

Graph Algorithms II



Agenda

01

Shortest path
คืออะไร, Algorithm

02

All pair shortest path
คืออะไร, Algorithm

03

Topological sort
คืออะไร, Algorithm

04

โจทย์
ตัวอย่างข้อสอบกราฟ

Shortest path

คืออะไร, Algorithm

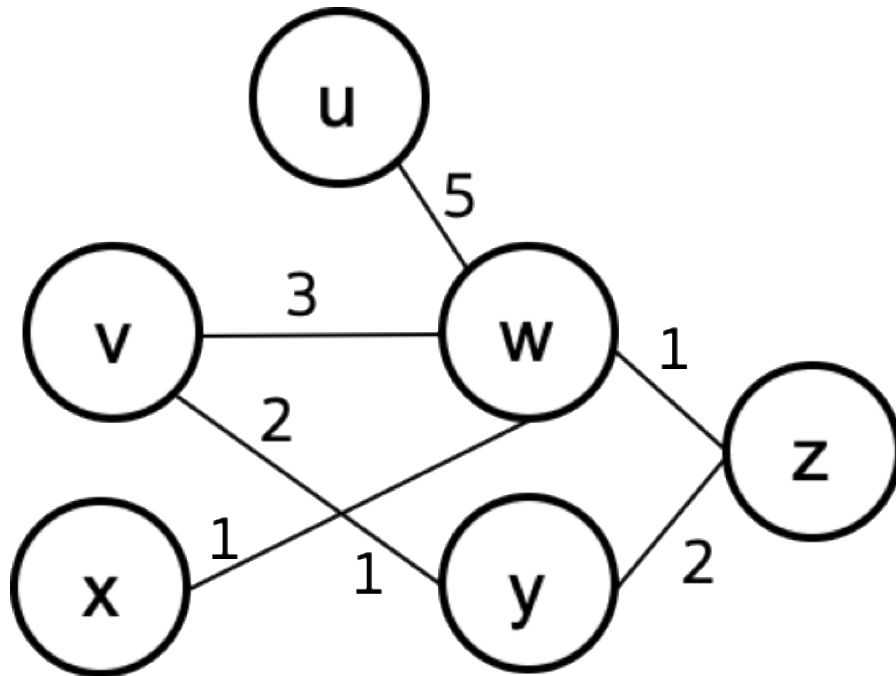


Shortest path of Graph

คืออะไร, Algorithm



- มีจุดเริ่มต้น (source) 1 จุด
- ถ้ามองว่า เส้นทางใด ไกลที่สุด / ใช้ระยะทางน้อยสุด / weight ต่ำที่สุด ในการไปถึง destination



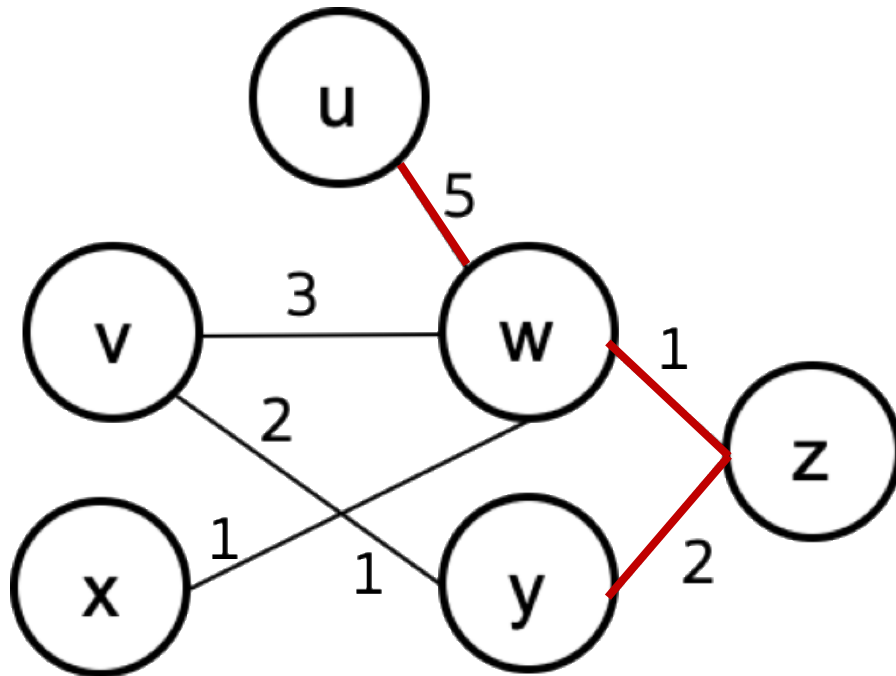
จงหาเส้นทางที่ใกล้ที่สุด จาก u ไป y

Shortest path of Graph

คืออะไร, Algorithm



- มีจุดเริ่มต้น (source) 1 จุด
- ถ้ามองว่า เส้นทางใด ไกลที่สุด / ใช้ระยะทางน้อยสุด / weight ต่ำที่สุด ในการไปถึง destination



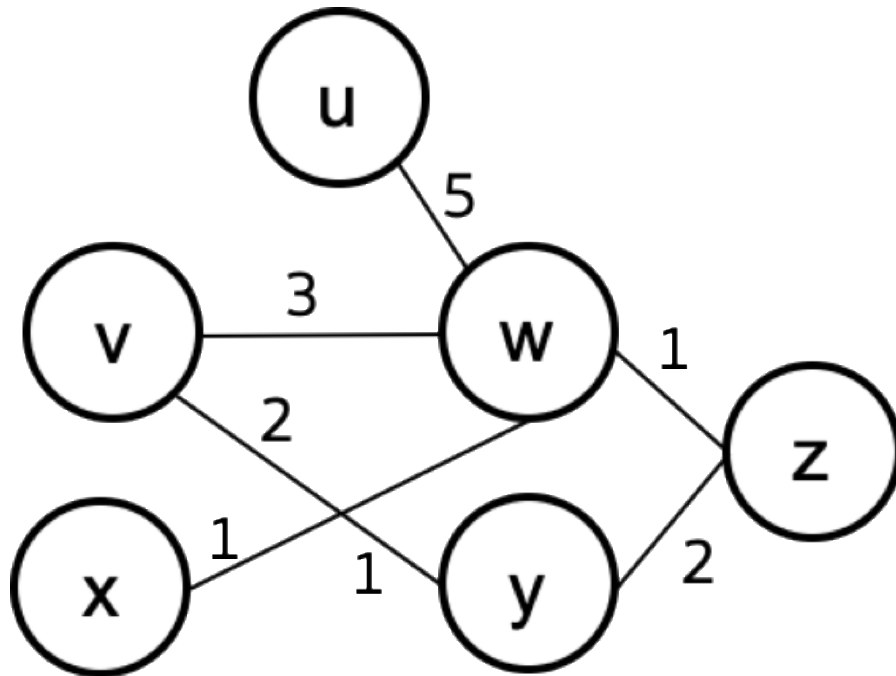
จงหาเส้นทางที่ใกล้ที่สุด จาก u ไป y

Shortest path of Graph

คืออะไร, Algorithm



- หรือ ถามว่า ระยะทางที่ใกล้ที่สุด จากจุดเริ่มต้น ไปยังทุกจุด เป็นเท่าใด (สร้าง Shortest Path Tree)



จงหาระยะทางที่สั้นที่สุด
ในการเดินไปให้ครบทุก node เมื่อเริ่มจาก u

u -> v -

u -> w -

u -> x -

u -> y -

u -> z -

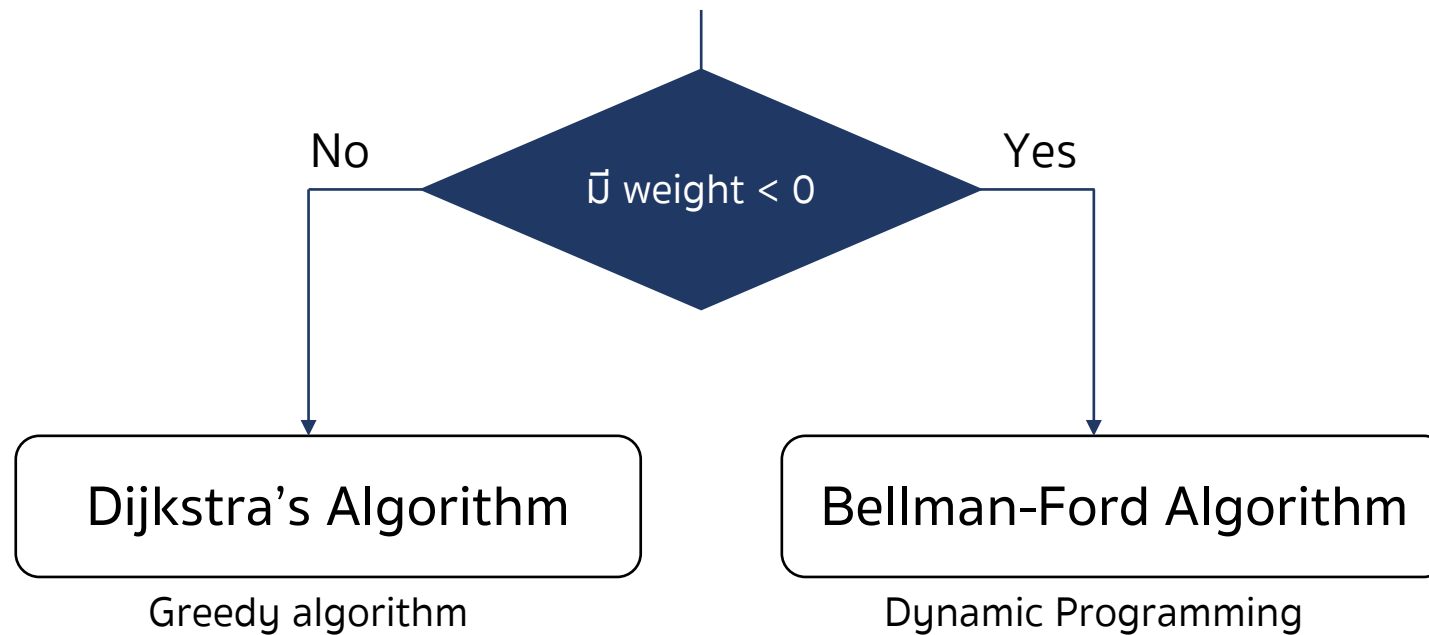
(ผลลัพธ์จากการสร้าง Shortest Path Tree)

Shortest path of Graph

คืออะไร, Algorithm

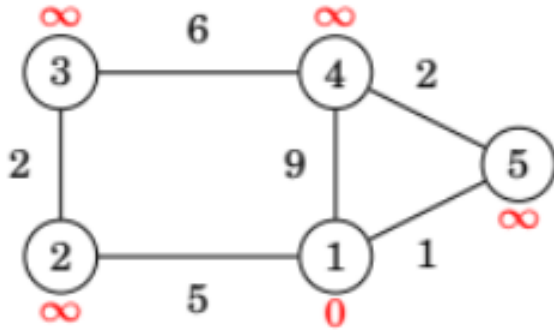


- วิธีการหา Shortest Path / Shortest Path tree

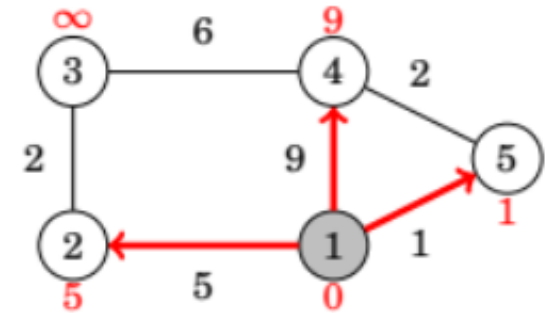


Shortest path of Graph

Dijkstra's Algorithm



ในแต่ละรอบ Dijkstra's algorithm จะเลือกโหนดที่ยัง
ไม่ถูกพิจารณาและมีระยะทางใกล้ที่สุด
ในตัวอย่างโหนดแรกเป็นโหนด 1 ซึ่งมีระยะทางเป็น 0

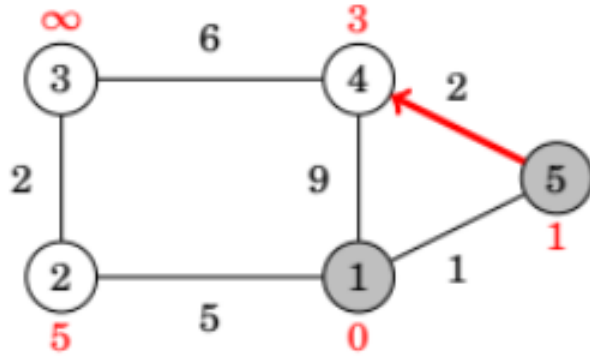


เมื่อโหนดถูกเลือก algorithm จะพิจารณาทุกเส้น
เชื่อมที่เริ่มต้นที่โหนดนั้นและลดระยะทาง

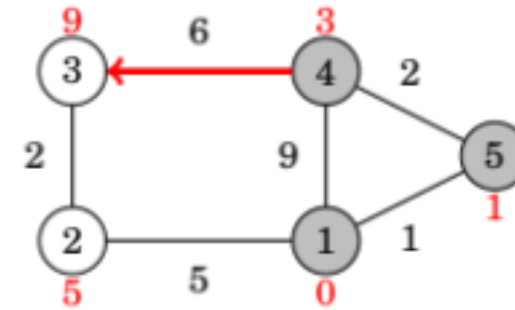
ในกรณีนี้เส้นเชื่อมจากโหนด 1 ลดระยะทางของโหนด
2, 4 และ 5 ซึ่งระยะทางปัจจุบันเป็น 5, 9 และ 1

Shortest path of Graph

Dijkstra's Algorithm



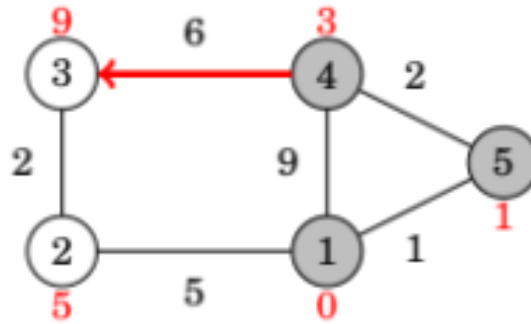
โหนดต่อไปที่จะประมวลผลคือโหนด 5 ที่มีระยะทางเป็น 1 ซึ่งจะไปลดระยะทางที่ไปยังโหนด 4 จาก 9 เป็น 3



หลังจากนั้นโหนดต่อไปคือโหนด 4 จะลดระยะทางจาก 3 เป็น 3

Shortest path of Graph

Dijkstra's Algorithm

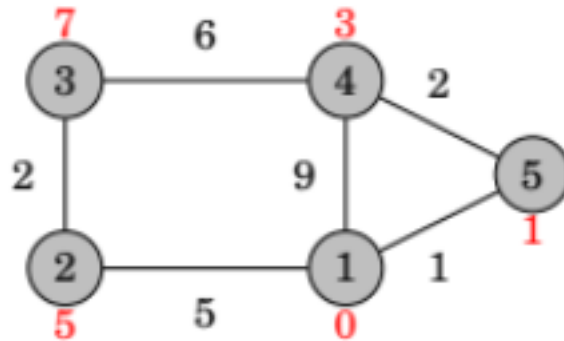


คุณสมบัติของ Dijkstra's algorithm คือเมื่อโหนดถูกเลือกแล้ว ระยะทางของมันจะถือว่าสิ้นสุดแล้ว

ตัวอย่างเช่นที่จุดนี้ ระยะทาง 0 1 และ 3 เป็นระยะทางสุดท้ายของโหนด 1, 5 และ 4

Shortest path of Graph

Dijkstra's Algorithm



หลังจากนั้น algorithm จะทำงานกับ 2 โหนดที่เหลือ ซึ่งได้ระยะทางสุดท้ายดังรูป

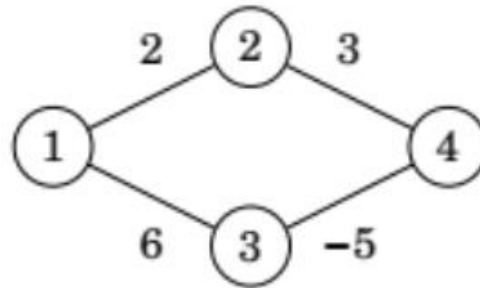
Shortest path of Graph

Dijkstra's Algorithm



ประสิทธิภาพของ Dijkstra's algorithm นั้นอยู่บนพื้นฐานของความจริงที่ว่าในกราฟไม่มีเส้นเชื่อมที่มีน้ำหนักเป็นลบ (negative edges)

ถ้ามีเส้นเชื่อมเป็นลบ algorithm อาจจะให้ผลลัพธ์ที่ผิดได้ ตัวอย่างเช่น



ระยะทางสั้นที่สุดจาก 1 ไปโหนด 4 คือ 1->3->4 ซึ่งความยาวเป็น 1 อย่างไรก็ตาม Dijkstra's algorithm หาเส้นทาง 1->2->4 ตามที่ค่าน้ำหนักน้อยที่สุดของเส้นเชื่อม

