Univalent Involutions

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In this paper I prove that isomorphisms are isomorphic to Univalent Involutions.

An isomorphism^[1] from a Category^[2] theoretic view is a morphism `f` with an inverse `g` such that:

$$f \cdot g \stackrel{<=>}{} id_{B}$$

$$g \cdot f \stackrel{<=>}{} id_{A}$$

$$f : A \rightarrow B$$

$$g : B \rightarrow A$$

An involution^[3] is a morphism 'h' such that:

$$h \cdot h \le id_T$$

 $h : T \rightarrow T$

Every involution is an isomorphism, but not every isomorphism is an involution.

It turns out that every isomorphism can be turned into a Univalent Involution:

 $h := (x : T) = if let some(a) = h'_{A}^{-1}(x) \{ h'_{B}(f(a)) \}$

The univalent involution `h` has the following normal paths (where `opt` is used as a functor):

$$h[h'_{A}^{-1} \rightarrow h'_{B}^{-1}] \le opt(f)$$

 $h[h'_{B}^{-1} \rightarrow h'_{A}^{-1}] \le opt(g)$

A Univalent Involution differs from ordinary involutions by the property it can be turned back into a heterogenous isomorphism, kind of like a tuple `(a, b)` can be turned into `a` and `b`:

Since equality in Intuitionistic Logic [4] using types is a tuple (f, g), this particular form of involution is thought to be univalent [5].

References:

[1]	"Isomorphism"
	Wikipedia
	https://en.wikipedia.org/wiki/Isomorphism

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 Wikipedia
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