

⑦ Maximize  $Z = 5x_1 + 3x_2$   
 subject to  $3x_1 + 5x_2 \leq 15$   
 $5x_1 + 2x_2 \leq 10$   
 $x_1, x_2 \geq 0$

Introducing slack variables, we rewrite the problem in standard form:

$$\begin{aligned} \text{Max } Z &= 5x_1 + 3x_2 + 0.x_3 + 0.x_4 \\ \text{s.t. } &3x_1 + 5x_2 + x_3 + 0.x_4 = 15 \\ &5x_1 + 2x_2 + 0.x_3 + x_4 = 10 \\ &x_j \geq 0 \quad j=1,2,3,4 \end{aligned}$$

Here,  $(a_3, a_4)$  form the initial basis.

Now,  
 $x_B$  = Basis variable  
 $B$  = Basis vector  
 $c_B$  = Coeff. of Basis variable  
 in objective f.

cm	B	$x_B$	b	$c_j$	5	3	0	0	Ratio	Operation	Remarks
a	$a_3$	$x_3$	15	3	5	1	0	0	$15/3$		$a_4 \rightarrow$ entering
$\leftarrow 0$	$a_4$	$x_4$	10	$\boxed{5}$	2	0	1	$\downarrow$	$10/5$	$a_4 \rightarrow$ departing	
	$z_j - c_j$			-5	-3	0	0			$\boxed{5} \rightarrow$ key element	
$\leftarrow 0$	$a_3$	$x_3$	0	0	$\frac{19}{5}$	1	- $\frac{3}{5}$	$\frac{9.5}{19}$	$R_2' = \frac{1}{5}R_2$	$a_2 \rightarrow$ entering	
8	5	$a_4$	$x_1$	2	1	$\frac{2}{5}$	0	$\downarrow$	$R_1' = R_1 - \frac{2}{5}R_2'$	$a_3 \rightarrow$ departing	
	$z_j - c_j$			0	-1	0	1			$\frac{10}{5} \rightarrow$ key element	
now	3	$a_2$	$x_2$	$\frac{45}{19}$	0	1	$\frac{5}{19}$	$-\frac{3}{19}$	$R_4' = 5R_4$		
is	5	$a_4$	$x_1$	$\frac{90}{19}$	1	0	$-\frac{2}{19}$	$\frac{5}{19}$	$R_2' = R_2 - \frac{1}{3}R_4'$		
at		$z_j - c_j$		0	0	$\frac{5}{19}$	1				

Here all  $z_j - c_j \geq 0$  and thus optimal conditions are achieved and optimal solns are

$$x_1 = \frac{20}{19}, \quad x_2 = \frac{95}{19} \quad Z_{\max} = \frac{100 + 135}{19} = \frac{235}{19} \text{ Ans.}$$

		$C_1$	$C_2$	$C_3$	$C_4$	$C_5$	$C_6$	Rate	Operation
$C_B$	$B$	$x_1$	$b$	$a_1$	$a_2$	$a_3$	$a_4$	$a_5$	
$C_1$	$a_1$	$x_2$	$b'_1$	<del><math>a_1'</math></del>	$a_2'$	$a_3'$	$a_4'$	$a_5'$	
$C_2$	$a_2$	$x_4$	$b'_2$	$0$	$a_{21}$	$a_{23}$	$a_{22}$	$a_{24}$	
$Z_j - C_j$									

~~Now~~ Since (1) is left  $\neq 0$ , so

by row operation we have to make

$$a_{11} \rightarrow 1$$

$$a_{21} \rightarrow 0$$

$$R'_1 = \frac{1}{a_{11}} R_1$$

$$R'_2 = a_{21} - a_{21} \times R'_1$$

Continue the process until all  $Z_j - C_j \geq 0$



Optimality Condition

To calculate departing variable

Consider the minimum ratio from the column of entering variable

i.e. For  $x_3$  if entry

$$\text{Min Ratio} = b_1/a_{11} \leftrightarrow b_2/a_{21}$$

Consider the minimum value

If  $\frac{b_1}{a_{11}}$  is min.