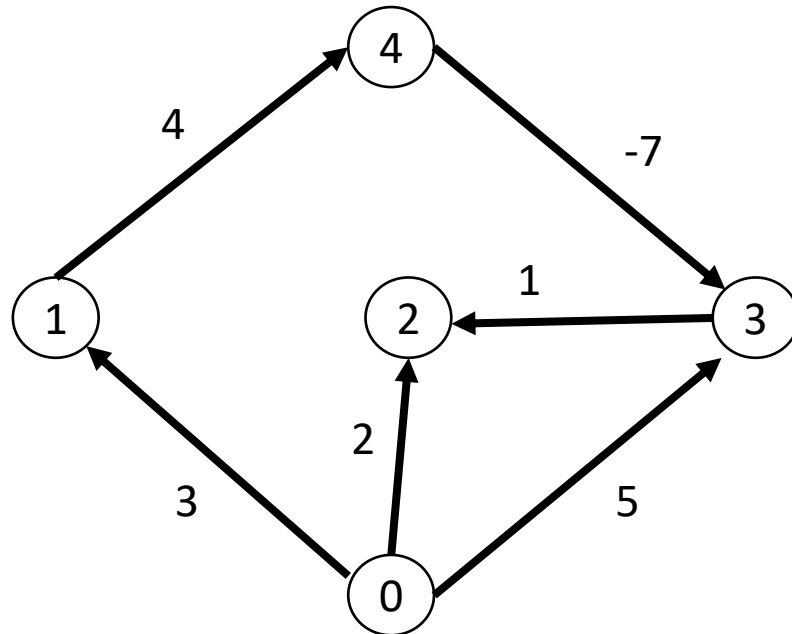
Shortest path of Graph

Dijkstra's Algorithm



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Shortest path of Graph

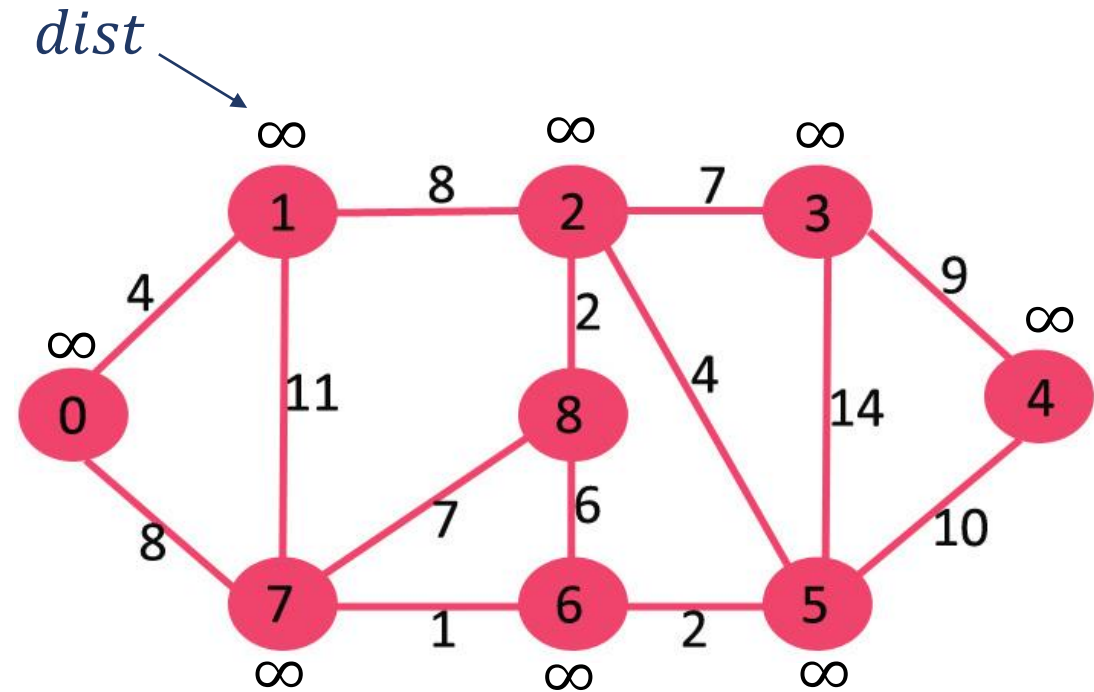
Dijkstra's Algorithm



- วิธีการสร้าง Shortest Path Tree (Dijkstra's Algorithm)

- Distance ที่สั้นที่สุด จาก src ถึงแต่ละ node
- ใช้ BFS with Priority queue

- Set up dist ทุกตัวเป็น INF
- เริ่มต้นจาก src $\rightarrow s = \text{src}$
- $\text{dist}[s] = 0$
- ถ้ายัง visited ไม่ครบทุกตัว
เลือก s ที่มี dist น้อยที่สุด
mark ว่า visit s แล้ว
For each node in $\text{Adj}(s)$
ถ้า $\text{dist}[s] + w(s \rightarrow \text{node}) < \text{dist}[\text{node}]$:
update $\text{dist}[\text{node}]$
- ค่า dist ที่ได้ของแต่ละ node คือ shortest distance from src



Shortest path of Graph

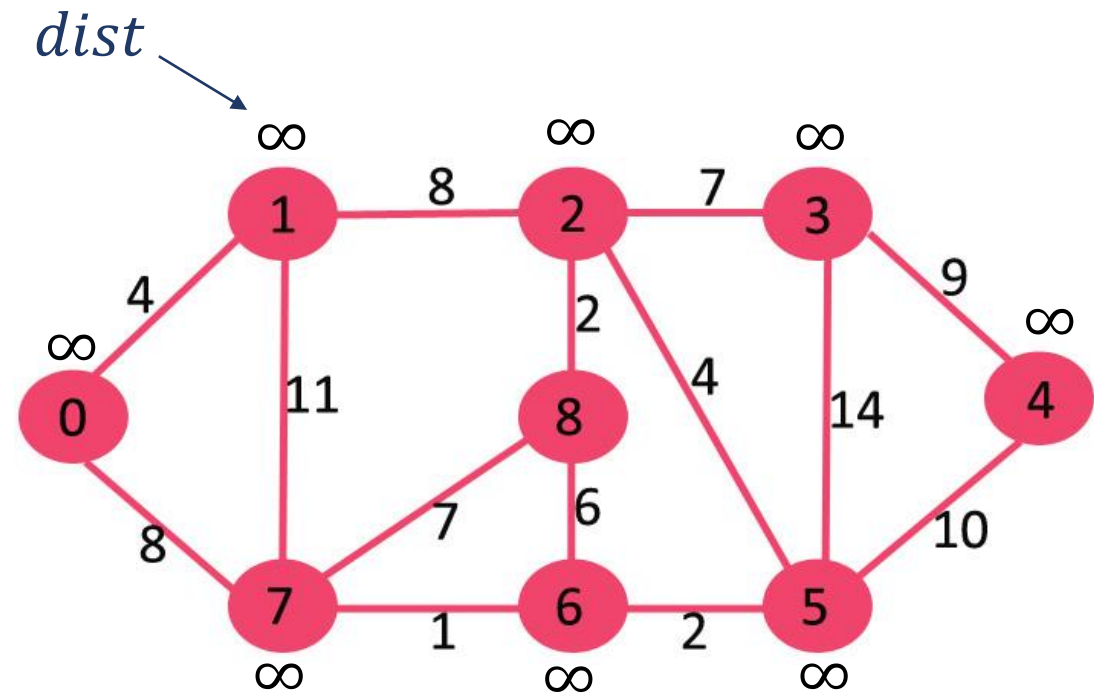
Dijkstra's Algorithm



- วิธีการสร้าง Shortest Path Tree (Dijkstra's Algorithm)

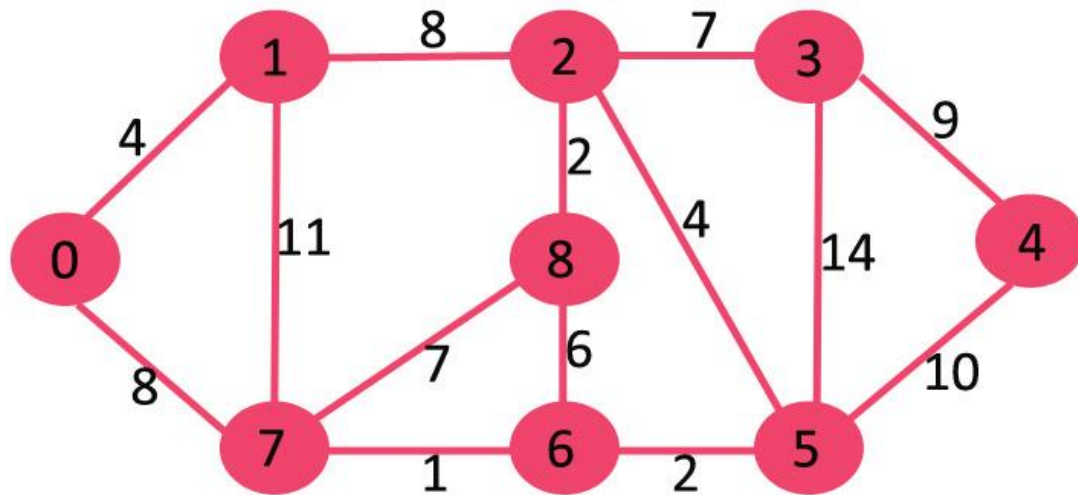
- Distance ที่สั้นที่สุด จาก src ถึงแต่ละ node
- ใช้ BFS with Priority queue

- เป็น BFS ยังไง?
 - ค่อยๆ อัปเดต dist ทีละขั้นตามแนวกว้าง
 - ใช้ Priority queue เพื่อเก็บ node ที่ต้องไปต่อแบบเรียงลำดับ dist จากน้อยไปมาก



Shortest path of Graph

Dijkstra's Algorithm



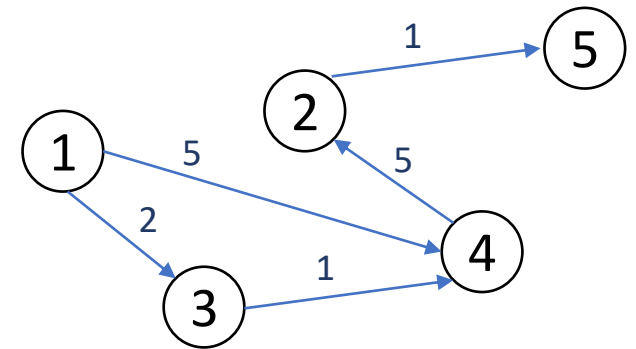
Let's code !!

Shortest path of Graph

Dijkstra's Algorithm



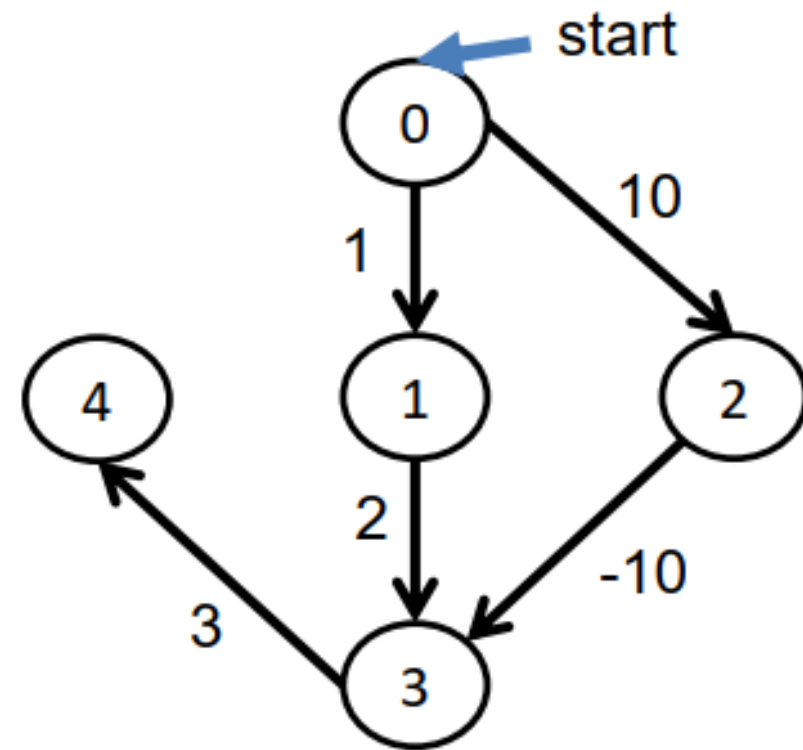
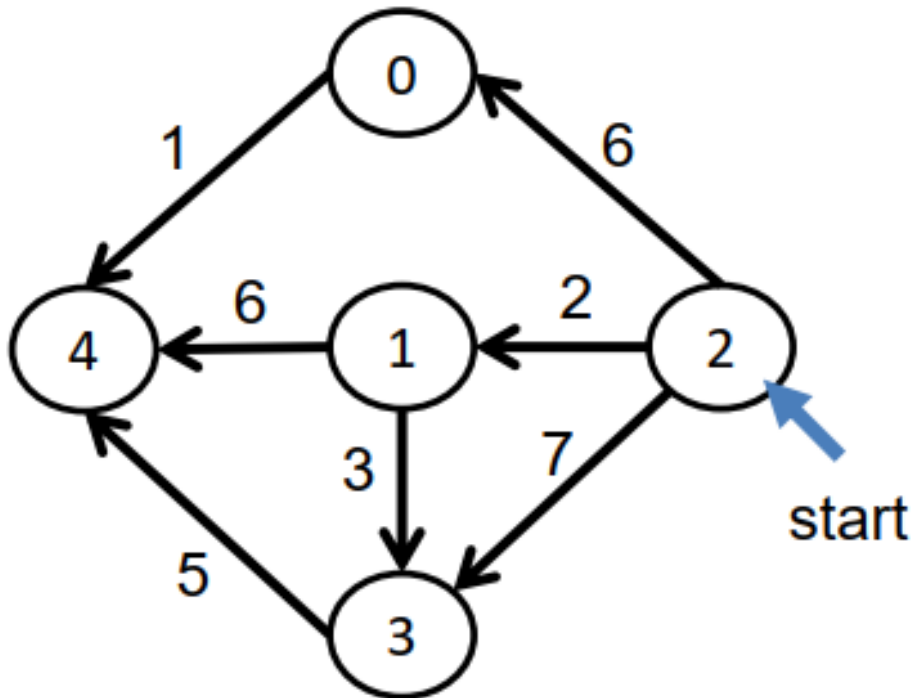
- แบบฝึกหัด
 - จงหาเส้นทางที่ใกล้ที่สุดในการเดินทางจากสถานที่เริ่มต้น ไปยังปลายทางที่กำหนด โดยมีกราฟเป็น input
- ข้อมูลนำเข้า (Input)
 - จำนวนสถานที่ในแผนที่ (N)
 - จำนวนเส้นทางที่เชื่อมต่อแต่ละสถานที่ (E)
 - สถานที่ต้นทาง สถานที่ปลายทาง
 - ต้นทาง ปลายทาง ระยะทาง ของ เส้นทางที่ 1
 - ต้นทาง ปลายทาง ระยะทาง ของ เส้นทางที่ 2
 - ...
 - ต้นทาง ปลายทาง ระยะทาง ของ เส้นทางที่ E
- ผลลัพธ์ (Output)
 - $V_1 V_2 V_3 \dots V_k$
(V_i แทนหมายเลขสถานที่ โดยค้นด้วยวิธีวนซ้ำ)
 - ระยะทางที่ใช้
- ตัวอย่าง input
 - 4
 - 4
 - 1 2
 - 1 4 5
 - 1 3 2
 - 3 4 6
 - 4 2 5
- ตัวอย่าง output
 - 1 3 4 2
 - 8



Shortest path of Graph

Dijkstra's Algorithm

ทดลองกราฟนี้



Shortest path of Graph

Bellman-Ford Algorithm



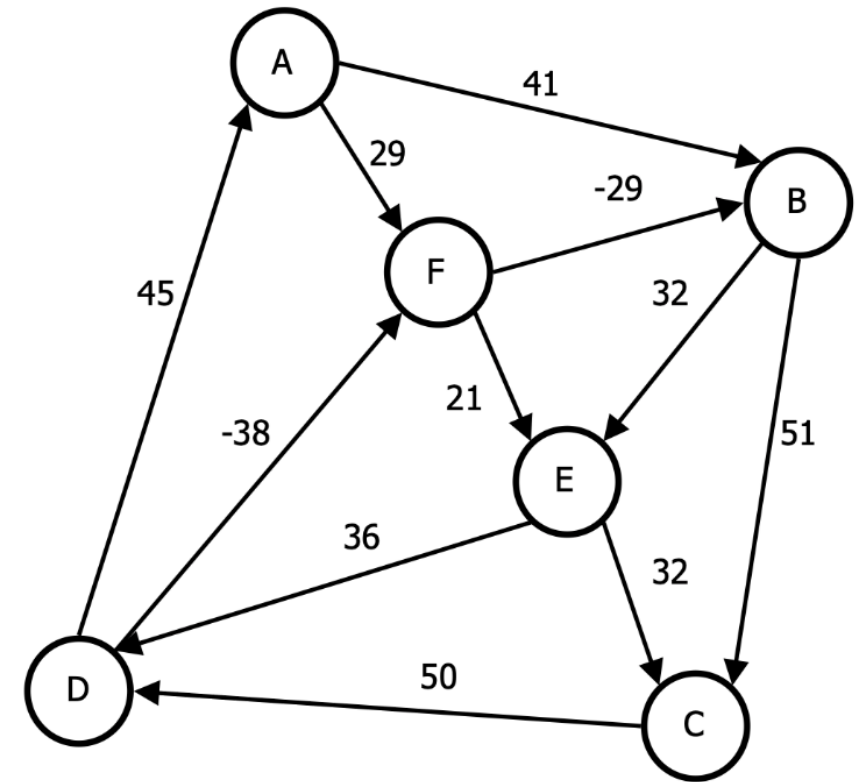
- ในกรณีที่กราฟมี weight เป็นค่าติดลบ ไม่สามารถรับประกันได้ว่า Dijkstra's จะถูก
- ให้ใช้ **Bellman-Ford Algorithm**
- (ต้องยอม ถึงแม้จะใช้เวลาเยอะกว่า)

Shortest path of Graph

Bellman-Ford Algorithm



- ตัวอย่างในชีวิตประจำวัน
 - แต่ละ **Node** แทนสินค้าทั้งหมดในรายการ
 - Edge** (u, v) แทนราคาเมื่อซื้อสินค้า u แล้วซื้อสินค้า v
 - เช่น เมื่อซื้อ A -> ซื้อ B จะเสียค่าใช้จ่าย 41 บาท
 - เมื่อซื้อ F -> ซื้อ B จะได้รับส่วนลด 29 บาท
 - การแก้ปัญหาว่า “ควรซื้อสินค้าแต่ละชนิดอย่างไร เพื่อให้ได้ราคาถูกที่สุด เมื่อซื้อสินค้า A เป็นชิ้นแรก”

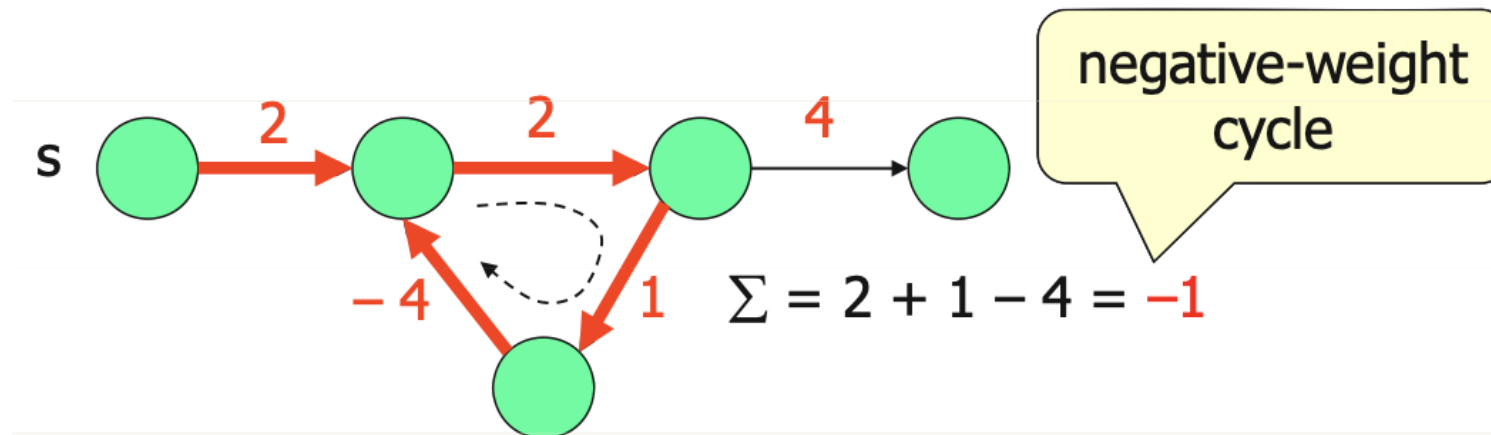


Shortest path of Graph

Bellman-Ford Algorithm

- Bellman-Ford Algorithm

- Input: source vertex
- Output:
 - If no negative cycle detected -> returns **Shortest Path Tree** (Shortest path to all vertices)
 - If negative cycle detected -> the shortest distance will not be calculated, the **cycle** is reported.



