Unique Universal Binary Relations

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In this paper I introduce unique universal binary relations, which simplifies abstract generalizations.

A binary relation is an ordered pair:

This can also be modeled as a predicate `p`:

$$p(a, b)$$
: bool

A *universal* binary relation can infer `p` from `b`:

$$role_of(b) = p$$

One says that `b` is assigned the "role" `p`. For any $q < \neg = p$, q(a, b) = false. A *unique* binary relation has the property that `b` can be inferred from p(a):

$$p(a) = b$$

Unique universal binary relations also permit multiple predicates, e.g. p(a) = b and q(a) = c.

Members of types are unique universal binary relations:

false : bool	type_of(false, bool)	type_of(false) = bool
true : bool	type_of(true, bool)	type_of(true) = bool

The role of `bool` is `type_of`. Every relation to `bool` is a type judgement.

The values `false` and `true` can also be assigned a role `value_of` to perform computations:

```
not(false) = true value_of(not(false), true) value_of(not(false)) = true
not(true) = false value_of(not(true), false) value_of(not(true)) = false
```

Now, look at the following law in Type Theory for Cartesian products:

$$\forall x : X, y : Y \{ (x, y) : (X, Y) \}$$

This law can be generalized over all Cartesian products, without quantification over all predicates:

type_of((not(false), not(true))) = (type_of(not(false)), type_of(not(true)))

$$\forall$$
 X, Y { role_of(X) == role_of(Y) => role_of((X, Y)) == role_of(X) } lift role
 \forall x, y, X, Y { (x, X) \land (y, Y) => ((x, y), (X, Y)) } lift role
 value_of((not(false), not(true))) = (true, false)