



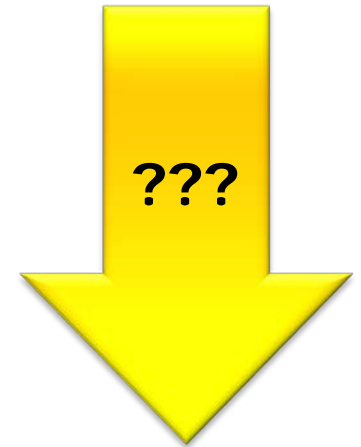
Reentrancy, and all the other things...

"Everything you were afraid to ask ..."

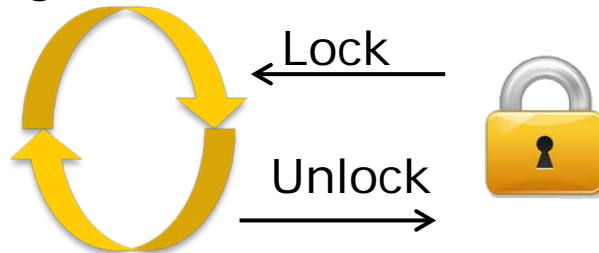
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Learning Goals

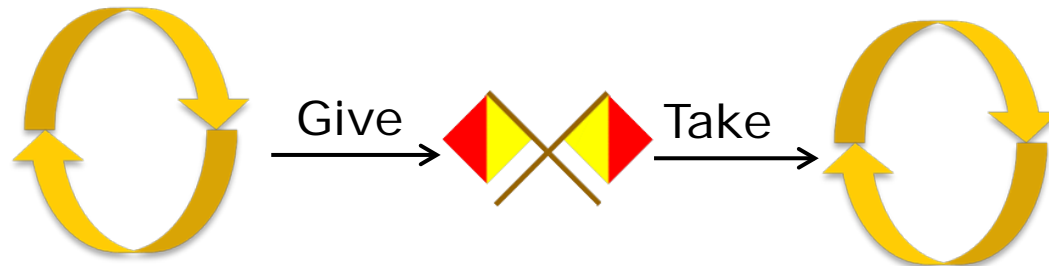
- Reentrancy
- Critical Section
- Thread Safe
- Semaphore
- Mutex
- Thread Safe
- FreeRTOS Implementation



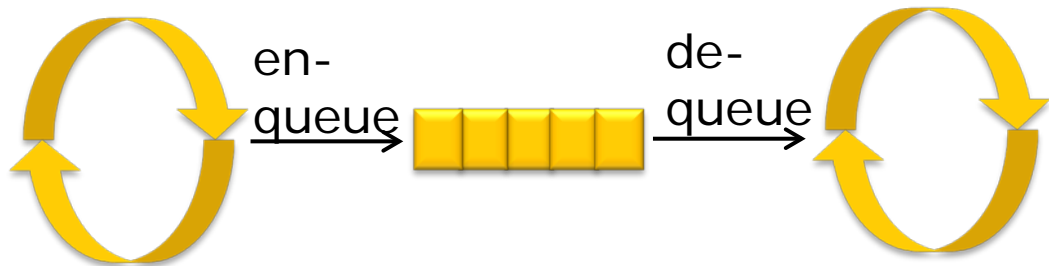
Synchronization Primitives



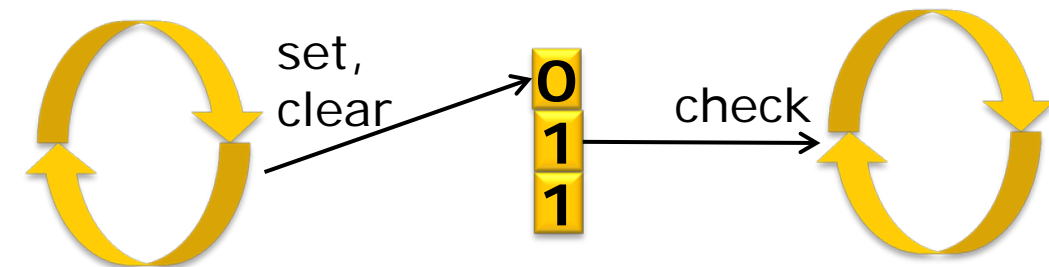
- **Mutex**
 - mutual exclusion



- **Semaphore**
 - Sync, Notify



- **Queues**
 - Communication



- **Flags**
 - Synchronization

Reentrancy

Behaviour of a programm

- Attribute of a program or subroutine
- Can be interrupted in the middle of execution
 - thread/task
 - interrupt
- **Reentrant**: Can be safely called (re-entered) by other thread/task or interrupt

```
int var;  
  
void decrement(void) {  
    if (var>0) {  
        var--;  
    }  
}
```

it's not reentrant, because can not interrupt the sequenz of this code



Critical Section

- Sequence of code
- Protected against concurrent execution
- Only **one** program flow is inside critical section
- Used to protect access to shared resource