คือ cycle ที่ weight รวมมีค่า < 0,
ซึ่งทำให้หาจุดสิ้นสุดไม่ได้ เพราะยังเดินในวง ค่ายังลดลงเรื่อย ๆ

Shortest path of Graph

Bellman-Ford Algorithm



- ขั้นตอนของ Bellman–Ford Algorithm

กำหนดให้ V = set of vertices, src = Vertex เริ่มต้น

1. Initial array **dist** size $|V|$ with **INF**, except for **dist[src] = 0**
2. Calculate shortest distance โดย ทำด้านล่างนี้ทั้งหมด $|V| - 1$ ครั้ง
 - แต่ละ **edge (u, v)**
if $dist[v] > dist[u] + \text{weight of edge}(u,v)$:
 $dist[v] = dist[u] + \text{weight of edge}(u,v)$

** สรุปคือ วนทุก edge -> แต่ละ edge เก็บค่าน้อยที่สุดไว้ที่ node ปลายทาง และวนต่อจนครบ -> ทำไป $|V| - 1$ รอบ

3. Report if there's negative cycle
 - แต่ละ **edge (u, v)**
if $dist[v] > dist[u] + \text{weight of edge}(u,v)$:
Report "Graph contains negative weight cycle"

** ข้อ 2. เป็นการการันตีแล้วว่านั้นคือเส้นทางที่สั้นที่สุด ถ้าตรงนี้ยังเจออีก แสดงว่าเจอ negative weight cycle

4. If there's **no** negative weight cycle:
 $dist[v] = \text{shortest path from src to } v$ เมื่อ $v \in V$

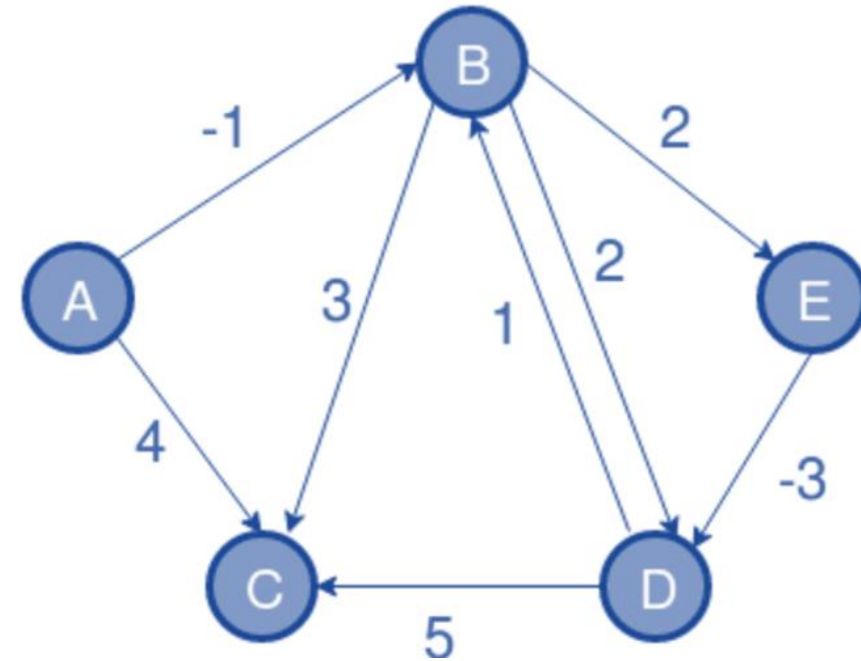
Shortest path of Graph

Bellman-Ford Algorithm



- ตัวอย่าง (src = A)

1. Initial array **dist** size $|V|$ with **INF**, except for **dist[src] = 0**
2. Calculate shortest distance
For $i = 0$ to $|V| - 1$:
For each **edge** (u, v):
if $\text{dist}[v] > \text{dist}[u] + w(\text{edge}(u,v))$:
 $\text{dist}[v] = \text{dist}[u] + w(\text{edge}(u,v))$
3. Report if there's negative cycle
For each **edge** (u, v):
if $\text{dist}[v] > \text{dist}[u] + \text{weight of edge}(u,v)$:
 "Graph contains negative weight cycle"
4. If there's **no** negative weight cycle:
 $\text{dist}[v] = \text{shortest path tree}$



$E = \{$

dist =

A	B	C	D	E

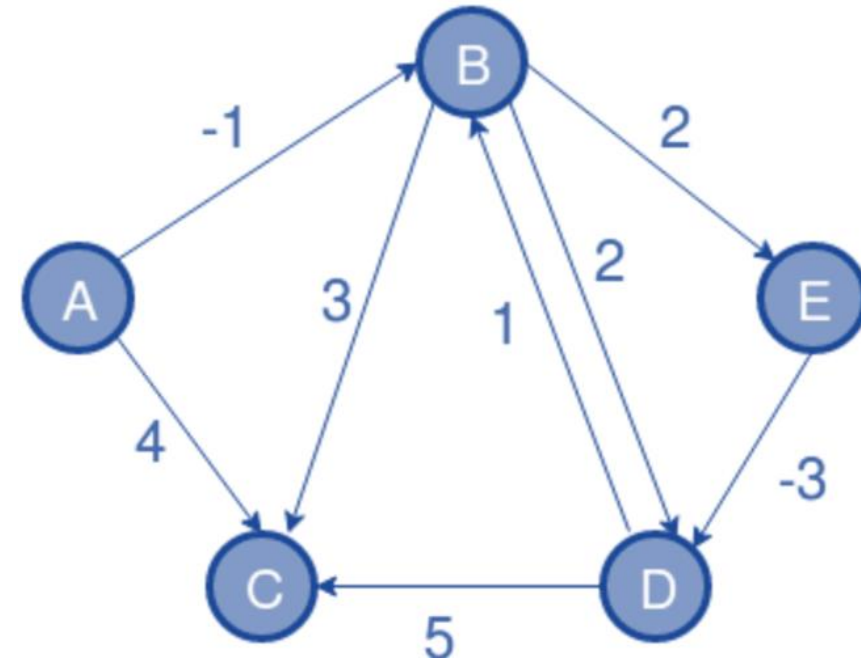
Shortest path of Graph

Bellman-Ford Algorithm



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 "Graph contains negative weight cycle"
4. If there's **no** negative weight cycle:
 $\text{dist}[v]$ = shortest path tree



$E = \{(A, B), (A, C), (B, C), (B, D), (B, E), (D, C), (D, B), (E, D)\}$

dist =

A	B	C	D	E

Shortest path of Graph

Bellman-Ford Algorithm



- ตัวอย่าง (src = A)

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2. Calculate shortest distance

For $i = 0$ to $|V| - 1$:

For each **edge** (u, v):

if $\text{dist}[v] > \text{dist}[u] + w(\text{edge}(u,v))$:

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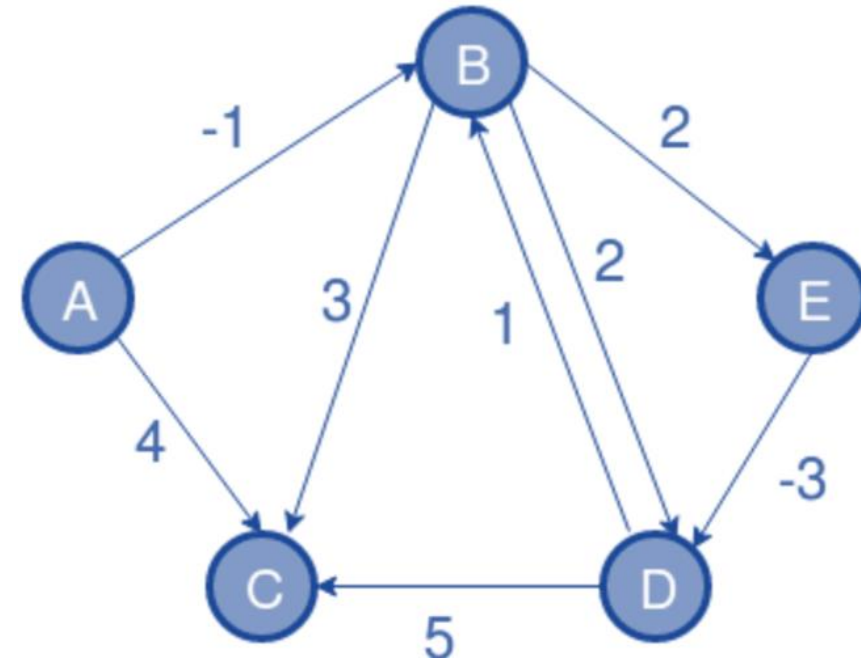
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$\text{dist}[v]$ = shortest path tree



$E = \{(A, B), (A, C), (B, C), (B, D), (B, E), (D, C), (D, B), (E, D)\}$

	A	B	C	D	E
dist =	0	INF	INF	INF	INF

Shortest path of Graph

Bellman-Ford Algorithm



- ตัวอย่าง (src = A)

1. Initial array **dist** size $|V|$ with **INF**, except for **dist[src] = 0**

2. Calculate shortest distance

For $i = 0$ to $|V| - 1$:

For each **edge** (u, v):

if $\text{dist}[v] > \text{dist}[u] + w(\text{edge}(u,v))$:

$\text{dist}[v] = \text{dist}[u] + w(\text{edge}(u,v))$

$i = 1$

3. Report if there's negative cycle

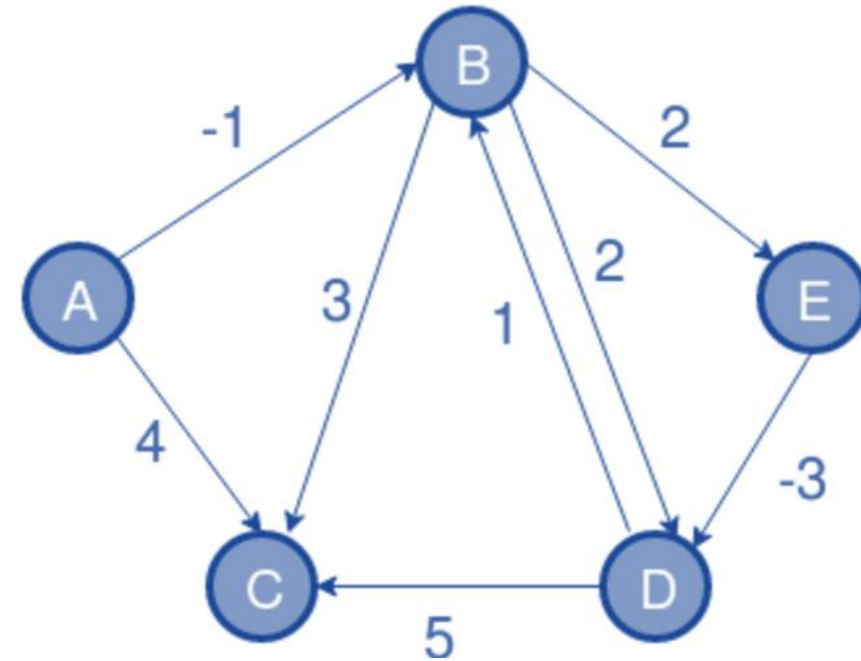
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	A	B	C	D	E
dist =	0	INF	INF	INF	INF

Shortest path of Graph

Bellman-Ford Algorithm



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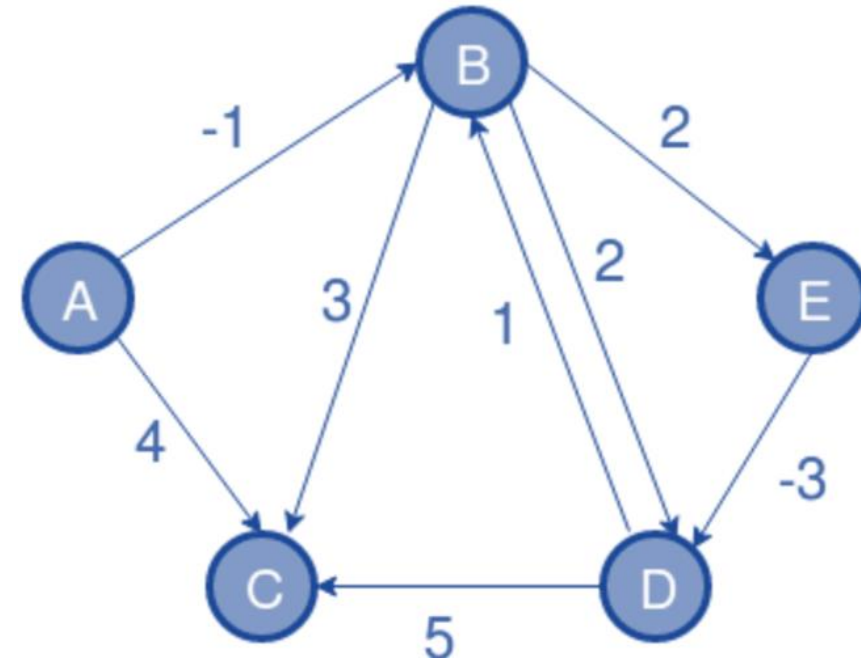
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$E = \{(A, B), (A, C), (B, C), (B, D), (B, E), (D, C), (D, B), (E, D)\}$

	A	B	C	D	E
dist =	0	INF	INF	INF	INF

Shortest path of Graph

Bellman-Ford Algorithm



- ตัวอย่าง (src = A)

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if $\text{dist}[v] > \text{dist}[u] + w(\text{edge}(u,v))$:

$\text{dist}[v] = \text{dist}[u] + w(\text{edge}(u,v))$

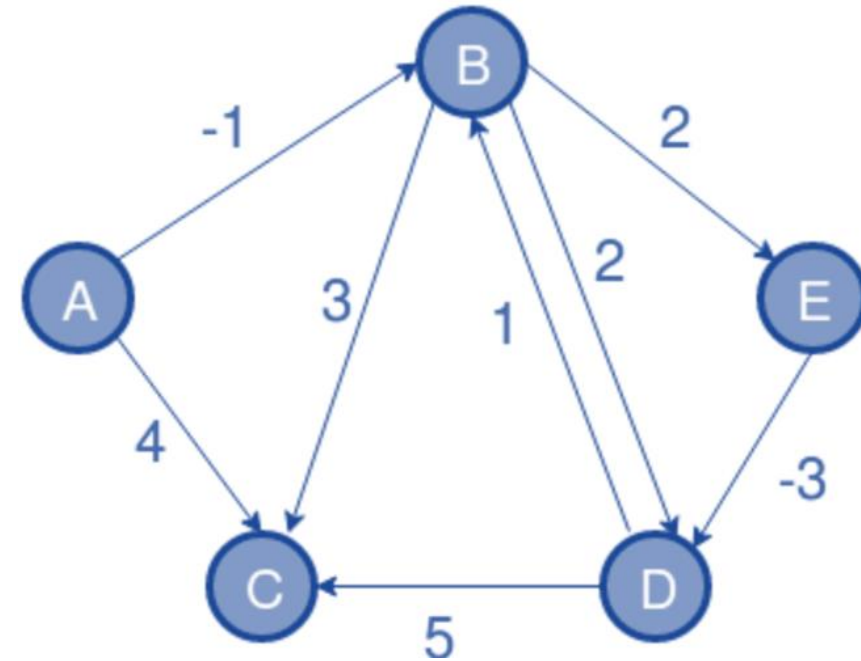
3. Report if there's negative cycle

For each **edge** (u, v):

if $\text{dist}[v] > \text{dist}[u] + \text{weight of edge}(u,v)$:

“Graph contains negative weight cycle”

4. If there's **no** negative weight cycle:
 $\text{dist}[v]$ = shortest path tree



$E = \{(A, B), (A, C), (B, C), (B, D), (B, E), (D, C), (D, B), (E, D)\}$

	A	B	C	D	E
dist =	0	-1	INF	INF	INF

Shortest path of Graph

Bellman-Ford Algorithm



- ตัวอย่าง (src = A)

1. Initial array **dist** size $|V|$ with **INF**, except for **dist[src] = 0**
2. Calculate shortest distance

For $i = 0$ to $|V| - 1$:

$i = 1$

For each **edge** (u, v):

if $\text{dist}[v] > \text{dist}[u] + w(\text{edge}(u,v))$:
 $\text{dist}[v] = \text{dist}[u] + w(\text{edge}(u,v))$

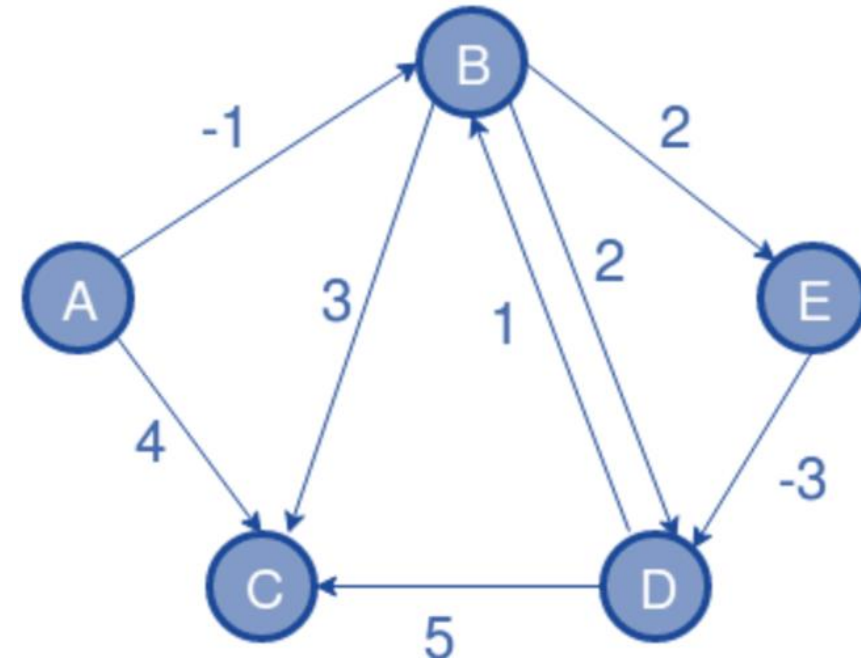
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For each **edge** (u, v):

if $\text{dist}[v] > \text{dist}[u] + \text{weight of edge}(u,v)$:

