

## Analysis: Year of More Code Jam

The setting of this problem is no doubt discrete probability. From the definition, the space consists of  $N^T$  equally likely possible outcomes. That can be, under our limits, as huge as  $10^{450}$ . Clearly, a naive approach is not feasible.

But let us do a little exercise in probability. Define the random variable  $X_i$  to be the number of contests on the  $i$ -th day, the quantity we want to compute is the average of  $X_i^2$ , i.e., the expectation  $E(\sum X_i^2)$ . By the linearity of the expectation, we have

$$E(\sum_{1 \leq i \leq N} X_i^2) = \sum_{1 \leq i \leq N} E(X_i^2).$$

So let us focus on the computation of the variable for a fixed day for the moment. Pick any  $i$ , and let  $X := X_i$ . Let us define more random variables. Define  $Y_j$  to be the *indicator* of whether the  $j$ -th tournament will have a contest on the  $i$ -th day. Clearly,  $X = \sum_j Y_j$ . So,

$$(*) \quad E(X^2) = E((\sum_{1 \leq j \leq T} Y_j)^2) = \sum_{1 \leq j \leq T} E(Y_j^2) + 2 \sum_{1 \leq j < k \leq T} E(Y_j Y_k).$$

We observe that each terms in the last expression is easy to compute. Being the indicator random variables, the  $Y$ 's take value 0 or 1. So

- $Y_j^2$  always has the same value as  $Y_j$ , and its expectation is just the probability that  $Y_j$  is 1, i.e., tournament  $j$  has a contest on day  $i$ .
- $Y_j Y_k$  is 1 if and only if both  $Y_j$  and  $Y_k$  are 1. The expectation is the probability that both the  $j$ -th and the  $k$ -th tournament has a contest on day  $i$ .

Let the input for the tournament  $j$ , i.e., the contest pattern be  $d_1=1, d_2, \dots, d_m$ . Denote  $D(i,j)$  be the number of the  $d$ 's for tournament  $j$  that are less than or equal to  $i$ . There are  $N$  choices for the starting date of a particular tournament. It is easy to see that the first probability above is  $D(i,j) / N$ ; while the second probability is  $D(i,j)D(i,k) / N^2$ .

So far we addressed the problem just for a single day  $i$ . We need to do this for every  $i$ . There are  $10^9$  of them. But notice that, as long as there is no input  $d_t = i$ ,  $D(i-1, j) = D(i, j)$  for all  $j$ . This means that the expectation for the  $i$ -th day is the same as the expectation for the  $(i-1)$ -th. There are at most  $T \max(M) \leq 2500$  such  $d$ 's in the input, so we need to compute (\*) for at most 2500 days. For our problem, it is good enough to realize that there is no  $d > 10000$ . So all the expectations after the 10000-th day are the same. We can just do the computation for the first 10000 days, and for the rest, a simple multiplication.

The last problem is the need for big integers. At the first glance we might have both numerators and denominators as big as  $10^{450}$ . But that is not the truth. Simply observe the above answer, which is a sum of various  $D(i,j) / N$  and  $D(i,j)D(i,k) / N^2$ . We actually proved that the denominator is never bigger than  $N^2$ . A careful implementation with 64-bit integers will be good enough.

For a further speed-up. The formula in (\*) involves computing  $O(T^2)$  terms. But if we do it from day 1, keep  $D(i,j)$  for each  $j$  and two more variables --  $S_1$  for the sum of all the  $D(i,j)$ 's, and  $S_2$  for the sum of  $D(i,j)^2$ 's, then we just need constant update time when we see an input  $d$ , and also constant computation time for each day we want to compute (\*).

