	0	0
0	0	1	0
0	1	0	1
0	1	1	1
1	0	0	0
1	0	1	0
1	1	0	1
1	1	1	0

K-Map



SoP form is:

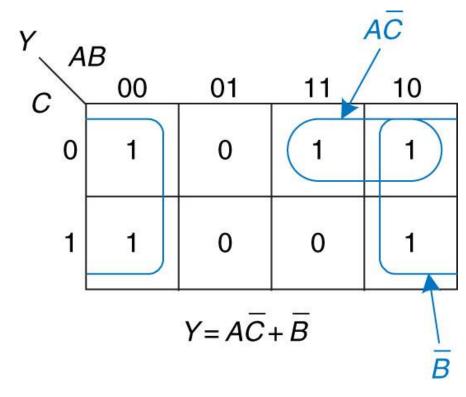
$$Y = \bar{A}B\bar{C} + \bar{A}BC + AB\bar{C}$$

K-map gives:

$$Y = \bar{A} B + B \bar{C}$$



Example



SoP:
$$Y = \overline{ABC} + \overline{ABC} + \overline{ABC} + \overline{ABC} + \overline{ABC} + \overline{ABC}$$

Using the K-map: Circle 1s.

Note that we can wrap around.

We cannot do a 3-by-1 rectangle.

Can still cover the 1s with only 2 rectangles.



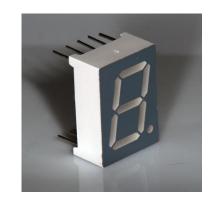
Overview of today's lecture

Karnaugh Maps: a more systematic way to simplify Boolean formulae

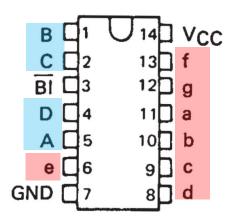
Example: 7-segment display driver



SN7449 BCD to 7-SEGMENT Decoder/Driver

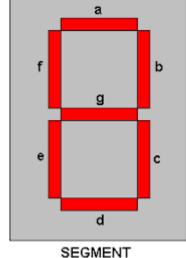


SN54LS49 . . . J OR W PACKAGE SN74LS49 . . . D OR N PACKAGE (TOP VIEW)



Binary Inputs	Decoder Outputs	7 Segment Display Outputs
D C B A	abcdefg	
0 0 0 0	1 1 1 1 1 1 0	0
0 0 0 1	0 1 1 0 0 0 0	1
0 0 1 0	1 1 0 1 1 0 1	2
0 0 1 1	1 1 1 1 0 0 1	3
0 1 0 0	0 1 1 0 0 1 1	4
0 1 0 1	1 0 1 1 0 1 1	5
0 1 1 0	1 0 1 1 1 1 1	6
0 1 1 1	1 1 1 0 0 0 0	7
1 0 0 0	1 1 1 1 1 1 1	8
1 0 0 1	1 1 1 1 0 1 1	9

$$D_3, D_2, D_1, D_0$$



SEGMENT IDENITIFICATION

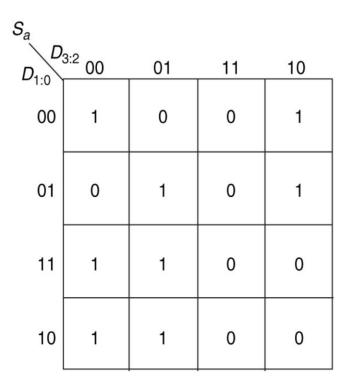


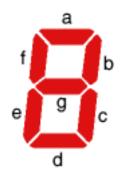


4 inputs D₃, D₂, D₁, D₀ (written D_{3:0}). 7 outputs.

Input represents 4-bit binary number.

Output should show corresponding decimal.

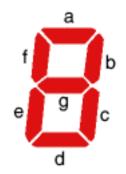


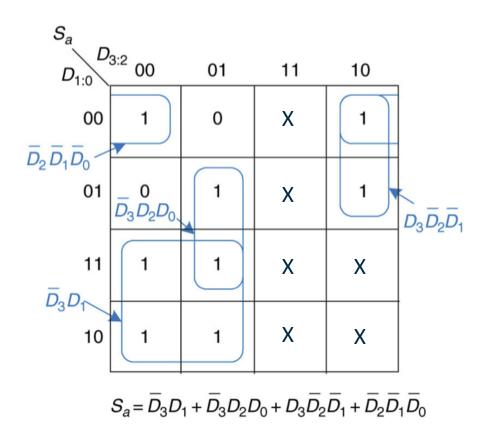


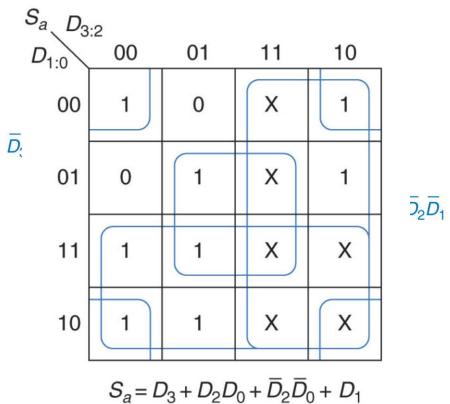
D_3	D ₂	D_1	D_0	S _a
0	0	0	0	1
0	0	0	1	0
0	0	1	0	1
0	0	1	1	1
0	1	0	0	0
0	1	0	1	1
0	1	1	0	1
0	1	1	1	1
1	0	0	0	1
1	0	0	1	1

