Propositional Logic Interpretation of Answered Modal Logic

by Sven Nilsen, 2020

In this paper I show that Answered Modal Logic can be interpreted, in part, with Propositional Logic. This model can only be fully developed under Higher Order Operator Overloading.

For `n` variables, the semantic model of canoncial expressions in Answered Modal Logic form an n-dimensional cube, where each dimension represents a variable in a modal set ` $\{! \diamondsuit, \neg! \diamondsuit, \Box\}$ `.

The corners of this n-dimensional cubes naturally can model propositions:

$\square X$	X is true
!≎X	X is false

The generalized Propositional Logic gate NOT is defined as following:

$$\begin{array}{ll} not(! \diamond X) = \square X & \neg ! \diamond X = \neg ! \diamond X \\ not(\square X) = ! \diamond X & \neg \square X = \neg \square X \\ not(\neg ! \diamond X) = \neg \square X & \neg \neg ! \diamond X = ! \diamond X \\ not(\neg \square X) = \neg ! \diamond X & \neg \neg \square X = \square X \end{array}$$

This is also consistent with:

$$not[\neg] \le not$$

To derive a full Propositional Logic interpretation, it is sufficent to construct a NAND gate. Since NOT is already defined, the remaining work is to construct a AND gate:

$and(! \diamond X, ! \diamond X) = ! \diamond X$	Carries over from Propositional Logic
and($! \diamond X$, $\neg ! \diamond X$) = $! \diamond X$	See proof D (next page)
$and(! \diamond X, \neg \Box X) = ! \diamond X$	See proof A (next page)
$and(! \diamond X, \Box X) = ! \diamond X$	Carries over from Propositional Logic
$and(\neg! \diamond X, ! \diamond X) = ! \diamond X$	See proof D (next page)
and $(\neg! \diamond X, \neg! \diamond X) = \neg! \diamond X$	See proof C (next page)
$and(\neg! \Diamond X, \neg \Box X) = \neg! \Diamond X$	See proof A (next page)
$and(\neg! \diamond X, \Box X) = \neg! \diamond X$	See proof B (next page)
$and(\neg \Box X, ! \Diamond X) = ! \Diamond X$	See proof A (next page)
and $(\neg \Box X, \neg! \Diamond X) = \neg! \Diamond X$	See proof A (next page)
$and(\neg \Box X, \neg \Box X) = \neg \Box X$	See proof C (next page)
$and(\neg \Box X, \Box X) = \neg \Box X$	See proof B (next page)
$and(\Box X, ! \Diamond X) = ! \Diamond X$	Carries over from Propositional Logic
and($\Box X$, $\neg! \diamond X$) = $\neg! \diamond X$	See proof B (next page)
$and(\Box X, \neg \Box X) = \neg \Box X$	See proof B (next page)
$and(\Box X, \Box X) = \Box X$	Carries over Propositional Logic

Notice that this is an operator on the functions of a variable `X`, which is permitted by using Higher Operator Overloading (HOOO) semantics.

The NAND gate is constructed by using `nand <=> not . and`.

For functions of different variables, one can not use the semantics of HOOO. However, in some cases there exists a model in Propositional Logic.

and(
$$\Box X$$
, $\Box Y$) Undefined, but has a model in propositional logic (corners of the cube) and($\Box X$, $\neg ! \diamond Y$) Undefined, no model

The proofs of AND are as following.

Since `and` is commutative:

and(
$$! \diamondsuit X$$
, $\neg \Box X$) = and($\neg \Box X$, $! \diamondsuit X$) and($\Box X$, $\neg \Box X$) = and($\neg \Box X$, $\Box X$)

The following cases are somewhat intuitive, since $\neg \Box X = \{! \diamond, \neg ! \diamond, \Box\} X$:

and(
$$! \diamond X$$
, $\neg \Box X$) = $! \diamond X$ Proof A and($\neg ! \diamond X$, $\neg \Box X$) = $\neg ! \diamond X$

The second case is a bit trickier:

- \therefore and($\square X, \neg \square X$) Proof B
- \therefore and($\Box X$)($\neg \Box X$)
- \therefore id($\neg \Box X$) Using `and(true)(x) => id(x)`
- ∴ ¬□X

The same trick can be used here:

and(
$$\square X$$
, $\neg ! \diamond X$) = $\neg ! \diamond X$

Two trivial cases are the following:

and
$$(\neg \Box X, \neg \Box X) = \neg \Box X$$
 Proof C and $(\neg ! \diamondsuit X, \neg ! \diamondsuit X) = \neg ! \diamondsuit X$

The only case left (two commutative) is the following, which I define using intuition:

and(
$$! \diamond X$$
, $\neg ! \diamond X$) = $! \diamond X$ Proof D

Q.E.D.