

Insertion Sort

Algorithm

Now we have a bigger picture of how this sorting technique works, so we can derive simple steps by which we can achieve insertion sort.

Step 1 – If it is the first element, it is already sorted. return 1;

Step 2 – Pick next element

Step 3 – Compare with all elements in the sorted sub-list

Step 4 – Shift all the elements in the sorted sub-list that is greater than the value to be sorted

Step 5 – Insert the value

Step 6 – Repeat until list is sorted

Pseudocode

procedure insertionSort(A : array of items)

 int holePosition

 int valueToInsert

 for i = 1 to length(A) inclusive do:

 /* select value to be inserted */

 valueToInsert = A[i]

 holePosition = i

 /*locate hole position for the element to be inserted */

 while holePosition > 0 and A[holePosition-1] > valueToInsert do:

 A[holePosition] = A[holePosition-1]

 holePosition = holePosition - 1

end while

/* insert the number at hole position */

A[holePosition] = valueToInsert

end for

end procedure

Selection Sort

Algorithm

Step 1 – Set MIN to location 0

Step 2 – Search the minimum element in the list

Step 3 – Swap with value at location MIN

Step 4 – Increment MIN to point to next element

Step 5 – Repeat until list is sorted

Pseudocode

procedure selection sort

list : array of items

n : size of list

for i = 1 to n - 1

/* set current element as minimum*/

min = i

/* check the element to be minimum */

for j = i+1 to n

if list[j] < list[min] then

```
        min = j;
    end if
end for
/* swap the minimum element with the current element*/
if indexMin != i then
    swap list[min] and list[i]
end if
end for
end procedure
```

Bubble Sort

Algorithm

```
begin BubbleSort(list)
    for all elements of list
        if list[i] > list[i+1]
            swap(list[i], list[i+1])
        end if
    end for
    return list
end BubbleSort
```

Pseudocode

```
procedure bubbleSort( list : array of items )  
  loop = list.count;  
  for i = 0 to loop-1 do:  
    swapped = false  
    for j = 0 to loop-1 do:  
      /* compare the adjacent elements */  
      if list[j] > list[j+1] then  
        /* swap them */  
        swap( list[j], list[j+1] )  
        swapped = true  
      end if  
    end for  
    /*if no number was swapped that means  
    array is sorted now, break the loop.*/  
    if(not swapped) then  
      break  
    end if  
  end for  
end procedure return list
```

Merge Sort

Algorithm

Step 1 – if it is only one element in the list it is already sorted, return.

Step 2 – divide the list recursively into two halves until it can no more be divided.

Step 3 – merge the smaller lists into new list in sorted order.

Pseudocode

```
procedure mergesort( var a as array )
```

```
  if ( n == 1 ) return a
```

```
  var l1 as array = a[0] ... a[n/2]
```

```
  var l2 as array = a[n/2+1] ... a[n]
```

```
  l1 = mergesort( l1 )
```

```
  l2 = mergesort( l2 )
```

```
  return merge( l1, l2 )
```

```
end procedure
```

```
procedure merge( var a as array, var b as array )
```

```
  var c as array
```

```
  while ( a and b have elements )
```

```
    if ( a[0] > b[0] )
```

```
      add b[0] to the end of c
```

```
      remove b[0] from b
```

```
    else
```

```
      add a[0] to the end of c
```

```
      remove a[0] from a
```

```
end if
end while
while ( a has elements )
add a[0] to the end of c
remove a[0] from a
end while
while ( b has elements )
add b[0] to the end of c
remove b[0] from b
end while
return c
end procedure
```

Quick Sort

Algorithm

Quick Sort Pivot selection Algorithm ()

- Step 1 – Choose the highest index value has pivot
- Step 2 – Take two variables to point left and right of the list excluding pivot
- Step 3 – left points to the low index
- Step 4 – right points to the high
- Step 5 – while value at left is less than pivot move right
- Step 6 – while value at right is greater than pivot move left
- Step 7 – if both step 5 and step 6 does not match swap left and right
- Step 8 – if $\text{left} \geq \text{right}$, the point where they met is new pivot

Quick Sort Algorithm ()

Step 1 – Make the right-most index value pivot

Step 2 – partition the array using pivot value

Step 3 – quicksort left partition recursively

Step 4 – quicksort right partition recursively

Pseudocode ()