Other inputs' output not specified





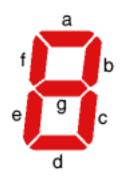


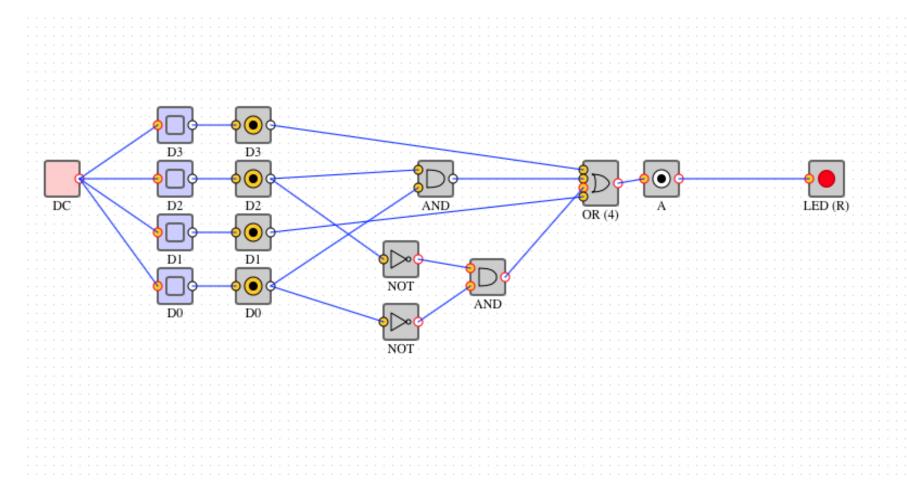


What about the unspecified inputs? Have so far assumed they are 0s.

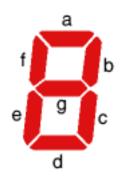
We could equally output 1s if it helped us reduce circuitry!

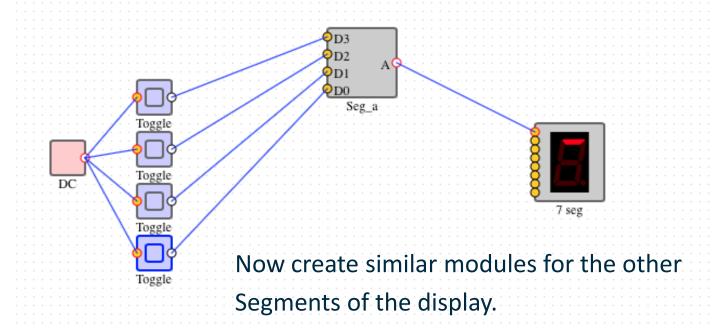




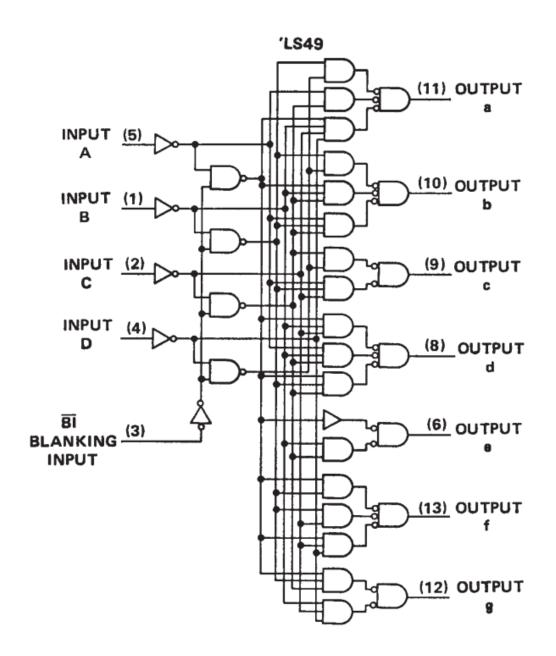














Note: Karnaugh Maps for PoS minimisation

1. Exact same creation of K-map

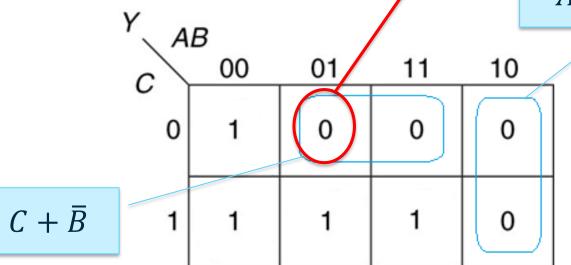
 $A + \overline{B} + C$

- 2. Circle the **zeros** in the map using the same rules:
 - you group 0s to produce minimum sum terms instead of grouping 1s to produce minimum product terms
 - Each cell corresponds to a sum term that is false in that cell

 $\bar{A} + B$

Resulting PoS:

$$(C + \overline{B}) \cdot (\overline{A} + B)$$





Summary

Karnaugh Maps: a more systematic way to simplify Boolean formulae

Example: 7-segment display driver

