A presentaion on A* search algorithm

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Introduction



Figure: Using Navigation Assistants

Definition

What is Graph Traversal?

Graph traversal (also known as graph search) refers to the process of visiting (checking and/or updating) each vertex in a graph.

Definition



Figure: Using Graph Traversal Algorithms

Definition

What is A* search?

A* search is a path-finding and graph traversal algorithm which is basically an evolution of Dijkstra's algorithm

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Why A* search?

As we have already mentioned, A* search is basically a modification of Dijkstra's algorithm.

■ But why do we need A* search?

Why A* search?

As we have already mentioned, A* search is basically a modification of Dijkstra's algorithm.

- But why do we need A* search?
- What are the benefits of A* over Dijkstra's algorithm or other such algorithms which validates the modification?

A* search Overview

To answer that question we first need to know how the A* search algorithm works

■ You see, Dijkstra's algorithm is a greedy algorithm and so it always immediately chooses the best possible path at every step

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- You see, Dijkstra's algorithm is a greedy algorithm and so it always immediately chooses the best possible path at every step
- Even if the chosen path strays away from the destination, it isn't smart enough to detect that
- That's why it needs to visit a lot more nodes to come to a conclusive decision



Figure: A* search vs Dijkstra

How does A* manage to be this efficient? HEURISTICS

What is Heuristics?

 Heuristics is an approach to problem solving, learning, or discovery



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- Heuristics is an approach to problem solving, learning, or discovery
- It employs a practical method not guaranteed to be optimal or perfect
- But it is sufficient for the immediate goals



Figure: Using Heuristics

Let's come back to this diagram again

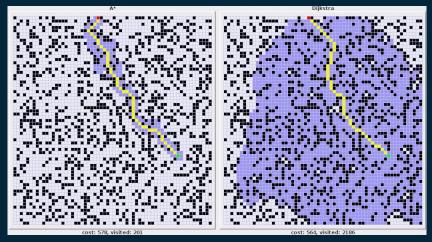


Figure: A* search vs Dijkstra

■ How is Heuristics used in A* search?

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- It can be of many different kinds



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- It can be of many different kinds
- For example, the Euclidean distance can be an obvious one

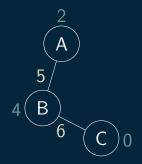
- How is Heuristics used in A* search?
- It can be of many different kinds
- For example, the Euclidean distance can be an obvious one
- The shorter the Euclidean distance of a node from the destination the more likely it is to be included in the optimal path



Figure: Using Euclidean Distance as a Heuristic measure

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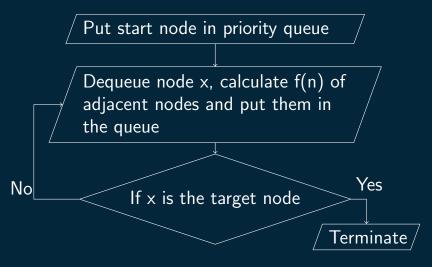
A* Search Algorithm



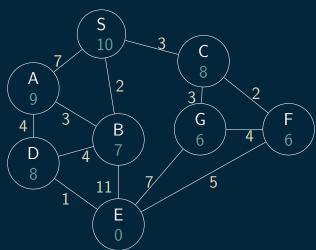
$$f(n) = d(n) + h(n)$$



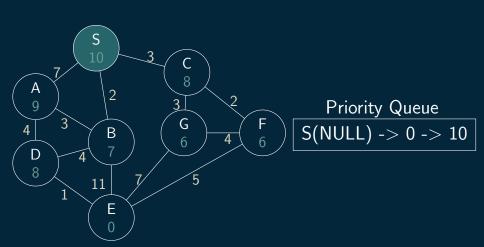
Flow Chart

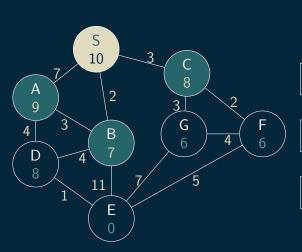


Graph



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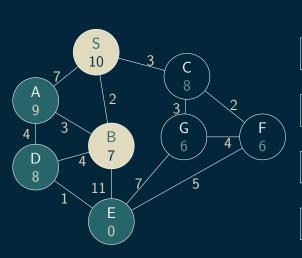




