

Path Sets

by Sven Nilsen, 2017

Sometimes a path^[1] is not unique, but has multiple solutions:

$$f[g] \Leftrightarrow \{h_0, h_1, \dots\}$$

This happens when the path is non-surjective^[2], for example:

$$\text{id}[\text{false_1}] \Leftrightarrow \{\text{false_1}, \text{id}\}$$

A non-surjective path collapses the domain^[3] of `h`, such that it becomes partial^[4]. All functions that contain the same partial function is member of the path set.

A path can also be non-existent, which corresponds to an empty set. This gives paths a natural semantics that corresponds to set theory.

If one thinks of variables as functions with 0 arguments, then it is easy to generalize function application^[5] to operate on sets of functions, where the result is also a set of functions:

$$\text{not}(\{\text{false}\}) \Leftrightarrow \{\text{true}\}$$

$$\text{not}(\{\text{false}, \text{true}\}) \Leftrightarrow \{\text{true}, \text{false}\}$$

This makes it possible to do function currying^[6] in a way that corresponds to paths:

$$\text{and}(\{\text{true}\}) \Leftrightarrow \{\text{id}\}$$

$$\text{and}(\{\text{false}\}) \Leftrightarrow \{\text{false_1}\}$$

$$\text{and}(\{\text{false}, \text{true}\}) \Leftrightarrow \{\text{false_1}, \text{id}\}$$

$$\text{and}(\{\text{false}, \text{true}\}) \Leftrightarrow \text{id}[\text{false_1}]$$

Actually, `id[false_1]` is not the only path:

$$\text{id}[\text{false_1}] \Leftrightarrow \{\text{false_1}, \text{not}, \text{id}, \text{true_1}\}[\text{false_1}]$$

$$\text{and}(\{\text{false}, \text{true}\}) \Leftrightarrow \{\text{false_1}, \text{not}, \text{id}, \text{true_1}\}[\text{false_1}]$$

Instead of writing every function of type `bool → bool`, one can just write:

$$\text{and}(\{\text{false}, \text{true}\}) \Leftrightarrow (\text{bool} \rightarrow \text{bool})[\text{false_1}]$$

This is only allowed because all functions in the set have the same path set by `false_1`. You take the intersection of the sets for all functions in `bool → bool`.

The set of all functions of type `bool → bool` have no partial function in common, so you can not write:

$$f[g] \leq (bool \rightarrow bool)$$

This is because there is no function you can construct with a path such that it is logically equivalent to all functions of type `bool → bool`. The type `bool` is non-empty, so the function space `bool → bool` is non-empty. The only case where you can do this is when the function space is empty.

The function space `bool → bool` has these properties:

$$\begin{aligned} (bool \rightarrow bool)[false_1] &\leq \{false_1, id\} \\ (bool \rightarrow bool)[not] &\leq \{\} \\ (bool \rightarrow bool)[id] &\leq \{\} \\ (bool \rightarrow bool)[true_1] &\leq \{id, true_1\} \end{aligned}$$

You can write sets of paths that have the same path sets:

$$(bool \rightarrow bool)[\{not, id\}] \leq \{\}$$

If path sets are not the same for a set of paths, then you take the intersection of the path sets:

$$(bool \rightarrow bool)[(bool \rightarrow bool)] \leq \{\}$$

Now, we construct a higher order function logically equivalent to `if`:

$$\begin{aligned} if : a \rightarrow a \rightarrow (bool \rightarrow a) \\ if = \lambda(x, y) = \lambda(c) = if\ c\ \{x\}\ \text{else}\ \{y\} \end{aligned}$$

Since `and` can be curried with `true` and `false`, we can construct `and` using `if`:

$$\begin{aligned} and(true) &\leq id \\ and(false) &\leq false_1 \\ \\ if(id, false_1) &\leq and \end{aligned}$$

Doing the same for `or`:

$$\begin{aligned} or(true) &\leq true_1 \\ or(false) &\leq id \\ \\ if(true_1, id) &\leq or \end{aligned}$$

It is easy to see that one can derive the symmetric paths of `if` by `not`:

$$\begin{aligned} and[not] &\leq or \\ if(id, false_1)[not] &\leq and[not] \\ if(id, false_1)[not] &\leq if(true_1, id) \end{aligned}$$

In general, the only interesting path by `bool → bool` is `not`, because the other 3 are trivial:

```
if(a, b)[false_1] <=> false_n
if(a, b)[not] <=> if(b[not], a[not])
if(a, b)[id] <=> if(a, b)
if(a, b)[true_1] <=> true_n
```

Any boolean function can be constructed by `if`. Its symmetric path by `not` can be simplified:

```
if(if(a0, a1), if(a2, a3))[not]
if(if(a2, a3)[not], if(a0, a1)[not])
if(if(a3[not], a2[not]), if(a1[not], a0[not]))
[not] if(if(a3, a2), if(a1, a0))
not(if(if(a3, a2), if(a1, a0)))
```

$a_i : \text{bool}$

Notice that it just reverses the order and inverts the output. You can easily do the same to any else-if expression by copying the expression, reversing the higher order arguments and inverting the output:

```
f := \ (a0, a1, a2) = \ (x0, x1) = if x0 { a0 } else if x1 { a1 } else { a2 }

f[not] := \ (a2, a1, a0) = \ (x0, x1) = !if x0 { a0 } else if x1 { a1 } else { a2 }

a_i : bool
```

A summary so far:

1. Paths form a set called “path set”
2. All boolean functions and their symmetric paths can be constructed with `if`

Now we want to study the cardinality sum of all symmetric and asymmetric path sets.

$$\sum_j \{ |f_j[g_{i \rightarrow n}]| \}$$

Creating a table for the function space `bool → bool`:

Column → Row	false_1	not	id	true_1
false_1	8	4	4	8
not	4	4	4	4
id	4	4	4	4
true_1	8	4	4	8

$$\sum_j \{ |f_j[g_{i \rightarrow n}]| \} = 4^2 \cdot 4 + 4 \cdot (8-4) = 4 \cdot 4 \cdot 4 + 4 \cdot 4 = 80$$

This was counted manually. In the future I might be able to generalize this further.

References:

- [1] “Normal Paths”
Sven Nilsen, 2019
https://github.com/advancedresearch/path_semantics/blob/master/papers-wip/normal-paths.pdf
- [2] “Surjective function”
Wikipedia
https://en.wikipedia.org/wiki/Surjective_function
- [3] “Domain of a function”
Wikipedia
https://en.wikipedia.org/wiki/Domain_of_a_function
- [4] “Partial function”
Wikipedia
https://en.wikipedia.org/wiki/Partial_function
- [5] “Function application”
Wikipedia
https://en.wikipedia.org/wiki/Function_application
- [6] “Currying”
Wikipedia
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