Marmara University - Faulty of Engineering

Department of Computer Engineering

CSE4219 Principles of Embedded System Design (Fall 2024)

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Arm Cortex M4 Problems

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Sections Of the Report: -

- <u>Section (1):</u> Problem (1) ARM Assembly Program for Repeated Digit Summation
- Section (2): Problem (2) ARM Assembly Program for Matrix Column Swap
- Section (3): Problem (3) ARM Assembly Program for Error Correcting Code (ECC)

All Inputs are given to programs as Data Memory Inputs - For Implementation Simplicity Purposes!!

This method is common in embedded systems, assembly language programming, and certain low-level software development contexts. It simplifies the process by avoiding complex input/output handling at runtime, enabling the program to fetch inputs directly from specified memory addresses in a predictable, structured way.

Section (1): Problem (1) - ARM Assembly Program for Repeated Digit Summation

In this assembly program, we aim to calculate the value of the function F(a,n)=a+aa+aaa+...F(a, n)=a+aa+aaa+... up to nnn terms based on the input arguments a and n. We begin by defining our data section, which holds the input values for a and n.

We initialize a register to hold the total sum and then load the values of a and n from memory into registers for further processing. A loop counter is initialized to zero to track the number of terms processed.

Inside the loop, we compute the new term by multiplying the accumulated sum by 10 and adding the current value of a. This approach builds the terms in the required sequence (e.g., 3, 33, 333) efficiently without needing to construct each term explicitly. We also check if the loop counter has reached n; if so, we exit the loop.

Once all terms have been processed, the final sum is stored in the designated register (r0). This structured approach allows for a clear and efficient calculation of the expression using basic arithmetic operations in assembly language.

Inputs & Outputs:

- 1) $A = 3, B = 5 \rightarrow \text{Output: } 37035 \text{ (In Hex.)}$
- 2) A = 45, $B = 3 \rightarrow Output: 459135 (In Hex.)$

```
INCLUDE core_cm4_constants.s
    INCLUDE stm321476xx_constants.s
    AREA myData, DATA, READWRITE ; define data section
input DCD 45,3; argumments a and n
    AREA P1, CODE
    EXPORT __main
    ALIGN
    ENTRY
 main PROC
    MOV r0, #0; total
    LDR r1, = input; load addr of input
    LDR r1, [r1]
    LDR r2, =input+4
    LDR r2, [r2]
    MOV r3, #0
    MOV r4, #0
    MOV r5, #1
    MOV r6, r1
    MOV r7, #10 ; temp reg to multiply and divide
count_digits
    CMP r6, #0
    BEQ loop; no mmore digits left begin main loop
    MUL r5, r5, r7; multp factor increased
    UDIV r6, r6, r7
    B count_digits
loop
```

```
CMP r2, r3
BEQ stop; if b equals counter branch to stop execution

MLA r4, r4, r5, r1; multiply with accumulate

ADD r0, r0, r4

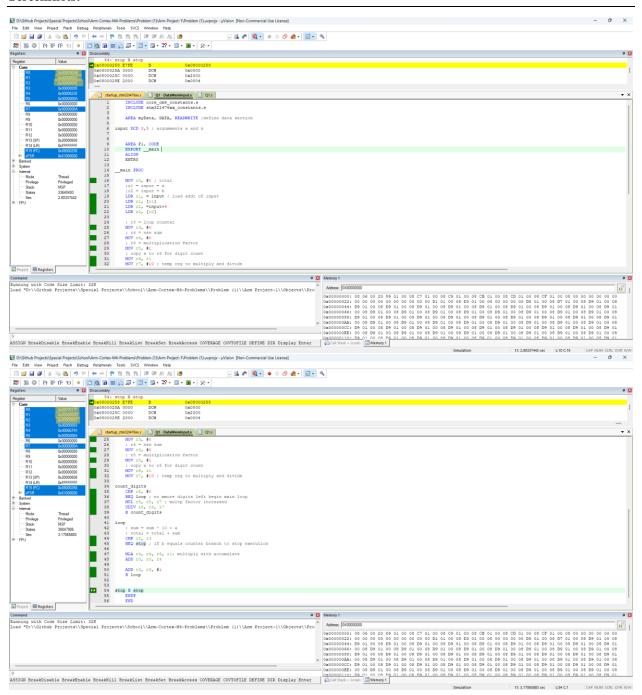
ADD r3, r3, #1
B loop

stop B stop

ENDP

END
```

Screenshots:



Section (2): Problem (2) - ARM Assembly Program for Matrix Column Swap

In this assembly program, our objective is to swap two columns in a 3x3 integer matrix. We start by defining the data section, which includes the matrix initialized with some integer values and a reserved block of memory to store the modified matrix after the swap.