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4. If there's **no** negative weight cycle:
 $\text{dist}[v]$ = shortest path tree



$E = \{(A, B), (A, C), (B, C), (B, D), (B, E), (D, C), (D, B), (E, D)\}$

	A	B	C	D	E
dist =	0	-1	INF	INF	INF

Shortest path of Graph

Bellman-Ford Algorithm



- ตัวอย่าง (src = A)

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2. Calculate shortest distance

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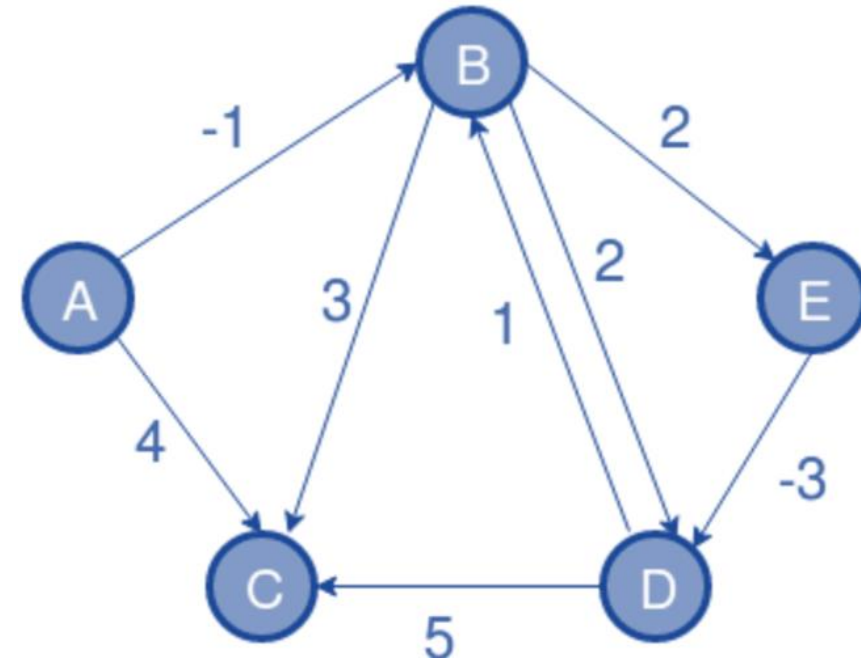
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4. If there's **no** negative weight cycle:
 $\text{dist}[v]$ = shortest path tree



$E = \{(A, B), (A, C), (B, C), (B, D), (B, E), (D, C), (D, B), (E, D)\}$

	A	B	C	D	E
dist =	0	-1	4	INF	INF

Shortest path of Graph

Bellman-Ford Algorithm



- ตัวอย่าง (src = A)

1. Initial array **dist** size $|V|$ with **INF**, except for **dist[src] = 0**
2. Calculate shortest distance

For $i = 0$ to $|V| - 1$:

For each **edge** (u, v):

if $\text{dist}[v] > \text{dist}[u] + w(\text{edge}(u,v))$:
 $\text{dist}[v] = \text{dist}[u] + w(\text{edge}(u,v))$

$i = 1$

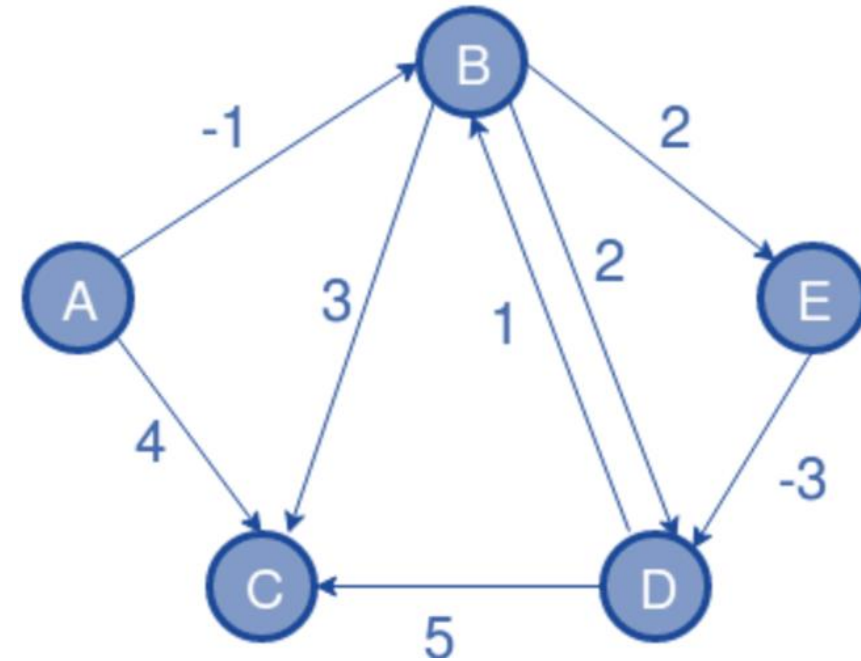
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4. If there's **no** negative weight cycle:
 $\text{dist}[v]$ = shortest path tree



$E = \{(A, B), (A, C), (B, C), (B, D), (B, E), (D, C), (D, B), (E, D)\}$

	A	B	C	D	E
dist =	0	-1	4	INF	INF

Shortest path of Graph

Bellman-Ford Algorithm



- ตัวอย่าง (src = A)

1. Initial array **dist** size $|V|$ with **INF**, except for **dist[src] = 0**
2. Calculate shortest distance

For $i = 0$ to $|V| - 1$:

$i = 1$

For each **edge** (u, v):

if $\text{dist}[v] > \text{dist}[u] + w(\text{edge}(u,v))$:

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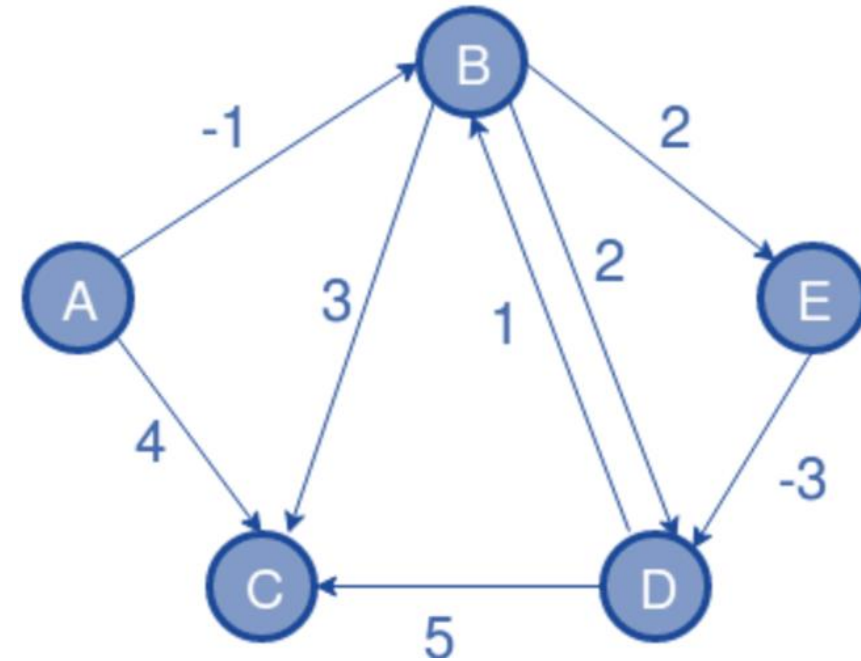
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 $\text{dist}[v]$ = shortest path tree



$E = \{(A, B), (A, C), (B, C), (B, D), (B, E), (D, C), (D, B), (E, D)\}$

	A	B	C	D	E
dist =	0	-1	2	INF	INF

Shortest path of Graph

Bellman-Ford Algorithm



- ตัวอย่าง (src = A)

1. Initial array **dist** size $|V|$ with **INF**, except for **dist[src] = 0**
2. Calculate shortest distance

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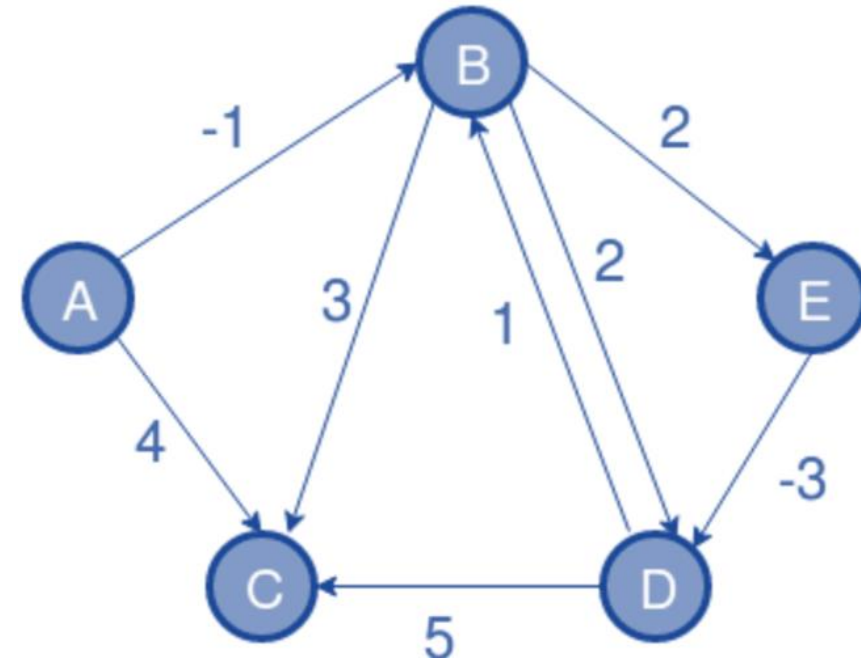
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	A	B	C	D	E
dist =	0	-1	2	INF	INF

Shortest path of Graph

Bellman-Ford Algorithm



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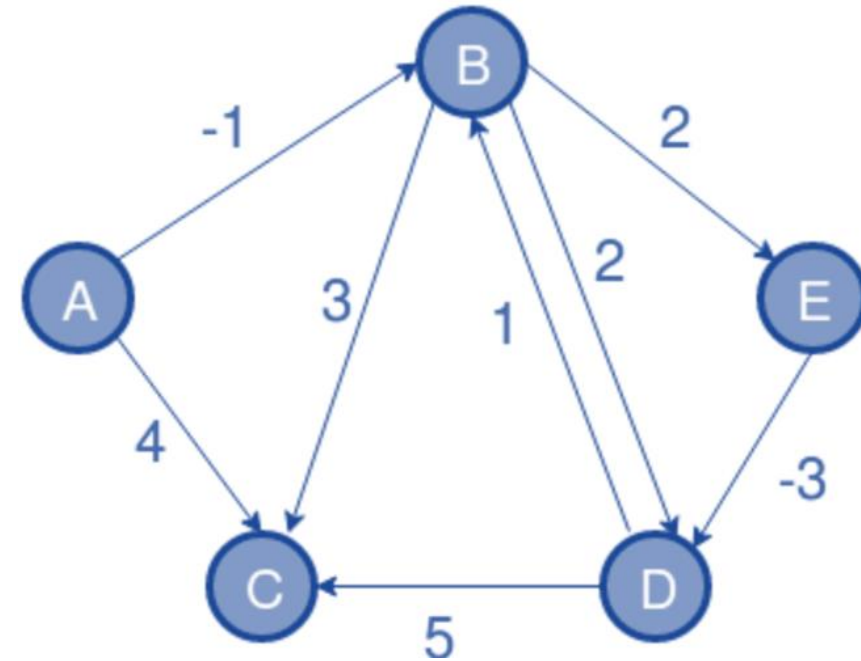
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4. If there's **no** negative weight cycle:
 $\text{dist}[v]$ = shortest path tree



$E = \{(A, B), (A, C), (B, C), (B, D), (B, E), (D, C), (D, B), (E, D)\}$

	A	B	C	D	E
dist =	0	-1	2	1	INF

Shortest path of Graph

Bellman-Ford Algorithm



- ตัวอย่าง (src = A)

1. Initial array **dist** size $|V|$ with **INF**, except for **dist[src] = 0**
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For $i = 0$ to $|V| - 1$:

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 $\text{dist}[v] = \text{dist}[u] + w(\text{edge}(u,v))$

$i = 1$

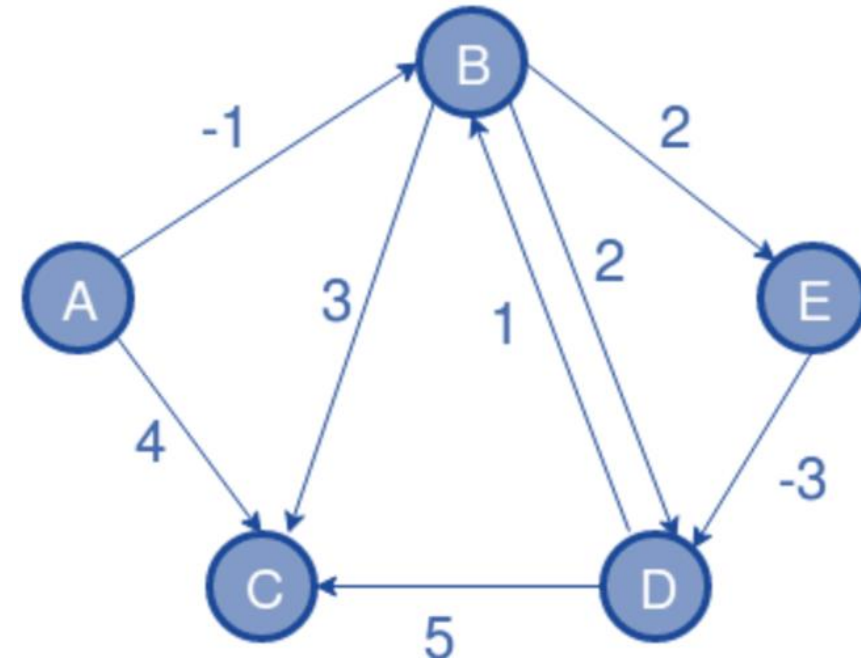
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	A	B	C	D	E
dist =	0	-1	2	1	INF

Shortest path of Graph

Bellman-Ford Algorithm



- ตัวอย่าง (src = A)

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For $i = 0$ to $|V| - 1$:

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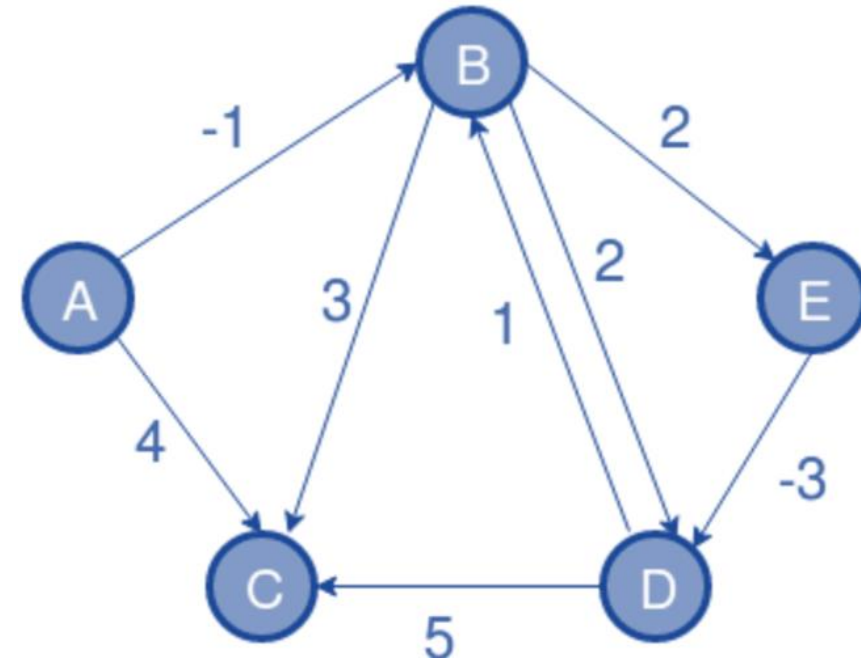
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	A	B	C	D	E
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Shortest path of Graph

Bellman-Ford Algorithm



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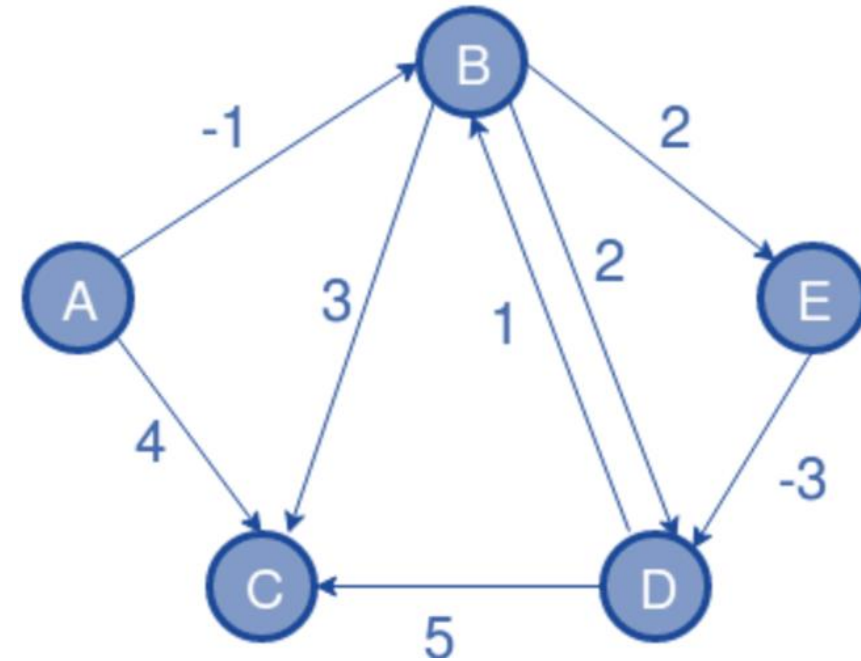
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	A	B	C	D	E
dist =	0	-1	2	1	1

