## **Analysis: The Next Number**

Let  $\mathbf{x}$  be our input. We want to find the next number  $\mathbf{y}$ . We denote L(s) to be the length of s, i.e., the number of digits in s.

Case 1. If all the digits in x are non-increasing, for example x = 776432100, then x is already the biggest one among the numbers in the list with L(x) digits. The next number, y, must have one more digit, that is, one more 0. Actually y must be the smallest one with L(x)+1 digits. To get y, we put the smallest non-zero digit in front, and all the other digits are put in non-decreasing order.

Case 2. Otherwise, x can be written as the concatenation x = ab, where b is the longest non-increasing suffix of x, so d, the last digit of a is smaller than the first digit of b. Let d' be the smallest among all the digits in b that are bigger than d.

Because b is non-increasing, x is the biggest number among those who has L(x) digits and starts with prefix a. The first L(a) digits of y must be bigger than a. The smallest we can do is to replace d with d', and then for the rest digits, we arrange them in non-decreasing order. Let us do another example. x = 134266530. Then a = 1342, b = 66530, d = 2, and d' = 3. The next number is y = 134302566.

In fact, we can unify the two cases above. Since the number of 0's is not restricted, we can just imagine there is one more 0 in the beginning of x, thus Case 1 is reduced to Case 2.

The above described is actually exactly the procedure to get the next permutation of a finite sequence in certain languages. Below is a solution that is essentially one line in C++. From the author:

```
deque<char> f;
...
f.push_front('0');
next_permutation(f.begin(), f.end());
if (f.front() == '0') f.pop_front();
```