Probabilistic Sub-Types

by Sven Nilsen, 2019

In this paper I introduce a notation for probabilistic sub-types.

Probabilistic sub-types is a generalization of notation for sub-types which truth value is `true` or `false`. Instead of using a boolean as a truth value, one uses probabilities. Formally, a probabilistic sub-type is a statement that has a truth value in the unit interval `[0, 1]` of real numbers:

```
[f]_p a : prob
prob := \(x : real) = (x \ge 0) \land (x \le 1)
```

The interpretation of sub-types which truth value is a boolean is referred to as "boolean sub-types". To keep probabilistic sub-types compatible with the notation for boolean sub-types, a sub-script `p` is written after the square bracket. This means that the sub-type is interpreted as a probability.

$$[f]_p a = (\exists_p f)(a)$$
 The probability that `f` returns `a` $\neg [f]_p a = 1 - (\exists_p f)(a)$ The probability that `f` does not return `a`

When there is no sub-script 'p' after the square bracket, the following interpretation is used:

[f] a
$$<=>$$
 [f]_p a > 0
[f] a = a: $[\mathbf{3}_p f]$ (> 0)

This interpretation does not terminate, but converges to some boolean value. If you are familiar with boolean type-checking in path semantics, you might know that double-existential paths `∃∃f` is one of four functions of type `bool → bool`. It is not possible to check the consistency of a statement without realizing the fact that this goes on forever but in a predictable and repeating pattern. Similarly, in probabilistic path semantics, the interpretation of a boolean sub-type goes on forever, but it is much more complex when defined in terms of the probabilistic existential path. By realizing the fact that consistency of probabilistic sub-types converges to checks predicted by boolean sub-types, one can stop execution and use the rules for checking boolean sub-types instead.

When mixing boolean sub-types with probabilistic sub-types, the boolean sub-types are used as domain constraints on the probabilistic existential path:

$$[g] b \wedge [f]_p a = (\exists_p f\{[g] b\})(a)$$

The above means the probability that `f` returns `a`, given that `g` returns `b`.

This can also be written in the familiar notation for conditional probability:

$$P([f] a | [g] b) = (\exists_p f\{[g] b\})(a)$$

Conditional probabilities share the semantics of mixing boolean sub-types and probabilistic sub-types.