

Equivalence Paths

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Assuming semantics of constrained functions, an equivalence path is an “auto-constraint” operator:

$$\sim := \lambda(f : A \rightarrow B) \rightarrow (A \rightarrow B) = f \{ \lambda(a : A) = |[f] f(a)| == 1 \}$$

This unary operator \sim is called “universal equivalence path”.

All structure preserving functions have themselves as equivalence paths:

$$\sim f \Leftrightarrow f \quad \text{if } f \text{ is structure perserving (} f \text{ got an inverse)}$$

Similarly, applying \sim twice has no effect:

$$\sim \sim f \Leftrightarrow \sim(\sim f) \Leftrightarrow f$$

There is a short version for accessing the two functions using \sim as a binary operator:

$$\begin{aligned} \forall \sim f \quad &\Leftrightarrow \quad f \sim 0 \quad \Leftrightarrow \quad \lambda(a : A) = |[f] f(a)| == 1 \\ \exists \sim f \quad &\Leftrightarrow \quad f \sim 1 \quad \Leftrightarrow \quad \lambda(b : B) = |[f] b| == 1 \end{aligned}$$

When a function f is constrained to $f \sim 0$, its existential path is $f \sim 1$:

$$\exists f \{ f \sim 0 \} \Leftrightarrow \exists \sim f \Leftrightarrow f \sim 1$$

If one knows $f \sim 1$, then one can construct $f \sim 0$ by composing f with $f \sim 1$:

$$f \sim 1 \cdot f \Leftrightarrow \forall \sim f \Leftrightarrow f \sim 0$$

One can think of $\sim f$ as crossing out all input-output pairs that intersect.

Using constrained functions, it is possible to fine-tune any desired equivalence path:

$\sim f$	$\sim f \{ (\neg = 3) \}$	$\sim f \{ (\neg = 2) \}$	$\sim f \{ (\neg = 2) \wedge (\neg = 5) \}$
$0 \rightarrow a$	$0 \rightarrow a$	$0 \rightarrow a$	$0 \rightarrow a$
$1 \rightarrow b$	$1 \rightarrow b$	$1 \rightarrow b$	$1 \rightarrow b$
$2 \rightarrow c$	$2 \rightarrow c$	$2 \rightarrow c$	$2 \rightarrow c$
$3 \rightarrow e$	$3 \rightarrow e$	$3 \rightarrow c$	$3 \rightarrow c$
$4 \rightarrow d$	$4 \rightarrow d$	$4 \rightarrow d$	$4 \rightarrow d$
$5 \rightarrow d$	$5 \rightarrow d$	$5 \rightarrow d$	$5 \rightarrow d$