

Analysis: Teach Me

Test set 1

We can solve this test set by considering all ordered pairs of employees. For an ordered pair (i, j) , we can check whether there is a skill that the i -th employee knows that the j -th employee does not know with a simple $O(C_i \times C_j)$ check using nested loops.

This solution runs in $O(5^2 \times N^2)$.

Test set 2

We can solve this test set by defining $m(i)$ as the number of employees who can mentor the i -th employee. If we can compute $m(i)$, the answer to the problem is the sum of all $m(i)$.

To compute $m(i)$, we can count the number of employees who **cannot** mentor the i -th employee instead. We can observe that the j -th employee cannot mentor the i -th employee if and only if the set of skills known by the j -th employee is a subset of the set of skills known by the i -th employee. Therefore, we would like to count the number of employees whose set of skills is a subset of the set of skills known by the i -th employee.

To count this, we can consider every subset of $\{A_{i1}, \dots, A_{iC_i}\}$. For a subset B , we can count the number of employees whose set of skills is exactly B . Taking the sum of all subsets gives us the number of employees whose set of skills knowledge is the subset of the set of skills known by the i -th employee. $m(i)$ is simply the number of employees subtracted by this value.

To be able to compute the number of employees whose set of skills is exactly a given set, we can preprocess the set of skills into an occurrences hashmap. In other words, we can keep a hashmap that takes a set of skills as its key and returns the number of employees who knows the exact same set of skills as its value.

For each employee i , we need to consider every subset of $\{A_{i1}, \dots, A_{iC_i}\}$. Since there can be up to 2^5 subsets, this solution runs in $O(2^5 \times N)$.