We initialize registers to hold the column indices to be swapped and load the matrix address and the destination memory address for the swapped result. The column indices are converted to the corresponding memory addresses based on the row-major storage of the matrix.

A loop is set up to iterate through the rows of the matrix. For each row, we calculate the memory addresses for the elements in the two columns that need to be swapped. The elements are then loaded from the source matrix, and the swapped values are stored in the destination memory.

This process repeats until all rows have been processed. Once the loop is complete, the program ends with the swapped matrix stored in the designated zeroed memory block. This implementation efficiently handles the column swap using basic arithmetic operations and memory addressing in assembly language.

Inputs & Outputs:

In the program the 2D Matrix is stored in Data Memory Section as a one-dimensional array, each array element is stored in one byte memory size resulting into $9 \times 4 = 36$ bytes of memory for each matrix before swap and after swap.

For Memory offset Calculation for each element in the two columns to be swapped. The following formula is used

$$(i * Num of Columns + j) * 4$$

for addressing each element in a specific row i and column j.

Matrix:

 $\begin{bmatrix} 1 & 2 & 3 \\ 1 & 2 & 3 \\ 1 & 2 & 3 \end{bmatrix}$

After Swapping:

 $\begin{bmatrix} 1 & 3 & 2 \\ 1 & 3 & 2 \\ 1 & 2 & 2 \end{bmatrix}$

Column 1 (2, 2, 2) and Column 2 (3, 3, 3) are swapped.

Before: (1, 2, 3) (1, 2, 3) (1, 2, 3) in Data Memory.

After $\underline{}$: (1, 3, 2) (1, 3, 2) (1, 3, 2) in Data Memory (Zero Memory Partition).

Program Source Code:

INCLUDE core_cm4_constants.s
INCLUDE stm321476xx_constants.s

AREA myData, DATA, READWRITE ;data section

Matrix DCD 1,2,3,1,2,3,1,2,3 ;row major each is a word for integer so total 36 bytes

zMem SPACE 36 ;reserving zeroed memory in data to store the modified matrix

```
AREA P2, CODE
    EXPORT __main
    ALIGN
    ENTRY
 main PROC
   MOV r0, #1; col1 = 2
   MOV r1, #2; col2 = 1
    LDR r2, =Matrix; get matrix address
    LDR r3, =zMem
    MOV r4, r0; get col 1 first element index in the matrix (j1)
    MOV r5, r1; get col 2 first element index in the matrix (j2)
    MOV r6, #0; row counter (i)
    MOV r8, #3; to multiply later for next element address
    ADD r11, r0, r1
    SUB r11, r8, r11
loop
   CMP r6, #3
    BEQ stop; if processed all rows exit
   MOV r7, r6
   MUL r7, r7, r8; row * 3
    ADD r4, r7, r0
    LSL r4, r4, #2
    ADD r5, r7, r1
    LSL r5, r5, #2
    ADD r12, r7, r11
```

```
LSL r12, r12, #2

; Load elements to swap

LDR r9, [r2, r4]

LDR r10, [r2, r5]

STR r10, [r3, r4]; now swap by storing to location of other column's element

STR r9, [r3, r5]

LDR r9, [r2, r12]; store remaining element

STR r9, [r3, r12]

ADD r6, r6, #1

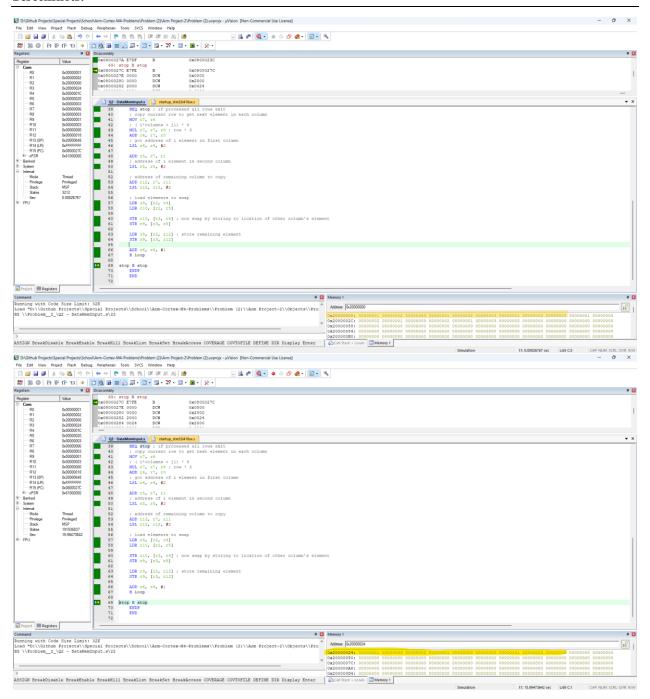
B loop

Stop B stop

ENDP

END
```

Screenshots:



Section (3): Problem (3) - ARM Assembly Program for Error Correcting Code (ECC)

In this assembly program, we start by defining our data section, which includes our 8-bit input data and a reserved zeroed memory location for storing the expanded result.

