

Analysis: Inconstant Ordering

Test Set 1

It is given that a valid string must start with A.

For $N = 1$, the only block B_1 should therefore start with a character greater than A. To keep building the alphabetically first valid string, we start B_1 with B which is just greater than A. We keep on incrementing the character as we move right to ensure that B_1 is strictly increasing. Notice that we will eventually end up with a prefix of BCDE...XYZ as B_1 . As $L_i \leq 25$, we will never run out of characters for this.

For $N = 2$, we construct the first block B_1 using the process described for $N = 1$. We now evaluate the possible ways to fill the strictly decreasing second block B_2 . We try to start B_2 with the smallest possible character to get the alphabetically first string. This implies that B_2 will end with the smallest character in the alphabet which is A. So, we start filling B_2 from the end with A and move to the left, incrementing the character.

Let b_1 be the last character in B_1 and b_2 be the first character in B_2 . If $b_1 > b_2$, then we will always get a valid string as it ensures that B_2 is strictly decreasing. On the other hand if $b_1 \leq b_2$, then the final string is invalid. For example, if $L_1 = 2$ and $L_2 = 4$, the resulting string from the above process will be ABCDCBA which is invalid.

To ensure that B_2 is always strictly decreasing, we update b_1 with the character just greater than b_2 . This will guarantee that the string is valid and alphabetically the first one. Again as $L_i \leq 25$, we will never run out of characters while filling B_2 . In the above example, ABCDCBA now transforms to a valid string ABEDCBA.

Test Set 2

We generalize the above solution for $N > 2$.

If N is even, we can divide the blocks into $N/2$ pairs as $[L_1, L_2][L_3, L_4] \dots [L_{N-1}, L_N]$ and solve the $N/2$ pairs independently using the same approach as for $N = 2$. We can do this as every such pair of blocks will end with A.

If N is odd, we solve for $N - 1$ blocks using the above approach and for the last block L_N , we treat it as the case when $N = 1$.

The time complexity is $O(L_1 + L_2 + \dots + L_N)$ for each test case.