Instant Quantum Partial Observations

by Sven Nilsen, 2020

In this paper I discuss the semantics of instant partial observations of quantum functions.

In the paper "Quantum Propagation", I mentioned that at any given instant, every outcome of a partial observation `g` is equally probable if `f` is semi quantum (all amplitudes have same length):

$$\begin{array}{ll} g \cdot f & \quad \text{`g` is a partial observation of a semi quantum function `f`} \\ f: () \to \mathbb{B}^n & \quad \text{`f` is semi quantum} & |\exists_{pc} f| <=> \exists_p f \\ g: \mathbb{B}^n \to \mathbb{B}^m & \quad m < n \end{array}$$

For simplicity, I will use m = 1 the rest of this paper.

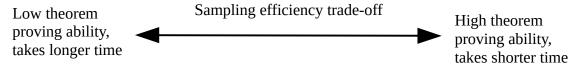
The intuition is that assuming some level of theorem proving abilities, one can construct a pair of complex probability amplitudes `a` and `b` in the expanded quantum propagation product:

Such that:

$$g(0) = g(1)$$
 $g(2) = g(3)$ $g(0) < g(2)$

There is no coincidence that this was in the paper about "Higher Order Non-Deterministic Diagrams". I am trying to understand how high the level of theorem proving ability is required to find a sampling algorithm, no matter which partial observation `g` that is chosen.

It is understood that this sampling algorithm is not trivial to write down in its most efficient form, but I assume that it will be unique in the sense that all output pairs will have same statistical properties. Under this assumption, there is a trade-off between theorem proving abilities and time:



Therefore, instant quantum partial observations belongs to a computational complexity class with respect to some oracle where one can take "every path" when interpretated as a statistical limit in itself. For `m` bits in the partial observation, it is not possible to extract more than `m` bits of information.