Shortest path of Graph

Bellman-Ford Algorithm



- ตัวอย่าง (src = A)

1. Initial array **dist** size $|V|$ with **INF**, except for **dist[src] = 0**
2. Calculate shortest distance

For $i = 0$ to $|V| - 1$:

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if $\text{dist}[v] > \text{dist}[u] + w(\text{edge}(u,v))$:
 $\text{dist}[v] = \text{dist}[u] + w(\text{edge}(u,v))$

$i = 1$

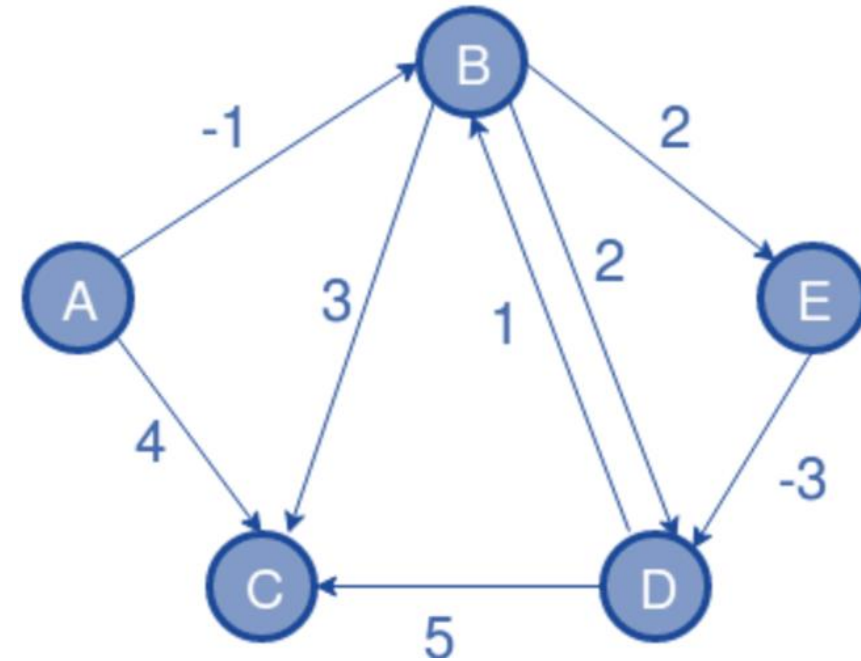
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$E = \{(A, B), (A, C), (B, C), (B, D), (B, E), (D, C), (D, B), (E, D)\}$

	A	B	C	D	E
dist =	0	-1	2	1	1

Shortest path of Graph

Bellman-Ford Algorithm



- ตัวอย่าง (src = A)

1. Initial array **dist** size $|V|$ with **INF**, except for **dist[src] = 0**
2. Calculate shortest distance

For $i = 0$ to $|V| - 1$:

For each **edge** (u, v):

if $\text{dist}[v] > \text{dist}[u] + w(\text{edge}(u,v))$:
 $\text{dist}[v] = \text{dist}[u] + w(\text{edge}(u,v))$

$i = 1$

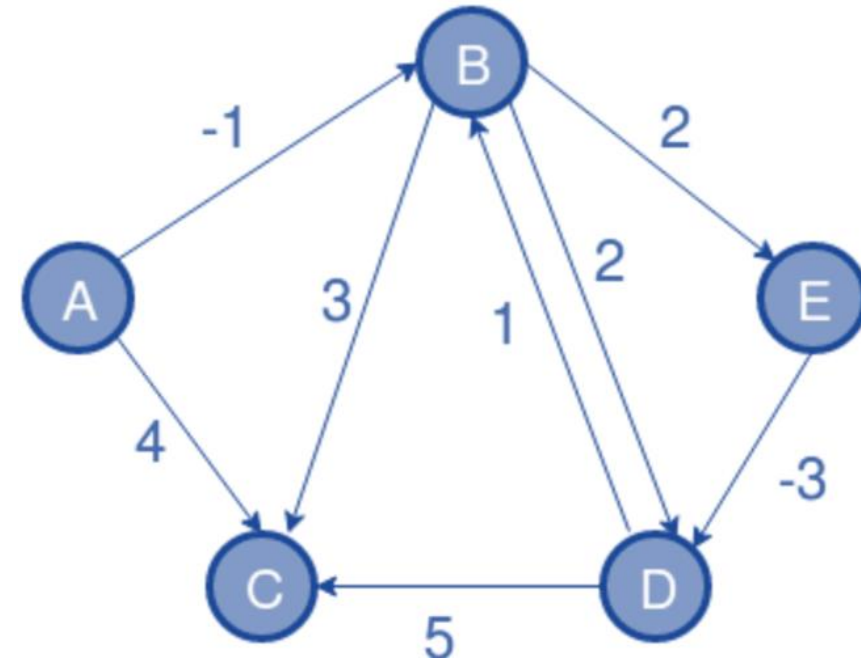
3. Report if there's negative cycle

For each **edge** (u, v):

if $\text{dist}[v] > \text{dist}[u] + \text{weight of edge}(u,v)$:

“Graph contains negative weight cycle”

4. If there's **no** negative weight cycle:
 $\text{dist}[v]$ = shortest path tree



$E = \{(A, B), (A, C), (B, C), (B, D), (B, E), (D, C), (D, B), (E, D)\}$

	A	B	C	D	E
dist =	0	-1	2	1	1

Shortest path of Graph

Bellman-Ford Algorithm



- ตัวอย่าง (src = A)

1. Initial array **dist** size $|V|$ with **INF**, except for **dist[src] = 0**
2. Calculate shortest distance

For $i = 0$ to $|V| - 1$:

$i = 1$

For each **edge** (u, v):

if $\text{dist}[v] > \text{dist}[u] + w(\text{edge}(u,v))$:

$\text{dist}[v] = \text{dist}[u] + w(\text{edge}(u,v))$

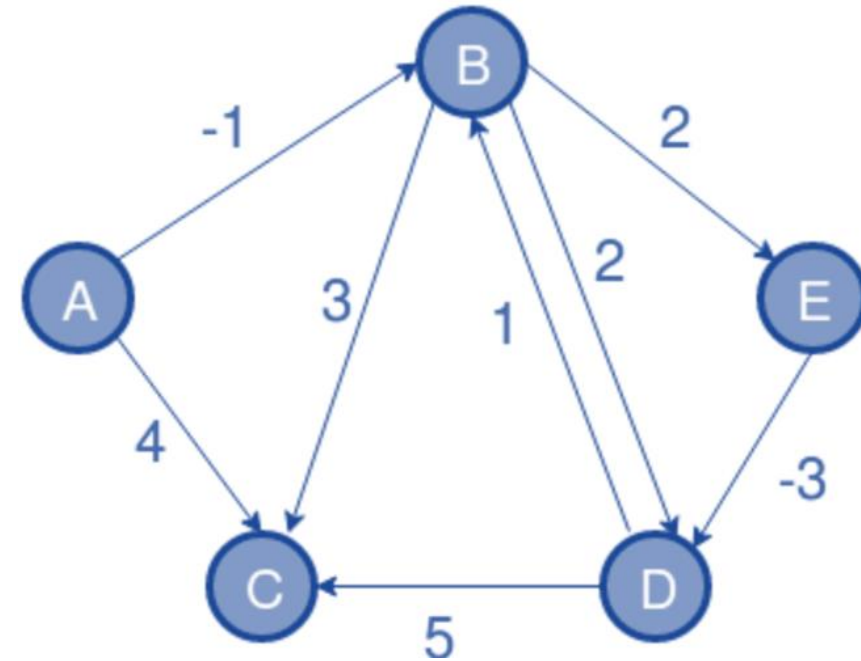
3. Report if there's negative cycle

For each **edge** (u, v):

if $\text{dist}[v] > \text{dist}[u] + \text{weight of edge}(u,v)$:

"Graph contains negative weight cycle"

4. If there's **no** negative weight cycle:
 $\text{dist}[v]$ = shortest path tree



$E = \{(A, B), (A, C), (B, C), (B, D), (B, E), (D, C), (D, B), (E, D)\}$

	A	B	C	D	E
dist =	0	-1	2	-2	1

Shortest path of Graph

Bellman-Ford Algorithm



- ตัวอย่าง (src = A)

1. Initial array **dist** size $|V|$ with **INF**, except for **dist[src] = 0**
2. Calculate shortest distance

For $i = 0$ to $|V| - 1$:

$i = 2$

For each **edge** (u, v):

if $\text{dist}[v] > \text{dist}[u] + w(\text{edge}(u,v))$:
 $\text{dist}[v] = \text{dist}[u] + w(\text{edge}(u,v))$

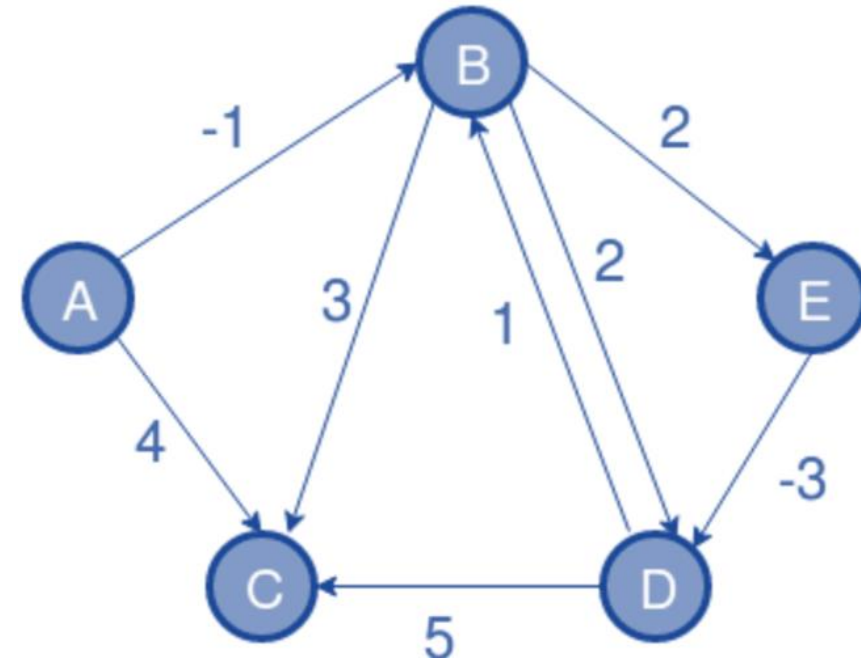
3. Report if there's negative cycle

For each **edge** (u, v):

if $\text{dist}[v] > \text{dist}[u] + \text{weight of edge}(u,v)$:

“Graph contains negative weight cycle”

4. If there's **no** negative weight cycle:
 $\text{dist}[v] = \text{shortest path tree}$



$E = \{(A, B), (A, C), (B, C), (B, D), (B, E), (D, C), (D, B), (E, D)\}$

	A	B	C	D	E
dist =	0	-1	2	-2	1

Shortest path of Graph

Bellman-Ford Algorithm



- ตัวอย่าง (src = A)

1. Initial array **dist** size $|V|$ with **INF**, except for **dist[src] = 0**
2. Calculate shortest distance

For $i = 0$ to $|V| - 1$:

$i = 2$

For each **edge** (u, v):

if $\text{dist}[v] > \text{dist}[u] + w(\text{edge}(u,v))$:

$\text{dist}[v] = \text{dist}[u] + w(\text{edge}(u,v))$

3. Report if there's negative cycle

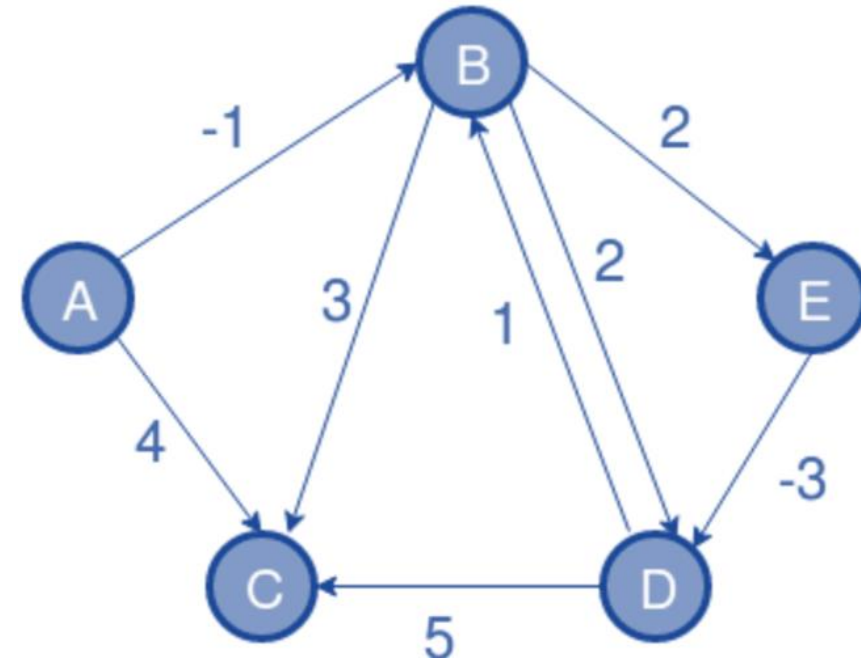
For each **edge** (u, v):

if $\text{dist}[v] > \text{dist}[u] + \text{weight of edge}(u,v)$:

"Graph contains negative weight cycle"

4. If there's **no** negative weight cycle:

$\text{dist}[v]$ = shortest path tree



$E = \{(A, B), (A, C), (B, C), (B, D), (B, E), (D, C), (D, B), (E, D)\}$

	A	B	C	D	E
dist =	0	-1	2	-2	1

Shortest path of Graph

Bellman-Ford Algorithm



- ตัวอย่าง (src = A)

1. Initial array **dist** size $|V|$ with **INF**, except for **dist[src] = 0**
2. Calculate shortest distance

For $i = 0$ to $|V| - 1$:

$i = 3$

For each **edge** (u, v):

if $\text{dist}[v] > \text{dist}[u] + w(\text{edge}(u,v))$:
 $\text{dist}[v] = \text{dist}[u] + w(\text{edge}(u,v))$

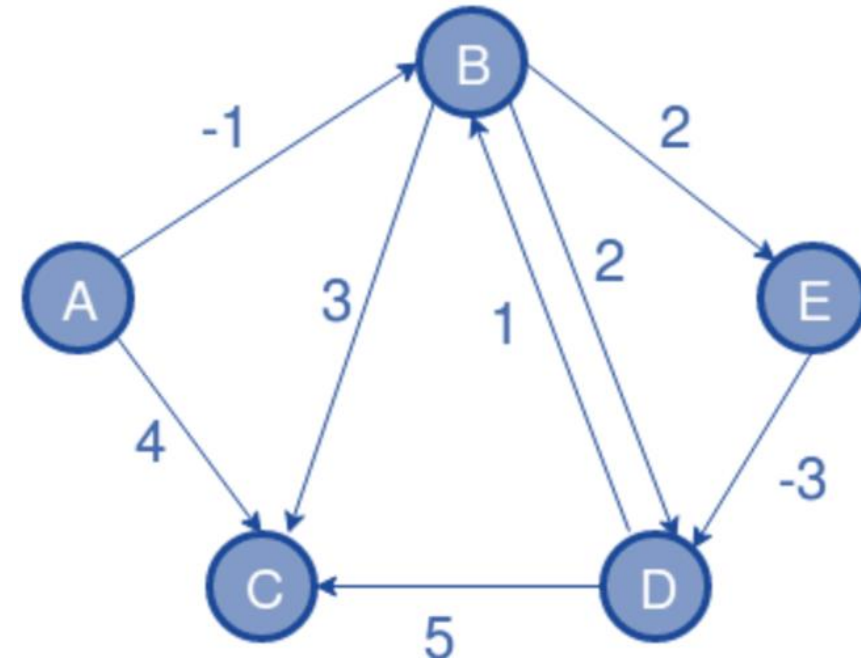
3. Report if there's negative cycle

For each **edge** (u, v):

if $\text{dist}[v] > \text{dist}[u] + \text{weight of edge}(u,v)$:

“Graph contains negative weight cycle”

4. If there's **no** negative weight cycle:
 $\text{dist}[v] = \text{shortest path tree}$



$E = \{(A, B), (A, C), (B, C), (B, D), (B, E), (D, C), (D, B), (E, D)\}$

	A	B	C	D	E
dist =	0	-1	2	-2	1

Shortest path of Graph

Bellman-Ford Algorithm



- ตัวอย่าง (src = A)

1. Initial array **dist** size $|V|$ with **INF**, except for **dist[src] = 0**
2. Calculate shortest distance

For $i = 0$ to $|V| - 1$:

$i = 4$

For each **edge** (u, v):

if $\text{dist}[v] > \text{dist}[u] + w(\text{edge}(u,v))$:

$\text{dist}[v] = \text{dist}[u] + w(\text{edge}(u,v))$

3. Report if there's negative cycle

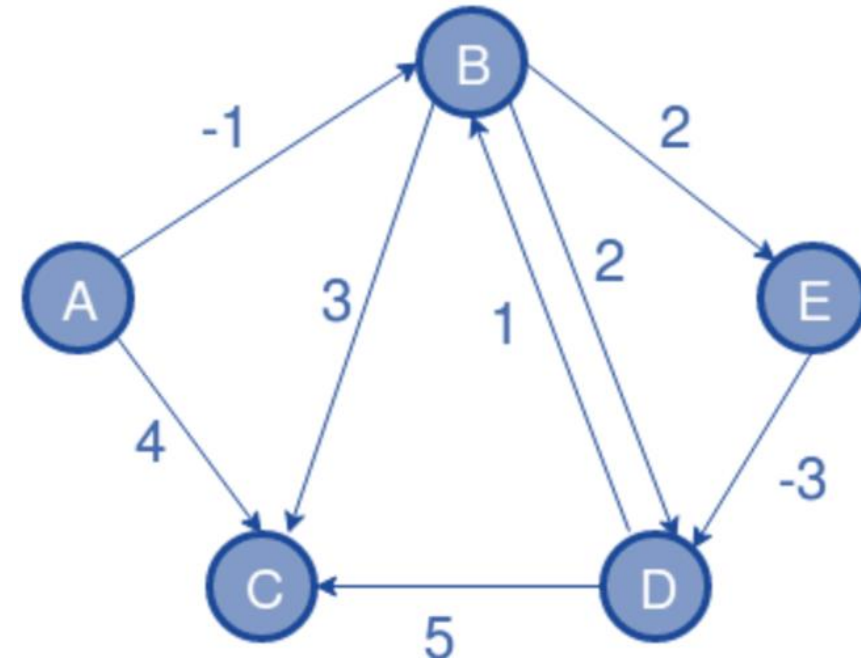
For each **edge** (u, v):

if $\text{dist}[v] > \text{dist}[u] + \text{weight of edge}(u,v)$:

