

Extracting Bits in Answered Modal Logic

by Sven Nilsen, 2020

In this paper I describe the extraction process of Answered Modal Logic in more detail.

The expression $\neg!\diamond X \vee \Box X$ can be encoded as following:

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The first bit tells whether $\neg!\diamond$ belongs to the set.

The second bit tells whether $\neg!\diamond$ belongs to the set.

The third bit tells whether \Box belongs to the set.

A special case is the empty expression:

000 `` empty expression

Handling empty expressions properly can be a bit tricky, since there is another special case:

111 $\neg\Box X$

When there are multiple variables, one can fill out the unmentioned variables:

$$\neg\Box X = \neg\Box X \wedge \neg\Box Y$$

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111

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The problem is that it is not allowed to remove the last term.

At least one term is needed to witness the fill operation:

$$\neg\Box X = \neg\Box Y$$

This is not equal to an empty expression:

$\neg\Box X \neg = ``$ Notice that $\neg =$ uses \neg not \neg

You can prove that that this rule fills out all the cases:

$\therefore \neg\Box X$

$\therefore !\diamond X \vee \neg!\diamond X \vee \Box X$

$\therefore (!\diamond X \wedge \neg\Box Y) \vee (\neg!\diamond X \wedge \neg\Box Y) \vee (\Box X \wedge \neg\Box Y)$

$\therefore (!\diamond X \wedge (!\diamond Y \vee \neg!\diamond Y \vee \Box Y)) \vee (\neg!\diamond X \wedge (!\diamond Y \vee \neg!\diamond Y \vee \Box Y)) \vee (\Box X \wedge (!\diamond Y \vee \neg!\diamond Y \vee \Box Y))$

$\therefore (!\diamond X \wedge !\diamond Y) \vee (!\diamond X \wedge \neg!\diamond Y) \vee (!\diamond X \wedge \Box Y) \vee (\neg!\diamond X \wedge !\diamond Y) \vee (\neg!\diamond X \wedge \neg!\diamond Y) \vee (\neg!\diamond X \wedge \Box Y) \vee (\Box X \wedge !\diamond Y) \vee (\Box X \wedge \neg!\diamond Y) \vee (\Box X \wedge \Box Y)$