

Avatar Logic to Set Theory

by Sven Nilsen, 2021

In this paper I introduce a method of translating Avatar Logic to Zermelo-Fraenkel Set Theory.

The axioms of Avatar Logic^[1] might be translated to Zermelo-Fraenkel Set Theory^[2]:

$$\begin{array}{ll} \mathbf{(a, b) \wedge b : p \wedge \text{uniq}(b)} & \mathbf{(a, q'(b)) \wedge q'(b) : p} \\ p(a, b) \wedge \exists! z \{ p(a, z) \} & p(a, \{\{q\}, \{q, b\}\}) \wedge \forall x \{ \exists! z \{ p(a, \{\{z\}, \{z, x\}\}) \} \} \end{array}$$

Translation must happen for every relation, otherwise it would require extending Second-Order Logic^[3] with tuples, roles and 1-avatars. Per relation requires only First-Order Logic^[4].

The translation uses Kuratowski's definition^[5] of an ordered pair $\{\{x\}, \{x, y\}\}$ for $x'(y)$. This representation is chosen because ordered pairs are not used as arguments in Avatar Logic.

Ordered pairs might also be used without $b : p$, but only to mean (a, b) as a binary relation.

The `uniq` predicate returns `true` for all atomic symbols, plus those 1-avatars that are optionally chosen to be behaving uniquely. Both axioms must be applied when the 1-avatar is unique.

In expanded form limited to quantifiers \forall, \exists , connectives $\Rightarrow, =, \in, \vee, \wedge$, negation \neg :

$$\begin{array}{l} \mathbf{(a, b) \wedge b : p \wedge \text{uniq}(b)} \\ p(a, b) \wedge \exists z \{ p(a, z) \wedge \neg \exists y \{ p(a, y) \wedge \neg(y = z) \} \} \\ \\ \mathbf{(a, q'(b)) \wedge q'(b) : p} \\ p(a, _k) \wedge \\ \exists r \{ \\ (r \in _k) \Rightarrow \forall c \{ (c \in _k) \Rightarrow ((c = q) \vee (c = r)) \} \wedge \\ \forall d \{ (d \in r) \Rightarrow ((d = q) \vee (d = b)) \} \\ \} \wedge \\ \forall x \{ \\ \exists z \{ p(a, _n) \wedge \neg \exists y \{ p(a, _m) \wedge \neg(y = z) \} \} \wedge \\ \exists r \{ \\ (r \in _n) \Rightarrow \forall c \{ (c \in _n) \Rightarrow ((c = z) \vee (c = r)) \} \wedge \\ \forall d \{ (d \in r) \Rightarrow ((d = z) \vee (d = x)) \} \\ \} \wedge \\ \exists r \{ \\ (r \in _m) \Rightarrow \forall c \{ (c \in _m) \Rightarrow ((c = y) \vee (c = r)) \} \wedge \\ \forall d \{ (d \in r) \Rightarrow ((d = y) \vee (d = x)) \} \\ \} \\ \} \end{array}$$

Notice that `p` might be written with uppercase letter in standard First-Order Logic notation. The variables starting with underscore e.g. `_n`, are introduced to bind the sub-expressions together.

References:

- [1] “Avatar Logic”
AdvancedResearch – Summary Page on Avatar Extensions
<https://advancedresearch.github.io/avatar-extensions/summary.html#avatar-logic>
- [2] “Zermelo-Fraenkel set theory”
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- [5] “Kuratowski’s definition – Ordered pair”
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