"Graph contains negative weight cycle"

4. If there's **no** negative weight cycle:

$\text{dist}[v]$ = shortest path tree



$E = \{(A, B), (A, C), (B, C), (B, D), (B, E), (D, C), (D, B), (E, D)\}$

	A	B	C	D	E
dist =	0	-1	2	-2	1

Shortest path of Graph

Bellman-Ford Algorithm

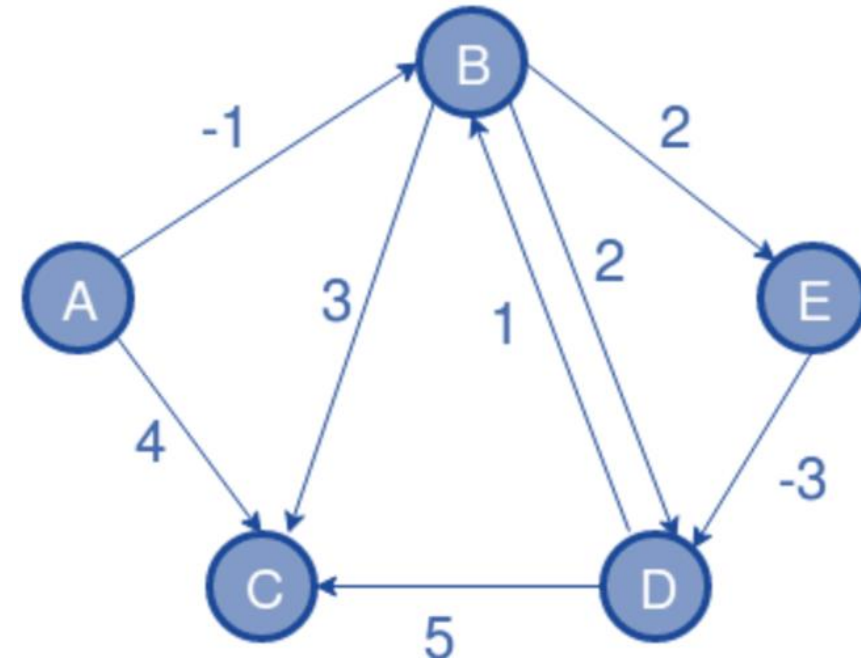


- ตัวอย่าง (src = A)

1. Initial array **dist** size $|V|$ with **INF**, except for **dist[src] = 0**
2. Calculate shortest distance
For $i = 0$ to $|V| - 1$:
For each **edge** (u, v):
if $\text{dist}[v] > \text{dist}[u] + w(\text{edge}(u,v))$:
 $\text{dist}[v] = \text{dist}[u] + w(\text{edge}(u,v))$

3. Report if there's negative cycle
For each **edge** (u, v):
if $\text{dist}[v] > \text{dist}[u] + \text{weight of edge}(u,v)$:
 "Graph contains negative weight cycle"

4. If there's **no** negative weight cycle:
 $\text{dist}[v]$ = shortest path tree



$$E = \{(A, B), (A, C), (B, C), (B, D), (B, E), (D, C), (D, B), (E, D)\}$$

	A	B	C	D	E
dist =	0	-1	2	-2	1

Shortest path of Graph

Bellman-Ford Algorithm



- ตัวอย่าง (src = A)

1. Initial array **dist** size $|V|$ with **INF**, except for **dist[src] = 0**
2. Calculate shortest distance

For $i = 0$ to $|V| - 1$:

For each **edge** (u, v):

if $\text{dist}[v] > \text{dist}[u] + w(\text{edge}(u,v))$:
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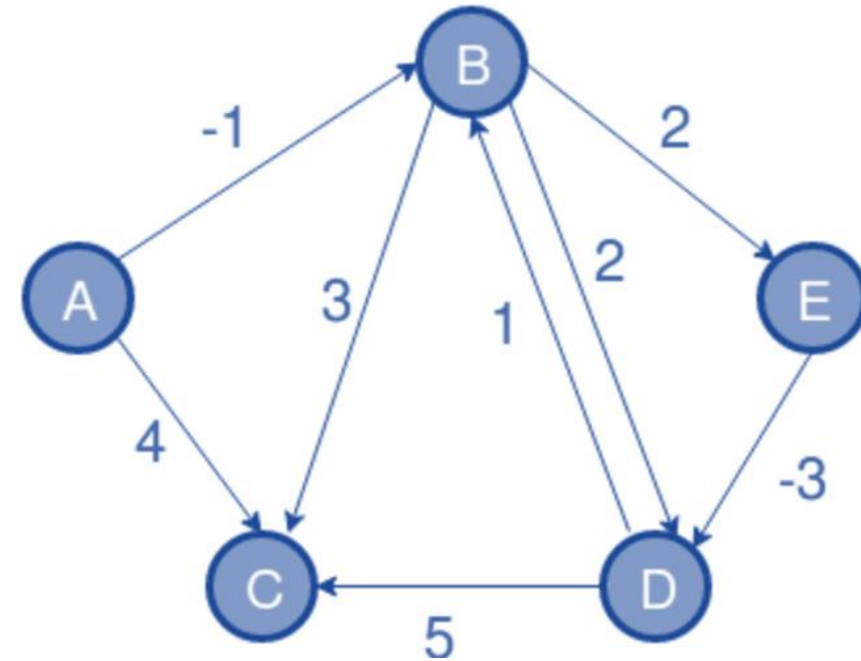
3. Report if there's negative cycle

For each **edge** (u, v):

if $\text{dist}[v] > \text{dist}[u] + \text{weight of edge}(u,v)$:

“Graph contains negative weight cycle”

4. If there's **no** negative weight cycle:
 $\text{dist}[v] = \text{shortest path tree}$

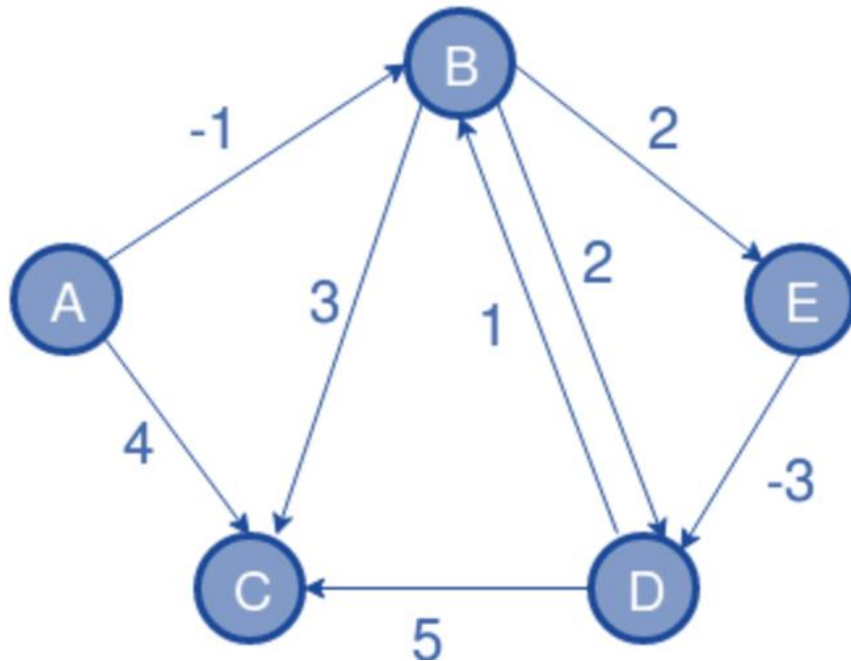


$E = \{(A, B), (A, C), (B, C), (B, D), (B, E), (D, C), (D, B), (E, D)\}$

	A	B	C	D	E
dist =	0	-1	2	-2	1

Shortest path of Graph

Bellman-Ford Algorithm



Let's code !!

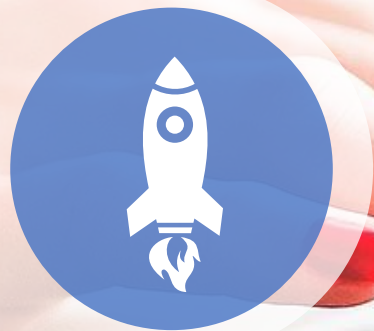
[06-Bellman-Ford.cpp]

adjMatrix <https://pastebin.com/e8gtvpad>

adjList <https://pastebin.com/S7ciAxii>

All-pair Shortest path

คืออะไร, Algorithm

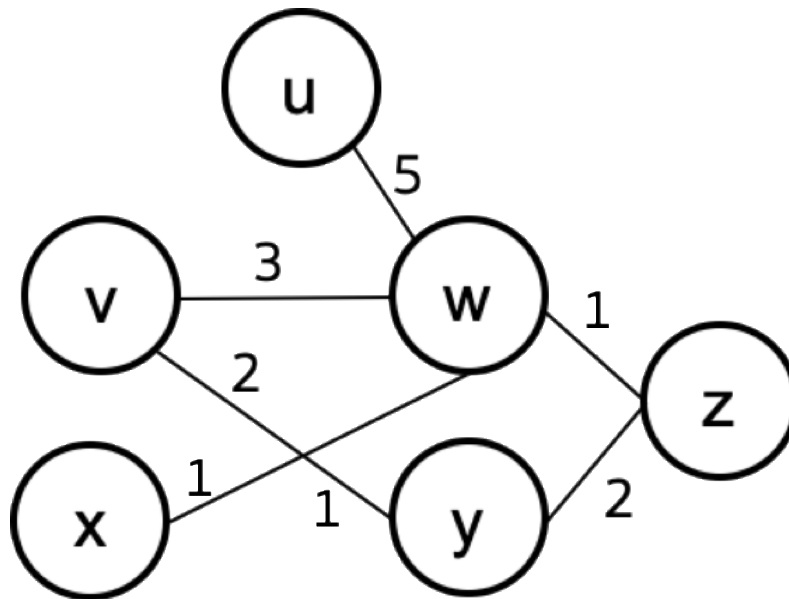


All-pair Shortest path of Graph

คืออะไร



- Shortest Path (แบบสั้น)
 - มีจุดเริ่มต้น (source) 1 จุด -> เส้นทางใดสั้นที่สุด
- All-pair Shortest Path
 - ทุก ๆ Node u, v ที่ไปถึงกันได้ -> หาเส้นทางที่ใกล้ที่สุด จาก u ไป v



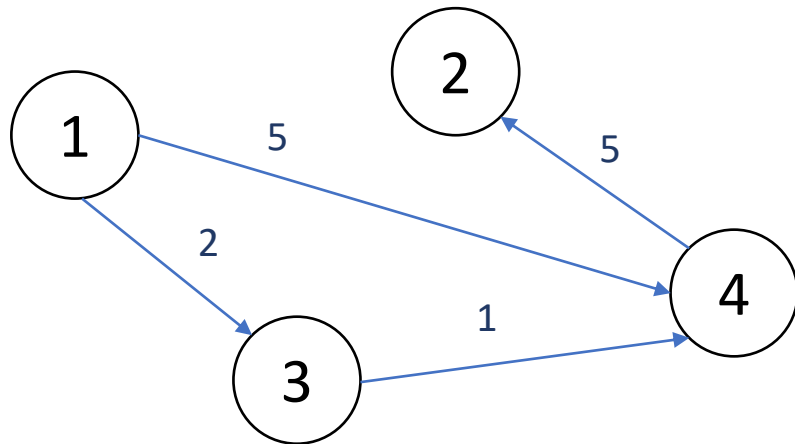
All-pair Shortest path of Graph

Floyd-Warshall Algorithm



- **Floyd-Warshall Algorithm**

- ทุก ๆ Node u, v ที่ไปถึงกันได้ \rightarrow ระยะทางที่ใกล้ที่สุด จาก u ไป v เป็นเท่าไรบ้าง



Node distance

1, 2

1, 3

1, 4

3, 2

3, 4

4, 2

All-pair Shortest path of Graph

Floyd-Warshall Algorithm



- ขั้นตอนของ Floyd-Warshall Algorithm

“ ในการหาเส้นทางจาก $i \rightarrow j$ ถ้ามีเส้นทางที่อ้อมไป ($i \rightarrow k \rightarrow j$) แล้วสั้นกว่า จะเก็บไว้เรื่อย ๆ ”

กำหนดให้ input graph $G = (V, E)$

1. Initialize array **dist** ขนาด $V \times V$
2. **dist** = **G**
 $\text{dist}[i][j] = 0$ if $i = j$ และ $\text{dist}[i][j] = \text{INF}$ if no (i, j) in E
3. for ($k=0; k < V; k++$):
 for ($i=0; i < V; i++$):
 for ($j=0; j < V; j++$):
 if ($\text{dist}[i][k] + \text{dist}[k][j] < \text{dist}[i][j]$)
 $\text{dist}[i][j] = \text{dist}[i][k] + \text{dist}[k][j]$

All-pair Shortest path of Graph

Floyd-Warshall Algorithm



กำหนดให้ input graph $G = (V, E)$

Initialize array **dist** ขนาด $V \times V$

dist = **G**

$\text{dist}[i][j] = 0$ if $i = j$

$\text{dist}[i][j] = \text{INF}$ if no (i, j) in E

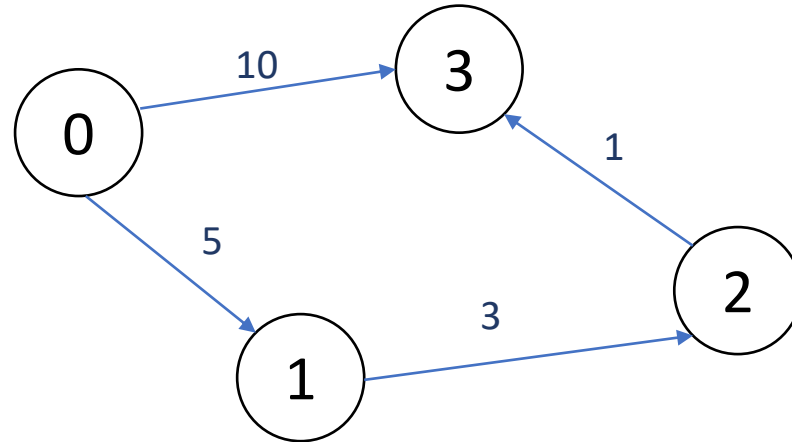
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 for ($i=0$; $i < V$; $i++$):

 for ($j=0$; $j < V$; $j++$):

 if ($\text{dist}[i][k] + \text{dist}[k][j] < \text{dist}[i][j]$)

$\text{dist}[i][j] = \text{dist}[i][k] + \text{dist}[k][j]$



All-pair Shortest path of Graph

Floyd-Warshall Algorithm



กำหนดให้ input graph $G = (V, E)$

Initialize array **dist** ขนาด $V \times V$

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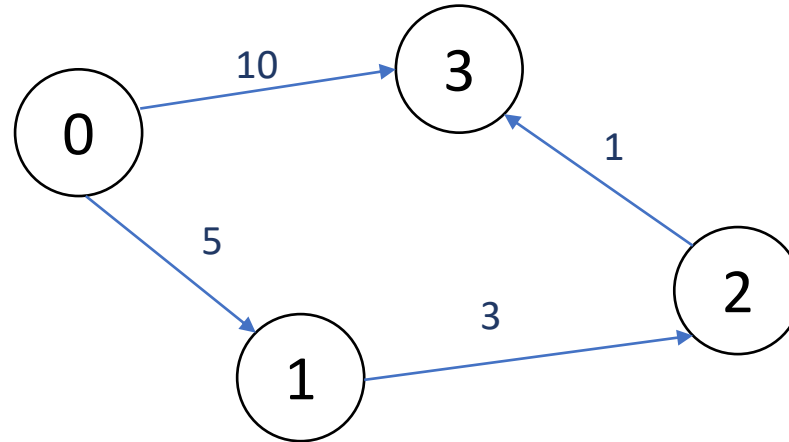
for ($k=0$; $k<V$; $k++$):

for ($i=0$; $i<V$; $i++$):

for ($j=0$; $j<V$; $j++$):

if ($\text{dist}[i][k] + \text{dist}[k][j] < \text{dist}[i][j]$)

$\text{dist}[i][j] = \text{dist}[i][k] + \text{dist}[k][j]$



dist				
	0	1	2	3
0				
1				
2				
3				

All-pair Shortest path of Graph

Floyd-Warshall Algorithm



กำหนดให้ input graph $G = (V, E)$

Initialize array **dist** ขนาด $V \times V$

dist = G

$\text{dist}[i][j] = 0$ if $i = j$

$\text{dist}[i][j] = \text{INF}$ if no (i, j) in E

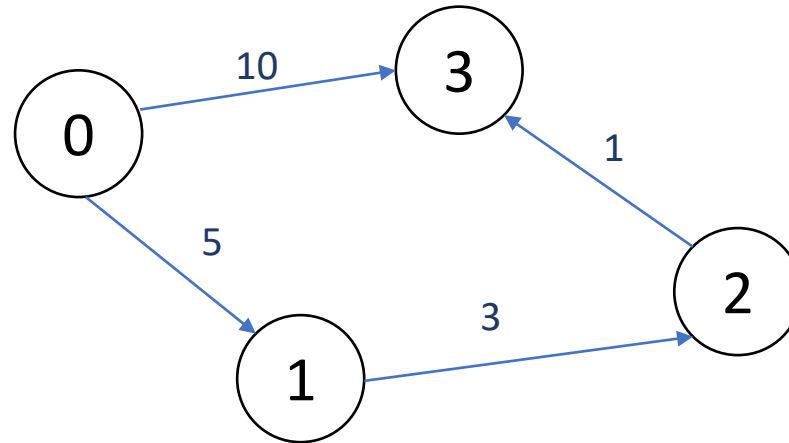
for ($k=0$; $k < V$; $k++$):

for ($i=0$; $i < V$; $i++$):

for ($j=0$; $j < V$; $j++$):

if ($\text{dist}[i][k] + \text{dist}[k][j] < \text{dist}[i][j]$)

