

# Structure-Preserving Functions

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Of all functions of type  $\text{bool} \rightarrow \text{bool}$ , there are only two structure-preserving functions:

$\text{not} := \text{if}(\text{false}, \text{true})$

$\text{id} := \text{if}(\text{true}, \text{false})$

$\text{if} := \lambda(a, b) = \lambda(x : \text{bool}) = \text{if } x \{ a \} \text{ else } \{ b \}$

These functions are called structure-preserving because when using in a path, the path set is always non-empty. All structure-preserving functions have inverses.

In order to be a structure preserving function, it must map from one type to another that contains the same or more number of elements.

$f : A \rightarrow B \wedge |A| \geq |B|$

However, this is not sufficient to define a structure-preserving function. A more accurate way is to think of the type  $B$  having some implicit sub-type that constrains it to those values returned by  $f$ . This constrained type must have the same number of elements as the input type:

$f : A \rightarrow B \wedge |A| = |f(A)|$

Another way of writing is using the existential path:

$f : A \rightarrow B \wedge |A| = |\exists f|$

This is a shorthand version for:

$f : A \rightarrow B \wedge |A| = |[\exists f] \text{ true}|$

Therefore, one has the following for any function of type  $A \rightarrow B$ :

$|f(A)| = |\exists f|$

$f(A) \subseteq B$

$f : A \rightarrow B$

To construct all structure-preserving functions of type  $A \rightarrow A$ , one can use the  $\text{id}$  function, take the partial function pairs  $(x, \text{id}(x))$  and then rearrange the outputs:

$(x_i, \text{id}(x_i)) \quad \neg \exists k: (\neg = i) \{ (x_k, \text{id}(x_k)) \}$

This means there exists  $|A|!$  number of structure-preserving functions of type  $A \rightarrow A$ .

If the existential path of a function is  $\text{true}_1$ , then the length of the existential path is equal to the length of the output type. The same is also true in the other direction:

$$(\exists f \Leftrightarrow \text{true}_1) \Leftrightarrow (|\exists f| == |B|)$$

$$f : A \rightarrow B$$

For example:

$$f : \text{bool} \times \text{bool} \rightarrow \text{bool} \times \text{bool} \wedge |\text{bool} \times \text{bool}| == |f(\text{bool} \times \text{bool})|$$

$$|\text{bool} \times \text{bool}| == |\exists f|$$

$$|\text{bool}| \cdot |\text{bool}| == |\exists f|$$

$$2 \cdot 2 == |\exists f|$$

$$4 == |\exists f|$$

$$\exists f \Leftrightarrow \text{true}_1$$

$$\text{Because } |\text{bool} \times \text{bool}| == 4 \text{ and } (\exists f \Leftrightarrow \text{true}_1) \Leftrightarrow (|\exists f| == |\text{bool} \times \text{bool}|).$$

There exists  $4! = 4 \cdot 3 \cdot 2 \cdot 1 = 24$  number of structure preserving functions of type  $\text{bool} \times \text{bool} \rightarrow \text{bool} \times \text{bool}$ .

For functions of type  $A \rightarrow B$  there is a way to count the number of structure-preserving functions:

$$|f| == |A|! \cdot \text{bin}(|B|, |A|) == |B|! / (|B| - |A|)!$$

$$f : A \rightarrow B \wedge |A| == |f(A)|$$

For example, for functions of type  $\text{bool} \rightarrow \text{bool} \times \text{bool}$ , there exists 12 structure-preserving functions:

$$|f| == |\text{bool}|! \cdot \text{bin}(|\text{bool} \times \text{bool}|, |\text{bool}|)$$

$$|f| == 2! \cdot \text{bin}(4, 2)$$

$$|f| == 2 \cdot 6$$

$$|f| == 12$$

$$f : \text{bool} \rightarrow \text{bool} \times \text{bool} \wedge |\text{bool}| == |f(\text{bool})|$$

The reason is that there are 2 structure-preserving functions of type  $\text{bool} \rightarrow \text{bool}$ . For each of these functions one can pick 2 values of type  $\text{bool} \times \text{bool}$  to use as output. The values can not be rearranged per function of type  $\text{bool} \rightarrow \text{bool}$ , because it would be the same as counting the same functions more than once.

From this way of counting structure-preserving functions of different types, there is an algorithm to construct all structure-preserving functions of type  $A \rightarrow B$ :

1. Construct all structure-reserving functions of type  $A \rightarrow A$ .
2. Find all ordered maps from  $B$  to  $A$ .
3. Replace outputs of 1) with 2) using the order of  $A$  to look up the value of  $B$ .