Quick Sort Pivot selection ()

function partitionFunc(left, right, pivot)

 leftPointer = left

 rightPointer = right - 1

 while True do

 while A[++leftPointer] < pivot do

 //do-nothing

 end while

 while rightPointer > 0 && A[--rightPointer] > pivot do

 //do-nothing

 end while

 if leftPointer >= rightPointer

 break

 else

 swap leftPointer, rightPointer

 end if

 end while


```
    swap leftPointer,right  
    return leftPointer  
end function
```

Quick Sort Pivot selection ()

```
procedure quickSort(left, right)  
    if right-left <= 0  
        return  
    else  
        pivot = A[right]  
        partition = partitionFunc(left, right, pivot)  
        quickSort(left,partition-1)  
        quickSort(partition+1,right)  
    end if  
end procedure
```


Heap Sort

Algorithm

HEAPSORT(A)

BUILD-HEAP(A)

for $i \leftarrow \text{length}[A]$ **downto** 2

do exchange $A[1] \leftrightarrow A[i]$

$\text{heap-size}[A] \leftarrow \text{heap-size}[A] - 1$

HEAPIFY(A, 1)

BUILD-HEAP(A)

$\text{heap-size}[A] \leftarrow \text{length}[A]$

for $i \leftarrow \lfloor \text{length}[A]/2 \rfloor$ **downto** 1

do **HEAPIFY**(A, i)

HEAPIFY (A, i)

$l \leftarrow \text{LEFT}(i)$

$r \leftarrow \text{RIGHT}(i)$

if $l \leq \text{heap-size}[A]$ and $A[l] > A[i]$

then $\text{largest} \leftarrow l$

else $\text{largest} \leftarrow i$

if $r \leq \text{heap-size}[A]$ and $A[r] > A[\text{largest}]$

then $\text{largest} \leftarrow r$

if $\text{largest} \neq i$

then exchange $A[i] \leftrightarrow A[largest]$

HEAPIFY($A, largest$)

Pseudo code

Heapsort(A) {

 BuildHeap(A)

 for $i \leftarrow \text{length}(A)$ downto 2 {

 exchange $A[1] \leftrightarrow A[i]$

 heapsize \leftarrow heapsize - 1

 Heapify($A, 1$)

 }

BuildHeap(A) {

 heapsize \leftarrow length(A)

 for $i \leftarrow \text{floor}(\text{length}/2)$ downto 1

 Heapify(A, i)

 }

Heapify(A, i) {

$le \leftarrow \text{left}(i)$

$ri \leftarrow \text{right}(i)$

 if ($le \leq \text{heapsize}$ and ($A[le] > A[i]$))

 largest $\leftarrow le$

 else

 largest $\leftarrow i$

```
if (ri<=heapsize) and (A[ri]>A[largest])
```

```
    largest <- ri
```

```
if (largest != i) {
```

```
    exchange A[i] <-> A[largest]
```

```
    Heapify(A, largest)
```

```
}
```

```
}
```

Linear Search

Algorithm

Linear Search (Array A, Value x)

Step 1: Set i to 1

Step 2: if $i > n$ then go to step 7

Step 3: if $A[i] = x$ then go to step 6

Step 4: Set i to $i + 1$

Step 5: Go to Step 2

Step 6: Print Element x Found at index i and go to step 8

Step 7: Print element not found

Step 8: Exit

Pseudo code

procedure linear_search (list, value)

for each item in the list

if match item == value

return the item's location

end if

end for

end procedure

Binary Search

Algorithm

- Step 1 – Start searching data from middle of the list.
- Step 2 – If it is a match, return the index of the item, and exit.
- Step 3 – If it is not a match, probe position.
- Step 4 – Divide the list using probing formula and find the new middle.
- Step 5 – If data is greater than middle, search in higher sub-list.
- Step 6 – If data is smaller than middle, search in lower sub-list.
- Step 7 – Repeat until match.

Pseudo code

Procedure binary_search

A \leftarrow sorted array

n \leftarrow size of array

x \leftarrow value to be searched

Set lowerBound = 1

Set upperBound = n

while x not found

if upperBound < lowerBound

EXIT: x does not exist.

set midPoint = lowerBound + (upperBound - lowerBound) / 2

if A[midPoint] < x

set lowerBound = midPoint + 1

if A[midPoint] > x

set upperBound = midPoint - 1

If $A[\text{midPoint}] = x$

EXIT: x found at location midPoint

end while

end procedure