Then  $x_3$  is departing

The element in the junction of entering and departing variable is called key element

here  $a_{11}$  will be key element

In next table  $x_3$  will be replaced by  $x_2$

$x_3 \leftarrow x_2$

Table

$C_B$	B	$x_B$	$Z_j$	$c_j$	$a_{1j}$	$a_{2j}$	$a_{3j}$	$a_{4j}$	Min Ratio	Operation	Remark
0	$a_{31}$	$x_3$	$b_1$	$c_1$	$a_{11}$	$a_{21}$	$a_{31}$	$a_{41}$	1	0	$b_1/a_{11}$
0	$a_{41}$	$x_4$	$b_2$	$c_2$	$a_{12}$	$a_{22}$	$a_{32}$	$a_{42}$	0	1	$b_2/a_{22}$
			$Z_j - c_j$		$-c_1 - c_2$						

Min Ratio is applicable only when the ratio is non-negative

$$\text{Ratio} = \frac{a_{ij}}{b_i}$$

~~Ratio~~

First calculate  $Z_j - c_j$

$$Z_j - c_j = \left( \sum a_{ij} x_i \right) - c_j$$

i.e.,  $a_{11}x_1 + a_{21}x_2 - c_1 = -c_1$   
 $a_{12}x_1 + a_{22}x_2 - c_2 = -c_2$

Consider the minimum most value

If  $-c_1$  is minimum then

i.e.,  $x_1$  " " " variable

WAT

$$\text{Max } z = c_1 x_1 + c_2 x_2$$

Subject to

$$a_{11}x_1 + a_{12}x_2 \leq b_1$$

$$a_{21}x_1 + a_{22}x_2 \leq b_2$$

$$x_1, x_2 \geq 0$$

$$\text{Max } z = c_1 x_1 + c_2 x_2 + 0.x_3 + 0.x_4$$

Subject to

$$a_{11}x_1 + a_{12}x_2 + a_{13}x_3 + 0.x_4 = b_1$$

$$a_{21}x_1 + a_{22}x_2 + 0.x_3 + x_4 = b_2$$

$$x_i \geq 0 \quad i=1, 2, 3, 4$$

 $x_B \rightarrow$  basis variables $B \rightarrow$  basis vector  $\begin{pmatrix} 1 \\ 0 \end{pmatrix}, \begin{pmatrix} 0 \\ 1 \end{pmatrix}$ or  $\begin{pmatrix} 1 \\ 0 \end{pmatrix}, \begin{pmatrix} 0 \\ 1 \end{pmatrix}, \begin{pmatrix} 0 \\ 0 \end{pmatrix}$  $c_B$ : coefficient of basis variableshere basis variables are  $x_3, x_4$ 

$$\text{So } c_B = 0, 0$$

Simplex Table

$$\begin{aligned} \text{Max } Z &= CX \\ \text{s.t. } AX &\leq (Z) b \end{aligned}$$

$$X \geq 0$$

Note:

For "Max" Prob

Case I Constraints must be  $\leq$  type

For "Min" Problem

Case II

it must be  $\geq$  type

Suppose for case I there is a constraint

$$3x_1 + 5x_2 \geq 7$$

Multiply by -ve sign

$$-3x_1 - 5x_2 \leq -7$$

It will be  $\leq$  type

Same for Case II also

Q3/ Three products are processed through three different operations. The times (in minutes) required per unit of each product, the daily capacity of the operations (in minutes per day) and the profit per sold for each product (in Rs) are as follows

operation	Time / min			operation capacity
	Product-I	Product-II	Product-III	
1	3	4	3	42
2	5	0	3	45
3	3	6	2	41
Profit	3	2	1	

The zero time indicates that the product does not require the given operation. The problem is to determine the optimum daily production for three products that maximise the profit.

Formulate the above problem as a linear programming prob.

If  $x_1, x_2, x_3$  are the number of units produced daily of products 1, 2, 3 then

$$\text{Max } Z = 3x_1 + 2x_2 + x_3$$

$$\text{s.t. } 3x_1 + 4x_2 + 2x_3 \leq 42$$

$$5x_1 + 3x_2 + 3x_3 \leq 45$$

$$3x_1 + 6x_2 + 2x_3 \leq 41$$

$$x_1, x_2, x_3 \geq 0$$

Q2

At a cattle breeding firm it is prescribed that food ration for one animal must contain at least 14, 22 and 1 units of nutrients A, B and C respectively. Two different kinds of fodder are available. Each unit weight of these two contains the following amounts of the three nutrients.

	Feeder 1	Feeder 2
Nutrient A	2	1
Nutrient B	2	3
Nutrient C	1	1

It is given that the costs of unit quantity of feeder 1 and 2 are 3 and 2 monetary units respectively. pose a linear programming prob in terms of minimizing the cost purchasing the feeders for the above cattle breeding firm.