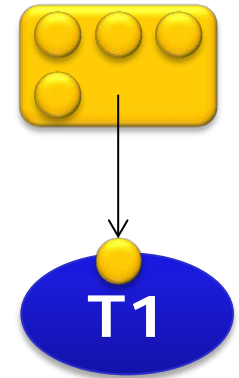
IMPLEMENTATION with: Disable Interrupts, EnterCritical(),...

```
int var;  
  
void decrement(void) {  
    if (var>0) {  
        var--;  
    }  
}
```

**Critical Section
Needed**

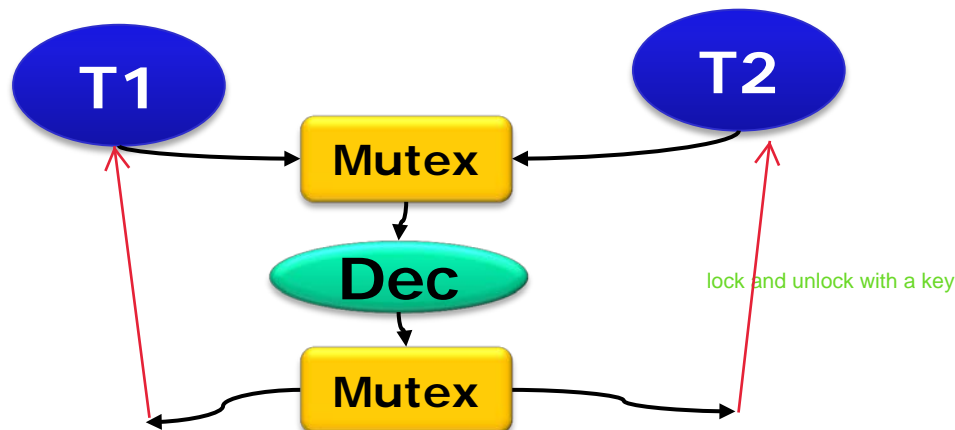
Semaphore

- Variable or abstract data type
- Used for synchronization
- Used to control access to a shared resource
 - By tasks, threads and interrupts
- **Can** be used to implement a Critical Section
- Binary and Counting semaphore
- Example **counting** semaphore
 - N Study Rooms
 - Students need to request and release room at front desk
 - Front desk decreases/increases number of available rooms
 - ➔ how many, not which room



Mutual Exclusion, Mutex

- Mutual Exclusion: Property of concurrency control to establish a critical section
- Establishes mutual exclusive execution of program sequence
- **Mutex**: abstract data type used for Mutual Exclusion
- Used for
 - Synchronization
 - Preventing race conditions



Thread Safe

- Attribute of a program or subroutine
- Guarantees safe execution by multiple **threads**
- Closely related to Reentrancy
- **Not** the same as Reentrancy in reentrancy we consider interrupts as well
 - Does not include the presence of interrupts
 - Thread might use Mutex to be thread-safe
 - Interrupt could run into Mutex (starves/blocks)
 - ➔ not safe!

```
int var;  
  
void decrement(void) {  
    LockMutex();  
    if (var>0) {  
        var--;  
    }  
    ReleaseMutex();  
}
```

typically we have a LockMutex() from ISR and a LockMutex() like that

Implementation in FreeRTOS

- Implemented as Queues with no data
- **Semaphore** does not implement priority inheritance
 - Binary, Counting binary (one flag one use)
 - Must not be returned
 - Used for critical sections and message passing
- **Mutex**
 - Binary (normal) and Recursive
 - Implements **priority inheritance**
 - MUST be returned
 - Used for critical sections

Quiz:

`taskDISABLE_INTERRUPTS()` -> just disable the interrupts, not designed for netsting, no context switch

`taskExitCritical` -> inside the nested section, allowed in a nested way

`vTaskSuspendAll()` -> stopps the scheduler from, context switch, really disable the scheduler, all the rest is running (lika a car in front of a red trafficlight)

How are context switch handled? What can trigger a context-switch -> interrupt SysTick, `yield()`, API call (`vTaskDelay()`)

Task context switch! -> ABKLÄREN

BASEPRI (M4) & PRIMASK (on our M0)

operating system uses `taskExitCritical`

Discussion: File System

- Discuss in Groups
 - File System, SPI bus to memory/SD Card
 - Multiple task using file system
 - open/write/read/close file(s)
- Identify Needs
 - Reentrancy, Critical Section,
 - Semaphore, Mutex

