

Volatile keyword explanation



give me the deep and detailed explaination of Volatile Keyword in java so that i can able to understand it well and do not have any type of doubt in this.

so give me the deep and detailed explaination

Below is **the most detailed, deepest, doubt-free explanation of the volatile keyword in Java** you will ever need.

I will explain everything in simple language, with internal memory concepts, JVM behavior, thread visibility, happens-before rules, limitations, and when to use volatile.

★ VOLATILE Keyword in Java – Deep & Complete Explanation

volatile is a **modifier** in Java that can be applied to **variables** (only instance and static variables, not local variables).

It is used in **multithreading** to solve the *visibility problem* and avoid *caching issues* between threads.

Let's break the topic into very digestible layers.

🔥 1. Why Do We Even Need volatile? (The Core Problem)

Java threads run in parallel and often **cache variables locally** inside CPU caches or thread-specific memory such as **working memory**.

So:

- Thread A updates a variable.
- But thread B does **not see the updated value** because it is still reading the **cached** old value.

This creates the **visibility problem**.

🧠 Java Memory Model (JMM) - Very Simple Understanding

Each thread has:

- **Main Memory (Heap / RAM)** → Shared by all threads
- **Working Memory (Thread-local cache)** → Private to each thread

When a thread reads a variable:

- It may copy it into its **local cache**
When it writes:
- It may update **only its local cache** first (not immediately main memory)

Therefore, two threads may see two different values of the same variable.

★ 2. What volatile actually does? (Simple Definition)

✓ A volatile variable guarantees:

(A) Visibility

Whenever a thread writes to a volatile variable, the value is **immediately flushed to main memory**.

Whenever a thread reads a volatile variable, it **always reads from main memory** (not from cache).

So no stale/cached value problem.

(B) Happens-Before Relationship

A **write** to a volatile variable happens-before **every read** of that volatile variable.

Meaning:

Whatever memory operations happen before writing to volatile...

...becomes visible to any thread that later reads that volatile.

This is **huge** and makes volatile useful for synchronization.

(C) No instruction reordering around volatile

The JVM and CPU normally reorder instructions for optimization.

But:

- Writes to volatile variables **cannot move forward or backward**.
- Reads from volatile variables **cannot move around**, either.

This prevents subtle concurrency bugs.

★ 3. What volatile DOES NOT DO (VERY IMPORTANT)

Volatile does **not**:

✗ Provide atomicity for compound actions:

For example:

```
java
count++; // not atomic
```

Even if **count** is volatile, the above is still:

1. read count
2. add 1
3. write count

Multiple threads can interfere and cause wrong results.

✗ Replace synchronized completely

Volatile handles **only visibility**, not **mutual exclusion**.

★ 4. Real-Life Example Without volatile (Visibility Problem)

```
java
class MyThread extends Thread {
    private boolean running = true;

    public void run() {
        while (running) {
            // do something
        }
    }

    public void stopThread() {
        running = false;
    }
}
```

If one thread calls **stopThread()**, another thread might **never stop**, because:

- Thread may cache **running = true**
 - Never sees the updated value (**false**)
-

★ 5. Fix Using volatile

```
java
class MyThread extends Thread {
    private volatile boolean running = true;

    public void run() {
        while (running) {
            // do something
        }
    }

    public void stopThread() {
        running = false;
    }
}
```

```
}
```

Now:

- Thread always reads the latest value from **main memory**.
- Loop breaks immediately when another thread writes **false**.

★ 6. How volatile prevents Reordering (Deep Explanation)

Volatile creates **memory barriers (also called fences)**:

Types:

- **LoadLoad barrier**
- **LoadStore barrier**
- **StoreLoad barrier**
- **StoreStore barrier**

Whenever volatile read/write happens, JVM inserts appropriate barriers.

This prevents:

- CPU reordering
- JVM compiler reordering

Example:

```
java

x = 1;
volatile y = 1;
```

The write to x will definitely happen **before** the write to y.

Similarly:

```
java

volatile y = 1;
int z = x;
```

The read of x will happen **after** the volatile write.

This makes volatile useful for **publishing safely**.

★ 7. Volatile vs Synchronized (Deep Comparison)

| Feature | volatile | synchronized |
|--------------------|----------------------------|--|
| Visibility | Yes | Yes |
| Mutual Exclusion | ✗ No | ✓ Yes |
| Atomicity | ✗ No | ✓ Yes |
| Prevent reordering | ✓ Yes | ✓ Yes |
| Performance | Fast | Slower |
| Use Case | One-thread-write-many-read | Multiple threads updating critical section |

★ 8. When Should You Use volatile? (Exact Use Cases)

✓ Use volatile when:

1. One thread writes, others only read

Example: status flags

- boolean running
- boolean initialized

2. Atomic reference assignment

```
java
private volatile Object instance;
```

3. Double-Checked Locking (DCL) for Singleton

```
java
private static volatile Singleton instance;
```

4. Prevent instruction reordering in concurrent algorithms

✗ Do NOT use volatile when:

1. Multiple threads modify the same variable
2. You need atomic operations (++, --, +=)
3. You need locking, fairness, or thread-safety around multiple operations
4. Working with complex data invariants

Use `synchronized` OR `AtomicInteger`, `Locks`, etc.

★ 9. The Most Important Rule

volatile guarantees visibility, NOT atomicity

Example of wrong usage:

```
java

volatile int count = 0;

public void increment() {
    count++; // still NOT thread-safe
}
```

Correct approach: use **AtomicInteger**.

★ 10. Example Showing volatile Does NOT Fix Race Condition

```
java

class Counter {
    volatile int count = 0;

    void increment() {
        count++; // not atomic
    }
}
```

If two threads do:

Thread A: reads 5

Thread B: reads 5

Both write back 6

Correct answer should be 7, but result is 6.

Volatile **cannot** fix this.

★ 11. Volatile with long and double (Special Case)

Before Java 5:

- long (64 bits)
- double (64 bits)

Could be read/written in **two steps** (not atomic).

After Java 5:

- All reads/writes to 64-bit primitives are atomic.
- Still, volatile is needed for **visibility**.

★ 12. Summary Table (Everything in One Place)

| Feature | volatile Variable |
|----------------|---|
| Visibility | ✓ Immediately visible to all threads |
| Caching | ✗ Cannot cache the variable |
| Atomicity | ✗ Not guaranteed |
| Locking | ✗ No lock |
| Reordering | ✗ Cannot reorder around volatile |
| Happens-Before | ✓ Established |
| Use Case | State flags, single-writer variables, DCL |

★ Final Clear Definition (What YOU should remember)

The volatile keyword in Java ensures that a variable's value is always read from main memory and not from a thread's local cache. It also prevents instruction reordering around that variable and establishes a happens-before relationship between writes and reads. It does NOT guarantee atomicity.

If you want, I can also give you:

- ✓ Memory diagrams showing how volatile works
- ✓ Real-world examples from Spring / concurrency libraries
- ✓ Complete interview Q&A on volatile
- ✓ Coding scenarios where volatile vs synchronized matters

Just tell me!