$\text{dist}[i][j] = \text{dist}[i][k] + \text{dist}[k][j]$



dist				
	0	1	2	3
0				
1				
2				
3				

All-pair Shortest path of Graph

Floyd-Warshall Algorithm



กำหนดให้ input graph $G = (V, E)$

Initialize array **dist** ขนาด $V \times V$

dist = G

$\text{dist}[i][j] = 0$ if $i = j$

$\text{dist}[i][j] = \text{INF}$ if no (i, j) in E

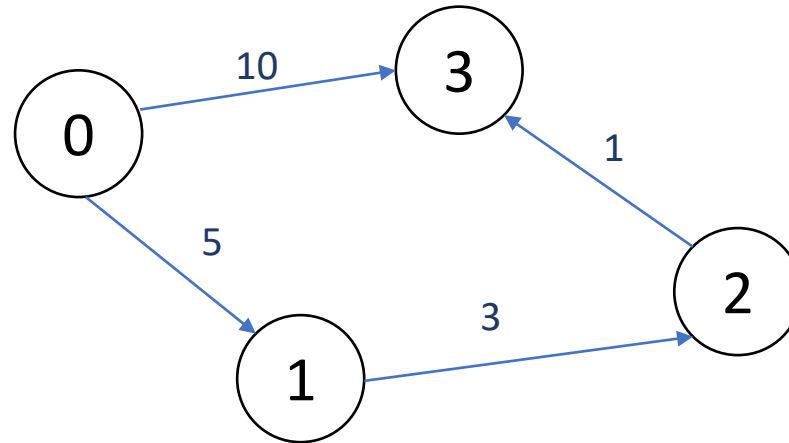
for ($k=0$; $k<V$; $k++$):

for ($i=0$; $i<V$; $i++$):

for ($j=0$; $j<V$; $j++$):

if ($\text{dist}[i][k] + \text{dist}[k][j] < \text{dist}[i][j]$)

$\text{dist}[i][j] = \text{dist}[i][k] + \text{dist}[k][j]$



dist

	0	1	2	3
0	0	5	INF	10
1	INF	0	3	INF
2	INF	INF	0	1
3	INF	INF	INF	0

All-pair Shortest path of Graph

Floyd-Warshall Algorithm



กำหนดให้ input graph $G = (V, E)$

Initialize array **dist** ขนาด $V \times V$

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$\text{dist}[i][j] = 0$ if $i = j$

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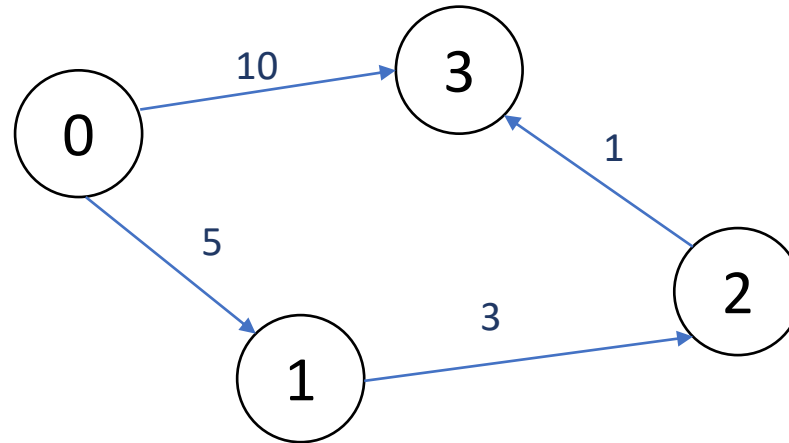
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if ($\text{dist}[i][k] + \text{dist}[k][j] < \text{dist}[i][j]$)

$\text{dist}[i][j] = \text{dist}[i][k] + \text{dist}[k][j]$



$i = 0$ $j = 0$ $k = 0$ \rightarrow

dist

	0	1	2	3
0	0	5	INF	10
1	INF	0	3	INF
2	INF	INF	0	1
3	INF	INF	INF	0

All-pair Shortest path of Graph

Floyd-Warshall Algorithm



กำหนดให้ input graph $G = (V, E)$

Initialize array **dist** ขนาด $V \times V$

dist = **G**

$\text{dist}[i][j] = 0$ if $i = j$

$\text{dist}[i][j] = \text{INF}$ if no (i, j) in E

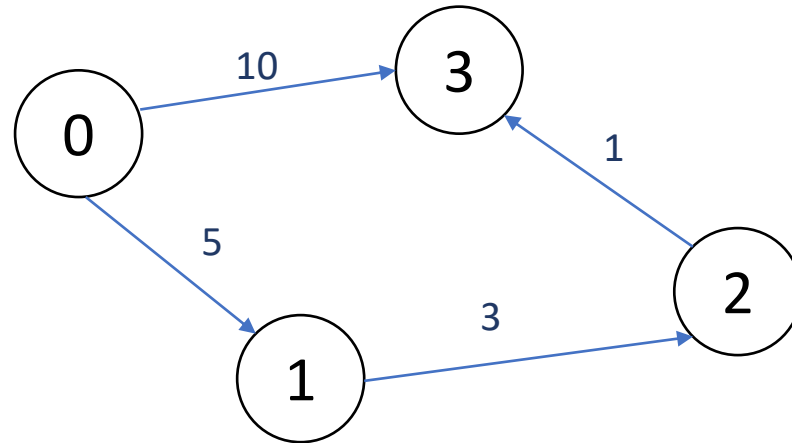
for ($k=0$; $k < V$; $k++$):

for ($i=0$; $i < V$; $i++$):

for ($j=0$; $j < V$; $j++$):

if ($\text{dist}[i][k] + \text{dist}[k][j] < \text{dist}[i][j]$)

$\text{dist}[i][j] = \text{dist}[i][k] + \text{dist}[k][j]$



dist

	0	1	2	3
0	0	5	INF	10
1	INF	0	3	INF
2	INF	INF	0	1
3	INF	INF	INF	0

$i = 0$ $j = 0$ $k = 0$ \rightarrow False

All-pair Shortest path of Graph

Floyd-Warshall Algorithm



กำหนดให้ input graph $G = (V, E)$

Initialize array **dist** ขนาด $V \times V$

dist = **G**

$\text{dist}[i][j] = 0$ if $i = j$

$\text{dist}[i][j] = \text{INF}$ if no (i, j) in E

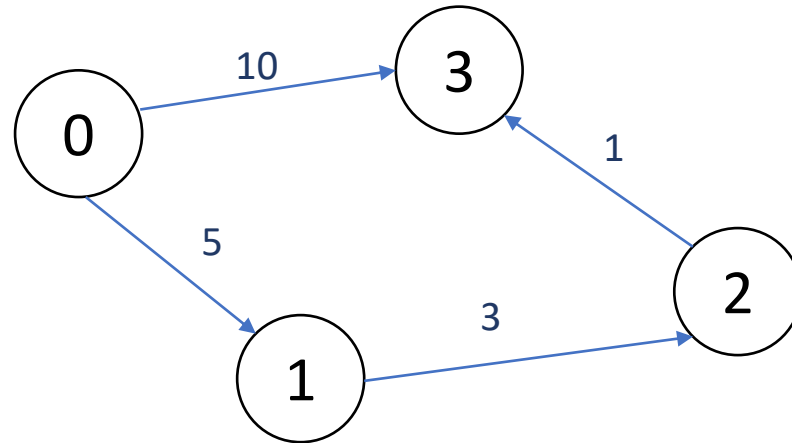
for ($k=0$; $k < V$; $k++$):

for ($i=0$; $i < V$; $i++$):

for ($j=0$; $j < V$; $j++$):

if ($\text{dist}[i][k] + \text{dist}[k][j] < \text{dist}[i][j]$)

$\text{dist}[i][j] = \text{dist}[i][k] + \text{dist}[k][j]$



$i = 0 \quad j = 1 \quad k = 0 \rightarrow$

dist

	0	1	2	3
0	0	5	INF	10
1	INF	0	3	INF
2	INF	INF	0	1
3	INF	INF	INF	0

All-pair Shortest path of Graph

Floyd-Warshall Algorithm



กำหนดให้ input graph $G = (V, E)$

Initialize array **dist** ขนาด $V \times V$

dist = **G**

$\text{dist}[i][j] = 0$ if $i = j$

$\text{dist}[i][j] = \text{INF}$ if no (i, j) in E

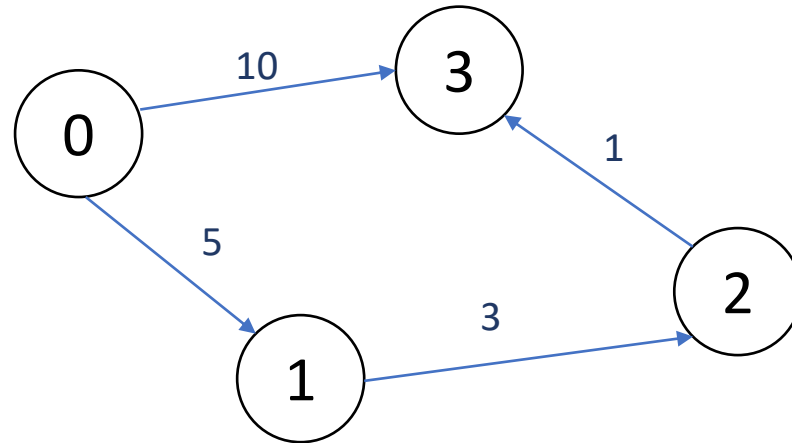
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for ($j=0$; $j < V$; $j++$):

if ($\text{dist}[i][k] + \text{dist}[k][j] < \text{dist}[i][j]$)

$\text{dist}[i][j] = \text{dist}[i][k] + \text{dist}[k][j]$



$i = 0$ $j = 2$ $k = 0$ \rightarrow

dist

	0	1	2	3
0	0	5	INF	10
1	INF	0	3	INF
2	INF	INF	0	1
3	INF	INF	INF	0

All-pair Shortest path of Graph

Floyd-Warshall Algorithm



กำหนดให้ input graph $G = (V, E)$

Initialize array **dist** ขนาด $V \times V$

dist = **G**

$\text{dist}[i][j] = 0$ if $i = j$

$\text{dist}[i][j] = \text{INF}$ if no (i, j) in E

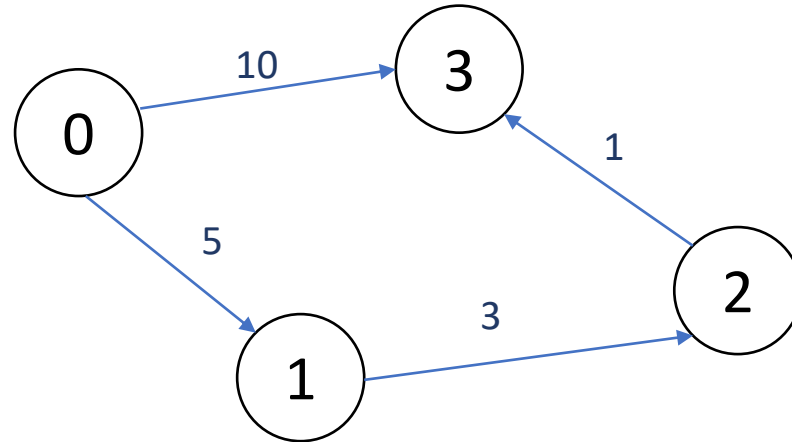
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if ($\text{dist}[i][k] + \text{dist}[k][j] < \text{dist}[i][j]$)

$\text{dist}[i][j] = \text{dist}[i][k] + \text{dist}[k][j]$



$i = 0$ $j = 3$ $k = 0$ \rightarrow

dist

	0	1	2	3
0	0	5	INF	10
1	INF	0	3	INF
2	INF	INF	0	1
3	INF	INF	INF	0

All-pair Shortest path of Graph

Floyd-Warshall Algorithm



กำหนดให้ input graph $G = (V, E)$

Initialize array **dist** ขนาด $V \times V$

dist = **G**

$\text{dist}[i][j] = 0$ if $i = j$

$\text{dist}[i][j] = \text{INF}$ if no (i, j) in E

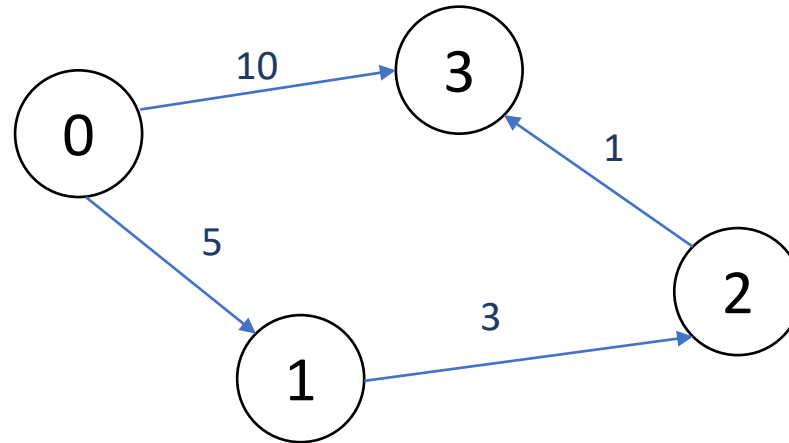
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for ($i=0$; $i < V$; $i++$):

for ($j=0$; $j < V$; $j++$):

if ($\text{dist}[i][k] + \text{dist}[k][j] < \text{dist}[i][j]$)

$\text{dist}[i][j] = \text{dist}[i][k] + \text{dist}[k][j]$



$i = 1$ $j = 0$ $k = 0$ \rightarrow

dist				
	0	1	2	3
0	0	5	INF	10
1	INF	0	3	INF
2	INF	INF	0	1
3	INF	INF	INF	0

All-pair Shortest path of Graph

Floyd-Warshall Algorithm



กำหนดให้ input graph $G = (V, E)$

Initialize array **dist** ขนาด $V \times V$

dist = **G**

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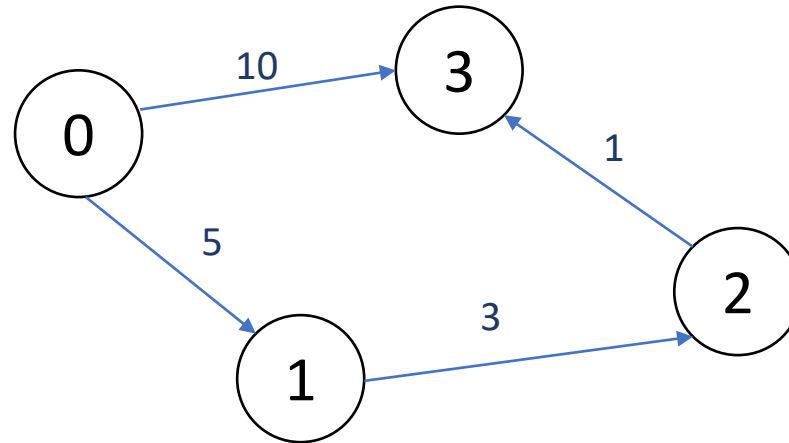
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$\text{dist}[i][j] = \text{dist}[i][k] + \text{dist}[k][j]$



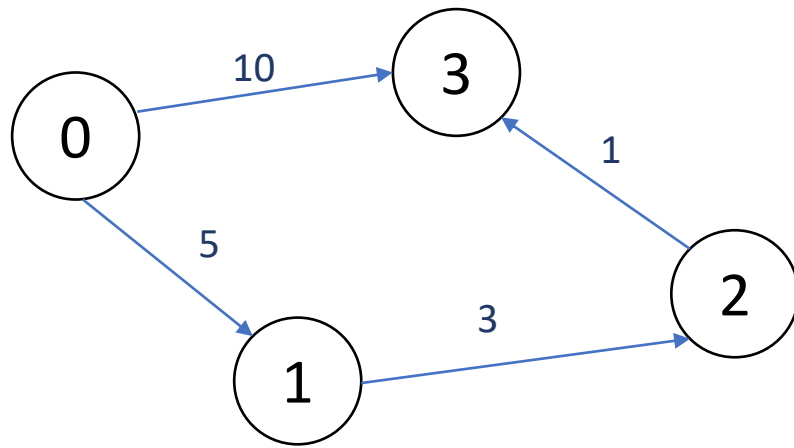
$i =$ $j =$ $k =$ \rightarrow

dist

	0	1	2	3
0	0	5	INF	10
1	INF	0	3	INF
2	INF	INF	0	1
3	INF	INF	INF	0

All-pair Shortest path of Graph

Floyd-Warshall Algorithm



Let's code !!