First, we load the 8-bit input value into a register and extract the bits into a 13-bit register by shifting the original value multiple times. This approach simplifies the process and avoids the need to access data memory repeatedly for bit manipulation.

Next, we apply a masking technique to isolate the bits relevant for each parity bit. For each required parity bit, we perform an AND operation between the current register (which contains the original data bits) and the mask corresponding to the parity bit position. This allows us to check only the bits that the parity bit is responsible for.

To compute the parity bits, we utilize an XOR operation to simulate even counting. By XORing each bit in the masked registers and shifting them right, we accumulate the count of 1s into a temporary register (r9). After processing all bits, we check the parity result stored in r9. If the count is odd, we set the corresponding parity bit to 1 by performing an OR operation with the expanded data register (r3). Finally, we store the complete 13-bit expanded data, including the original data bits and the calculated parity bits, into our reserved memory location.

This method ensures an efficient calculation of the error-correcting codes while maintaining clarity and simplicity in the implementation.

Program Source Code:

```
INCLUDE core_cm4_constants.s
   INCLUDE stm321476xx_constants.s

AREA myData, DATA, READWRITE ; data section

input DCB 0xB3 ; 8 bit data input
zMem SPACE 4  ; reserving zeroed memory in data to store the expanded data

AREA P3, CODE
   EXPORT __main
   ALIGN
   ENTRY

__main PROC

LDR r0, =input
   MOV r1, #0
   LDRB r1, [r0] ; load input data to a register
   LDR r2, =zMem

MOV r3, #0 ; 13-bit extended
   ; extract bits and store inside r3
   ; bit 7 to 12
```

```
MOV r4, r1, LSR #7
ORR r3, r3, r4, LSL #12
MOV r4, r1, LSR #6
AND r4, r4, #1; mask LSB
ORR r3, r3, r4, LSL #11
MOV r4, r1, LSR #5
AND r4, r4, #1; mask LSB
ORR r3, r3, r4, LSL #10
MOV r4, r1, LSR #4
AND r4, r4, #1; mask LSB
ORR r3, r3, r4, LSL #9
MOV r4, r1, LSR #3
AND r4, r4, #1; mask LSB
ORR r3, r3, r4, LSL #7
MOV r4, r1, LSR #2
AND r4, r4, #1; mask LSB
ORR r3, r3, r4, LSL #6
MOV r4, r1, LSR #1
AND r4, r4, #1; mask LSB
ORR r3, r3, r4, LSL #5
MOV r4, r1
AND r4, r4, #1; mask LSB
ORR r3, r3, r4, LSL #3
LDR r4, =0x0AA8; p1
LDR r5, =0x0CC8; p2
LDR r6, =0x10E0 ; p4
LDR r7, =0x1E00; p8
LDR r8, =0x1FFE; p0
AND r4, r3, r4; its for checking p1
AND r5, r3, r5; p2
AND r6, r3, r6; P4
AND r7, r3, r7 ; p8
AND r8, r3, r8; p0
MOV r9, #0; to use in xor to check even parity
```

```
parity1
    ANDS r10, r4, #1
    EOR r9, r9, r10
    LSR r4, r4, #1
    CMP r4, #0; check if checked all bits
    BNE parity1
    ORR r3, r3, r9, LSL #1
    MOV r9, #0; to use in xor to check even parity
parity2
    ANDS r10, r5, #1
    EOR r9, r9, r10
    LSR r5, r5, #1
    CMP r5, #0; check if checked all bits
    BNE parity2
   ORR r3, r3, r9, LSL #2
   MOV r9, #0; to use in xor to check even parity
parity4
    ANDS r10, r6, #1
    EOR r9, r9, r10
    LSR r6, r6, #1
    CMP r6, #0; check if checked all bits
    BNE parity4
   ORR r3, r3, r9, LSL #4
   MOV r9, #0; to use in xor to check even parity
parity8
    ANDS r10, r7, #1
    EOR r9, r9, r10
    LSR r7, r7, #1
    CMP r7, #0; check if checked all bits
    BNE parity8
    ORR r3, r3, r9, LSL #8
```

```
; next parity check for p2
   MOV r9, #0; to use in xor to check even parity
parity0
   ANDS r10, r8, #1
   EOR r9, r9, r10
   LSR r8, r8, #1
   CMP r8, #0; check if checked all bits
   BNE parity0
   ; if data contain odd number r9 is 1 and store one into p1 (bit 1)
   ORR r3, r3, r9, LSL #0

   ; store expanded data 13 bit
   STR r3, [r2]

stop B stop
   ENDP
   ENDP
   ENDP
```

Screenshots:

