

## Analysis: Minimum Sort

In this problem, we are given only one tool to get information about our input: the minimum query. Moreover, the query becomes cheaper when done over long distances. This means a sorting algorithm that normally spends running time finding minimums would align well with our needs. One well known such algorithm is [Selection sort](#).

In Selection sort, we search for the minimum of the full list, put that in its correct position (the beginning), and continue solving recursively searching for the minimum of a suffix of the list, and moving that minimum to the beginning. This can be implemented with  $N - 1$  pairs of the minimum query and (possibly) a swap operation as follows:

```
for i := 1 to N-1
  j = minimum_index_of(i, i+1, ..., N)
  if i != j:
    swap(i, j)
```

The minimum queries above are over ranges of sizes  $N, N - 1, \dots, 2$ . This means the total cost is exactly

$$\sum_{i=2}^N \left\lceil \frac{10^8}{i} \right\rceil.$$

For  $N = 100$ , this is 418737795, which is less than the limit of  $6 \times 10^8$ .