$$B(S) -> 2 -> 9$$

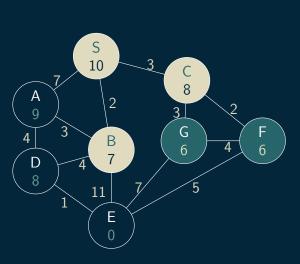
$$C(S) -> 3 -> 11$$

$$A(S) -> 7 -> 16$$



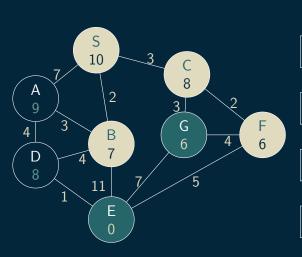
$$A(B) -> 5 -> 14$$

$$D(B) -> 6 -> 14$$



$$G(C) -> 6 -> 12$$

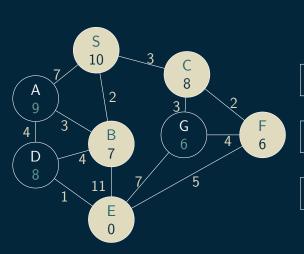
$$A(B) -> 5 -> 14$$



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$$G(C) -> 6 -> 12$$

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$$D(B) -> 6 -> 14$$

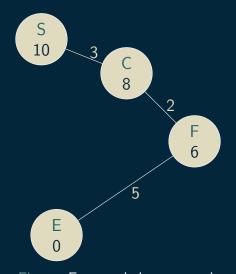


Figure: Expected shortest path

Exact Heuristics

Using exact distance from any node n to goal node as h(n)

Precompute distances between each pair of nodes

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- Precompute distances between each pair of nodes
- Time-consuming

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Exact Heuristics

Using exact distance from any node n to goal node as h(n)

- Precompute distances between each pair of nodes
- Time-consuming
- Equal to Euclidean distance if there is no other node on the direct path

Manhattan Distance

$$h(n) = | n.x - g.x | + | n.y - g.y |$$

- Suitable for 2D grids
- Movement is allowed only in x or y axis (4 directions)

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Diagonal Distance

$$h(n) = \max\{| n.x - g.x |, | n.y - g.y |\}$$

- Suitable for games like chess
- Both axial and diagonal movements are allowed (8 directions)

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Euclidean Distance

$$h(n) = \sqrt{(n.x-g.x)^2 + (n.y-g.y)^2}$$

- Suitable for graphs
- Movement in any direction is allowed

Implementation

Implementation details that can affect the performance of A*

Breaking ties in deciding minimum element of priority queue

Implementation

Implementation details that can affect the performance of A*

- Breaking ties in deciding minimum element of priority queue
- Alternative heap implementations of priority queue

Breaking Ties

How to break ties in priority queue?

- lowest heuristic cost
 - Prefer the nodes closest to the goal

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Breaking Ties

How to break ties in priority queue?

- lowest heuristic cost
 - Prefer the nodes closest to the goal
- cross product of s-g and n-g vectors
 - Prefer paths on the start-goal straight line
- LIFO like DFS
 - Prefer new insertions to old insertions



Which heap to use?

Binary Heap

VS

Fibonacci Heap

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Which heap to use?

Operation	Delete-min	Insert	Decrease- key
Binary Heap	$\Theta(\log n)$	$\mathcal{O}(\log n)$	$\mathcal{O}(\log n)$
Fibonacci Heap	Amortized $\mathcal{O}(\log n)$	Θ(1)	Amortized $\Theta(1)$

Table: Time-complexity comparisons

Which heap to use?

Binary Heap

VS

Fibonacci Heap

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Time Complexity

Branching Factor (b)

Average number of neighbours at each node

Time Complexity

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Average number of neighbours at each node

Depth of Solution (d)

Number of nodes in the shortest path

Time Complexity

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Average number of neighbours at each node

Depth of Solution (d)

Number of nodes in the shortest path

Time Complexity : $\mathcal{O}(b^d)$ or $\mathcal{O}(\mid E\mid)$

A* search





Figure: Tower Defense

- Finding shortest path in games
 - mostly single source, single destination

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- Graph Traversal
 - originally designed for this purpose



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 - informed search algorithm for agents

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- Graph Traversal
 - originally designed for this purpose
- Artificial Intelligence
 - informed search algorithm for agents
- Navigation
 - Google Maps & other digital maps

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THE END

Thank You! Feel free to ask any questions