

# Complexity of Path Semantical Logic

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*In this paper I derive the time complexity of Path Semantical Logic for level 1 and 0.*

Path Semantical Logic separates propositions into levels, such that an equality between two propositions in level  $N+1$ , propagates into equality between uniquely associated propositions in level  $N$ .

In the paper “Faster Brute Force Proofs”<sup>[1]</sup>, I suggested a method to speed up brute force propositional logic separating propositions into levels 1 and 0. To show the potential for improved performance, I provided a table with fractions<sup>[2]</sup> as a function of  $|F|$  (level 1) and  $|X|$  (level 0). However, I did not provide a method to calculate this fraction directly without counting. Since the denominator of the fraction is simply  $2^{|F|+|X|}$ , only a formula for computing the numerator is missing.

The formula for computing the numerator is also a measure of the time complexity for level 1 and 0.

There are two conditions that determines whether a case should be counted:

1. No more than one zero in  $F$
2. All  $X$ 's are equal

These two conditions are combined using logical OR.

For simplicity, I will refer to the first condition as  $O(F)$  and the second as  $O(X)$ .

It is simpler to analyze the individual conditions and use the following relationship:

$$O(F \vee X) = O(F) + O(X) - O(F \wedge X)$$

The simplest term is  $O(X)$ , because there are only two solutions  $00000\dots 0$  and  $11111\dots 1$ . For each solution, one enumerates the entire  $F$ :

$$O(X) = 2 \cdot 2^{|F|} = 2^{1+|F|}$$

To compute  $O(F)$ , one enumerates the entire  $X$  when there are zero or one  $0$ 's in  $F$ :

1111...1      0111...1      1011...1      1101...1      ...

$$O(F) = (1 + |F|) \cdot 2^{|X|}$$

The third term  $O(F \wedge X)$  is the intersection, which is simply a product of  $2$  and  $1 + |F|$ :

$$O(F \wedge X) = 2 \cdot (1 + |F|)$$

Finally, putting this together and simplifying  $O(F \vee X)$  to  $O(|F|, |X|)$ :

$$O(|F|, |X|) = 2^{1+|F|} + (1 + |F|) \cdot (2^{|X|} - 2)$$

## References:

- [1] “Faster Brute Force Proofs”  
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[https://github.com/advancedresearch/path\\_semantics/blob/master/papers-wip/faster-brute-force-proofs.pdf](https://github.com/advancedresearch/path_semantics/blob/master/papers-wip/faster-brute-force-proofs.pdf)
  
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