

Geometric Paths

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In this paper I introduce the concept of geometric paths, used to analyze path-connected spaces. I also introduce the concept of anageometric paths, a generalization of geometric paths.

A geometric path `f` is a higher order function constructed by an axiom function `f₁`:

$$f_1(a, b) \iff \text{continuum}(a, b)$$

$$f_1(b, c) \iff \text{continuum}(b, c)$$

$$f_1(_, _) \iff \text{false}_1$$

Example of an axiom that defines conditional continuum between some points. All other points are disconnected.

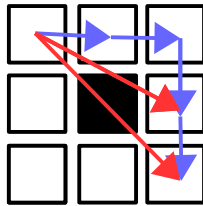
$$f := \lambda(a : T, c : T) = \lambda(g : T \rightarrow \text{bool}) = f_1(a, c)(g) \vee \exists b : g \{ f(a, b)(g) \wedge f(b, c)(g) \}$$

$$f : T \times T \rightarrow (T \rightarrow \text{bool}) \rightarrow \text{bool}$$

$$\text{continuum} := \lambda(a : T, b : T) = \lambda(g : T \rightarrow \text{bool}) = g(a) \wedge g(b)$$

It is only permitted to use `continuum` or `false₁` for any pair of inputs in `f₁`.

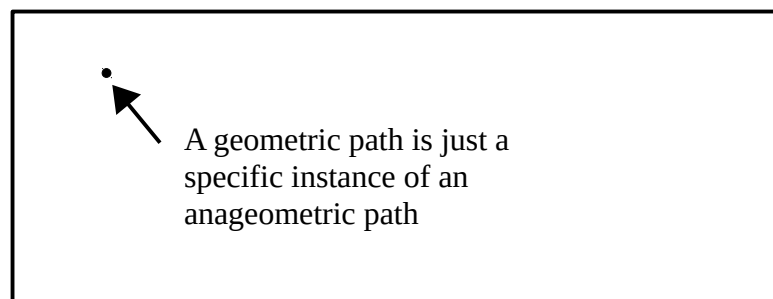
The geometric path can be thought of as connecting a space by composing small arrows:



When there exist a path of **blue** arrows, then the space is connected by “filling in” with **red** arrows. The space is undirected, such that all arrows can be flipped around.

A geometric path is a *higher order unlabeled undirected graph* for any function of type `T → bool`.

Each geometric path is part of an infinite family of similar functions, called “anageometric path”:



By anageometric path we mean the infinite family of functions similar to a geometric path, constructed by same axiom

The rest of this paper explains how an anageometric path is constructed.

Assume one has a property `g` that defines some sub-type of 3 variables `a`, `b` and `c`:

$$a : [g] a' \quad b : [g] b' \quad c : [g] c'$$

$$g : T \rightarrow \text{bool}$$

All possible interpretations of the order “abc” can be given descriptive names:

abc => a'b'c'

000	external
001	tip
010	bridge
011	tail
100	head
101	wall
110	neg tip
111	internal

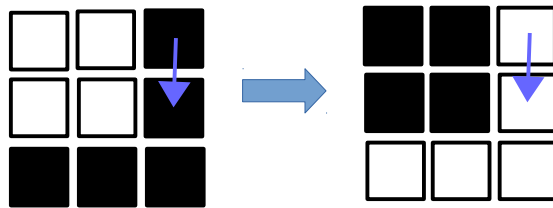
A geometric path is the interpretation “internal”, since the property `g` must hold for all points.

It is common to write the geometric path the following way:

$$f_{111} \quad \text{Using the interpretation code `111` in the anageometric path}$$

By swapping `g` with `¬g`, one gets the interpretation “external” (`¬g` <=> not · g`):

$$f_{000} := \lambda(a : T, b : T) = \lambda(g : T \rightarrow \text{bool}) = f_{111}(a, b)(\neg g)$$

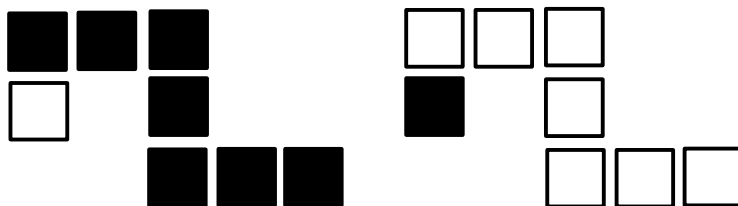


The paths in these two spaces are very different, even the relation between the spaces is simple.

A “head” can be thought of as a path starting inside but then immediately goes outside:

$$f_{100} := \lambda(a : T, c : T) = \lambda(g : T \rightarrow \text{bool}) = \neg g(a) \wedge f_1(a, b)(\text{true}_1) \wedge \exists b : g \{ f_{000}(b, c)(g) \}$$

$$f_{011} := \lambda(a : T, b : T) = \lambda(g : T \rightarrow \text{bool}) = f_{100}(a, b)(\neg g)$$



The “neg tip” and “tip” are constructed in a similar way:

$$\begin{aligned} f_{110} &:= \lambda(a : T, c : T) = \lambda(g : T \rightarrow \text{bool}) = \neg g(c) \wedge \exists b : g \{ f_{000}(a, b)(g) \wedge \neg f_{000}(b, c)(\text{true}_1) \} \\ f_{001} &:= \lambda(a : T, b : T) = \lambda(g : T \rightarrow \text{bool}) = f_{110}(a, b)(\neg g) \end{aligned}$$

The “wall” and “bridge” are constructed the following way:

$$\begin{aligned} f_{101} &:= \lambda(a : T, c : T) = \lambda(g : T \rightarrow \text{bool}) = g(a) \wedge g(c) \wedge \exists b : \neg g \{ f_1(a, b)(\text{true}_1) \wedge f_1(b, c)(\text{true}_1) \} \\ f_{010} &:= \lambda(a : T, b : T) = \lambda(g : T \rightarrow \text{bool}) = f_{101}(a, b)(\neg g) \end{aligned}$$

There is another way of constructing a “wall”, without using `f₁`:

$$\begin{aligned} f_{101} &:= \lambda(a : T, c : T) = \lambda(g : T \rightarrow \text{bool}) = \\ &\quad \exists b : g \{ f_{100}(a, b)(g) \wedge f_{110}(a, b)(g) \wedge f_{100}(b, c)(g) \wedge f_{001}(b, c)(g) \} \end{aligned}$$

One then generalizes this construction to an interpretation code of arbitrary length:

- Topological equivalent binary numbers up to two bits
- A single bit is interpreted as “wall” or “bridge” respectively

For example:

$$\begin{aligned} 00000100000 &\Rightarrow 00100 \\ 00010100111 &\Rightarrow 001010011 \end{aligned}$$

This topological transform can be described by the following L-system using two rules:

$$\begin{aligned} 000 &\Rightarrow 00 && \text{if code is longer than 3} \\ 111 &\Rightarrow 11 && \text{if code is longer than 3} \end{aligned}$$

The anageometric path is the **infinite family of functions** from this construction:

$$f_x \quad \text{where `x` is some interpretation code of the anageometric path}$$

For example, to ask whether some space `f` filled with `g` connects `a` to `b` by the following:

$$(a, b) : (f_{1101001} g) \quad \Leftrightarrow \quad (a, b) : [f_{1101001} g] \text{ true}$$

Here, `a` starts inside the space, then after an internal path goes through a wall, a bridge, an external path and finally inside again to `b`:



With no access to `f₁`, it is not possible to construct the anageometric path, because one can not find two points that are only locally connected but not connected through any longer path.