Soln

Let  $x_1, x_2$  are unit weights of feeder 1,

$$\therefore \text{Min } Z = 3x_1 + 2x_2$$

$$\text{s.t } 2x_1 + 2x_2 \geq 14$$

$$2x_1 + 3x_2 \geq 22$$

$$x_1 + x_2 \geq 1$$

$$x_1, x_2 \geq 0$$

## Formulation of LPP

→ A factory is engaged in manufacturing two products A and B which involve lathe work, grinding and assembling. The cutting, grinding and assembling times required for one unit of A are 2, 1 and 1 hours respectively and for one unit of B are 3, 1 and 3 hours respectively. The profits on each unit of A and B are Rs 200 and Rs 300 respectively.

Assuming that 300 hours of lathe time, 300 hours of grinding time and 240 hours of assembling time, are available, pose a linear programming problem in terms of maximizing the profit.

Let  $x_1$  = no. of units of A  
 $x_2$  = " " " " " B

∴ Maximize the Profit

$$\therefore \text{Max } Z = 2x_1 + 3x_2$$

S.T.

$$2x_1 + 3x_2 \leq 300$$

$$x_1 + x_2 \leq 300$$

$$x_1 + 3x_2 \leq 240$$

$$x_1, x_2 \geq 0$$

	Cutting	Grind	Assem
A	2	1	1
B	3	1	3
hours	300	300	240

## Formulation of Linear Programming Problem

The structure is

maximize | minimize → Objective  
function

subjected to a set of constraints placed either in the form of equality / inequality

The variables are known as decision variables

$$x_i \geq 0$$

Example if there are two decision variables say  $x_1, x_2$

$$\text{Then } \max z = 5x_1 + 3x_2$$

$$\text{subject to } 5x_1 - x_2 \geq 10$$

$$x_1 + 3x_2 \geq 20$$

$$x_1, x_2 \geq 0$$

The matrix representation is given by

$$\max z = cx$$

$$\text{where } Ax \geq b$$

$$x_1, x_2 \geq 0$$

$$\text{where } c = \begin{bmatrix} 5 & 3 \end{bmatrix}$$

$$x = \begin{bmatrix} x_1 \\ x_2 \end{bmatrix}$$

$$A = \begin{bmatrix} 5 & -1 \\ 1 & 3 \end{bmatrix}, b = \begin{bmatrix} 10 \\ 20 \end{bmatrix}$$

Point A is the intersection of lines  $2x+y=2$  and  $x+3y=2$

$$08, 2x+y-2=x+3y-2$$

$$08, x=2y$$

$$2(2y)+y=2$$

$$08, y=0.4$$

$$\therefore x=2-(3 \times 0.4)$$

$$= 0.8$$

The feasible region is ABCA. The co-ordinates of extreme points of the feasible region are,

$$A=(0.8, 0.4), B=(1, 0), C=(2, 0)$$

Objective function -  $\max Z = 3x-y$

$$Z|_{(0.8, 0.4)} = 3 \times 0.8 - 0.4 = 2.0$$

$$Z|_{(1, 0)} = 3.0$$

$$Z|_{(2, 0)} = 6.0$$

Therefore, the maximum value of  $Z=6$  which occurs at point  $C(2, 0)$ , the solution to the given problem is,  $x=2, y=0$  and  $\max Z=6$ .

2014

(2) Solve graphically the following LPP -

$$\text{Minimize } Z = 3x_1 + 2x_2$$

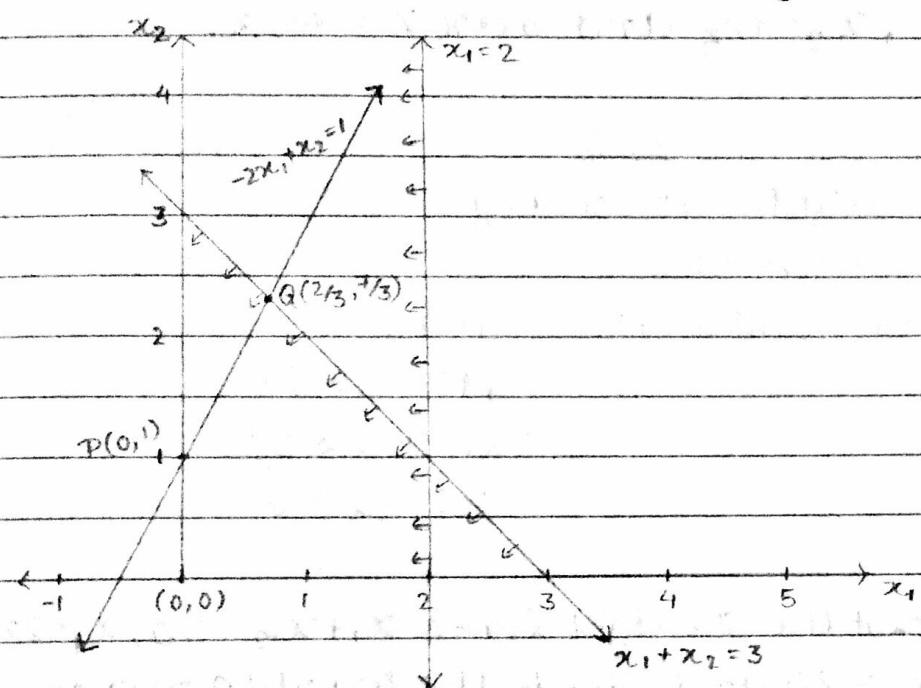
$$\text{Subject to Constraints } -2x_1 + x_2 \leq 1,$$

