## **Univalent Involutions**

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*In this paper I prove that isomorphisms are isomorphic to Univalent Involutions.* 

An isomorphism<sup>[1]</sup> from a Category<sup>[2]</sup> theoretic view is a morphism `f` with an inverse `g` such that:

$$f \cdot g \stackrel{\text{def}}{=} id_{B}$$

$$g \cdot f \stackrel{\text{def}}{=} id_{A}$$

$$f : A \rightarrow B$$

$$g : B \rightarrow A$$

An involution<sup>[3]</sup> is a morphism `h` such that:

$$h \cdot h \le id_T$$
  
 $h : T \rightarrow T$ 

Every involution is an isomorphism, but not every isomorphism is an involution.

It turns out that every isomorphism can be turned into a Univalent Involution:

 $h := (x : T) = if let some(a) = h'_{A}^{-1}(x) \{ h'_{B}(f(a)) \}$ 

The univalent involution `h` has the following normal paths:

$$h[h'_{A}^{-1} \rightarrow h'_{B}^{-1}] \le f[some]$$
  
 $h[h'_{B}^{-1} \rightarrow h'_{A}^{-1}] \le g[some]$ 

A Univalent Involution differs from ordinary involutions by the property it can be turned back into a heterogenous isomorphism, kind of like a tuple `(a, b)` can be turned into `a` and `b`:

Since equality in Intuitionistic Logic [4] using types is a tuple (f, g), this particular form of involution is thought to be univalent [5].

## **References:**

[1]	"Isomorphism"
	Wikipedia
	https://en.wikipedia.org/wiki/Isomorphism

- [2] "Category (mathematics)"
  Wikipedia
  <a href="https://en.wikipedia.org/wiki/Category">https://en.wikipedia.org/wiki/Category</a> (mathematics)
- [3] "Involution (mathematics)"
  Wikipedia
  https://en.wikipedia.org/wiki/Involution (mathematics)
- [4] "Intuitionistic logic"
  Wikipedia
  https://en.wikipedia.org/wiki/Intuitionistic\_logic
- [5] "Homotopy type theory"
  Wikipedia
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