[07-Floyd-Warshall.cpp]

<https://pastebin.com/WMgwUzVy>

All-pair Shortest path of Graph

Floyd-Warshall Algorithm



- Dijkstra's Algorithm vs. Floyd-Warshall Algorithm

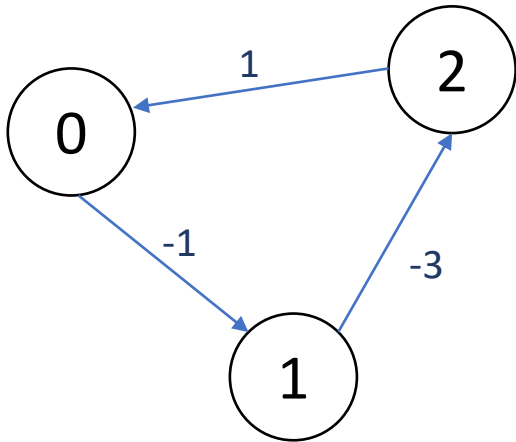
Dijkstra's	Floyd-Warshall
Shortest Path (Single source)	All-pair Shortest Path
$O(E \log V) \rightarrow \text{all nodes} = O(V^3 \log V)$	$O(V^3)$
Not Support negative weights	Support negative weights (if no negative weight cycle detected)

All-pair Shortest path of Graph

Floyd-Warshall Algorithm



- How to detect **negative weight cycle** using Floyd-Warshall Algorithm
 - ปกติแล้ว $\text{dist}[i][j] = 0$ เสมอ เมื่อ $i == j$
 - ถ้าพบว่า $\text{dist}[i][j] < 0$ แสดงว่าพบ **negative weight cycle**



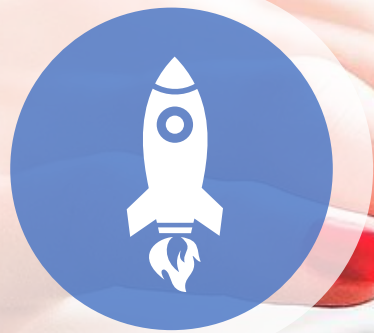
Sum of edges in this cycle = $-1 + -3 + 1 = -3$
(จึงเป็น Negative Weight Cycle)

dist			
	0	1	2
0	0	-1	INF
1	INF	0	-3
2	1	INF	0

$i =$ $j =$ $k =$
if ($\text{dist}[i][k] + \text{dist}[k][j] < \text{dist}[i][j]$)
 $\text{dist}[i][j] = \text{dist}[i][k] + \text{dist}[k][j]$

Topological sort

คืออะไร, Algorithm



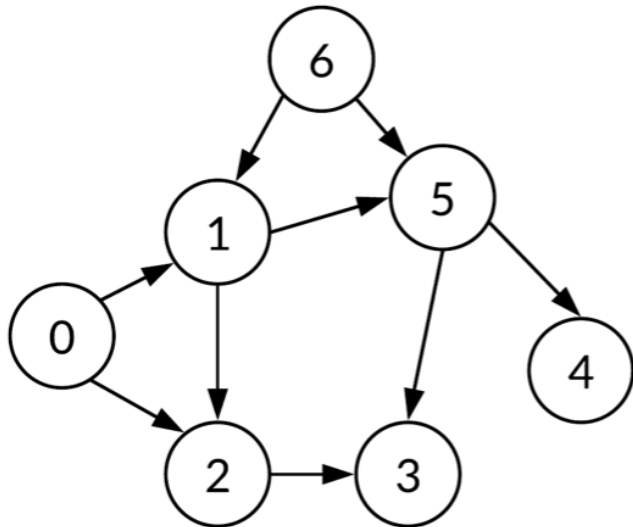
Topological sort of Graph

คืออะไร, Algorithm

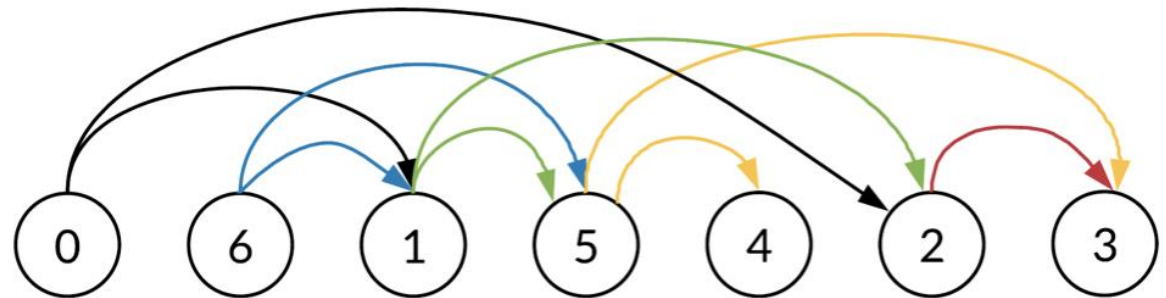


- จาก Directed Acyclic Graph (DAG) $G = (V, E)$
- Directed Acyclic Graph คือ กราฟแบบมีทิศทาง ที่ไม่มี cycle อยู่ข้างใน
- Topological sort of G คือ ลำดับของ Vertex ที่ทุก ๆ $\text{edge}(u, v)$ u จะมาก่อน v เสมอ

Unsorted graph



Topologically sorted graph

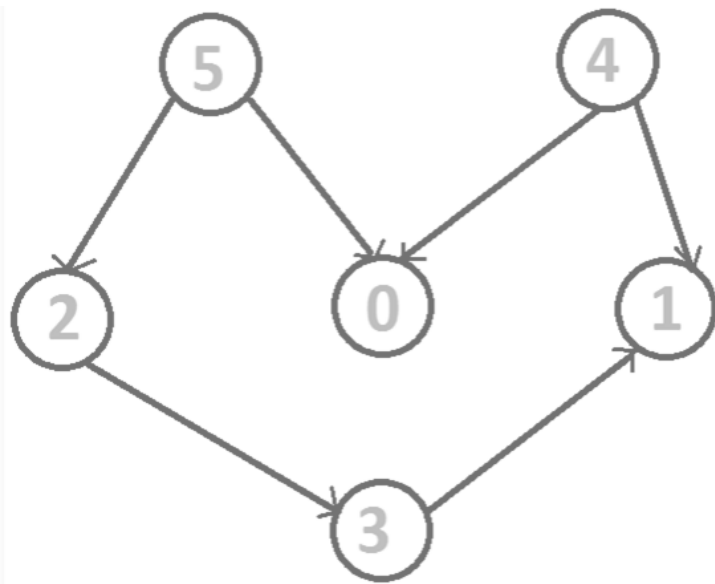


Topological sort of Graph

คืออะไร, Algorithm



- ข้อใดเป็น Topological sort ของกราฟด้านล่างบ้าง?



☐ 5 4 2 3 1 0

☐ 4 5 2 3 1 0

☐ 5 2 3 1 4 0

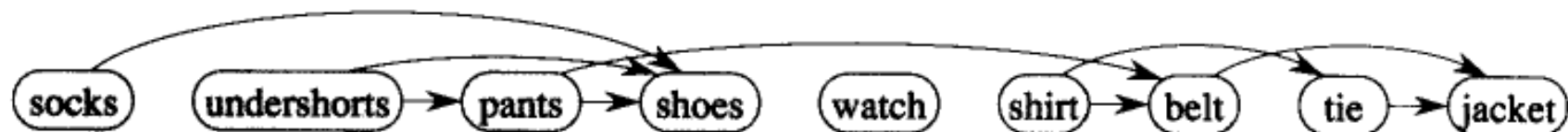
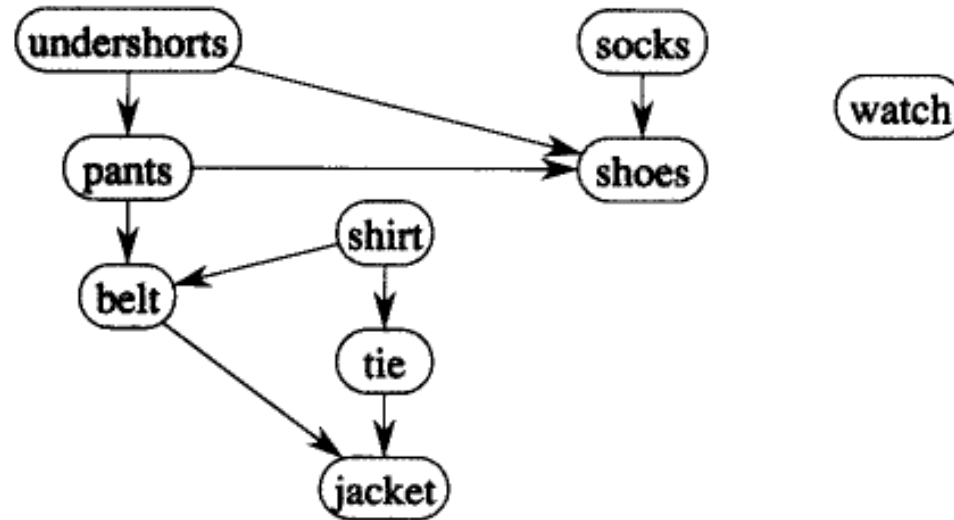
☐ 4 2 3 1 5 0

Topological sort of Graph

คืออะไร, Algorithm



- ประโยชน์

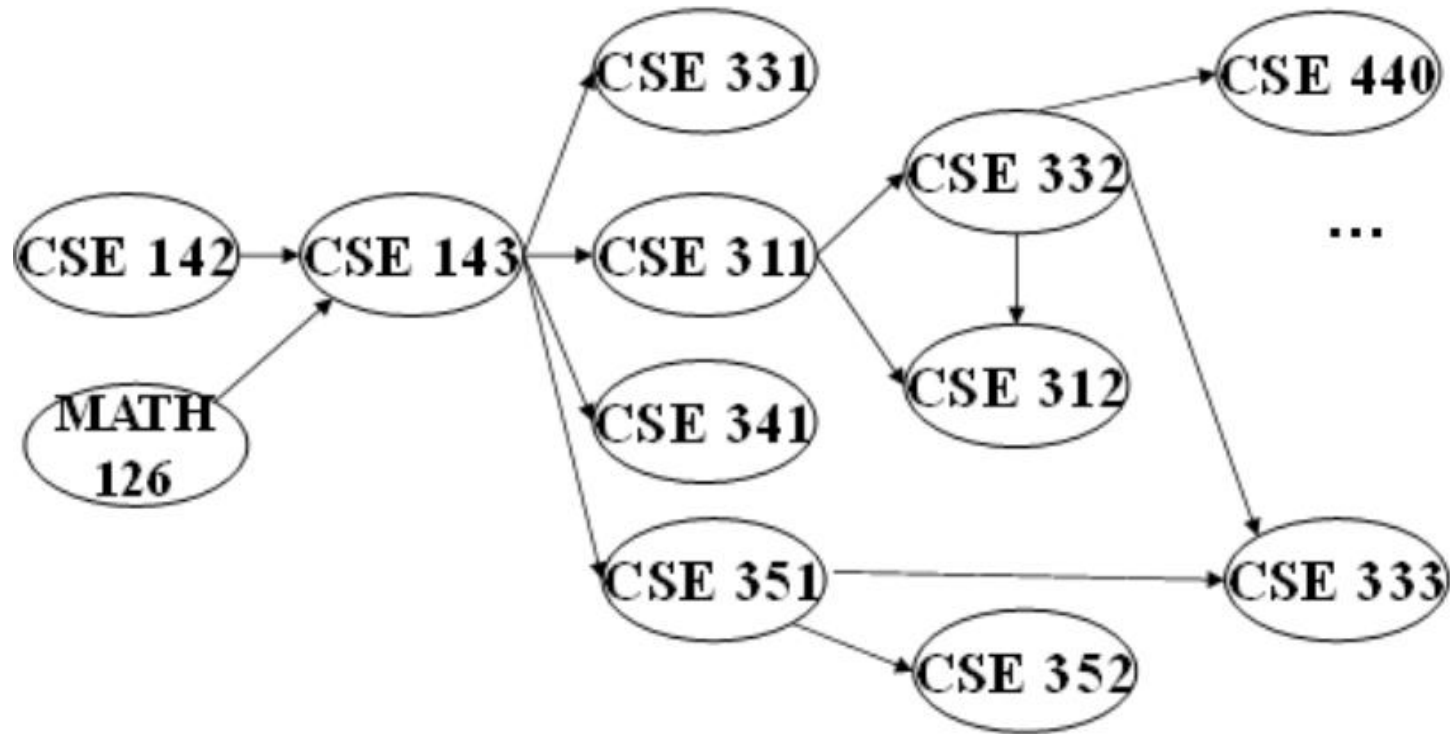


Topological sort of Graph

คืออะไร, Algorithm



- ประโยชน์



Topological sort of Graph

คืออะไร, Algorithm



- วิธีหา Topological sort “ประยุกต์จาก DFS แต่เพิ่ม stack เพื่อเก็บ node ที่เป็นลำดับก่อนหน้าให้ออกมาก่อน ”

DFS:

initialize all in **visited[V]** = false

```
//call recursive function  
DFS_util(source, visited);
```

DFS_util(v, visited):

```
show v;  
mark visited[v] = true;  
For each node in adjacency(v):  
    if visited[node] == false:  
        DFS(node, visited);
```

TopologicalSort:

initialize all in **visited[V]** = false
empty stack S

```
//call recursive function for each Vertex v  
for each v:  
    if visited[v] != true  
        TopologicalSortUtil(v, visited, S)  
  
// print from stack
```

TopologicalSortUtil(v, visited):

```
visited[v] = true // ** Not print yet  
For each node in adjacency(v):  
    if visited[node] == false:  
        TopologicalSortUtil(node, visited, S)  
  
push v to stack S
```

Topological sort of Graph

คืออะไร, Algorithm



TopologicalSort:

initialize all in **visited[V]** = false
empty stack **S**

//call recursive function for each Vertex **v**
for each **v**:

if **visited[v] != true**

TopologicalSortUtil(v, visited, S)

// print from stack

TopologicalSortUtil(v, visited):

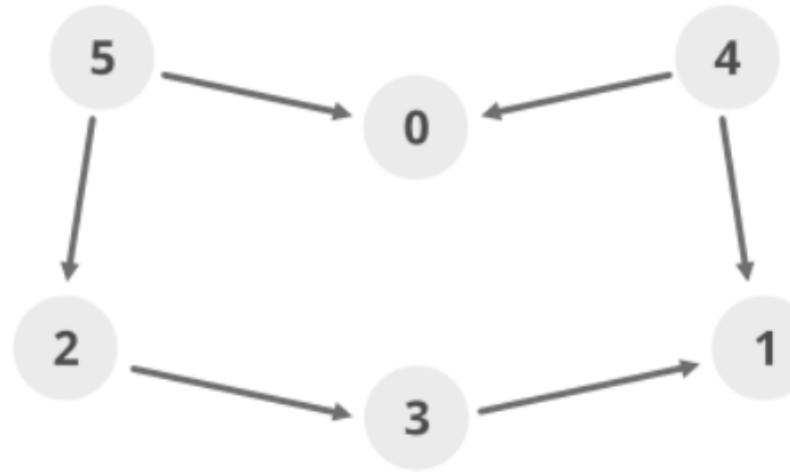
visited[v] = true

for each **node** in **adjacency(v)**:

if **visited[node] == false**:

TopologicalSortUtil(node, visited, S)

push **v** to stack **S**



Adja cent list (G)

- 0 →
- 1 →
- 2 → 3
- 3 → 1
- 4 → 0, 1
- 5 → 2, 0

visited =

0	1	2	3	4	5
F	F	F	F	F	F

stack S = []

Topological sort of Graph

คืออะไร, Algorithm



TopologicalSort:

initialize all in **visited[V]** = false
empty stack **S**

//call recursive function for each Vertex **v**
for each **v**:

if **visited[v] != true**

TopologicalSortUtil(v, visited, S)

// print from stack

TopologicalSortUtil(v, visited):

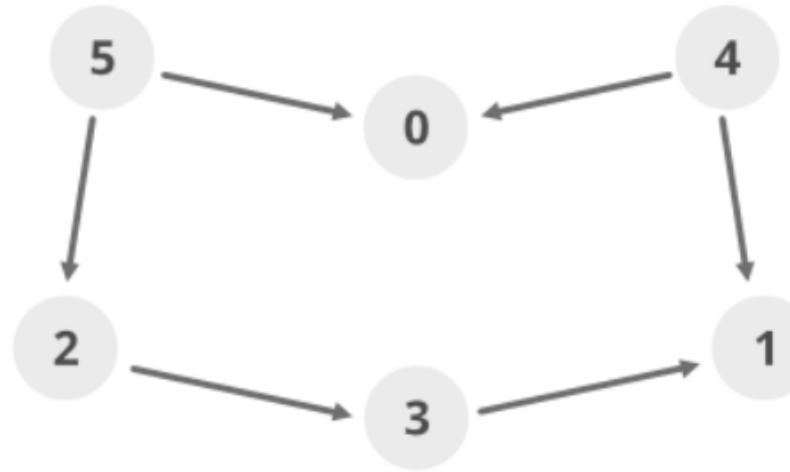
visited[v] = true

for each **node** in **adjacency(v)**:

if **visited[node] == false**:

TopologicalSortUtil(node, visited, S)

push **v** to stack **S**



Adja cent list (G)

0 →
1 →
2 → 3
3 → 1
4 → 0, 1
5 → 2, 0

“ 5 4 2 3 1 0 ”



Thank you

ສູ້ໆ ນະຄະທຸກຄນ

ຄລິບຢ້ອນກລັງ

<https://youtu.be/auAJWyOwU9c>

<https://youtu.be/x45Qrj6PAwE>



Appendix

Shortest path of Graph

Dijkstra's Algorithm



```
#include <bits/stdc++.h>
using namespace std;

typedef pair<int, int> ii;
typedef vector<int> vi;
typedef vector<ii> vii;
#define INF 1000000000

int main() {
    int V, E, s, u, v, w;
    vector<vii> AdjList;
    scanf("%d %d %d", &V, &E, &s);

    AdjList.assign(V, vii());
    // assign blank vectors of pair<int, int>s to AdjList
    for (int i = 0; i < E; i++) {
        scanf("%d %d %d", &u, &v, &w);
        AdjList[u].push_back(ii(v, w));
    }
    // directed graph
}
```

Shortest path of Graph

Dijkstra's Algorithm



```
vi dist(V, INF);
dist[s] = 0; // INF = 1B to avoid overflow
priority_queue< ii, vector<ii>, greater<ii> > pq;
pq.push(ii(0, s)); // ^to sort the pairs by increasing distance from s

while (!pq.empty()) { // main loop
    ii front = pq.top();
    pq.pop(); // greedy: pick shortest unvisited vertex
    int d = front.first, u = front.second;
    if (d > dist[u]) continue; // this check is important, see the explanation
    for (int j = 0; j < (int)AdjList[u].size(); j++) {
        ii v = AdjList[u][j]; // all outgoing edges from u
        if (dist[u] + v.second < dist[v.first]) {
            dist[v.first] = dist[u] + v.second; // relax operation
            pq.push(ii(dist[v.first], v.first));
        }
    }
}

for (int i = 0; i < V; i++)
    printf("SSSP(%d, %d) = %d\n", s, i, dist[i]);

return 0;
}
```

Shortest path of Graph

Bellman Ford Algorithm



```
#include <algorithm>
#include <cstdio>
#include <vector>
#include <queue>
using namespace std;

typedef pair<int, int> ii;
typedef vector<int> vi;
typedef vector<ii> vii;
#define INF 1000000000

int main() {
    int V, E, s, u, v, w;
    vector<vii> AdjList;
    scanf("%d %d %d", &V, &E, &s);

    AdjList.assign(V, vii());
    for (int i = 0; i < E; i++) {
        scanf("%d %d %d", &u, &v, &w);
        AdjList[u].push_back(ii(v, w));
    }
}
```

Shortest path of Graph

Bellman Ford Algorithm



```
// Bellman Ford routine
vi dist(V, INF);
dist[s] = 0;
for (int i = 0; i < V - 1; i++) // relax all E edges V-1 times, overall O(VE)
    for (int u = 0; u < V; u++) // these two loops = O(E)
        for (int j = 0; j < (int)AdjList[u].size(); j++) {
            ii v = AdjList[u][j]; // we can record SP spanning here if needed
            dist[v.first] = min(dist[v.first], dist[u] + v.second); // relax
        }
```

```
bool hasNegativeCycle = false;
for (int u = 0; u < V; u++) // one more pass to check
    for (int j = 0; j < (int)AdjList[u].size(); j++) {
        ii v = AdjList[u][j];
        if (dist[v.first] > dist[u] + v.second) // should be false
            hasNegativeCycle = true; // but if true, then negative cycle exists!
    }
```

```
printf("Negative Cycle Exist? %s\n", hasNegativeCycle ? "Yes" : "No");
```

```
if (!hasNegativeCycle)
    for (int i = 0; i < V; i++)
        printf("SSSP(%d, %d) = %d\n", s, i, dist[i]);
```

```
return 0;
```

```
/*
// Graph has negative weight, but no negative cycle
5 5 0
0 1 1
0 2 10
1 3 2
2 3 -10
3 4 3

// Graph with negative cycle exists
3 3 0
0 1 1000
1 2 15
2 1 -42
*/
```