$$x_1 + x_2 \leq 3,$$

$$x_1 \leq 2,$$

$$x_1, x_2 \geq 0$$

Firstly, we plot the straight lines  $-2x_1 + x_2 = 1$ ,  $x_1 + x_2 = 3$  and  $x_1 = 2$  and shade the feasible region as shown in figure -



Point Q is the intersection of lines  $-2x_1 + x_2 = 1$  - (1)

and  $x_1 + x_2 = 3$  - (2)

From (2)  $x_2 = 3 - x_1$  and putting in equation (1) we get,

$$-2x_1 + 3 - x_1 = 1$$

$$-3x_1 = -2$$

$$x_1 = 2/3$$

$$\therefore x_2 = 3 - x_1$$

$$= 3 - 2/3$$

$$= 7/3$$

$$\text{So, } Q(2/3, 7/3)$$

Point B - Intersection of  $3x_1 + 8x_2 = 24$   
and  $x_1 + x_2 = 4$

$$\therefore 3x_1 + 8(4 - x_1) = 24$$

$$3x_1 + 32 - 8x_1 = 24$$

$$-5x_1 = -8$$

$$x_1 = \frac{8}{5} = 1.6$$

Therefore,  $x_2 = 2.4$

$$B(2.4, 1.6)$$

Point C - Intersection of  $10x_1 + 7x_2 = 35$

$$\text{and } x_1 + x_2 = 4$$

$$\text{or, } 10x_1 + 7(4 - x_1) = 35$$

$$3x_1 = 7$$

$$x_1 = \frac{7}{3}$$

$$\therefore x_2 = 4 - \frac{7}{3} = \frac{5}{3}$$

$$C\left(\frac{5}{3}, \frac{7}{3}\right)$$

The feasible region is OABCDO

The coordinates of extreme points of the feasible

region are O(0,0), A(0,3), B( $\frac{8}{5}, \frac{12}{5}$ ), C( $\frac{5}{3}, \frac{7}{3}$ ),

D( $\frac{7}{2}, 0$ ).

Objective function  $\rightarrow Z = 5x_1 + 7x_2$

$$Z|_{(0,0)} = 0, Z|_{(0,3)} = 21, Z|_{\left(\frac{8}{5}, \frac{12}{5}\right)} = 5 \cdot \frac{8}{5} + 7 \cdot \frac{12}{5} = \frac{124}{5},$$

$$Z|_{\left(\frac{5}{3}, \frac{7}{3}\right)} = 5 \cdot \frac{5}{3} + 7 \cdot \frac{7}{3} = \frac{74}{3}, Z|_{\left(\frac{7}{2}, 0\right)} = \frac{35}{2}$$

Since, max value of  $Z$  occurs at B( $\frac{8}{5}, \frac{12}{5}$ ) which  
is  $Z = \frac{124}{5}$ , the solution to the given problem is  
 $x_1 = \frac{8}{5}, x_2 = \frac{12}{5}$  and  $\max Z = \frac{124}{5}$ .

Since min values of  $Z$  occurs at  $y(15, 1.25)$  which is 205, the solution to the given problem is  $x_1 = 15$ ,  $x_2 = 1.25$  and  $\min Z = 205$ .

2013

(2) Solve the LPP by graphical method -

$$\text{Maximize } Z = 5x_1 + 7x_2$$

$$\text{Subject to } x_1 + x_2 \leq 4$$

$$3x_1 + 8x_2 \leq 24$$

$$10x_1 + 7x_2 \leq 35$$

$$x_1, x_2 \geq 0$$

We plot the straight lines and shade the feasible region as shown in the figure.

$$x_1 + x_2 = 4 \rightarrow (0, 4) \text{ and } (4, 0)$$

$$3x_1 + 8x_2 = 24 \rightarrow (0, 3) \text{ and } (8, 0)$$

$$10x_1 + 7x_2 = 35 \rightarrow (0, 5) \text{ and } (7/2, 0)$$

